# CS 354 - Machine Organization & Programming Thursday, September 12, 2019

Waitlisted? Complete the form at: https://forms.gle/CRvL1oR8i9Bymvyo6

Jim Skrentny, 5379 CS, skrentny@cs.wisc.edu

Course website: https://canvas.wisc.edu/courses/154937

Project p1 (3%): DUE at 10 pm on Monday, September 23rd

**Exam Conflicts:** report any by tomorrow using form at: https://forms.gle/6TwXssFmUCh7o8GS8

TA lab consulting & PM drop-in hours: are scheduled, see links on course front page

PM BYOL: Tonight 6 pm 1207 CS

Linux Workshop: Friday 5:30 pm 1240 CS

#### **Last Time**

C Program Structure C Logical Control Flow Recall Variables Meet Pointers

#### **Today**

Practice Pointers
Recall 1D Arrays
1D Arrays and Pointers
Passing Addresses

#### **Next Time**

Pointers Get More Interesting

#### Read:

K&R Ch. 7.8.5: Storage Management

K&R Ch. 5.5: Character Pointers and Functions

K&R Ch. 5.6: Pointer Arrays; Pointers to Pointers

**See:** Piazza post for web alternatives to K&R readings

## **Practice Pointers**

→ How many pointer variables are declared in the code above? ι.

To avoid, use: int \*p1, \*p2;

 $\rightarrow$  What is p1's value?

Danger: pi is considered uninstitutived It has whatever bit parties are in the mean allocated to 17.

 $\rightarrow$  Write the code to initialize p1 with address 0.

profer pl = NULL; or (Pl = 0; Pl = Ux0)

→ Given i below, write the code to make p1 point to i.

int i = 22;DI= & 1;

→ Write the code to display p1's pointee's value.

print [ ("% i m", \*p1) 5

 $\rightarrow$  What happens if the code above executes when p1 is NULL?

Donger: program crash due to segmentation fault.

→ Write the code to display p1's value?

printf (" %p \n", P1);

→ Is it useful to know a pointer's exact value?

usually not for basic nemory dragram.

→ What does the code below do?

int \*\*q = &p1;

## Recall 1D Arrays

What? An array variable is

Can change.

A consposed unit of storage having parts of elements whose values

Acless using identifier and indexing to get to a particular element.

Majorentod

· Allocated as a continions fixed size block of memory

Why?

· For storing a collection of desta. If the same type, with fast access to its ellements.

· Because it is much easier then declaring individual variables for each items in collection.

How?

→ How many integer elements have been allocated memory?

→ Where in memory was the allocation made? Twk

→ Write the code that gives the element at index 1 a value of 11.

→ Draw a basic memory diagram showing array a.

Danger: Elements are uninitialized by default.

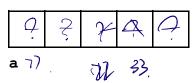
Is associated with the stanting address of the array. \* In C, the identifier for an array

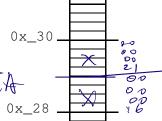
it is not a seperate ran.

# 1D Arrays and Pointers

#### Given:

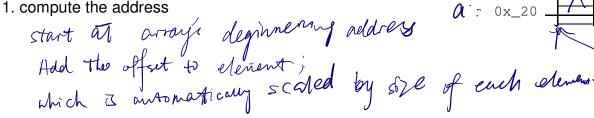
void someFunction(){ int a[5]; a[2] = 22;





**Address Arithmetic** 

 $*a[i] = *(\alpha + i)$ 



E fast.

PM

2. dereference the computed address to access the element

→ Write address arithmetic code to give the element at index 3 a value of 33.

 $\rightarrow$  Write address arithmetic code equivalent to a [0] = 77;

# **Using a Pointer**

→ Write the code to create a pointer b having the address of array a above.  $\hat{M} + P = A$ 

→ Write the code that uses p to give the element in a at index 4 a value of 44.

\* In C, pointers and arrays are closely related.

## Passing Addresses

### Recall Call Stack Tracing:

mannally excuting codes with funes in a that winic the waching.

call get a box (Stack Frame) its mem storing params, Local warmings

running and those belows are suspended waiting for their chance to return.

What is output by the code below?

```
void f(int pv1, int *pv2, int *pv3, int pv4[]) {
int lv = pv1 + *pv2 + *pv3 + pv4[0];
      = 11;
 - pv1
  *pv2 = 22;
  *pv3 = 33;
  pv4[0] = lv;
pv4[1] = 44;
int main(void) {
 \int int lv1 = 1, lv2 = 2;
  int *lv3;
 int 1v4[] = \{4,5,6\}; after
[1v3 = 1v4 + 2]
[ f(lv1, &lv2, lv3, lv4);
  printf("%i,%i,%i\n",lv1,lv2,*lv3);
  printf("%i,%i,%i\n",lv4[0],lv4[1],lv4[2]);
  return 0;
            out put: 1,22,33
13,44,35
```

DV3 main 204 lv3 lu2 LVI

. Callee

raller

Pass-by-Value

scalars: scalar parameter gets a copy of the argument

param gets a copy of its APPR ARG

Changing a parameter

does not change its argument. Just changes the collee's copy of the address.

\* BUT passing an address

regime the other to trust the adder.

now the called know where the ary is and can change it. CS 354 (F19): L3 - 5