IPv6 Specification

簡介: Some changes from IPv4 to IPv6:

- **Expanded Addressing Capabilities** (RFC 1884)
 - from 32 bits to 128 bits (more level and nodes)
 - improve multicast routing ("scope" field)
 - "anycast address": send a packet to any one of a group of nodes
- Header Format Simplification
 - —limit bandwidth cost
- **Extensions and Options**
- —more flexibility
- Flow Labeling Capability
- -define traffic "flow" to support special QoS
- **Authentication and Privacy Capabilities**
- —authentication, data integrity, and data confidentiality

IPv6 Header Format:

32 bits

Version	Prio.	Flow Label					
	Payload	Length	Next Header	Hop Limit			
Source address							
	Destination address						

IPv6 Extension Headers:

Example:

1	
IPv6 header	TCP Header + data
Next Header = TCP	

IPv6 header Next Header =	Routing header Next Header = TCP	TCP Header + data
Routing		

IPv6 header	Routing header	Fragment header	Fragment of TCP
Next Header =	Next Header =	Next Header $=$ TCP	Header + data
Routing	Fragment		

⇒ <u>Hop-by-Hop Options</u>, <u>Routing</u>, <u>Fragment</u>, <u>Destination Options</u>, <u>Authentication</u> (RFC 1826), <u>Encapsulating Security Payload</u> (RFC 1827)

⇒除了 hop-by-hop options 外, 其他的 extension headers 不可以被在傳送路徑中的其他非 destination 的 nodes 檢查與處理

⇒each extension header 的內容與語意決定是否要處理下一個 header, 因此, extension header must processed strictly in Order

⇒Extension Header Order:

IPv6

Hop-By-Hop Options header (0)

Destination Options header (60)

Routing header (43)

Fragment header (44)

Authentication header

Encapsulating Security Payload header

Destination Options header

Upper-layer header

Note: TCP (6), *UDP*(17), *Nothing*(59)

⇒錯誤處理: ICMP (RFC 1885)

⇒Each extension header occurs at most once, except for Destination Options header which occurs at most twice

Options:

⇒Hop-by-Hop Options header and Destination Options header carry a variable number type-length-value (TLV) encoded "options"

⇒Option format:

⇒Option type:

highest-order two bits 表示當無法辨識此 option type 時所要採取的行動

- 00—skip this option and continue processing the header
- 01 —discard the packet
- 10 —discard the packet and send ICMP message to source
- 11—discard the packet and if Destination address was not a multicast address send ICMP message to source

third-highest-order bit 表示在傳送的過程中 Option data 是否被改變。

- 0—Option data doesn't change en-route
- 1 —Option data may change en-route

Alignment:

Notation xn+y: option type must appear at an integer multiple of x bytes from the start of the header, plus y bytes.

Padding options:

Pad1 option: used to insert one byte of padding into the option area

One byte, special option format: no length and value field

PadN option: used to insert more than one byte

		<u> </u>
1	Opt Len	Opt data = $0 \dots$

Example: 4n+3 (hop-by-hop options)

1 1 1 1 /						
Next Header	Hdr Ext Len =1	Pad1 Option=0 Option type=Y				
Opt Data Len=7	1-Byte field	2-byte filed				
4-byte field						
Pad4 Option=1	Opt Len=2	0	0			

Hop-by-Hop Options header:

Next header	Hdr Ext len					
Options						
	(multiple of 8 bytes)					

Jumbo Payload option (4n+2) – hop-by-hop option

•	`	,			1 1		
					194	O	pt Data Len=4
		Jumbo F	aylo	oad	length		

- ⇒Used to send packets with payload longer than 65535 bytes.
- ⇒Payload length in IPv6 header must be zero
- ⇒Jumbo payload options must not be used in a packet that carries a Fragment header
- \Rightarrow Links' MTU must be greater than 65575 = 65535+40
- ⇒功用: 如 the special handling by routers (source desires)

Routing Header

Next header	Hdr Ext Len	Routing Type	Segment Left				
Type-specific data							

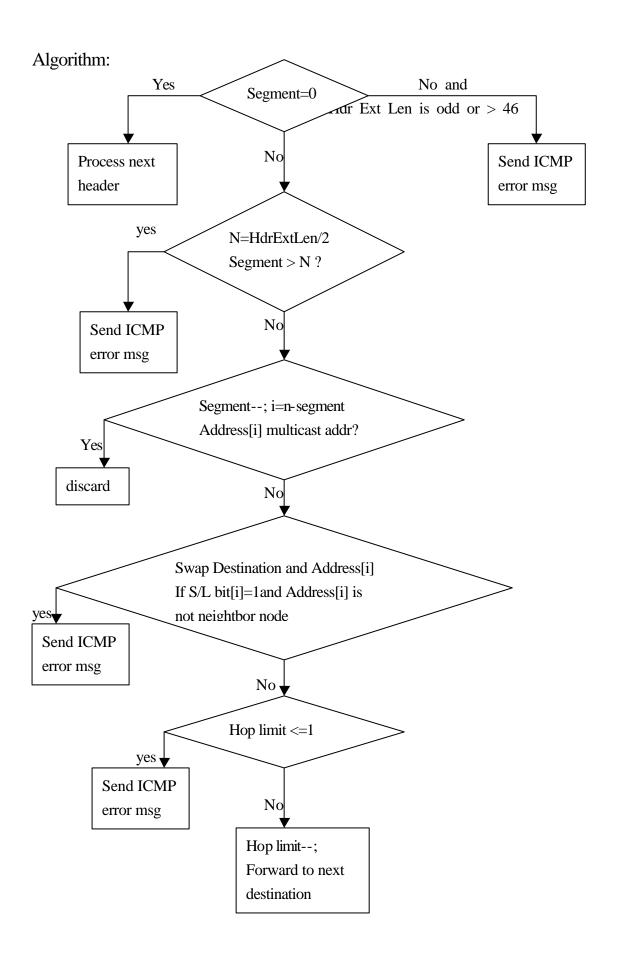
Type 0 Routing header:

Next header	Hdr Ext Len	Routing Type=0	Segment Left				
Reserved	Strict/Loose Bit Map						
Address[1]							

Address[n]

[⇒]Segment Left: Maximum legal number = 23

^{⇒24-}bits bit-map, number 0 to 23, left-to-right, indicate whether or not the next destination address must be a neighbor: 1–strict, 0–loose



Example:

Source: S, Destination: D Intermediates: I1, I2, I3

As the packet travels from S to I1

Source address = S Hdr Ext Len = 6
Destination=I1 Segment left = 3
(if bit 0 of Bit Map is 1, Address[1]=I2
S and I1 must be neighbors; Address[2]=I3

Address[3]=D

Address[3]=D

Address[3]=I3

As the packet travels from I1 to I2

This is checked by S)

Source address = S Hdr Ext Len = 6
Destination = I2 Segment left = 2
(if bit 1...) Address[1]=I1
Address[2]=I3

As the packet travels from I2 to I3

Source address = S Hdr Ext Len = 6
Destination = I3 Segment left = 1
(if bit 2). Address[1]=I1
Address[2]=I2
Address[3]=D

As the packet travels from I3 to D

 $Source \ address = S \\ Destination = D \\ (if \ bit \ 3 \ .). \\ Segment \ left = 0 \\ Address[1] = I1 \\ Address[2] = I2$

Fragment Header:

⇒ fragmentation in IPv6 is performance only by Source

Next Header	Reserved	Fragment offset	Res	M		
Identification						

⇒Every packet must have different identification, increase each time when a packet must be fragmented

original packet:

Unfragme	ntable Part	Fragmentable Part
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- ⇒Unfragmentable Part: IPv6 header and extension headers that must be processed by nodes en-route to the destination
- ⇒Fragmentable Part: extension header that need be processed only by the final destination, plus upper–layer header and data

original packet:

Unfragmentable	First	Second	• • • • •	Last
part	fragment	fragment		fragment

Fragment packets:

Unfragmentable	Fragment	First
part	Header	fragment

.

Unfragmentable	Fragment	Last
part	Header	fragment

Each packet is composed of:

- (1) Unfragmentable part: payload length = fragment packet only
- (2) Fragment header: Fragment offset, M flag, Identification
- (3) Fragment
- ⇒Reassemble packet from fragment, packets must have the same <u>Source</u> <u>Address</u>, <u>Destination Address</u>, and <u>Fragment Identification</u>

- ⇒Error conditions:
- (1)insufficient fragments are received to complete reassembly of a packet within 60 seconds
- (2)flag M = 1 and payload length is not a multiple of 8 bytes
- (3) payload length of fragment > 65535

Destination Options Header

⇒used to carry optional information that need be examined only by destination

Next header	Hdr Ext Len	
_		
	optio	ıs

⇒two way to encode optional destination information: Destination Options Header, or as a separate extension header (fragment/authentication)

Packet size issues:

- ⇒MTU of every link must >= 576 bytes,如果無法提供,則必須靠下層來做 fragmentation and reassembly
- ⇒可以利用 Path MTU Discovery[RFC-1191] to discover and take advantage of path with MTU greater than 576 bytes
- ⇒最簡單的方法就是限制 packet 的 size <=576 bytes
- \Rightarrow a node must not send fragments that reassemble to a size greater than 1,500 byte.
- ⇒IPv6 packet is sent to an IPv4 destination.
 - Source may receive ICMP message.
 - Subsequent packets to 528 (payload=576-40-8) and include Fragment header, so that the IPv6-to-IPv4 translating router can know.

Flow Labels

⇒A flow: a sequence of packets which the source desires special handling (by hop-by-hop option or control protocol)

identified by combination of source address and non-zero flow label(randomly and uniformly chosen), as a hash key used by router.

All packets belonging to the same flow must have the same <u>source</u> address, <u>destination</u> address, <u>priority</u>, and flow label

⇒The router may cache the information of next-hop interface, how to queue the packet based on its priority, etc...

Priority:

0—7: provide congestion control

8—15: no back off in response to congestion, real-time packets

8 : most willing to discard under congestion (e.g. high-fidelity video traffic)

15: least willing ...(e.g. low-fidelity audio traffic)

Upper-layer Protocol Issues

⇒Upper-layer Checksums – Pseudo-header

⇒Maximum Packet Lifetime – don't enforce maximum packet lifetime

 \Rightarrow Upper-layer payload length –

IPv4: max packet size – 40 (20 for IP header, 20 for TCP header)

IPv6: max packet size – 60