

## Assignment 3: Algorithm Analysis

1. (5 points) Give a big-Oh notation, in terms of  $n$ , of the running time of the algorithm shown in the following code fragment:

```
/** Returns the sum of the prefix sums of given array. */
public static int example3(int[ ] arr) {
    int n = arr.length, total = 0;
    for (int j = 0; j < n; j++) // loop from 0 to n-1
        for (int k = 0; k <= j; k++) // loop from 0 to j
            total += arr[j];
    return total;
}
```

2. (5 points) Consider the problem of matrix multiplication, where the input includes two matrices,  $A$  (of dimension  $n \times k$ ) and  $B$  (of dimension  $k \times m$ ), the output is matrix  $C$  (of dimension  $n \times m$ ). Recall that each element  $C[i][j]$  in matrix  $C$  is the sum of the element-wise products of the  $i$ -th row of  $A$  and the  $j$ -th column of  $B$ . Consider the following pseudocode for matrix multiplication. Express the running time of the algorithm in big-Oh notation with the above dimension parameters.

Algorithm matProduct( $A$ ,  $B$ ):

```
for (i = 0:n-1) do
    for (j = 0:m-1) do
        C[i][j] = 0
        for (h = 0:k-1) do
            C[i][j] += A[i][h] * B[h][j]
```

3. (a) (10 points) The following pseudocode shows a recursive algorithm for computing  $2^n$  for any integer  $n \geq 0$ . It is based on the recurrent relationship:  $2^n = 2^{n-1} + 2^{n-1}$ . Only

consider addition as a primitive operation. Use the substitution method to show the running time of this algorithm in big-Oh notation.

```
Algorithm recurPower(n):  
    if (n == 0) then  
        return 1  
    else  
        // addition as a primitive operation  
        return recurPower(n-1) + recurPower(n-1)
```

- (b) **(5 points)** The following pseudocode shows a non-recursive algorithm for computing  $2^n$ . Show the running time of this algorithm in big-Oh notation. Here, you can consider multiplication as a constant-time primitive operation.

```
Algorithm iterPower(n):  
    if (n == 0) then  
        return 1  
    else  
        power = 1  
        for (i = 0:n-1) do  
            // consider multiplication as a primitive  
            operation  
            power = power * 2  
        return power
```

4. **(10 points)** Suppose there is a recursive algorithm and its running time  $T(n)$  satisfies a recurrence relationship  $T(n) = 2T(n/2) + n$  when the input size  $n > 1$ , and  $T(1) = 1$ . Use the substitution method to show the running time  $T(n)$  in the big-Oh notation.