

# Functional and Logic Programming

*Bachelor in Informatics and Computing Engineering*  
2024/2025 - 1<sup>st</sup> Semester

## Introduction to Prolog

# Agenda

- Facts and Rules
- Queries
- How Prolog works

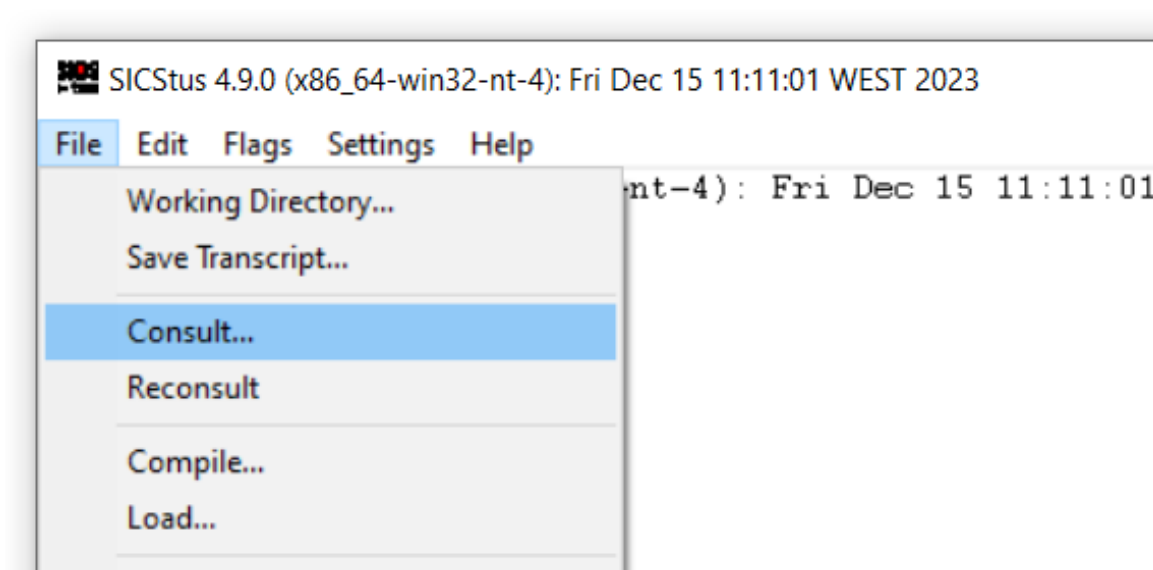
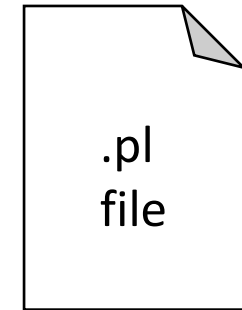
# Prolog

- Prolog is the most widely used logic programming language
  - There are some language dialects, such as Edinburgh Prolog, and also a standardization - ISO Prolog
    - See Péter Szabó & Péter Szeredi, Improving the ISO Prolog Standard by Analyzing Compliance Test Results. Proc. of the 2006 Int'l Conf. on Logic Programming, pp 257-269
- There are several Prolog systems, both free and commercial
  - Some of the most popular are SICStus Prolog and SWI-Prolog
  - Another notable system is YAP, developed at DCC, FCUP
    - See Körner et al., “Fifty Years of Prolog and Beyond”. Theory and Practice of Logic Programming 22 (6), pp. 776-858, 2022

In this course, we'll be using SICStus Prolog v 4.9  
(link to installer and keys available in Moodle)

# Prolog

- Write your code in a text file with a .pl extension
  - Use the text editor of your choice
- In SICStus, load it using the *File -> Consult...* menu
  - Or call directly on the SICStus console: `| ?- consult('path/to/file.pl').`



Alternatively, you can use SPIDER  
(SICStus Prolog IDE, based on Eclipse)

# Facts

- Facts express a relation that is true
  - You can (kind of) interpret them as lines in a database table

Statements end with a period

male(homer) .	% homer is a male
female(marge) .	% marge is a female
father(homer, bart) .	% homer is the father of bart
mother(marge, bart) .	% marge is the mother of bart

Arguments between parentheses and separated by commas

Predicate (relation) names start with lowercase letter

## Semantics

- The semantics (interpretation) needs to be defined and shared

```
father(homer, bart).      % homer is the father of bart
father(homer, bart).      % the father of homer is bart
```

- This inherent ambiguity only highlights the importance of using appropriate and descriptive names as well as code comments

```
% single-line comment
```


```
/* multi-line
   comment */
```

Naming conventions and code comments represent a part of the evaluation of the practical assignment

# Rules

- Rules allow for the deduction of new knowledge from existing knowledge (facts and other rules)
  - Rules are expressed in the form of Horn Clauses:
    - Head :- Body

```
grandfather(X, Y) :- father(X, Z), parent(Z, Y).    % X is the grandfather of Y
                                                    % if X is the father of Z
                                                    % and Z is a parent of Y
```



%multiple definitions of a rule with the same head: rule one **or** rule two **or**...

```
parent(X, Y) :- father(X, Y).    % X is a parent of Y if X is the father of Y
parent(X, Y) :- mother(X, Y).   % X is a parent of Y if X is the mother of Y
```

## Disjunction

- Disjunction can also be expressed with the ; operator

```
parent(X, Y):- father(X, Y).      % X is a parent of Y if X is the father of Y
parent(X, Y):- mother(X, Y).      % X is a parent of Y if X is the mother of Y
```

% is equivalent to

```
parent(X, Y):- father(X, Y) ; mother(X, Y).
```

- The disjunction operator (;) should be used sparingly
  - Always use parentheses to clarify



# Rules

- Rules have both a declarative and a procedural interpretation
  - Declarative interpretation

<code>grandfather(X, Y) :-</code>	<code>% X is the grandfather of Y</code>
<code>    father(X, Z),</code>	<code>% if X is the father of Z</code>
<code>    parent(Z, Y) .</code>	<code>% and Z is a parent of Y</code>

- Procedural interpretation

<code>grandfather(X, Y) :-</code>	<code>% to solve grandfather(X,Y)</code>
<code>    father(X, Z),</code>	<code>% first solve father(X, Z)</code>
<code>    parent(Z, Y) .</code>	<code>% and then parent(Z, Y)</code>
	<code>% (solve = execute)</code>

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# Rules

- The head of a rule can have 0 or more arguments

```
parent(X, Y):- father(X, Y).      % X is a parent of Y if X is the father of Y

father(X):- father(X, Y).         % X is a father if he is the father of some Y

fathers:- father(X, Y).           % fathers is true if there is a(t least one)
                                  % father/child relation
```

A rule with no arguments is a good entry point to a program

# Prolog Programs

- A Prolog program is a finite set of predicates
  - Predicates use facts and rules to express knowledge as relations
  - Relations are generalizations of functions
    - Usually more versatile, usable in multiple directions
- A computation is a proof of a goal from a program
  - Using [a form of] SLD resolution with a unification algorithm
- A **correct** program does not allow the deduction of unwanted facts
- A **complete** program allows the deduction of everything intended

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# Terms

- Everything in Prolog is a *term*, which can be a *constant*, a *variable* or a *compound term*
- **Constants** represent elementary objects
  - **Numbers**
    - **Integers** (e.g., 4, -8) (bases other than decimal can also be used, e.g., 8'755)
    - **Floats** (e.g., 1.5, -1.6) (also supports exponent, e.g., 23.4E-2)
  - **Atoms**
    - Start with lower-case letter (e.g., john\_doe, johnSmith42)
    - String within single quotes (e.g., 'John Doe', 'John Smith 42')

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# Terms

- **Variables** act as placeholders for arbitrary terms
  - Start with a capital letter (e.g., Variable1)
  - Start with an underscore (e.g., \_Var2)
  - Single underscore (\_) (anonymous variable)
- **Compound terms** are comprised of a *functor* and *arguments* (which are terms)
  - The functor is characterized by its *name* (an atom) and *arity* (the number of arguments), usually represented as *name/arity*
  - E.g., point/2 represents a functor named *point* with two arguments
    - point(4, 2) is a possible instance of point/2, and so is point(foo, point(3, bar))

# Variables in Programs

- Variables are universally instantiated in logic programs

```
plus(0, S, S).           % 0 is the neutral element of addition
mult(1, V, V).           % 1 is the neutral element of multiplication
```

```
human(Homer).            % everything is human
father(homer, Bart)       % homer is the father of everything
```

```
grandfather(X, Y) :- father(X, Z), parent(Z, Y).
```

Variables occurring only in the body of a rule can be seen as existentially quantified

We need to be careful when using variables with facts

## Coding Efficiency Considerations

- Use implicit unification instead of additional variables

```
change_player(X, Y) :- X = 1, Y = 2.  
change_player(X, Y) :- X = 2, Y = 1.
```

Should instead be written as

```
change_player(1, 2).  
change_player(2, 1).
```

- Always place input arguments before output arguments
  - SICStus indexes predicates by their first argument

= is the unification operator (kind of '[possibly] equal');  
\= (not unifiable) can be interpreted as 'can't be equal'

## Coding Style Considerations

- Although white space and code indentation are meaningless, there are some coding style guidelines you should consider following, to increase code readability:
  - Indent the code consistently
  - Put each sub-goal on a separate, indented line
  - Use human-readable names for predicates and variables
  - Try to limit the length of code lines and number of lines per clause
  - ...

See Covington et al. (2012). Coding Guidelines for Prolog. Theory and Practice of Logic Programming, 12(6): 889-927



# Agenda

- Facts and Rules
- **Queries**
- How Prolog works

# Queries

- Computations in Prolog start with a question, which has two possible answers:
  - Yes (possibly with answer substitution - variable binding)
  - No
- The attempt to prove the question right/wrong (is it a consequence of the program?) produces the computations

```
| ?- male(homer).  
yes
```

```
| ?- father(homer, bart).  
yes
```

```
| ?- female(marge).  
yes
```

```
| ?- father(marge, bart).  
no
```

## Variables in Queries

- Queries can include variables
  - Variables are existentially quantified in queries
- A variable starting with an underscore is a '*don't care*'

```
| ?- father(X, bart).  
X = homer ?  
yes
```

```
| ?- father(_X, bart).  
yes
```

```
| ?- male(_).  
yes
```

```
| ?- male(X).  
X = homer ?  
yes
```

```
| ?- male(X).  
X = homer ? ;  
X = bart ? n  
no
```

If satisfied with the answer, just hit enter

If you want another answer, type 'n', 'no' or ';'

## Variables and Compound Queries

- Queries can be more complex, combining goals
- Variables are used to glue together the different goals
  - Underscore alone (`_`) is the exception

```
| ?- male(X), parent(X, bart).  
X = homer ? ;  
no
```

```
| ?- male(_X), parent(_X, bart).  
yes
```

```
| ?- male(_X), parent(Y, bart).  
Y = homer ? ;  
Y = marge ? ;  
Y = homer ? ;  
Y = marge ? ;  
no
```

Why the duplicates?  
Just wait a few slides!

## Closed World Assumption

- Assumption that everything that is true is known to be true (i.e., is represented as a clause in the program)
- Therefore, everything that cannot be deduced from the clauses in the program is assumed to be false

```
| ?- male(donald) .  
no
```

- Requires attention to make sure everything we want to deduce can be deduced from the program clauses

# Horn Clauses

- Everything in Prolog is expressed as a Horn Clause

- Rules are complete horn clauses (head :- body)

```
parent(X, Y) :- father(X, Y) .
```

- Fact are horn clauses where the body is always true (just the head)

```
male(homer) :- true           ⇔           male(homer) .
```

- Queries are horn clauses without a head (just the body)

```
| ?- father(X, bart) .
```

# Predicates

- A **predicate** is a set of clauses for the same functor
  - **Clauses** are either facts or rules
- Functors with the same name but different arity refer to different predicates

```
father(X) :- father(X, Y) .      % X is a father
                                   % if X is the father of some Y
```

# Documentation

- Documentation should include a **mode declaration** for each argument:
  - + (input): the argument is instantiated when the predicate is called
  - - (output): the argument is not instantiated in the predicate call
  - ? (in/out): the argument can be instantiated or not

```
% square(+number, -square)
```

```
% calculates the square  
% of a given number
```

```
% parent(?parent, ?child)
```

```
% parent/child relation
```

- One of the most powerful properties of Prolog is its versatility



# Prolog Versatility

- The versatility of Prolog can be seen in most predicates:
  - For instance, parent/2 allows:
    - Confirming that two given people are parent/child
    - Obtaining the children of a given person
    - Obtaining the parents of a given person
    - Obtaining all parent/child pairs
- In most other languages, we would need to implement four different functions to achieve this, or include extra logic to test instantiation

# Prolog and Relational Algebra

- A Prolog program can be seen as a database
  - Facts represent tables
  - Rules represent views
- Prolog can be used to implement all relational algebra operations, like union, cartesian product, projection, selection, ...

# Relational Operations

- Union

$$r\_union\_s( X_1, \dots, X_n ):- r( X_1, \dots, X_n ).$$

$$r\_union\_s( X_1, \dots, X_n ):- s( X_1, \dots, X_n ).$$

- Cartesian product

$$r\_times\_s( X_1, \dots, X_m, X_{m+1}, \dots, X_{m+n} ):- r( X_1, \dots, X_m ), s( X_{m+1}, \dots, X_{m+n} ).$$

- Projection

$$r\_1\_3( X_1, X_3 ):- r( X_1, X_2, X_3 ).$$

- Selection

$$r\_1( X_1, X_2, X_3 ):- r( X_1, X_2, X_3 ), X_2 > X_3.$$

- Intersection

$$r\_inters\_s( X_1, \dots, X_n ):- r( X_1, \dots, X_n ), s( X_1, \dots, X_n ).$$

- Join

$$r\_join\_s( X_1, X_2, X_3 ):- r( X_1, X_2 ), s( X_2, X_3 ).$$

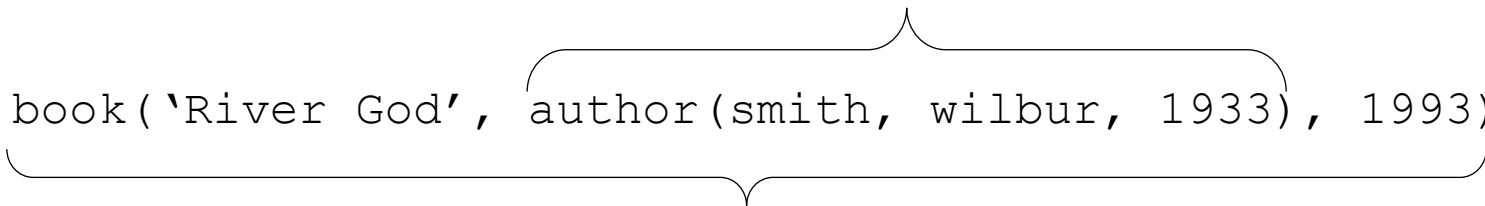
- Difference

$$r\_minus\_s( X_1, \dots, X_n ):- r( X_1, \dots, X_n ), \setminus + s( X_1, \dots, X_n ).$$

# Prolog and Relational Algebra

- Complex terms vs ‘normalized’ facts

`has(john, book('River God', author(smith, wilbur, 1933), 1993)).`



```
author(a37, smith, wilbur, 1933).  
book(b521, 'River God', 1993).  
author(a37, b521).  
person(p432, john).  
has(p432, b521).
```

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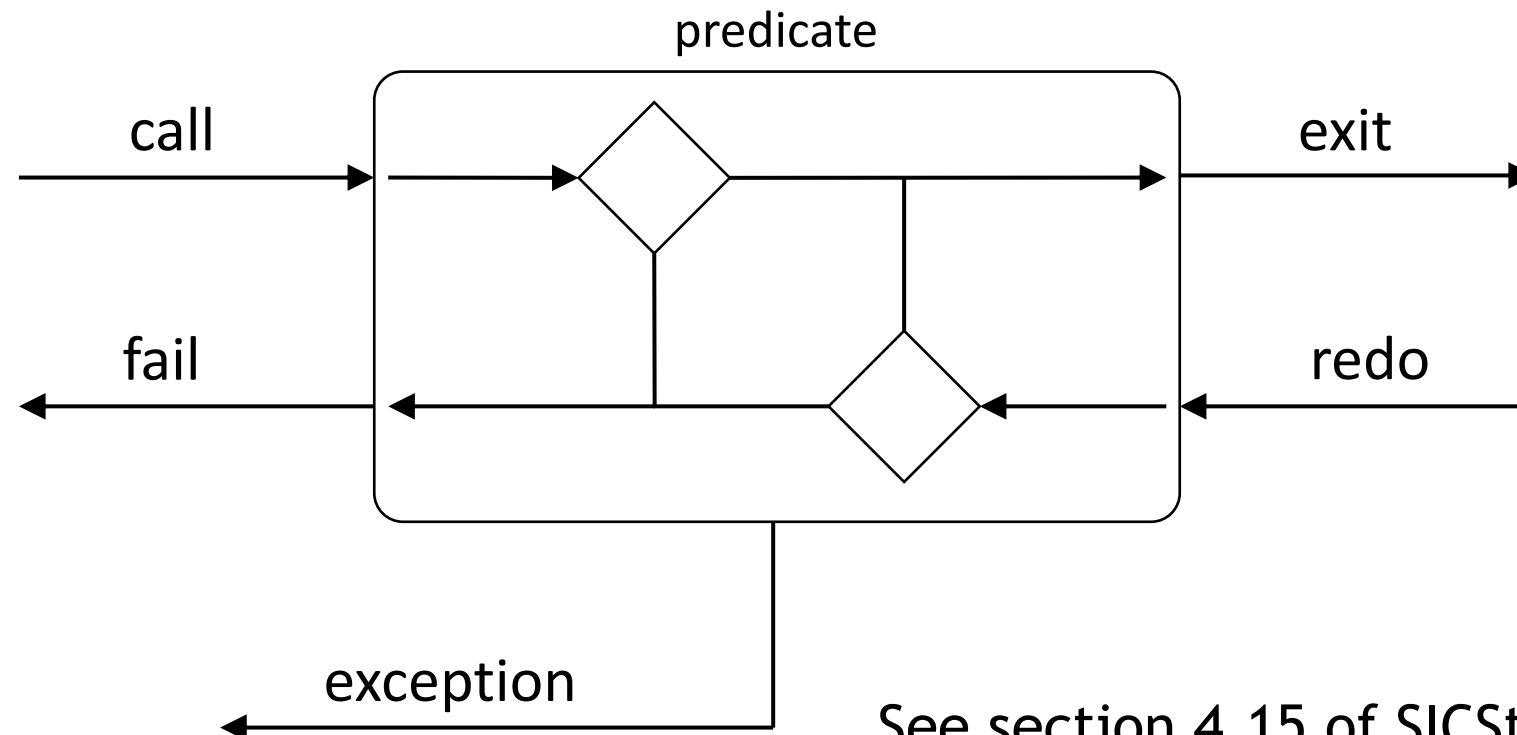
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# How Prolog Works

- Prolog's mechanics work
  - Top to bottom
    - The order of clauses is important
  - Left to right
    - In rules, prove sub-goals in left-to-right order
  - With backtracking
    - If a sub-goal fails, go back to previous decision point

## The Prolog Box Model

- Each call to a goal can be modelled as a four-gate box model



See section 4.15 of SICStus's Manual  
for more information on exceptions

## Tracing

- Trace mode allows us to follow the computations step by step
  - Can be activated from the menu Flags -> Debugging -> trace
  - Or in the code, by calling *trace*
    - Disable it by calling *notrace*

```
foo(bar) :- fubar(bar, baz),  
            trace,                % activate trace mode  
            qux(baz),             % the call to qux will be traced  
            notrace,             % deactivate trace mode  
            quux(bar) .
```

See section 5 of SICStus's Manual for  
more information on Trace and Debugging



# Tracing

- Trace message format:

N S InvID Depth Port: Goal ?

- N (only visible at Exit ports) indicates that the goal call may backtrack to find alternative solutions
- S indicates the existence of a spypoint
- InvID (Invocation ID) is a unique identifier for each goal (can be used to match messages from the various ports)
- Depth is an indication of the general call depth
- Port is one of Call, Exit, Redo, Fail or Exception
- Goal is the current goal of the computation

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## Additional Readings

- Prolog
  - Leon Sterling and Ehud Shapiro (1994). The Art of Prolog. The MIT Press (2<sup>nd</sup> ed). ISBN: 978-0262691635
  - Krzysztof R. Apt (1996). From Logic Programming to Prolog. Prentice Hall. ISBN: 978-0132303682
  - Patrick Blackburn, Johan Bos and Kristina Striegnitz (2006). Learn Prolog Now! College Publications. ISBN: 978-1904987178
  - Ivan Bratko (2011). Prolog Programming for Artificial Intelligence. Addison Wesley (4<sup>th</sup> ed). ISBN: 978-0321417466
  - Max Bramer (2013). Logic Programming with Prolog. Springer (2<sup>nd</sup> ed). ISBN: 978-1447154860

## Q & A

