

COMP9024 Assignment 4

This is the design manual of COMP9024 Assignment 4.

In this assignment, a digraph is used to represent a bus network and several functions are implemented to explore its feature.

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Data Structures

Graph Structures

- [BusStops](#)

This structure stores bus stop names and bus stop serial numbers, according to assignment spec bus stop names are no longer than 20 characters.

```
typedef struct bus_stop{
    char name[20];          // bus name is no more than 20 characters
    int id;                 // serial number of a bus stop
}BusStop;
```

Go to [appendix](#) to see its relationship with other structures.

- [Edge](#)

An edge links two vertices with a given distance. In this structure we only need to specify the destination vertex and distance, because [Edge](#) is designed to be the subclass of [vertex](#) so the source vertex is just the vertex itself.

Unfortunately circular dependency exists between [Edge](#) and [Vertex](#) (they mutually define each other !), to resolve this issue it is important to declare : `struct vertex;` before we define [Edge](#)

```
struct vertex;           // declare vertex beforehand
typedef struct edge{
    struct vertex *dst;   // destination vertex
```

```
int distance;          // the length of the edge
}Edge;
```

Go to [appendix](#) to see its relationship with other structures

- **Vertex**

A vertex must hold information of a bus stop as well as its corresponding edges. In **Vertex** we used **List** to store edges (named **edgesQueue**) as well as reversed edges (named **reversed_edgesQueue**). Reversed edges tell which vertices can reach this vertex and how far they are from this vertex, **reversed_edgesQueue** is useful in **StronglyConnectivity** where we need to apply dfs reversely. Variables **distanceFromStartVertex**, **previousVertexInShortesPath**, **positionInHeap** are useful in Dijkstra's Algorithm. **isMark** is a book-keeping variable to ensure that in BFS or DFS no vertex is accessed more than once.

```
typedef struct vertex{
    BusStop *thisBusStop;
    struct list *edgesQueue;
    struct list *reverse_edgesQueue;
    struct vertex *previousVertexInShortestPath;
    int distanceFromStartVertex;
    int positionInHeap;          // the position of the vertex in binary
    heap
    int isMarked;
}Vertex;
```

Go to [appendix](#) to see its relationship with other structures

- **Graph**

A graph is just a collection of vertices and edges. Since edges are already stored in their corresponding vertice, there's need to explicitly declare them in **Graph** again. Maximum number of vertices, current number of vertices and edges are saved in **Graph** for book-keeping.

```
typedef struct graphRep{
    Vertex **vertices; // store an array of vertex pointer
    int size;          // max number of vertices that can be store in
    graph
    int nV;            // number of vertices currently stored in graph
    int nE;            // number of edges currently stored in graph
}GraphRep;
typedef GraphRep *Graph;
```

Go to [appendix](#) to see its relationship with other structures

All structures discussed in this section are defined in **MyGrap.h**

List Structures

- **ListNode**

List is mainly used in storing edges of graph. Hence a **ListNode** contains a pointer that points to an edge.

```
typedef struct list_node{
    struct edge *thisEdge;
    struct list_node *next;
}ListNode;
```

Go to [appendix](#) to see its relationship with other structures

- **List**

In addition to the first node of the list, last node of the list is also cached in **List**, which makes it possible to append a node to **List** in $O(1)$ time, regardless of how long the list is.

```
typedef struct list{
    ListNode *firstNode;
    ListNode *lastNode;
    int size;
}List;
```

Go to [appendix](#) to see its relationship with other structures

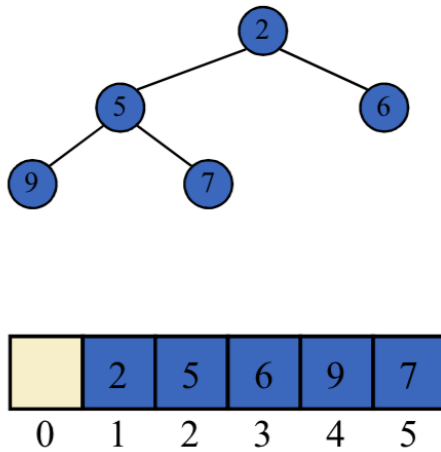
All structures discussed in this section are defined in **MyList.h**.

Heap Structure

To implement Dijkstra's algorithm a heap structure is indispensable. The basic idea is to store every vertex in a heap, ordered by its distance from the start point. For each run we take the smallest vertex out of the binary heap and compute relaxation until the heap is empty.

```
typedef struct binary_heap{
    struct vertex **node; // the vertex pointers of graph are all stored
in an array
    int n;                // current number of nodes
    int size;              // total number of nodes
}BinaryHeap;
```

Inside the binary heap, vertices are stored in an array. Below example depicts how an array represents a binary heap:



Notice that array at position 0 is unused.

Go to [appendix](#) to see its relationship with other structures

Algorithms

- Depth-first Search (DFS)
 - used in **StronglyConnectivity** to search all bus stops that can be reached
- Breadth-first Search (BFS)
 - used in **reachableStops**
- Kosaraju's Algorithm
 - used in **maximalStronlyComponents**
- Dijkstra's Algorithm
 - used in **TravelRoute** to find shortest path

Time Complexity Analysis

Function name	Time complexity
StronglyConnectivity	$O(M+N)$
maximalStronlyComponents	$O(M+N)$
reachableStops	$O(M+N)$
TravelRoute	$O((M+N)*\log N)$

Appendix

Structures Dependencies

Diagram below depicts the relationship between different structures, if structure A encloses structure B then structure B is the subclass of structure A (aka, there is a B declared in A).

Graph

Vertex

Bus stop

List

ListNode

Edge

Vertex

BinaryHeap

Vertex