Virtual Memory: Systems

15-213/18-213/14-513/15-513: Introduction to Computer Systems 18th Lecture, March 26, 2019

s bits b bits

index offset

Address of word:

data begins at this offset

CT

Review of Symbols

Basic Parameters

- N = 2ⁿ: Number of addresses in virtual address space
- M = 2^m: Number of addresses in physical address space
- P = 2^p : Page size (bytes)

Components of the virtual address (VA)

TLBI: TLB index

TLBT: TLB tag

VPO: Virtual page offset

VPN: Virtual page number

TLBT — TLBI — TL

0 1 2

E = 2^e lines per set

S = 2^s sets

valid bit

Components of the *physical address* (PA)

PPO: Physical page offset (same as VPO)

PPN: Physical page number

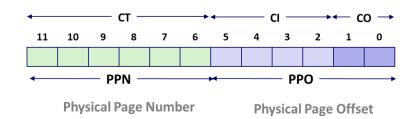
CO: Byte offset within cache line

CI: Cache index

CT: Cache tag

(bits per field for our simple example)

B = 2^b bytes per cache block (the data)



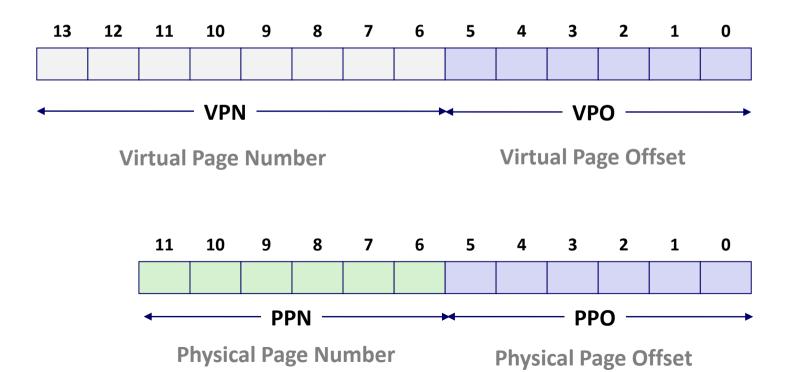
Today

- Simple memory system example
- Case study: Core i7/Linux memory system
- Memory mapping

Simple Memory System Example

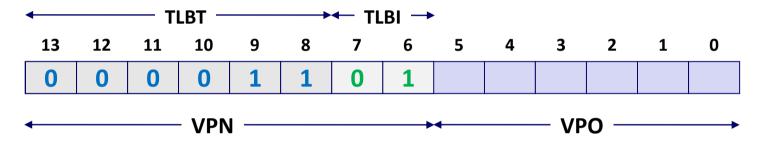
Addressing

- 14-bit virtual addresses
- 12-bit physical address
- Page size = 64 bytes



Simple Memory System TLB

- 16 entries
- 4-way associative



VPN = 0b1101 = 0x0D

Translation Lookaside Buffer (TLB)

Set	Tag	PPN	Valid									
0	03	-	0	09	0D	1	00	-	0	07	02	1
1	03	2D	1	02	_	0	04	_	0	0A	_	0
2	02	_	0	08	_	0	06	-	0	03	_	0
3	07	_	0	03	0D	1	0A	34	1	02	-	0

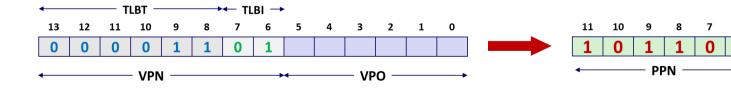
Simple Memory System Page Table

Only showing the first 16 entries (out of 256)

VPN	PPN	Valid
00	28	1
01	1	0
02	33	1
03	02	1
04	_	0
05	16	1
06	_	0
07	_	0

VPN	PPN	Valid
08	13	1
09	17	1
0A	09	1
ОВ	_	0
OC	-	0
0 D	2D	1
0E	11	1
OF	0D	1

 $0x0D \rightarrow 0x2D$



Simple Memory System Cache

- 16 lines, 4-byte cache line size
- Physically addressed

Direct mapped

	11	10	9	0	,	О	/ >	4	3	2	1	N.
1	1	0	1	1	0	1	1	0	1	0	0	1
	_		_	_		_						
			— РР	NI					— рр	<u> </u>		
			FF	IN		'			rr	U		

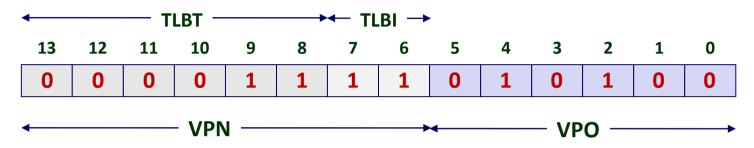
Idx	Tag	Valid	В0	B1	B2	В3
0	19	1	99	11	23	11
1	15	0	-	-	-	_
2	1B	1	00	02	04	08
3	36	0	_	_	_	_
4	32	1	43	6D	8F	09
5	0D	1	36	72	F0	1D
6	31	0	_	_	_	_
7	16	1	11	C2	DF	03

Idx	Tag	Valid	В0	B1	B2	В3
8	24	1	3A	00	51	89
9	2D	0	_	_	-	_
Α	2D	1	93	15	DA	3B
В	0B	0	-	_	_	_
С	12	0	-	-	-	_
D	16	1	04	96	34	15
Е	13	1	83	77	1B	D3
F	14	0	_	_	_	_

13

Address Translation Example

Virtual Address: 0x03D4



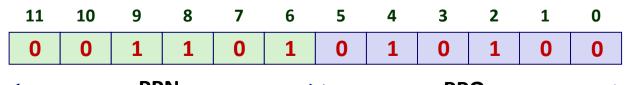
VPN **0x0F**

TLBI 0x3 TLBT 0x03 TLB Hit? Y Page Fault? N PPN: 0x0D

TLB

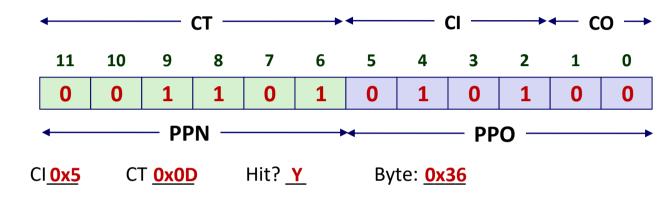
3	Set	Tag	PPN	Valid									
	0	03	_	0	09	0D	1	00	_	0	07	02	1
	1	03	2D	1	02	_	0	04	_	0	0A	_	0
	2	02	_	0	08	_	0	06	_	0	03	_	0
	3	07	_	0	03	0D	1	0A	34	1	02	_	0

Physical Address



Address Translation Example

Physical Address



Cache

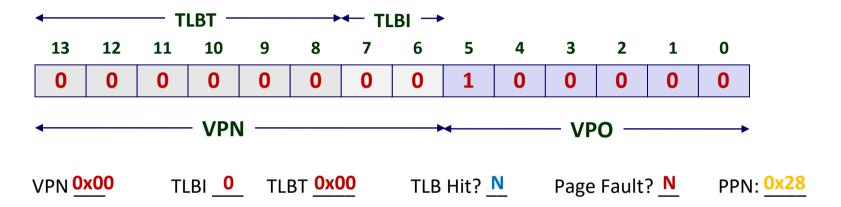
CO <u>0</u>

Idx	Tag	Valid	В0	B1	B2	В3
0	19	1	99	11	23	11
1	15	0	_	-	_	_
2	1B	1	00	02	04	08
3	36	0	_	_	_	_
4	32	1	43	6D	8F	09
5	0D	1	36	72	F0	1D
6	31	0	_	_	_	_
7	16	1	11	C2	DF	03

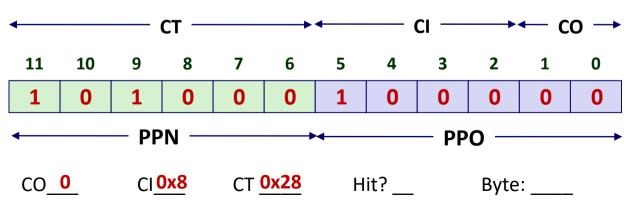
Idx	Tag	Valid	В0	B1	B2	В3
8	24	1	3A	00	51	89
9	2D	0	-	_	_	_
Α	2D	1	93	15	DA	3B
В	0B	0	_	_	_	_
С	12	0	_	_	_	_
D	16	1	04	96	34	15
E	13	1	83	77	1B	D3
F	14	0	_	_	_	_

Address Translation Example: TLB/Cache Miss

Virtual Address: 0x0020



Physical Address



Page table

rage	abic	
VPN	PPN	Valid
00	28	1
01	_	0
02	33	1
03	02	1
04	_	0
05	16	1
06	_	0
07	_	0

Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

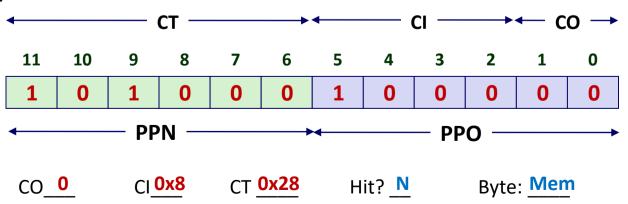
Address Translation Example: TLB/Cache Miss

Cache

ldx	Tag	Valid	В0	B1	B2	В3
0	19	1	99	11	23	11
1	15	0	_	-	_	_
2	1B	1	00	02	04	08
3	36	0	_	_	_	_
4	32	1	43	6D	8F	09
5	0D	1	36	72	F0	1D
6	31	0	_	_	_	_
7	16	1	11	C2	DF	03

Idx	Tag	Valid	В0	B1	B2	В3
8	24	1	3A	00	51	89
9	2D	0	-	_	_	_
Α	2D	1	93	15	DA	3B
В	0B	0	_	_	_	_
С	12	0	-	-	-	_
D	16	1	04	96	34	15
Е	13	1	83	77	1B	D3
F	14	0	_	_	_	_

Physical Address



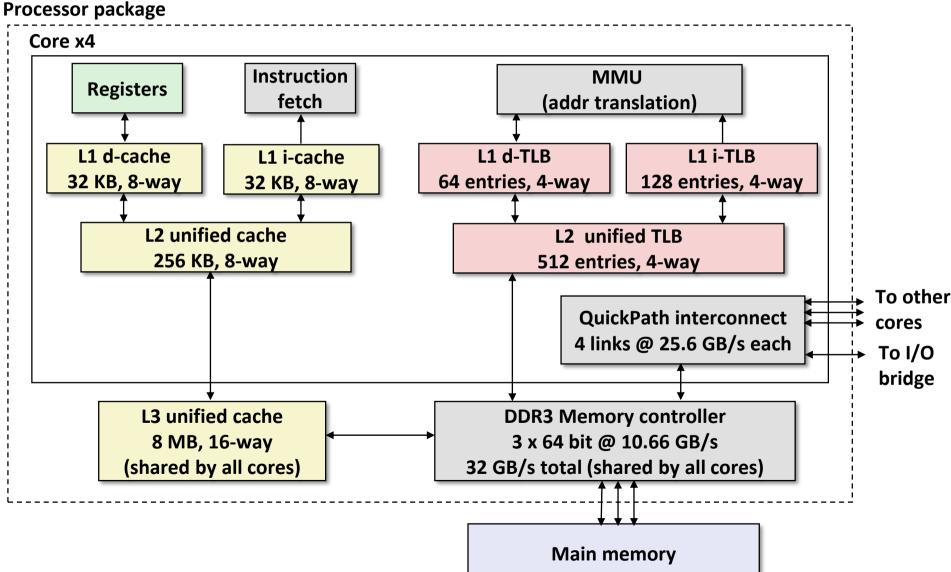
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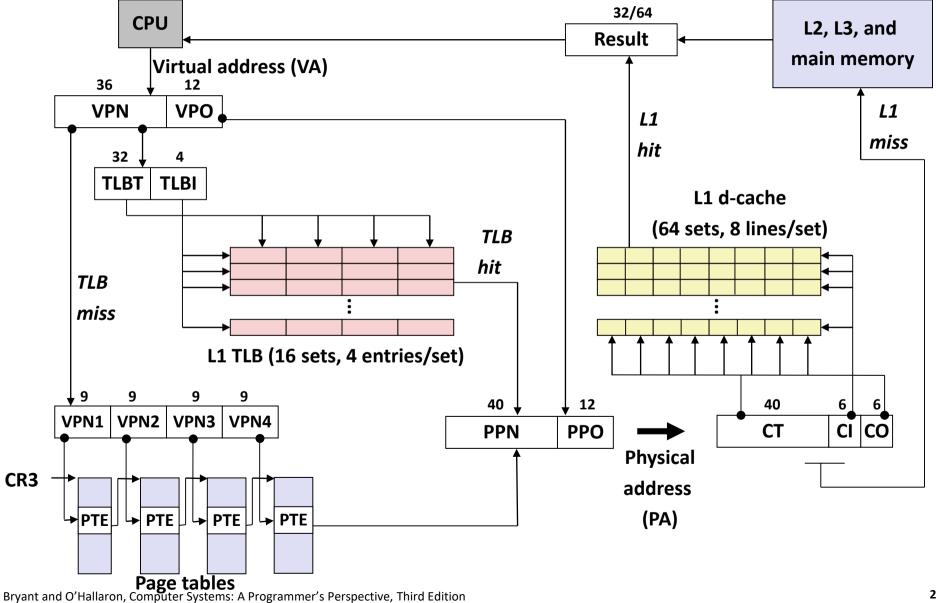
Intel Core i7 Memory System

Drocossor nackago

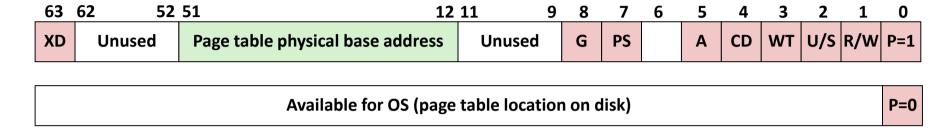
Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



End-to-end Core i7 Address Translation



Core i7 Level 1-3 Page Table Entries



Each entry references a 4K child page table. Significant fields:

P: Child page table present in physical memory (1) or not (0).

R/W: Read-only or read-write access access permission for all reachable pages.

U/S: user or supervisor (kernel) mode access permission for all reachable pages.

WT: Write-through or write-back cache policy for the child page table.

A: Reference bit (set by MMU on reads and writes, cleared by software).

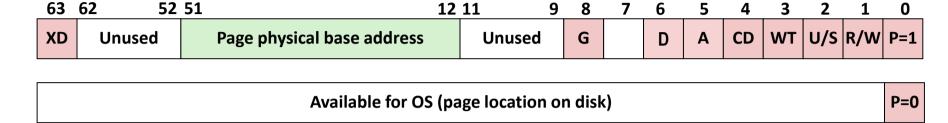
PS: Page size either 4 KB or 4 MB (defined for Level 1 PTEs only).

Page table physical base address: 40 most significant bits of physical page table address (forces page tables to be 4KB aligned)

XD: Disable or enable instruction fetches from all pages reachable from this PTE.

G: Process global – pages readable by all processes, subject to access restrictions

Core i7 Level 4 Page Table Entries



Each entry references a 4K child page. Significant fields:

P: Child page is present in memory (1) or not (0)

R/W: Read-only or read-write access permission for child page

U/S: User or supervisor mode access

WT: Write-through or write-back cache policy for this page

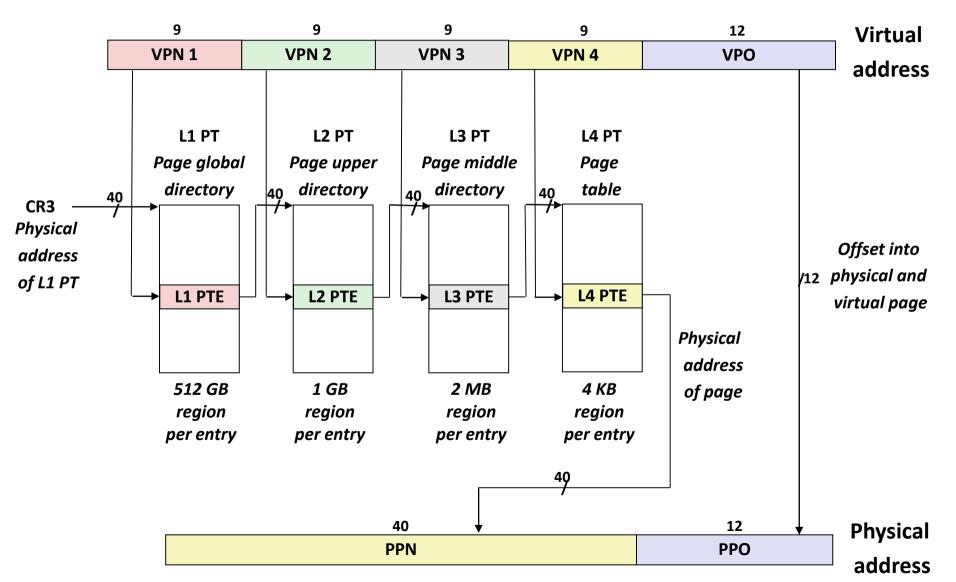
A: Reference bit (set by MMU on reads and writes, cleared by software)

D: Dirty bit (set by MMU on writes, cleared by software)

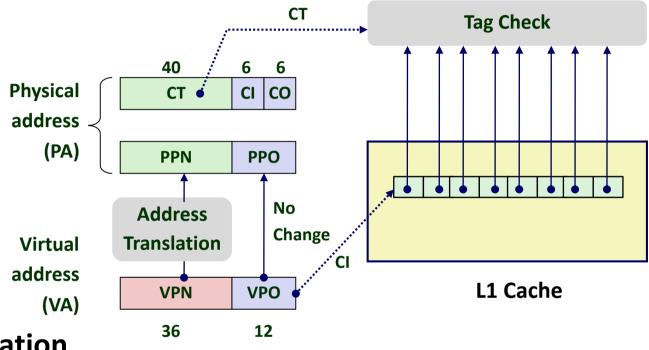
Page physical base address: 40 most significant bits of physical page address (40 out of 52 bits; 12 remaining bits address the page & force 4KB alignment)

XD: Disable or enable instruction fetches from this page.

Core i7 Page Table Translation



Cute Trick for Speeding Up L1 Access



Observation

- Bits that determine CI identical in virtual and physical address
- Can index into cache while address translation taking place
- Generally we hit in TLB, so PPN bits (CT bits) available next
- "Virtually indexed, physically tagged"
- Cache carefully sized to make this possible