

Overview

- nodes

- transaction (tx)

"Full node" store blockchain

→ □ → □ → □ → □ → □ → □ →

tx
/ |
⋮

read write
1 token burn 1 token

有沒有改變合約狀態?

- Smart Contract
(flexible txs, ERC20, ERC721)
read, write
Token NFT

- Verify S.C code need Flatten.

- bytecode to EVM.

Blockchain Award

- Bitcoin: Proof-of-Work, $H(\text{pre-hash}, \text{tx's root}, \text{nonce}) < t$

- Proof-of-Stake, (tx's root) 不用 Nonce

哪邊放多少錢, 就有多少機會賺更多錢

破冰

1. Blockchain 中 Smart Contract 代表什麼?

Ans. Token 表達的一種形式。

2. 為什麼要做 Flatten?

Ans. 減少依賴, 確保合約獨立運行。

3. 100% 在 POS 的機率分配, 顏色相同要一樣分配

A B C D E F G
1 10 11 10 15 20 3
4 14 14 14 25 25 4

使用先折半分配

4. 為什麼省電?

Ans. 加入質押, 要過效期才可取回

大戶(交易所)一起質押, 交易所一起挖

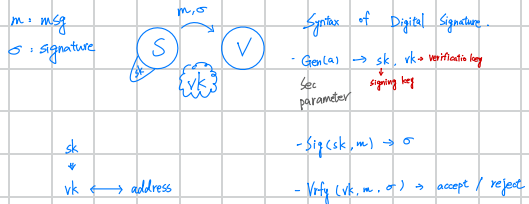
5. 隨機哪裡來(用來決定誰挖到, 不可預測)

Ans. 使用 node 產生時間毫秒值, 看數值的接近

Digital Signature

1. tx generation by ^{address} ① who hold coins.

2. verification of tx really from ①



Signature

- Authentication: Verify m from S.
- Integrity: m is not modified.
- Non-repudiation: cannot withdraw the sig.

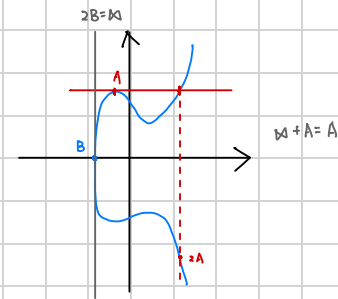
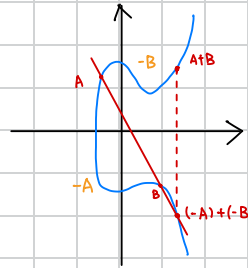
下週中

Signature 橢圓曲線密碼學

$$y^2 = x^3 + ax + b \mod P$$

$$x, y \in \mathbb{Z}_P$$

$$a, b \in \{0, \dots, P-1\}$$



$$y^2 = x^3 + ax + b \mod P$$

P, Q: a point in E (group defined by a EC)

- P is generator $E = \{P, 2P, \dots, (P-1)P\}$

- Give a point Q in E, $Q = xP$

EC hardness: Give (P, Q), hard to find $x \in \mathbb{Z}_P$ s.t. $Q = xP$
public generator

Elliptic Curve Digital Signature Algorithm (ECDSA)

- Gen(n): Let P is a prime, E is EC over $\mod P$

P is the generator of E , $sk = \alpha \in \mathbb{Z}_P$, $vk(E, P, Q = \alpha P)$

- Sign(sk, m): 1. random $r \in \mathbb{Z}_P$

$\rightarrow (\sigma_1, \sigma_2)$ 2. compute $\sigma_1 = rP$

3. compute $u = H(\sigma_1, \alpha P, m)$

$$(\sigma_1, \sigma_2) \quad \sigma_2 = r^{-1}(H(m) + \alpha \cdot u)$$

- Vrfy(vk, m, σ): $\sigma_2(\sigma_1) = r^{-1}(H(m) + \alpha \cdot u) \cdot (rP)$

$$= (H(m) + \alpha u) P$$

$$= H(m)P + u(\alpha P) \rightarrow vk$$