* High Performance Computing *
Shiva Savan
BE B-20 * Unit Test 1 Question 2(a) * Non Uniform Memory Acress - It allows memory access to every processor without any restrictions. A block of is attached to the processor and all blocks memory can be accessed through the path provided by the use interconnected network. - A process has divided path to the black of memory Eq. If accessing Memory block Mem! from CPUI will be much faster than accessing block Mem 2 from CPUI. It has significant implication ie. If we map address carefully it may be possible to keep most of the information by a processor in the block attached to it required by a processor in the block attached to it.

)

- Therefore, the CPU can access that number diretly
and reducing the contention for the common box. Since the time to access a memory location depends on whether it is attached to invoking CPV or not.
Since the time to access a memory location depends
on whether it is attached to invoking CPU or not.
This model is called NumA.
Question 2 (b)
* Unilor Memory Access
- lu this parle surces arte parle sinch to anne
the main man of the
* Uniform Memory Access - In this, each processor gets equal priority to access the main memory of the microprospor is called as
in fact, memory arcess of each processor is identical
in UMA platform. as shown below:
and the contract of the contra
Mem Mem Mem
Interespection network Bus
Interepreted for network DUS
CPU CPU CPU
CPU CPU 1 CPU
Fig. UMA
- The programming is much casive easier in such platfam because of availability of global nemony space in the system.
because of availability of global nemony space in the
- The coding for read only just ction among the no. of programs' running in different processes processors is not at all seen by the programmers. - This happens because coding for such platforms is
programs' running in different processes processors is
not at all seen by the grag ranners.
- This happens because coding for sub plathers in
Jan Jan Jan Jan

Jimilar to coding usually done in social program for single microprocessor.

The interaction among the read/write operation are handled requires the use of mutual eachsian or some other took for synchorization. Quostion 3 (a) * Principle of Parallel Agt Algorithm Design:

- An algorithm provides step by step solution of the given problem. An algorithm accepts from the user and based upon input. after performing the defined computations, it provides the output.

- The basic steps in the design of the algorithm an:

Pat Partitioning of overall at computation into Smaller computation into different processors. Question 3 (b)

* Decomposition, tasks & Dependancy graphs:

• Decomposition:

- The overall computation can be partioned juto number of small size computation 30 that these computation execute in parallel.

- The decomposition deals with the approaches of partioning the overall computation into sub-problems.

When a computation is divided into many small tasks it is referred as fine-grained composition.

- In a program, the basic unit of computation is referred as a task & is controlled by OS.

- In contrast of the program, the task are unit of computation based upon that the overall computation is decomposed. · Park- Dependany Graph

- It is directed as acyclic graph, typically a graph is a collection of nodes & edges, the task-dependany graph also contains nodes & edges.

- The node in this graph is a task whereas edges.

between any two nodes represent dependancy between them. Eg. There is adje that exists between 2 notes TiaTz.

If T2 is to be executed, it must be after Ti. Question 5 (a) * One to All Broad cost Communication One to all broadcast is the operation in which a single process sends identical data to all other processes.

(ii) Parallel algorithms always needs this operation.

(iii) Consider that data of size M is to be sent to all the processes. The processes.

- Initially always source process has the date.

- After termination of algorithm there will be copy of rist initial data with each process.

- Processor copies of data will be generated where as processor P is the number of processor as almost in Hg.

	(and The
M one to All M	
One to All M One to All M All tone	(P-1)
All tone	
Fig. One to AU & All to one L	1 . 0
Fig. One to All to one book	Keduction,
Question 5(b)	
* All to All Broadcast & Reduction. O All to All broadcast is a generation of broadcast in which all P nodes simult broadcast.	
1) All to All proadcast in a growtier	1 - + 1
broadcast in which all P wide in	one to all
broadcast.	laneously iniale a
(i) A process sends M- word massage to ever	. 4
but different processor man homed at di	y other process
(in) It is used to in matrix me stime	planent message
metria multiplication & matrix - vactor	H. 12 L.V
in The dual of all to all modution.	Surpucation.
is the distinction of an all-to- one	Latia V mode
(ii) A process sends M-word message to ever but different processes may broadcast di (iii) It is used to in matrix operation matrix multiplication & matrix-vector n (iv) The dual of all to all reduction, in is the destination of an all-to-one w	a Crion.
Au to Au	1,7-1
	,
Broadcast M, M,	M
M. M. Mer, All to All Mo M.	N.
M. M. Mer, All to All Mo M. (1) (P-1) Reduction	(P-1)
- One approach to perform all to all broaders.	deast is performing
one to all broaders.	1 1
- This approach may take I times to comp	lete communication
- Commication like can be used more	efficiently by simular
- only performing all to P one to all brown	adcast. V
- This approach may take I times to comp - Communication link can be used more - cousty performing all to P one to all brown - By this there will be contention of all no the same path at the same time into a	ressages traversing
the same path at the same three into a	Single message
$\times - \times - \times$	CV V