

Functional and Logic Programming

Bachelor in Informatics and Computing Engineering
2024/2025 - 1st Semester

Prolog

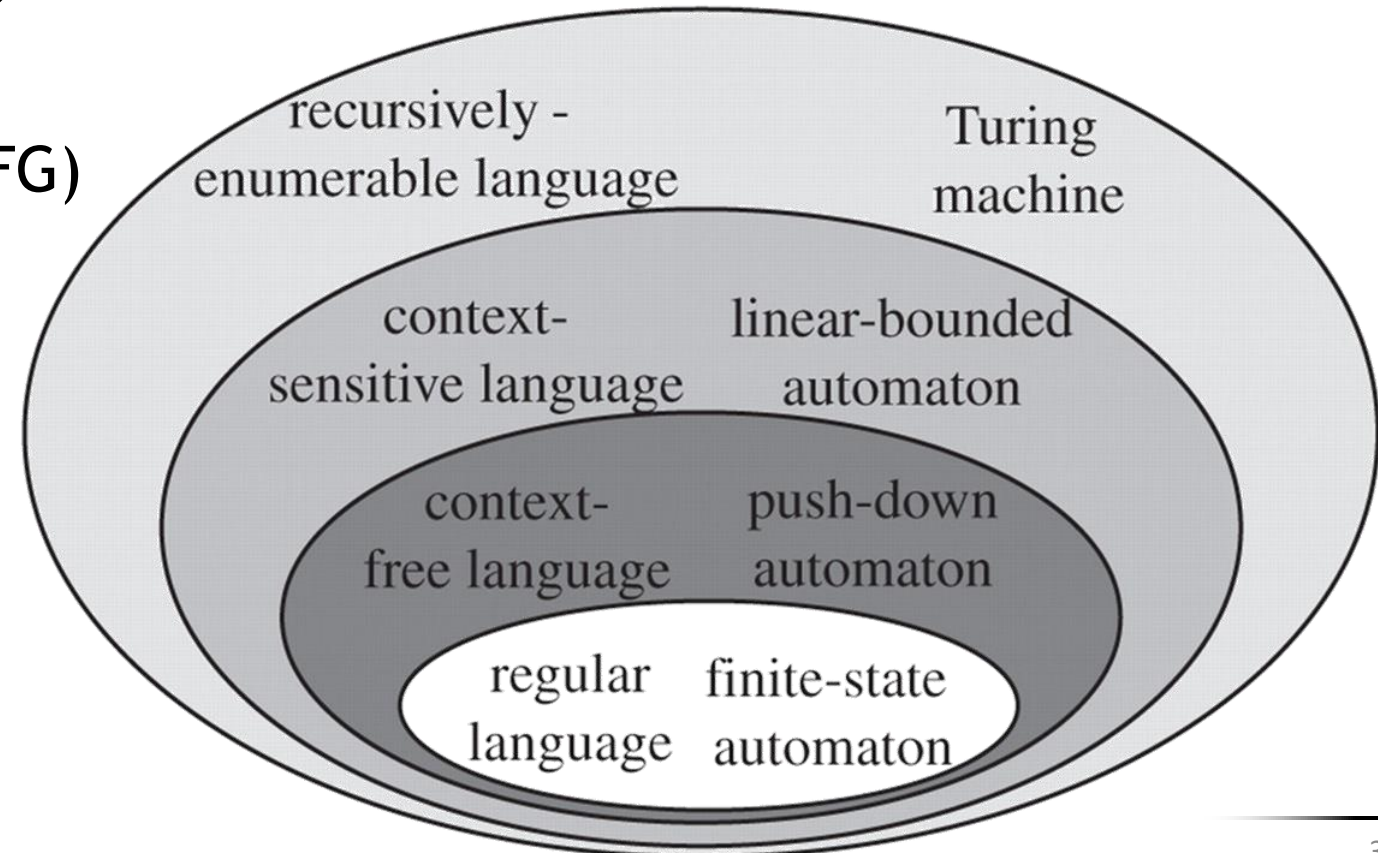
Applications and Libraries

Agenda

- Computational Models
- Incomplete Data Structures
 - Difference Lists
- Syntactic sugar
- Statistics
- SICStus Libraries

Computational Models

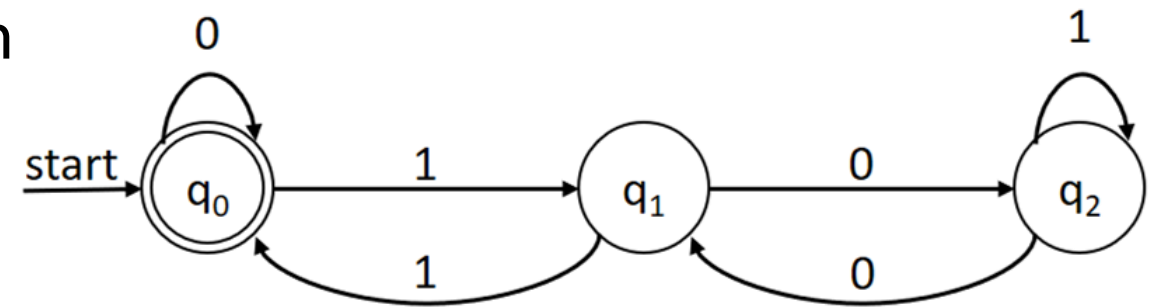
- Remembering Theory of Computation...
 - Finite Automata (DFA / NFA)
 - Pushdown Automata (PDA)
 - Context-Free Grammars (CFG)
 - Turing Machines (TM)



Computational Models

- We can easily create a Prolog program to emulate DFAs / NFAs
 - Generic solver uses graph search

DFA = $\langle Q, \Sigma, \delta, I, F \rangle$



```

accept(Str) :-
    initial(State),
    accept(Str, State).

accept([], State) :-
    final(State).
accept([S|Ss], State) :-
    delta(State, S, NState),
    accept(Ss, NState).
  
```

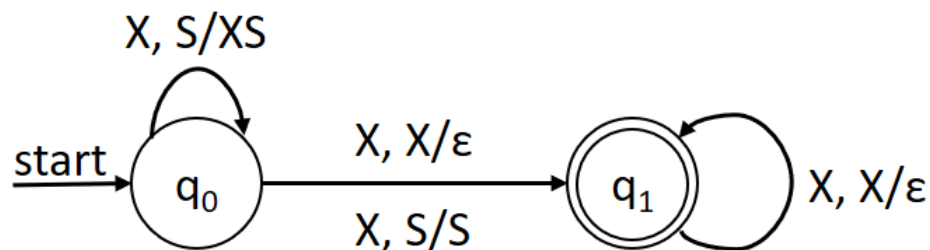
```

initial(q0).
final(q0).
final(q2).
delta(q0, 0, q0).
delta(q0, 1, q1).
delta(q1, 0, q2).
delta(q1, 1, q0).
delta(q2, 0, q1).
delta(q2, 1, q2).
  
```

Computational Models

- The same kind of logic can be used to emulate PDAs

$\text{PDA} = \langle Q, \Sigma, \Gamma, \delta, q_0, Z_0, F \rangle$



`initial(q0).`

`final(q1).`

`delta(q0, X, Stack, q0, [X|Stack]).`

`delta(q0, X, Stack, q1, Stack).`

`delta(q0, X, [X|Stack], q1, Stack).`

`delta(q1, X, [X|Stack], q1, Stack).`

`accept(Str) :- initial(State), accept(Str, State, []).`

`accept([], State, []) :- final(State).`

`accept([S|Ss], State, Stack) :-`

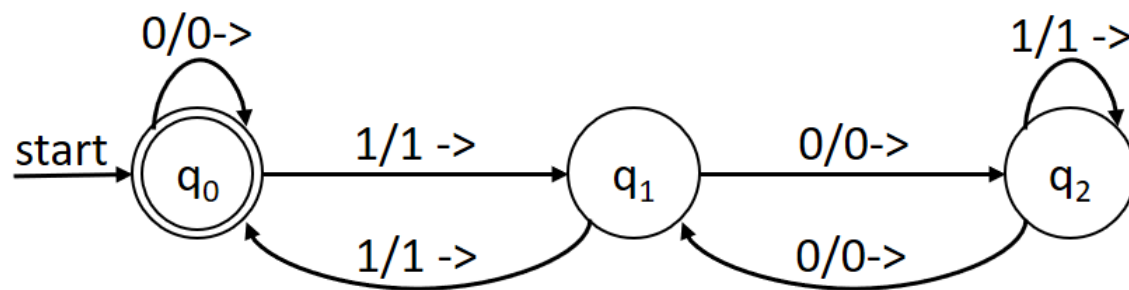
`delta(State, S, Stack, NewState, NewStack),`

`accept(Ss, NewState, NewStack).`

Computational Models

- The same kind of logic can also be used to emulate TMs

$$TM = \langle Q, \Sigma, \Gamma, \delta, q_0, B, F \rangle$$



```

tm(Str) :- initial(State),
             append(Str, [empty], StrEmpty),
             tm([empty], StrEmpty, State).

```

```

tm(Left, [S|Right], State) :-
    delta(State, Left, [S|Right], NewState, NewLeft, NewRight), !,
    tm(NewLeft, NewRight, NewState).

```

```

tm(_, _, State) :- final(State).

```

```

initial(q0).
final(q0).
delta(q0,L,[0|R],q0,[0|L],R).
delta(q0,L,[1|R],q1,[1|L],R).
delta(q1,L,[0|R],q2,[0|L],R).
delta(q1,L,[1|R],q0,[1|L],R).
delta(q2,L,[0|R],q1,[0|L],R).
delta(q2,L,[1|R],q2,[1|L],R).

```

Computational Models

- We can also easily emulate CFGs

CFG = $\langle V, T, P, S \rangle$

$S \rightarrow \varepsilon$

$S \rightarrow X$

$S \rightarrow XSX$

```
accept(Str) :- s(Str).
```

```
s([]).
```

```
s([X]).
```

```
s([X|SX]) :- append(S, [X], SX), s(S).
```

Computational Models

- The definition of CFGs can be simplified using DCGs (Definite Clause Grammars)
 - It uses a syntax similar to the specification of grammar rules
 - It can be used both to recognize and to generate strings

```
pal --> [].
pal --> [_].
pal --> [S], pal, [S].
```

```
| ?- phrase(pal, "not a pal").
no
| ?- phrase(pal, "abba").
true ? ;
no
| ?- phrase(pal, "madamimadam").
true ?
yes
| ?- phrase(pal, X).
X = [] ? ;
X = [_A] ? ;
X = [_A, _A] ? ;
X = [_A, _B, _A] ? ;
X = [_A, _B, _B, _A] ? ;
X = [_A, _B, _C, _B, _A] ? ;
X = [_A, _B, _C, _C, _B, _A] ? ;
X = [_A, _B, _C, _D, _C, _B, _A] ?
yes
```


Computational Models

- Verifications can be made as extensions to the grammar rules

```
palb --> [].
```

```
palb --> [S], { [S] = "0"; [S]="1" }.
```

```
palb --> [S], palb, [S], { [S] = "0"; [S]="1" }.
```

```
| ?- phrase(palb, "abba").  
no  
| ?- phrase(palb, "01x10").  
no  
| ?- phrase(palb, "01010").  
true ?  
yes  
| ?- phrase(palb, "00").  
true ?  
yes
```

Computational Models

- More complex rules can be used

`expr(Z) --> term(X), "+", expr(Y), {Z is X + Y}.`

`expr(X) --> term(X).`

`term(Z) --> num(X), "*", term(Y), {Z is X * Y}.`

`term(Z) --> num(Z).`

`num(X) --> [D], num(R), {"0"=<D, D=<"9", X is (D-"0")*10 + R}.`

`num(X) --> [D], {"0"=<D, D=<"9", X is D-"0"}.`

```
| ?- phrase(expr(X), "2+4").
X = 6 ?
yes
| ?- phrase(expr(X), "6+4*3").
X = 18 ?
yes
| ?- phrase(expr(X), "12*4+16").
X = 64 ?
yes
```

Agenda

- Computational Models
- Incomplete Data Structures
 - Difference Lists
- Syntactic sugar
- Statistics
- SICStus Libraries

Incomplete Data Structures

- Incomplete data structures increase efficiency by allowing ‘partial’ or ‘incomplete’ structures to be specified and incrementally constructed during runtime
 - This is achieved by maintaining a free variable as the final element of the structure, as opposed to a constant (such as [] for lists or null)
 - Changes to the incomplete structure can be made by [partially] instantiating the ending variable, thus not requiring the use of an extra output argument

Incomplete Data Structures

- Implementation of a dictionary using incomplete lists

```
lookup(Key, [ Key-Value | Dic ], Value).
lookup(Key, [ K-V | Dic ], Value):-
    Key \= K,
    lookup(Key, Dic, Value).
```

- When *Key* is present, *Value* is verified/returned
- When *Key* is not present, the new *Key-Value* pair is added to the dictionary

```
| ?- Dic = [x-1, y-2, z-3 | _R], lookup(y, Dic, V).
Dic = [x-1,y-2,z-3|_R],
V = 2 ?
yes
| ?- Dic = [x-1, y-2, z-3 | _R], lookup(w, Dic, 4).
Dic = [x-1,y-2,z-3,w-4|_A] ?
yes
```

Incomplete Data Structures

- Dictionary implemented with incomplete binary search tree

```
lookup(Key, dtnode(Key-Value, _L, _R), Value).
```

```
lookup(Key, dtnode(K-_V, L, _R), Value):-
```

```
    Key < K, lookup(Key, L, Value).
```

```
lookup(Key, dtnode(K-_V, _L, R), Value):-
```

```
    Key > K, lookup(Key, R, Value).
```

```
| ?- dtree(DTree).
```

```
DTree = dtnode(3-b,dtnode(1-(a),_A,_B),dtnode(7-d,dtnode(5-c,_C,_D),dtnode(9-(e),_E,_F))) ?
```

```
yes
```

```
| ?- dtree(DTree), lookup(5, DTree, V).
```

```
DTree = dtnode(3-b,dtnode(1-(a),_A,_B),dtnode(7-d,dtnode(5-c,_C,_D),dtnode(9-(e),_E,_F))),
```

```
V = c ?
```

```
yes
```

```
| ?- dtree(DTree), lookup(4, DTree, g).
```

```
DTree = dtnode(3-b,dtnode(1-(a),_A,_B),dtnode(7-d,dtnode(5-c,dtnode(4-g,_C,_D),_E),dtnode(9-(e),_F,_G))) ?
```

```
yes
```

Difference Lists

- While lists are widely used, some common operations may not be very efficient, as is the case of appending two lists
 - Linear on the size of the first list
- Idea: increase efficiency by ‘also keeping a pointer to the end of the list’
 - This is accomplished by using difference lists
 - We can use any symbol to separate the two parts of the difference list
 - With this representation, we can have an incomplete list (when the second list is not instantiated)

$$X = [1, 2, 3]$$

$$X = [1, 2, 3, 4, 5, 6] \setminus [4, 5, 6]$$

$$X = [1, 2, 3, a, b, c] \setminus [a, b, c]$$

$$X = [1, 2, 3] \setminus []$$

$$X = [1, 2, 3 \mid T] \setminus T$$

Difference Lists

- We can now append two (difference) lists in constant time
 - To append $X \setminus Y$ with $Z \setminus W$, simply unify Y with Z

```
append_dl (X\Y, Y\W, X\W) .
```

- Note that the two lists must be compatible - the tail of the first list must either be uninstantiated or be equal to the second list

```
| ?- append_dl( [a, b, c | Y ]\Y, [d, e, f | W]\W, A) .  
Y=[d,e,f|W]  
A=[a,b,c,d,e,f|W]\W
```

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Syntactic Sugar

- Do loops can be used instead of iteration predicates

`(Iterators do Body)`

- There are several iterators available (see section 4.2.3.5)
 - *foreach(Elem, List)*
 - *for(Iterator, MinExpression, MaxExpression)*
 - *count(Iterator, MinExpression, Max)*
 - *fromto(First, In, Out, Last)*
 - ...

Syntactic Sugar

```
| ?- foreach(X, [1,2,3,4]) do Y is X*2, write(X-Y),nl.  
1-2  
2-4  
3-6  
4-8  
yes  
| ?- Min is 3, (for(X, Min, 2*Min) do Y is X*2, write(X-Y),nl).  
3-6  
4-8  
5-10  
6-12  
Min = 3 ?  
yes  
| ?- Min is 3, (count(X, 2*Min, 9) do Y is X*2, write(X-Y),nl).  
6-12  
7-14  
8-16  
9-18  
Min = 3 ?  
yes
```

Syntactic Sugar

- More than one Iterator can be used
 - They are iterated synchronously (first element of each, second of each, ...)

```
| ?- foreach(Elem, [2,4,6,8]), foreach(Res, List) do Res is Elem*3.  
List = [6,12,18,24] ?  
yes  
| ?- M is 3, (for(X, 2*M, 3*M), count(Iter,1,Iterations) do Y is X*2, write(X-Y),nl).  
6-12  
7-14  
8-16  
9-18  
M = 3,  
Iterations = 4 ?  
yes  
| ?- M is 3, (for(X, 2*M, 3*M), count(Iter,1,Iterations), foreach(Y,List) do Y is X*2).  
M = 3,  
Iterations = 4,  
List = [12,14,16,18] ?  
yes
```

Syntactic Sugar

- Parameters can be used to access external variables from within the body of the iterator
 - This is also required for nested iterators

```
| ?- foreach(X, [1,2]) do foreach(Y, [1,2]) do write(X-Y),nl.  
_2351-1  
_2603-2  
_2921-1  
_3173-2  
yes  
| ?- foreach(X, [1,2]) do foreach(Y, [1,2]), param(X) do write(X-Y),nl.  
1-1  
1-2  
2-1  
2-2  
yes
```

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- **Statistics**
- SICStus Libraries

Statistics

- Execution statistics can be obtained using the *statistics/0* or *statistics/2* predicates
 - *statistics/0* prints statistics related to memory usage, execution time, garbage collection and others (counting from session start)
 - *statistics(?Keyword, ?Value)* obtains values (or lists of values) for several available statistics
 - See section 4.10.1.2 for a full list of available keywords and respective details

```
| ?- statistics(runtime, [Before|_]), fib(30,F), statistics(runtime, [After|_]),  
Time is After-Before.  
Before = 16849,  
F = 832040,  
After = 19207,  
Time = 2358 ?  
yes
```

```
| ?- statistics.
memory (total)      787611008 bytes
  global stack      443817856 bytes:      7112 in use, 443810744 free
  local stack       183558176 bytes:      368 in use, 183557808 free
  trail stack       73793272 bytes:       64 in use, 73793208 free
  choice stack      73793768 bytes:       560 in use, 73793208 free
  program space     12647872 bytes:    11145360 in use, 1502512 free
program space breakdown:
  compiled code      3376112 bytes
  JIT code           2590352 bytes
  sw_on_key          1535184 bytes
  try_node           869248 bytes
  predicate          818400 bytes
  aatree             656208 bytes
  atom               515072 bytes
  interpreted code   333376 bytes
  incore_info        252464 bytes
  atom table         98336 bytes
  miscellaneous      46752 bytes
  SP_malloc          32064 bytes
  int_info           9936 bytes
  FLI_stack          5456 bytes
  BDD hash table     3168 bytes
  module             1840 bytes
  numstack           1056 bytes
  source info        176 bytes
  foreign resource    160 bytes
7279 atoms (343192 bytes) in use, 33547152 free
No memory resource errors

0.656 sec. for 13 global, 43 local, and 11 choice stack overflows
15.370 sec. for 79 garbage collections which collected 3866624680 bytes
0.000 sec. for 0 atom garbage collections which collected 0 atoms (0 bytes)
0.000 sec. for 0 defragmentations
0.000 sec. for 412 dead clause reclamations
0.000 sec. for 0 dead predicate reclamations
0.485 sec. for JIT-compiling 1097 predicates
29.365 sec. runtime
=====
45.391 sec. total runtime
734307.664 sec. elapsed time
yes
```


Statistics

```
measure_time(Keyword, Goal, Before, After, Diff):-  
    statistics(Keyword, [Before|_]),  
    Goal,  
    statistics(Keyword, [After|_]),  
    Diff is After-Before.
```

```
| ?- measure_time(runtime, fib(30,F), Before, After, Time).  
F = 832040,  
Before = 21973,  
After = 24333,  
Time = 2360 ?  
yes  
| ?- measure_time(total_runtime, fib(30,F), Before, After, Time).  
F = 832040,  
Before = 37110,  
After = 41141,  
Time = 4031 ?  
yes
```

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SICStus Libraries

- SICStus has several (~55) libraries, with different purposes
 - Providing common data structures, such as sets and ordered sets, bags, queues, association lists, trees, or graphs, among others
 - Promoting interoperability, with functionalities such as
 - Parsing and writing information in CSV, JSON or XML format
 - Connecting with databases
 - Connecting with Java or .Net applications
 - Sockets and web programming
 - Providing Object-Oriented abstraction
 - ...

Aggregate Library

- The *aggregate* library provides operators for SQL-like queries
 - Results can be aggregated using sum, count, min, max, ...

```
| ?- aggregate(count, Child^parent(Person, Child), NChildren), NChildren >1.  
Person = cameron,  
NChildren = 2 ? ;  
Person = claire,  
NChildren = 3 ?  
yes
```

```
| ?- aggregate( sum(Dur), O^D^C^T^flight(O, D, Company, C, T, Dur), _TotalDur),  
               aggregate( count, O^D^C^T^Dur^flight(O,D,Company, C, T, Dur), _Count),  
               AvgDuration is _TotalDur/_Count.  
Company = iberia,  
AvgDuration = 108.33333333333333 ? ;  
Company = lufthansa,  
AvgDuration = 165.0 ? ;  
Company = tap,  
AvgDuration = 122.0 ? ;  
no
```

CLPFD Library

- The *clpfd* library provides one of the best constraint programming solvers and library for integers
 - Very good for puzzles, and combinatorial optimization problems
 - Example: solve the 3x3 magic square

```
| ?- L = [A,B,C,D,E,F,G,H,I], domain(L, 1, 9), all_distinct(L),  
    A+B+C#=Sum, D+E+F#=Sum, G+H+I#=Sum,  
    A+D+G#=Sum, B+E+H#=Sum, C+F+I#=Sum,  
    A+E+I#=Sum, C+E+G#=Sum, labeling([], L).
```

```
L = [2,7,6,9,5,1,4,3,8],
```

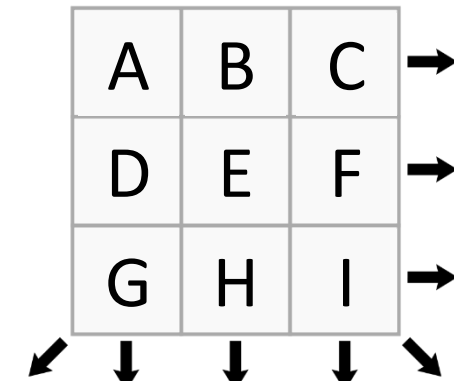
```
A = 2,
```

```
...
```

```
I = 8,
```

```
Sum = 15 ?
```

```
yes
```



CLPFD Library

- Another example: schedule seven resource-consuming tasks so they finish as quickly as possible and such that no more than a maximum is consumed at any given time

Task	Duration	Energy Consumption
1	16	2
2	6	9
3	13	3
4	7	7
5	5	10
6	18	1
7	4	11

Maximum instantaneous energy consumption: 13

CLPFD Library

```
schedule(Ss, End):-
    length(Ss, 7), domain(Ss, 1, 30),
    length(Es, 7), domain(Es, 1, 50),
    buildTasks(Ss, [16,6,13,7,5,18,4], Es, [2,9,3,7,10,1,11], Tasks),
    maximum(End, Es),
    cumulative(Tasks, [limit(13)]),
    labeling([minimize(End)], [End|Ss]).
```

```
buildTasks([], [], [], [], []).
buildTasks([S|Ss], [D|Ds], [E|Es], [C|Cs], [task(S, D, E, C, 0)|Ts]):-
    buildTasks(Ss, Ds, Es, Cs, Ts).
```

```
| ?- schedule(Starts, End).
Starts = [1,17,10,10,5,5,1],
End = 23 ?
yes
.
```

Q & A

