

Recurrences

Recurrence relation

Definition: Given a sequence of numbers $a_0, a_1, a_2, a_3, \dots$, a **recurrence relation** is an equation that expresses a_n in terms of one or more previous terms in the sequence for all $n \geq n_0$ where $n, n_0 \in \mathbb{Z}^+$.

= inductive def of a sequence
= recursive def of a sequence

a_0 = initial value.

(initial condition)

Example: compound interest

Suppose you deposit \$5,000 at 3% interest (compounded annually). How much money will you have after 30 years?

$$P_n = \text{amount after } n \text{ years} = P_{n-1} + 0.03 \times P_{n-1} = 1.03 P_{n-1} \text{ for } n > 0$$

$$P_0 = 5000$$

$$P_1 = 1.03 \cdot P_0 = 1.03 \times 5000$$

\vdots

$$P_n = 1.03^n \times 5000 = (1.03)^n P_0$$

← formula for P_n only including n & P_0 initial value.

Example: vending machine

Suppose a vending machine accepts only \$1 coins, \$1 bills, and \$5 bills. **Find a recurrence relation for the number of ways to deposit n dollars into the vending machine**, where the order in which the coins/bills are deposited matters.

Let $a_n = \# \text{ ways to deposit } n \text{ dollars into vending machine.}$

$$a_0 = 1$$

$$a_1 = 2$$

$$a_2 = 4 \quad [c, c, b] \quad (b, c) \quad (c, b) \quad (c, c)]$$

three $a_3 = 8$. 2 choices for 1st dollar, 2 choices for 2nd dollar, 2 choices for 3rd \$.

$$a_4 = 16$$

$$a_5 = 1\$5 \text{ or } 5\$1 = 1 + 2^5 = 33 = a_0 + a_2 + a_4.$$

$$a_6 = a_1 \quad \text{deposit } \$5 \text{ 1st}$$

$$+ a_5 \quad \text{deposit } \$1 \text{ bill first}$$

$$+ a_5 \quad \text{deposit } \$1 \text{ coin first}$$

\vdots

$$\text{solution } a_0 = 1, a_1 = 2, a_2 = 4$$

$$a_3 = 8, a_4 = 16,$$

$$a_n = a_{n-5} + 2 \cdot a_{n-1} \text{ for } n \geq 5$$

$$a_n = \underset{\substack{\uparrow \\ \text{1st is } \$5 \text{ bill}}}{a_{n-5}} + \underset{\substack{\uparrow \\ \text{1st is } \$1 \text{ bill}}}{a_{n-1}} + \underset{\substack{\uparrow \\ \text{1st is } \$1 \text{ coin}}}{a_{n-1}} = a_{n-5} + 2 \cdot a_{n-1}$$

1st is \$5 bill 1st is \$1 bill 1st is \$1 coin.

Solving a recurrence relation

solving the recurrence = coming up with a formula for P_n with only n and constants in it.

Solving linear homogeneous recurrence relations of degree 2

Given a sequence $\{a_n\}$ described by the recurrence relation $a_n = c_1 a_{n-1} + c_2 a_{n-2}$ where $c_1, c_2 \in \mathbf{R}$ and where $r^2 - c_1 r - c_2 = 0$ has 2 distinct roots r_1 and r_2

The solution to the recurrence relation is $a_n = \alpha r_1^n + \beta r_2^n$ for $n = 0, 1, 2, \dots$ where α and β are constants.
 \uparrow characteristic equation

Example:

Solve the recurrence relation $a_n = 5a_{n-1} - 6a_{n-2}$ for $n \geq 2$, $a_0 = 1$, $a_1 = 0$

$$\begin{aligned} c_1 &= 5 & c_2 &= -6 \\ r_1 &= 3 & r_2 &= 2 \end{aligned} \quad \begin{aligned} &\text{characteristic equation} \rightarrow r^2 - 5r + 6 = 0 \\ &(r^2 - c_1 r - c_2) & (r-3)(r-2) = 0 \end{aligned}$$

Solution looks like: $a_n = \alpha \cdot 3^n + \beta \cdot 2^n$ (what α & β)? \leftarrow use the initial condition $(a_0=1, a_1=0)$

Using $a_0 = 1$: $1 = \alpha \cdot 3^0 + \beta \cdot 2^0 = \alpha + \beta \rightarrow \alpha = 1 - \beta$

$a_1 = 0$: $0 = \alpha \cdot 3^1 + \beta \cdot 2^1 = 3\alpha + 2\beta$ \leftarrow $\beta = 3$
 $\alpha = -2$

Solution is $a_n = -2 \cdot 3^n + 3 \cdot 2^n$

Double-check answer:

n	$a_n = 5a_{n-1} - 6a_{n-2}$	$a_n = -2 \cdot 3^n + 3 \cdot 2^n$
0	1	$(-2) \cdot 3^0 + 3 \cdot 2^0 = 1$
1	0	$-2 \cdot 3^1 + 3 \cdot 2^1 = 0$
2	$5a_1 - 6a_0 = -6$	$-2 \cdot 3^2 + 3 \cdot 2^2 = -18 + 12 = -6$
3	$5a_2 - 6a_1 = -30$	$-2 \cdot 3^3 + 3 \cdot 2^3 = -54 + 24 = -30$

Recurrences (continued)

Example: Towers of Hanoi

The Towers of Hanoi is a game that consists of n disks and three poles: A (the source), B (the destination), and C (the spare). The disks are of different sizes and each has a hole in the middle so that it may be placed on a pole. The game starts with all the disks on pole A arranged by size with the largest disk on the bottom. The goal of the game is to move all the disks to pole B , with the following restrictions:

- only one disk at a time may be moved from one pole to another pole
- a disk can only be placed on top of a larger disk

Disks can be placed on pole C during the transfer, as long as the above restrictions are met.

How many moves does it take to solve the Towers of Hanoi with n disks?

Step 1) Find a recurrence relation for the number of moves.

Let $m(n)$ = # moves to solve with n disks

1 disk \Rightarrow 1 move

n disks \Rightarrow move $n-1$ disks from A to C ($m(n-1)$ moves)

move 1 disk from A to B (1 move)

move $n-1$ disks from C to B ($m(n-1)$ moves)

Recurrence relation: $m(1) = 1$

$$m(n) = m(n-1) + 1 + m(n-1) = 1 + 2 \cdot m(n-1) \text{ for } n > 1.$$

Step 2) Find a pattern and propose a solution.

n	$m(n)$
1	$1 = 2^1 - 1$
2	$1 + 2m(1) = 3 = 2^2 - 1$
3	$1 + 2m(2) = 7 = 2^3 - 1$
4	$1 + 2m(3) = 15 = 2^4 - 1$
5	$1 + 2m(4) = 31 = 2^5 - 1$

Proposed Solution
 $m(n) = 2^n - 1 \quad (\forall n \geq 1)$

$$m(1) = 1$$

$$m(n) = 1 + 2m(n-1) \text{ for } n > 1$$

Step 3) Prove the solution works using induction.

$$P(n) = m(n) = 2^n - 1 \quad \text{show } P(n) \text{ holds } \forall n \in \mathbb{N}^+$$

Base case: show $P(n)$ holds

$m(1) = 1$ by def of recurrence relation

$$2^1 - 1 = 1 \quad \text{so } m(1) = 2^1 - 1 \text{ \& } P(1) \text{ holds.}$$

Inductive step: show $P(k) \Rightarrow P(k+1)$

Assume $P(k)$ holds, i.e. $m(k) = 2^k - 1$ (IH)

Now consider $m(k+1)$

$$m(k+1) = 1 + 2 \cdot m(k) \text{ using recurrence relation}$$

$$= 1 + 2(2^k - 1) \text{ using (IH)}$$

$$= 1 + 2^{k+1} - 2$$

$$= 2^{k+1} - 1 \quad \text{so } P(k+1) \text{ holds.}$$

\therefore by induction $m(n) = 2^n - 1 \quad \forall n \in \mathbb{N}^+$

□

Recurrences and program analysis

Program analysis goal: analyze the **complexity** of algorithms

= how # of operations performed by algorithm grows a function of the input size of sequences → # of ops performed in input on size of 1, 2, 3...
 recursive programs → recurrence relations.

Given the following program specification:

Input: $a \in \mathbb{Z}, b \in \mathbb{Z}^+$

b: input size.

Output: $a^b = a \cdot a \cdot \dots \cdot a$ (i.e., a multiplied b times)

count multiplication

consider power 3 (3, 4)

Consider the following 3 implementations:

m_1 procedure $power1(a, b)$

current multiplications

- (1) if $b = 1$ then return a
- (2) else return $a \cdot power1(a, b - 1)$

b = input size

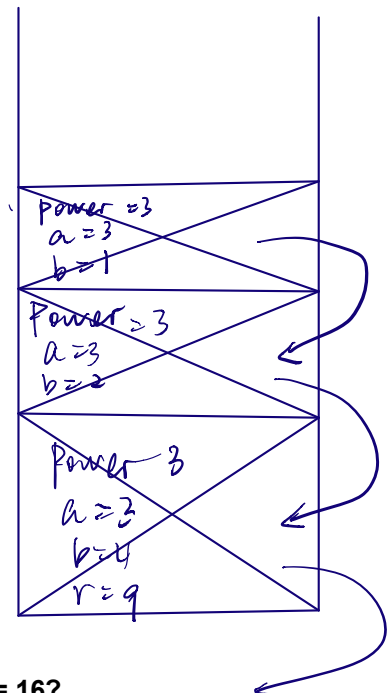
m_2 procedure $power2(a, b)$

- (1) if $b = 1$ then return a
- (2) else if b is odd then return $a \cdot power2(a, \lfloor b/2 \rfloor) \cdot power2(a, \lfloor b/2 \rfloor)$
- (3) else return $power2(a, \lfloor b/2 \rfloor) \cdot power2(a, \lfloor b/2 \rfloor)$

*floor $\lfloor x \rfloor$
 = largest integer K
 s.t. $K \leq x$*

m_3 procedure $power3(a, b)$

- (1) if $b = 1$ then return a
- (2) else $r = power3(a, \lfloor b/2 \rfloor)$
- (3) if b is odd then return $a \cdot r \cdot r$
- (4) else return $r \cdot r$



How many multiplications are done by each implementation when $a = 5$ and $b = 16$?

Let $m_i(n) = \# \text{ numbers done in power } i \text{ with input size } n = b$ Return 81

$$power(5, 16) : m_1(16) = 15$$

$$power(5, 16) : m_2(16) = 15$$

$$power(5, 16) : m_3(16) =$$

$$m_1(16) = 1 + m_1(15) = 15$$

$$\hookrightarrow 1 + m_1(14) = 14$$

$$\hookrightarrow 1 + m_1(13) = 13$$

$$\dots \hookrightarrow 1 + m_1(1) = 1$$

$\hookrightarrow 0 \text{ multiplications}$

$$m_2(16) = 1 + 2 \cdot m_2(8) = 5$$

$$m_3(16) = m_3(8) + 1 = 4$$

$$m_2(8) = 1 + 2 \cdot m_2(4) = 7$$

$$m_3(8) = m_3(4) + 1 = 3$$

$$m_2(4) = 1 + 2 \cdot m_2(2) = 3$$

$$m_3(4) = m_3(2) + 1 = 2$$

$$m_2(2) = 1 + 2 \cdot m_2(1) = 1$$

$$m_3(2) = m_3(1) + 1 = 1$$

$$m_2(1) = 0 \text{ multiplication}$$

$$m_3(1) = 0 \text{ multiplication.}$$