${\tt concordance=TRUE}$

Problem Set 5

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1 problem 2

```
1 = (-1)^0 * 1.0 * 2^0, that is, S = 0, d = 0, e = 1023
 2 = (-1)^0 * 1.0 * 2^1, that is, S = 0, d = 0, e = 1024
 3 = (-1)^0 * 1.1 * 2^1, that is, S = 0, d = 1, e = 1024
 d and e adds on for increasing integers.
 to store 2^{5}3 - 2, S = 0, dis
 times followed by 0), e = 1075
 e = 1075
 and when increasing the lieast significant bit by 1, we get
 So we can't represent 2^53 + 1.
 with the case of 2^54, e = 1023 + 54,
 so increasing the least significant bit which is the 52th bit of d, would result in an increas of 2^{(-52+54)}
 we can see in R, bits of 2^53 - 1 and 2^53 is different but 2^53 and 2^53 + 1 is the same.
```

2 problem 3

2.1 a)

from the following code we can see it is faster to copy a large vector of integers than a numeric vector of the same length.

```
#integer case
x<- c(1:1e8)
object.size(x)
## 400000040 bytes
# let y point to x so when change one element of the vector a whole new copy is made.
v <- x
# note that if we use x[4] \leftarrow 50 the time is much longer
\# since type of x changes from integer to double.
system.time(x[4] \leftarrow 50L)
##
      user system elapsed
      0.08
             0.05
                     0.13
##
# numeric case
x1<- rnorm(1e8)
v1 <- x1
object.size(x1)
## 80000040 bytes
system.time(x1[4] <- 50)
##
      user system elapsed
  0.12 0.14 0.26
```

2.2 b)

the time to take a subset of size k=n/2 from an integer vector is a little bit faster.

```
x<- c(1:1e8)
x1<- rnorm(1e8)
system.time(sub <- x[1:5e7])

## user system elapsed
## 0.29 0.08 0.37

system.time(sub1 <- x1[1:5e7])

## user system elapsed
## 0.28 0.13 0.41</pre>
```

3 problem 4

3.1 a)

Because we should consider both the time for each computation and the the amount of communication that needs to happen. If we have very many tasks and each one takes little time, the communication overhead of starting and stopping the tasks will reduce efficiency, which will be the case if we do n individual column-wise computations.

3.2 b)

the amount of memory used:

A: each of the p workers will use $n^2 for X$, n * m for Y and n * m for result, $So(n^2 + 2nm)pintotal$.

B:each of the p workers will use n*m for X, n*m for Y and m*m for result, So $(m^2 + 2nm)pintotal$.

A uses more memory than B

the communication cost:

A:p task is needed and each task need to pass n*n for X, n*m for Y and n*m for result,

So $C1 = (2mn+n^2)pintotal$.

B:p*p task is needed and each task need to pass m*n for X, n*m for Y and m*m for result,

So $C2 = (2mn+m^2)p^2 intotal$.

substract these two, and substitute p=n/m:

C1-C1 = $(2mn+n^2)*n/m - (2mn+m^2)*(n/m)^2 = n^2(1-n/m) < 0$

So A takes less Commication cost than B.

So B is better for minimizing memory use and A for minimizing communication