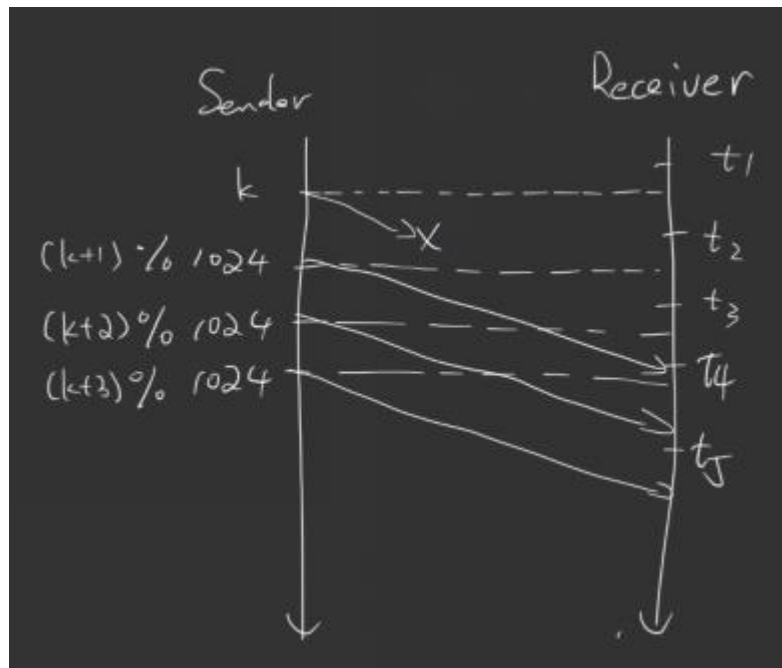


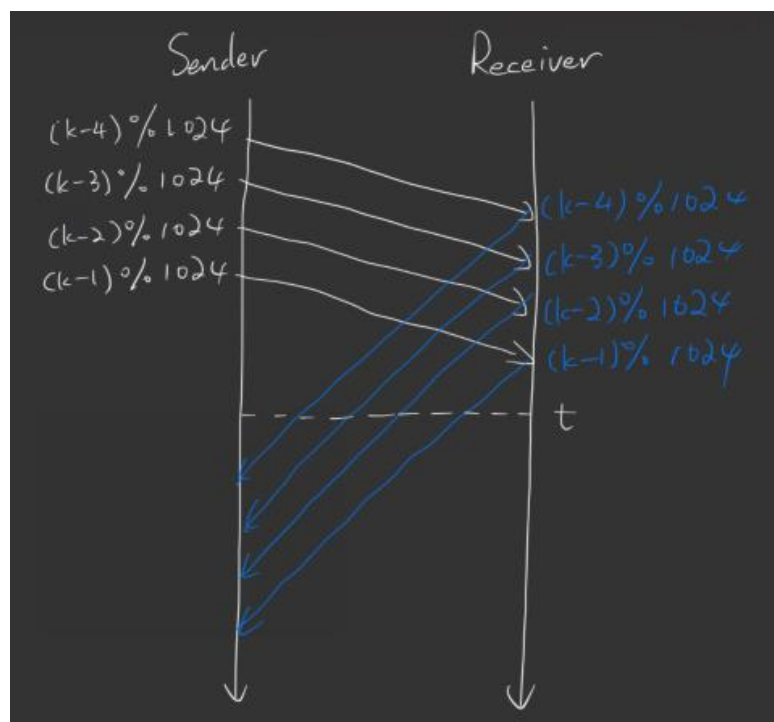
Q1:

a.



As we can see in the picture. At time  $t_1$ ,  $t_2$ ,  $t_3$ ,  $t_4$ ,  $t_5$ , the next in-order packet that the receiver is expecting all has a sequence number of  $k$ . However, at time  $t_1$ , the send base of the sender is  $k$ , the next sequence number is  $k$ . At time  $t_2$ , the send base of the sender is  $k$ , the next sequence number is  $(k + 1) \% 1024$ . At time  $t_3$ , the send base of the sender is  $k$ , the next sequence number is  $(k + 2) \% 1024$ . At time  $t_4$ , the send base of the sender is  $k$ , the next sequence number is  $(k + 3) \% 1024$ . At time  $t_5$ , the send base of the sender is  $k$ , the next sequence number is  $(k + 4) \% 1024$ .

b.



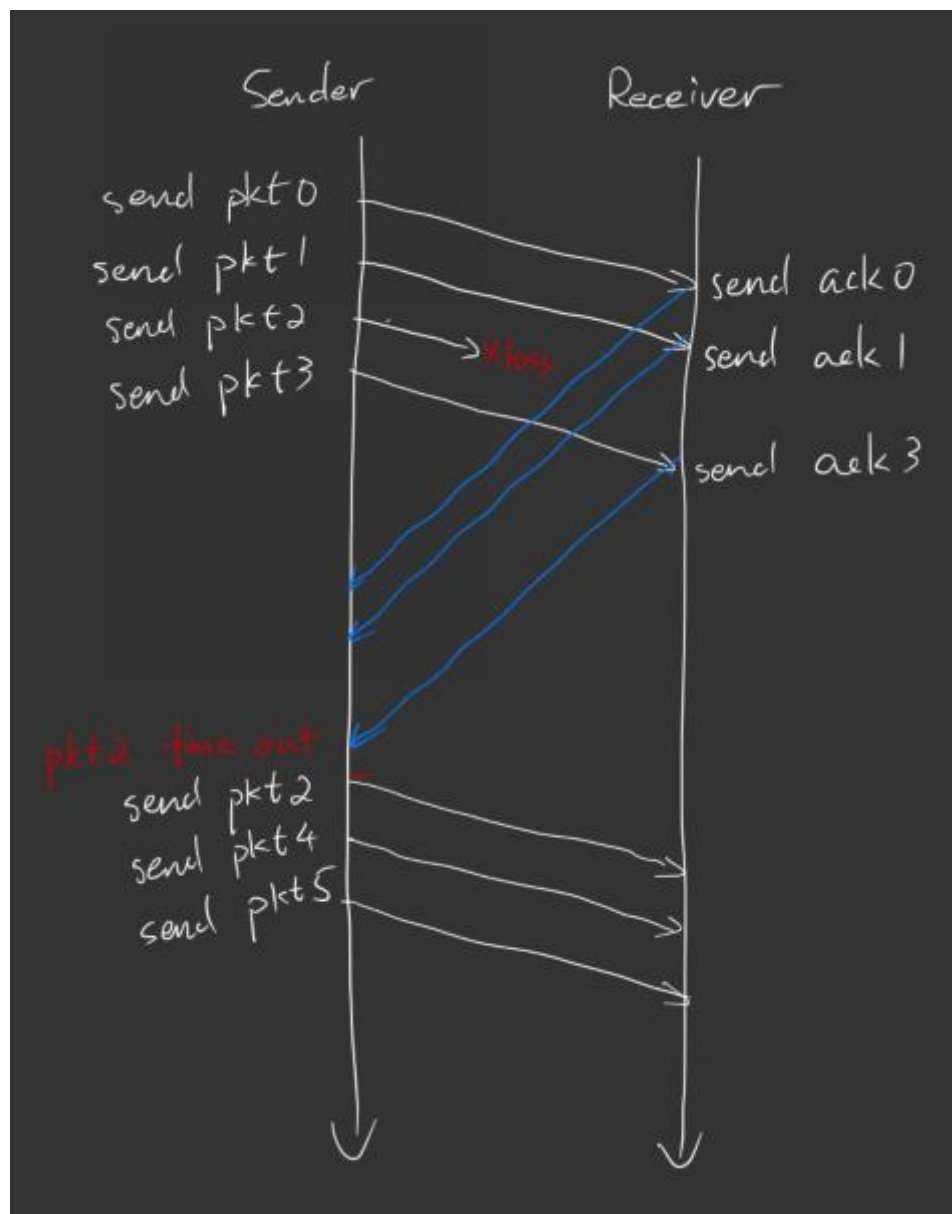
As we can see in the picture. At time  $t$ , the next in-order packet that the receiver is expecting all has a sequence number of  $k$ . In this situation, there are four ACK messages are currently being propagating back to the sender which are ACK message for  $(k - 4) \% 1024$ ,  $(k - 3) \% 1024$ ,  $(k - 2) \% 1024$ ,  $(k - 1) \% 1024$ . These are all possible values of the ACK field in all possible messages currently propagating back to the sender at time  $t$ .

## Q2:

The selective repeat protocol is used because the ack3 is sent when receiver receives pkt3. If it is Go-Back-N protocol, the receiver will send ack1 instead of ack3.

a. The position of sender window is 0 to 3. The position of receiver window is 2 to 5.

b.

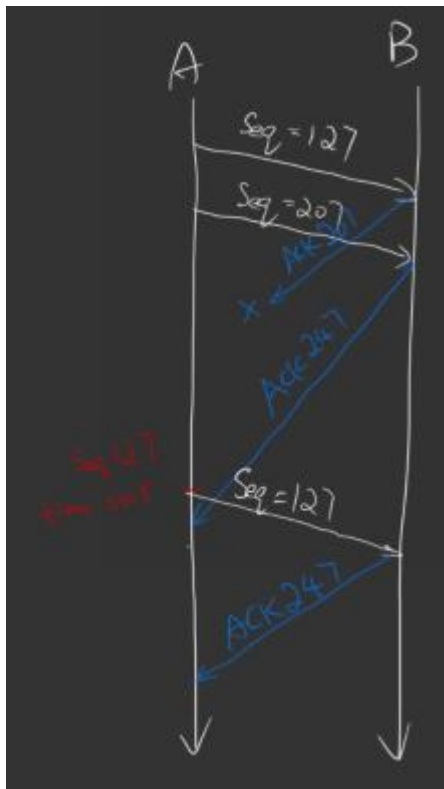


**Q3:**

- a.  $110 - 90 = 20$  bytes
- b. It will be acknowledgment 110.

**Q4:**

- a. The sequence number is  $127 + 80 = 207$   
The source port is 302 and the destination port is 80.
- b. The acknowledgment number is  $127 + 80 = 207$   
The source port is 80 and the destination port is 302.
- c. The acknowledgment number is 127, the source port is 80 and the destination port is 302.
- d.



**Q5:**

- a. from 1st to 6th and 23rd to 26th transmission round
- b. from 6th to 23rd transmission round
- c. A triple duplicate ACK.
- d. A timeout.
- e. 32
- f. 21
- g. 14
- h. 7th transmission round
- i. congestion window size =  $8 / 2 + 3 = 7$   
ssthresh =  $8 / 2 = 4$

j.  $ssthresh = 42 / 2 = 21$

congestion window =  $1 * 2 * 2 = 4$  segments

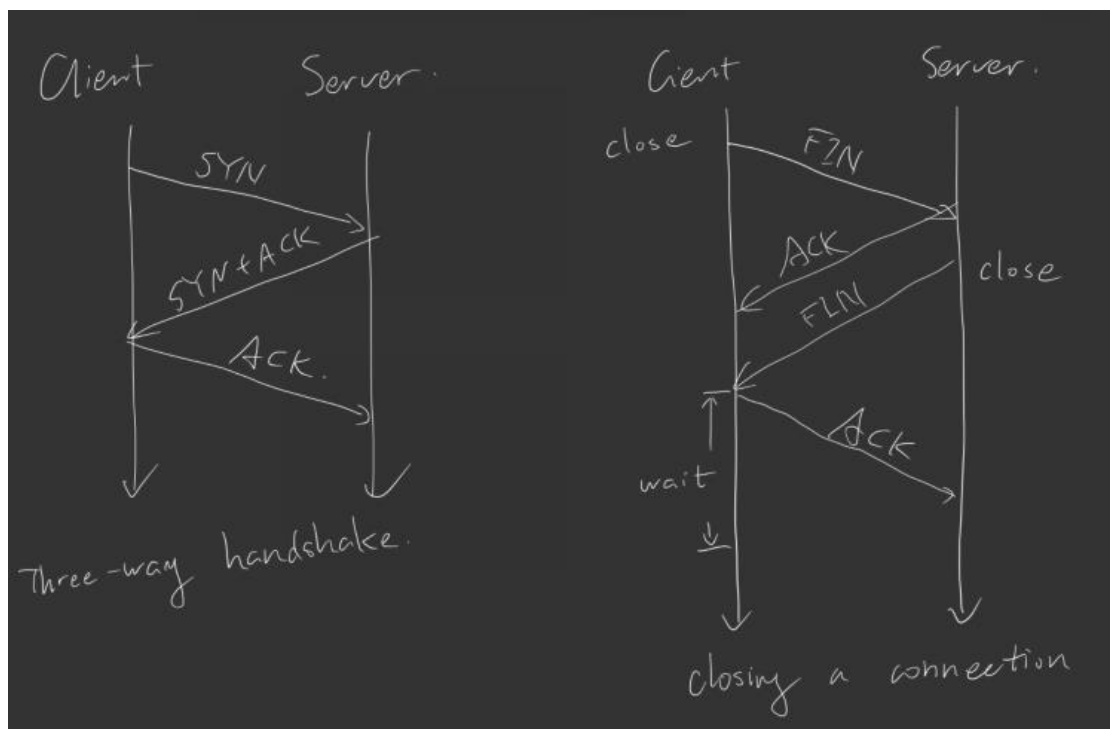
k.  $1 + 2 + 4 + 8 + 16 + 32 = 63$  packets

**Q6:**

a.  $2^{(4 * 8)} * 536 = 2Tb$

b.  $2^{(4 * 8)} * (536 + 66) * 8 / (155 * 1024 * 1024) = 127266 \text{ s} = 2121 \text{ min} = 35 \text{ h}$

**Q7:**



**Q8:**

Sequence number = 65530

Acknowledgement number =  $81920 + 10240 = 92160$

**Q9:**

Since  $rwnd$  is bigger than  $cwnd$ , the number of bytes that can be sent is determined by  $cwnd$ .

So  $3000 - 2000 = 1000$  more bytes can be sent.

**Q10:**

Transmission#	Congestion Window	Threshold	Receive Window
1	4Kb	12Kb	32Kb
2	8Kb	12Kb	32Kb
3	16Kb	12Kb	32Kb
4	20Kb	12Kb	32Kb
5	24Kb	12Kb	32Kb
6	28Kb	12Kb	32Kb
7	32Kb	12Kb	32Kb
8	4Kb	16Kb	32Kb
9	8Kb	16Kb	32Kb
10	16Kb	16Kb	32Kb
11	20Kb	16Kb	32Kb
12	24Kb	16Kb	32Kb
13	28Kb	16Kb	32Kb
14	32Kb	16Kb	32Kb
15	36Kb	16Kb	32Kb
16	30Kb	18Kb	32Kb
17	34Kb	18Kb	32Kb
18	4Kb	17Kb	32Kb
19	8Kb	17Kb	32Kb
20	16Kb	17Kb	32Kb
21	32Kb	17Kb	32Kb
22	36Kb	17Kb	32Kb
23	4Kb	18Kb	32Kb
24	8Kb	18Kb	32Kb

The sender attempts to send 8Kb data on the 24th transmission.