CS 354 - Machine Organization & Programming Thursday, September 26, 2019

Midterm Exam - Thursday, October 3rd, 7:15 - 9:15 pm

- Lec 1 (2:30 pm): room 3650 of Humanities
- Lec 2 (4:00 pm): room B10 of Ingraham Hall
- UW ID required
- closed book, no notes, no electronic devices (e.g., calculators, phones, watches)
- see "Midterm Exam 1" on course site Assignments for topics

Project p2A (3%) DUE: 10 pm, Monday, September 30th Project p2B (3%) DUE: 10 pm, Monday, October 7th

TIP: Use blank outlines to study for the exam.

Homework hw1 (1.5%) DUE TOMORROW: 10 pm, Friday, September 27th

Homework hw2 (1.5%) DUE: 10 pm, Wednesday, October 2nd

Last Time

Array Caveats
Command-line Arguments
Meet Structures
Nested Structures and Arrays of Structures
Passing Structures
Pointers to Structures

Today

Pointers to Structures (from last time) Standard & String I/O and stdio.h File I/O and stdio.h Copying Text Files

Three Faces of Memory Virtual Address Space C's Abstract Memory Model Where Do I Live?

Next Time

Globals and Static Locals
Linux Processes and Address Spaces
------ END of Exam 1 Material ----The Heap & Dynamic Memory Allocators
Read: B&O 9.9.4 - 9.9.5

Standard and String I/O in stdio.h Library

Standard I/O

Standard Output (console) putchar one char at a time. puts one string at a time. int printf(const char *format_string, comma-separated-list-of-vars) returns number of characters written, or a negative if error chars to display, format spenfiers And special chars it in format string format specifiers such %i, %d, %f, %p, %s Each forat specifiers much have a matching variable. Standard Input (keyboard) getchar gets int scanf(const char *format string, comma-separated-list-of-var-addrs) returns number of inputs successfully matched and assigned format string cheraiters to skip and format specifiers
touch forat specifiers much have a northery variable address whitespace lipit Separater (spaces, tabs, newlines) leading whitespace is skipped. Standard Error (console) void perror(const char *str) preferel for displaying error messages should be tig enough. String I/O int sprintf(char *str, const char *format, ...) prints formatted output to str int sscanf(const char *str, const char *format, ...) reads formatted input from str.

File I/O in stdio.h Library

```
Standard I/O Redirection
       $ a. out < input filerare = output filerane.
   File I/O
      File Output
          fputc/<del>putc</del>
          fputs
          int fprintf(FILE *stream, const char *format, ...)
             returns number of characters written, or a negative if error
      File Input
          fgetc/<del>getc</del>, ungetc
          fgets
          int fscanf(FILE *stream, const char *format, ...)
             returns non-negative value, or EOF if error
   File Pointers and Descriptors (3)
stdin, stdout, stderr predefined.

Cosole keybourd, window

printf ("text") = fprintl (stdont, "text");
   Opening and Closing
          FILE *fopen(const char *filename, const char *mode)
             returns open file pointer, or NULL if access problem
          int fclose(FILE *stream)
             returns 0, or EOF if error
```

Copying Text Files

```
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char *argv[]) {
   if (argc != 3) {
       fprintf(stderr, "Usage: copy inputfile outputfile\n");
       exit(1);
   }
   FILE *ifp, *ofp;
   ifp = . fopen ( or gv [ 1] / " r" );
if (ifp == NULL) {
       fprintf(stderr, "Can't open input file %s!\n", argv[1]);
      exit(1);
   ofp = fopen ( argr [2], "w");
if (ofp == NULL) {
      fprintf(stderr, "Can't open output file %s!\n", argv[2]);
exit(1); fclose (ifp);
   const int bufsize = 257; // worning: assume line & x5b chers. char buffer[bufsize];
while (facts (buffer, buffize, ifp 7! = NVLL)

fput (buffer, fp))
    folse (ifp);
flore (ofp);
   return 0;
```

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Three Faces of Memory

to enable multiple processing to run concurrently on a single markine. * A key OS goal is process: is a running pragram. **Process View = Virtual Memory** Goal: provide a single view of memory that 13 consistent for all processing Stack virtual address space (VAS): Musion by Os that each process has gone continuous mem "standbox" virtual address: Simulated adds that a process generate during execution Heap Data Code Hardware View = Physical Memory Goal: keep the CPV busy for speed and efficiency **CPU** L0: Registers physical address space (PAS): installed memory (in CPU) L1: Cache Note: PAS << VAS L2: Cache organized as a multi-Level Hierarhy L3: Cache that ensure fragmency used datal as (SRAM) L4: Main Memory (DRAM) physical address: actural address used to L5: Local Secondary Storage (HDD,SSD,DVD) access machine memory L6: Network Storage (NAS, Cloud) System View = Illusionist (CS 537) Goal: make memory sharable secure, Process 1 and efficiently divided arong processes Main Memory pages: mem blocks managed by of. Process 2 <u>os</u> page table: VAS Page OS structure that maps virtural pages Table Secondary to physical pages, which enables Storage 1 mapping vitual adds to physical adds Process 3 Secondary VAS Storage 2

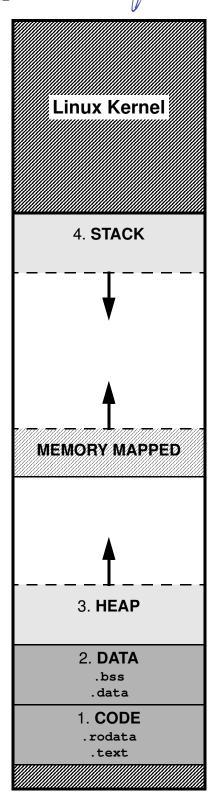
A Process' Virtual Address Space (IA-32/Linux)

32-bit Processor = 32-bit Addresses

recall: <u>byte addressability</u>: each address accesses 1 byte max addressable bytes: $2^{32} = 4,294,967,296 = 4GB$

Kernel: mem residents part of os

0xC0000000 1100000000000000000000000000 ntwal velde spher



 $0x08048000\\0000\underline{1000}0000\underline{0100}1000\underline{0000}0000\underline{0000}\\0x00000000$

C's Abstract Memory Model

1. CODE Segment Contains: the program .text section instruction .rodata section litteral string, switch Jump Table. Lifetime: entire program's execution Initialization: From exe File. Binary Instruction and data Access: Read Only sharable 2. DATA Segment Contains: VAR: Variables that are either globals or locals Lifetime: entire program's execution Initialization: .data section from code // initialized to none zero value. I not moved .bss section of // uninitialized or initialized to zero. Access: Read / Write **3. HEAP** (AKA Free Store) Contains: mem, allocated & fresel ding execution. Lifetime: managed by programer Initialization: none by default. Access: read/write 4. STACK (AKA Auto Store) Contains: mem in stack frames that auto allocated and fixed during execution.

stack frame (AKA activation record)

contain: nok-static, local, parame, temporary rariables,... Lifetime: from declaraction to end of suspe, Initialization: more by default

Access: read / wite

Where do I live?

· Doita

BSS

```
it initialize with non zero
#include <stdio.h>
#include <stdlib.h>
int gus = 14; Data Section
int guy; . BSS section STACK
int madison(int pam) {
    static int max = 0; BSS section of Data
    int meg[] = \{22, 44, 88\};
    int *mel = &pam;
    max = gus - -;
    return max + meg[1] + *mel;
}
int *austin(int *pat){
    static int amy = 33; . Duta sether of Deta
    int *ari = malloc(sizeof(int),*44);
    gus--; . Data section of Put
    *ari = *pat;
    return ari;
}
int main(int argc, char *argv[]) {
    int vic[] = {33,66,99}; Stuk
    int *Wes = malloc(sizeof(int));
    *wes = 55;
                  Heap
    guy = 66;
    free (wes);
    wes = vic;
    wes[1] = madison(quy);
    wes = austin(&gus);
    free (wes);
    printf("Where do I live?");
   return 0;
}
* Arrays, structs, pointers can live in data, heap and/or stack
* Pointer variables can store any address but you will get
    A seg fautt on most platforms of your dereference an address that is outside of your mem segments.
```