Graph Applications

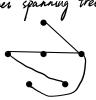
Spanning trees

, F = E.

Definition: Let G = (V, E) be a graph and T = (V, F) be a connected subgraph of G. If T is a tree, we call it a **spanning tree** of G. another spanning tree.

Example





Theorem: Every connected graph *G* has a spanning tree.

Proof (outline): STAT T= G

If I is a tree, we are done. otherwise, find a simple cycle in G

pick any edge in that cycle & remove it to obtain Ti.

Pick any edge in that cycle & remove it to obtain Ti.

Ti is still connected and Ti hers same vertices as G.

(So Ti is a subgraph of G) Sur T=T, & repeat the process

Removing edges will it? When I have 1814 edges
I will be a spanning tree

How can we find a spanning tree?

In zypook 1 Cs 400, depth-fint-traversal & breath-first-traversal.

Graph coloring

Definitions

coloring: an assignment of a color to each vertex in a graph so that no two adjacent vertices have the same color

k-coloring: a coloring that uses k colors

chromatic number of a graph G is the fewest number of colors needed for a coloring, i.e., the smallest k for which G has a k-coloring

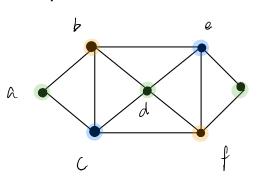
denoted X(b) (X is the greek chi)

• $\chi(G) = 1$

Theorem: A graph is > colorable if and only if it is bipartite.

Examples

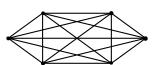
Example 1: What is the chromatic number of this graph?



XCG) 73 since a,b, c requires 3 colong

what if add an eelge $\{a, g\} \Rightarrow g$ require a new color.

Example 2 : K_6 , C_5 , C_8

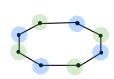


Kb = complete graph in b vertices

X(kb) = b since every vertex 13 connected to every other vertex.

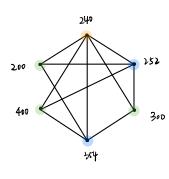


C5 = cycle with 5 vertices x(C5) = 3



C₈ = cycle with 8 vertices X(C₈) = 2

Example 3: Scheduling problem



How many different exam times one required so these are not any conflicts?

In general, colorings me not neversarily unique enth color = one exam time

50 XCG) 13 the fewers #10f exam time
required

know 1(G) 73,
XCG)=3