



Team 7

Computer Theoretical Project





Algorithm Selected:

The chosen algorithm for analysis is the A* algorithm, implemented in the transportation optimization system to find the fastest emergency routes from any location to the nearest hospital. A* is a heuristic-based shortest path algorithm that balances efficiency and optimality, making it ideal for real-time navigation in dynamic environments like urban road networks.



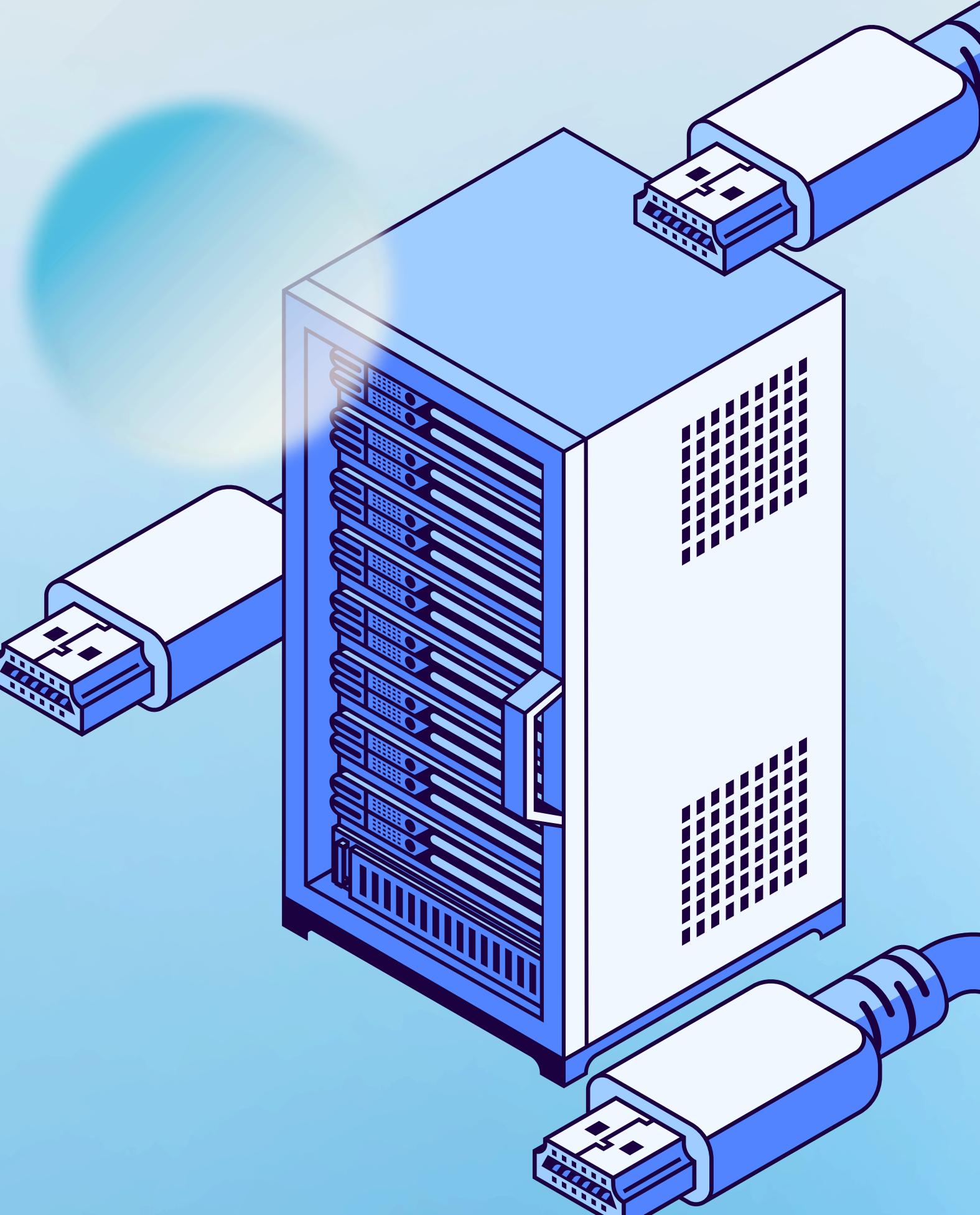
Introduction

Problem: Find fastest emergency routes in Cairo's traffic.

A*: Heuristic-based shortest path algorithm.

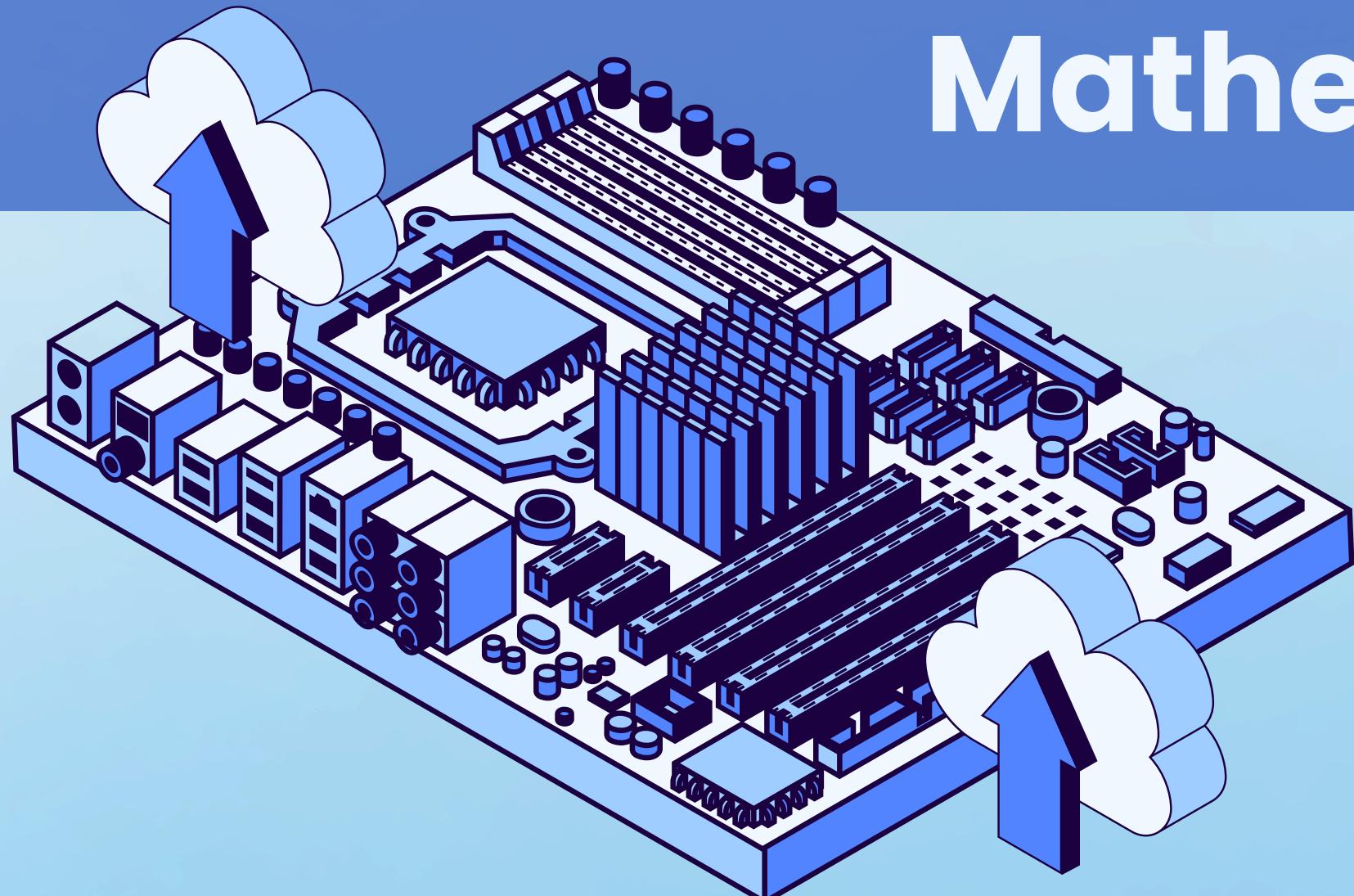
Why A*: Balances speed and optimality using heuristics.

Application: Finds fastest emergency routes to hospitals.





Mathematical Foundations :-



A* Formula: $f(n) = g(n) + h(n)$

$g(n)$: Actual cost from start to node n . $h(n)$: Heuristic estimate to goal (Euclidean distance in Cairo).

Heuristic: $h(n,t) = \sqrt{(X_n - X_t)^2 + (Y_n - Y_t)^2}$ (Euclidean).

Admissible and consistent \rightarrow optimal path.
Optimal Path :-

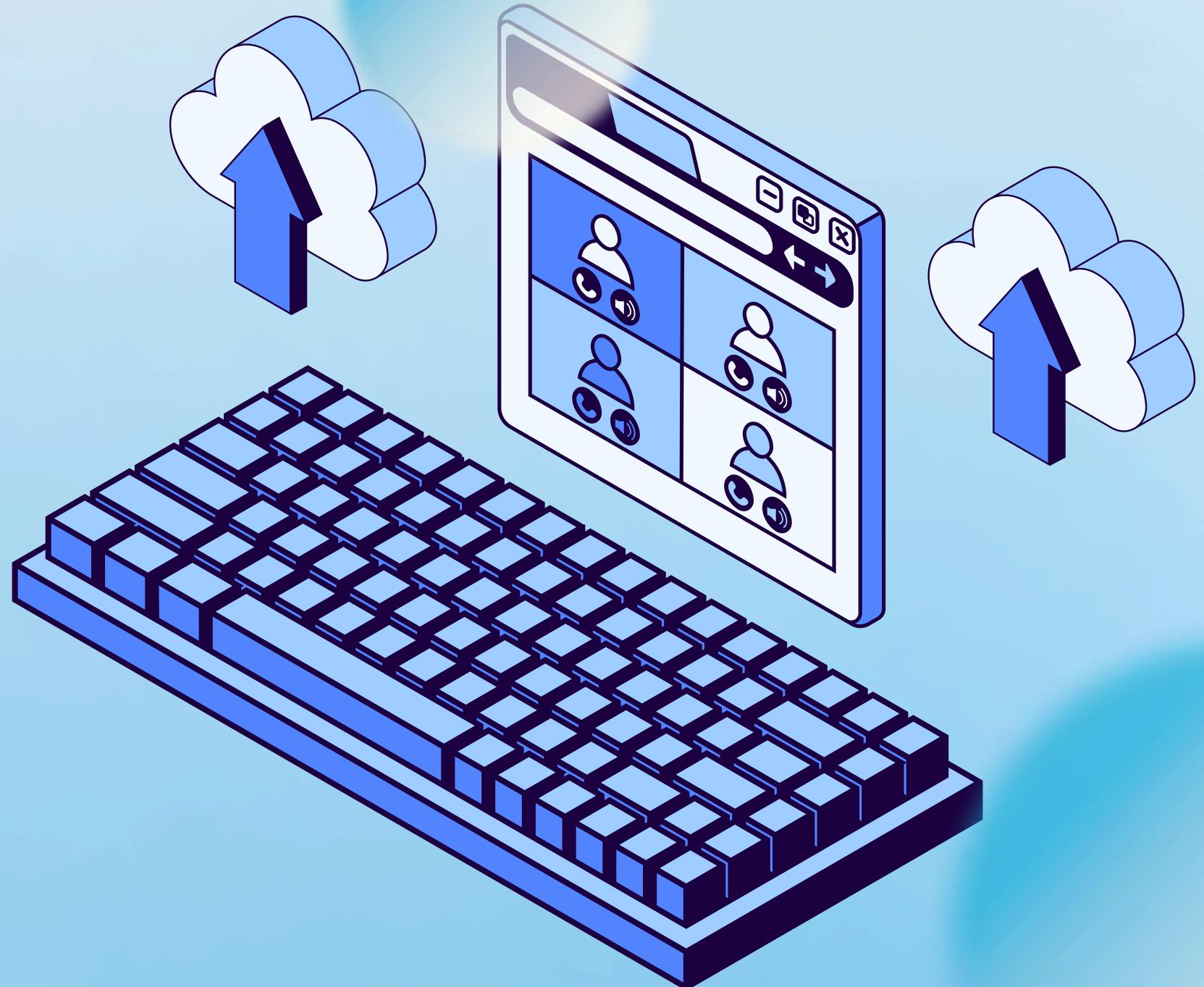
Admissible Heuristic : $h(n) \leq h^*(n)$
Euclidean distance \leq road distance.

Consistent Heuristic: $h(n) \leq c(n,n') + h(n')$
Satisfies triangle inequality.



Complexity Analysis :-

- n: number of nodes,
- m : number of edges. Time Complexity (Worst Case): $O(m + n \log n)$ (binary minheap, priority queue).
- Space Complexity: $O(n)$ (queue, dictionaries). for storing openSet, gScore, fScore, and cameFrom.



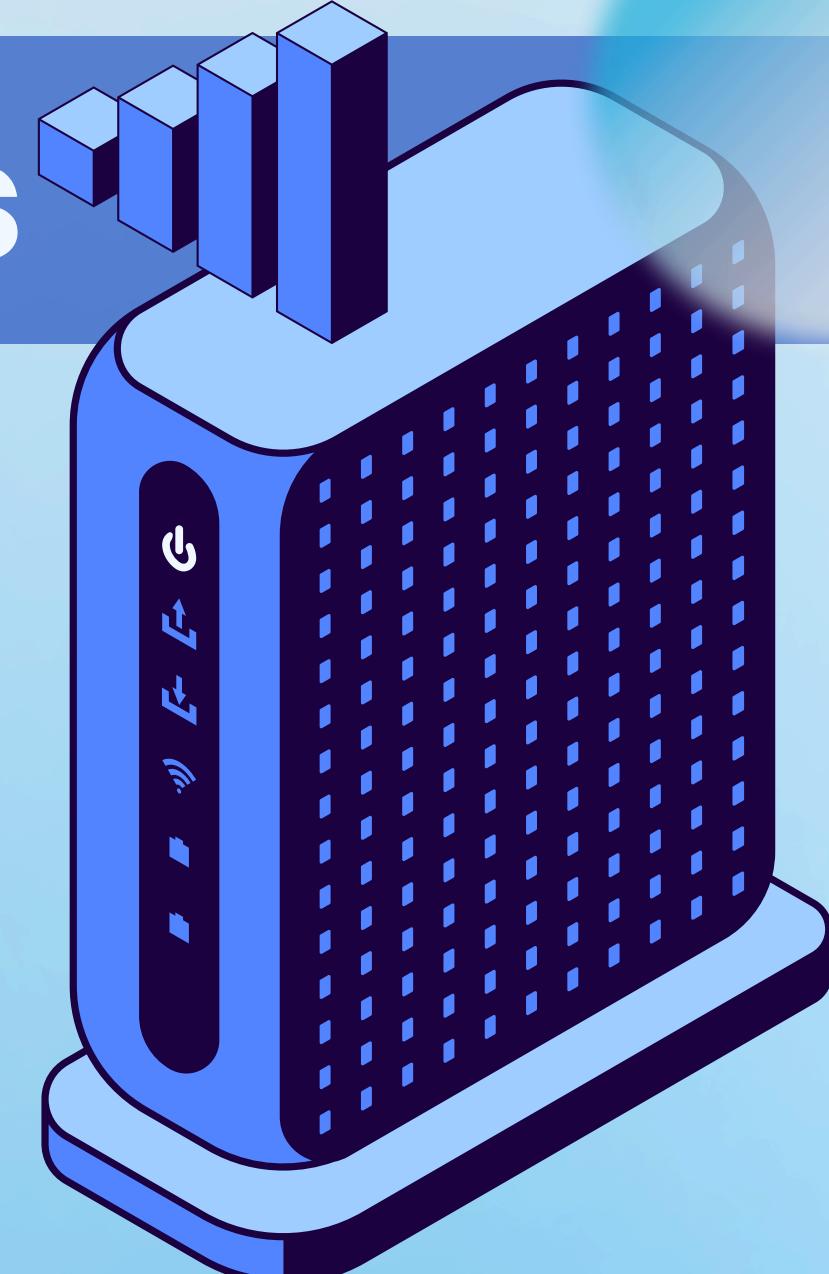


Implementation and Modifications

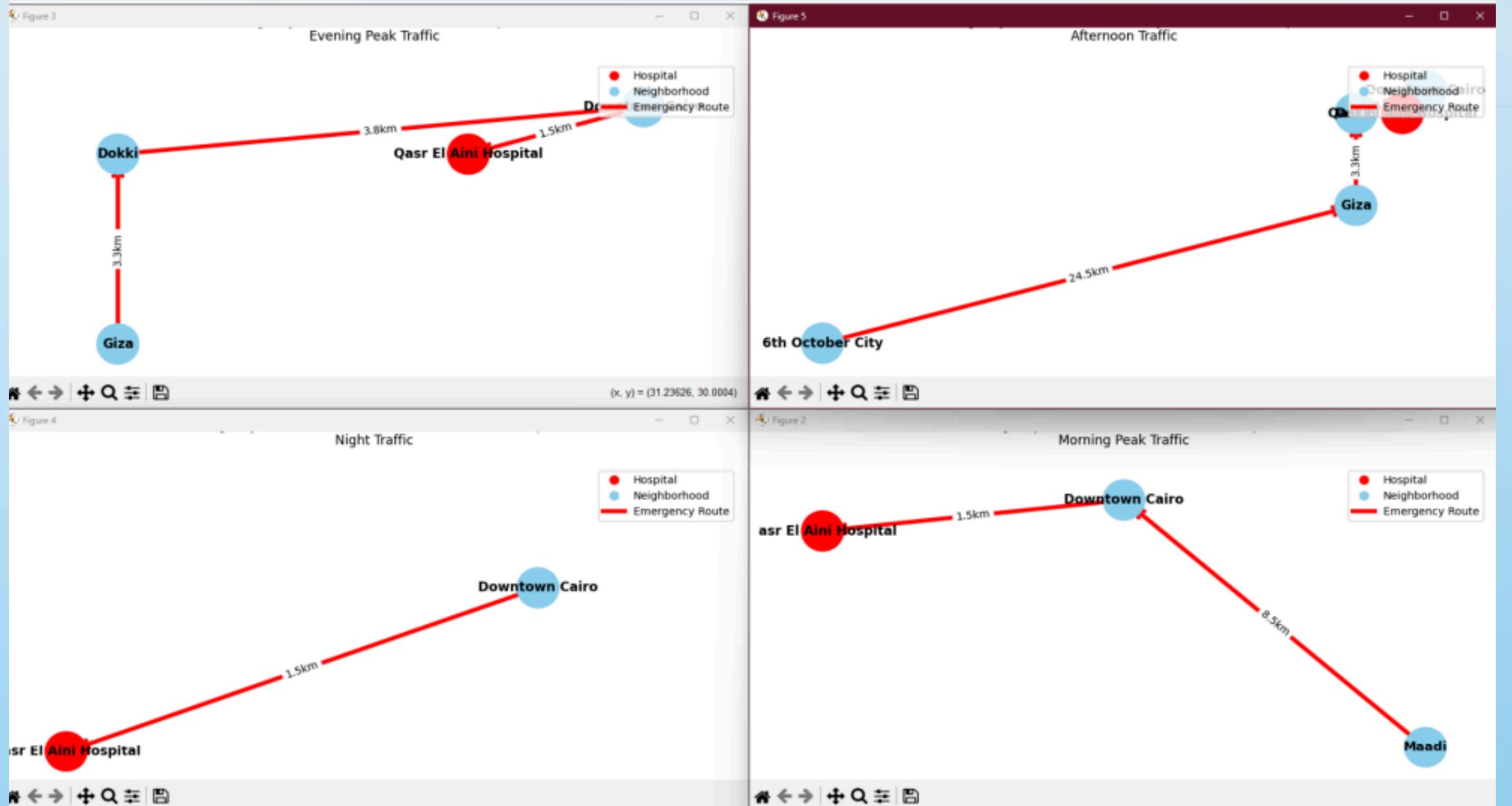
we modified A* as follows :-

- Edge Weights were made time-dependent :
- Heuristic was Euclidean distance to nearest hospital:
- Graph Preprocessing ensured all roads are bidirectional, and road conditions and traffic levels are normalized.
- Traffic-Aware Modifications:
- Implementation of A*:

```
• def get_traffic_factor(time_of_day, capacity, flow): utilization = flow / capacity return 1.0 + (0.8 if time_of_day == 'morning_peak' else 0.3) * utilization
• function A_star(start, goal): openSet := {start} cameFrom := empty map gScore[start] := e fScore[start] := h(start, goal) while openSet is not empty: current := node in openSet with lowest fScore if current == goal: return reconstruct_path(cameFrom, current) remove current from openSet for each neighbor of current: tentative_gScore := gScore[current] + d(current, neighbor) if tentative_gScore < gScore[neighbor]: cameFrom[neighbor] := current gScore[neighbor] := tentative_gScore fScore[neighbor] := gScore[neighbor] + h(neighbor, goal)if neighbor not in openSet: add neighbor to openSet
```

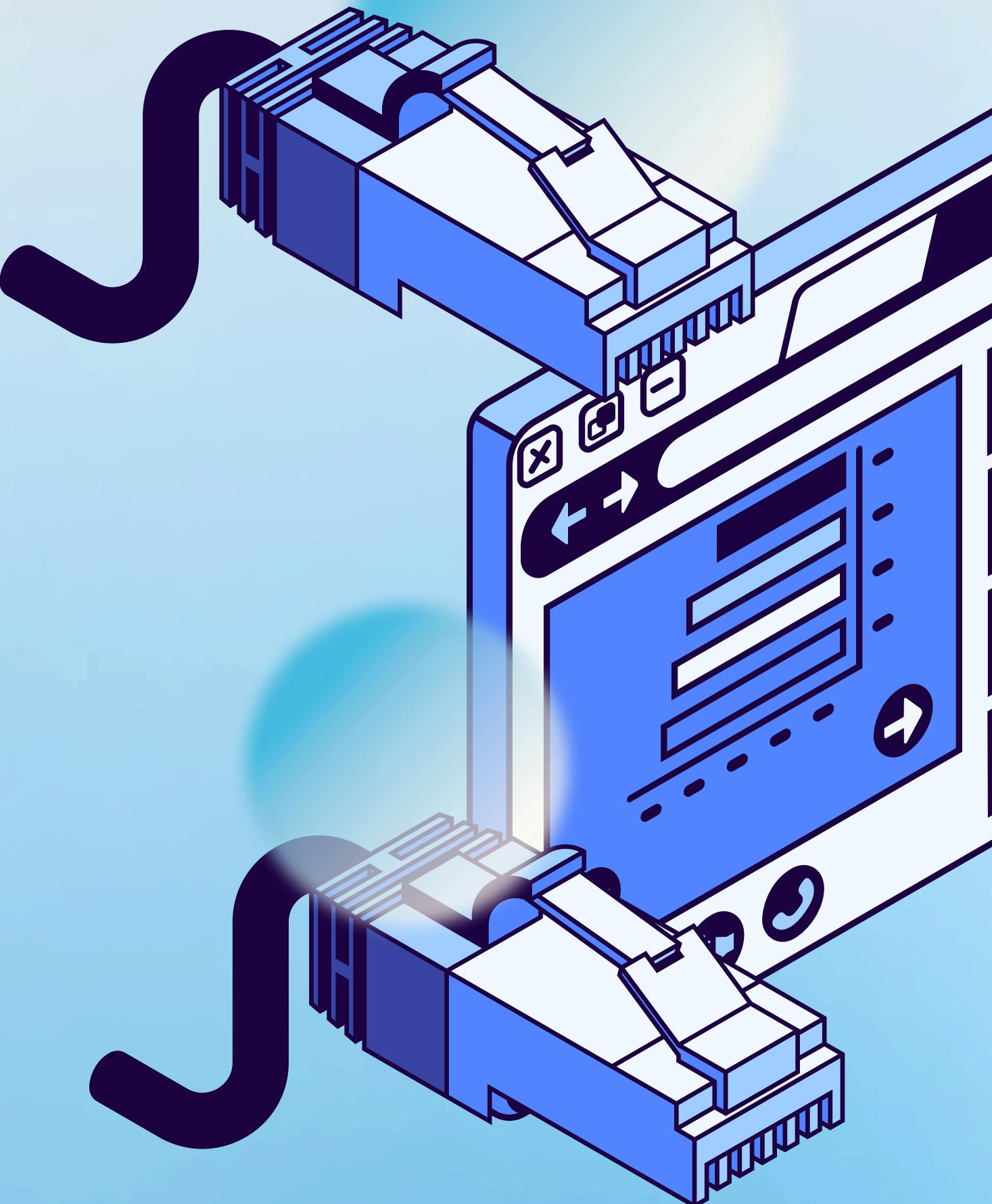


Mathematical Foundations :-



We evaluated A* in different real-time traffic scenarios using our Greater Cairo graph dataset. These Test Cast Traffic Scenarios :-





■ Morning Traffic (bottom right) :

The route from Maadi to Qasr El Aini Hospital was dynamically chosen via Downtown Cairo, selected the least-cost path (10 km total).

■ Afternoon Traffic (top right):

A longer emergency route of 24.5 km was computed from 6th October City to Giza, then to Qasr El Aini Hospital.

■ Evening Traffic (top left):

In a more central scenario, A* routed traffic from Giza to Dokkī (3.3 km) and then to the hospital (total 8.6 km).

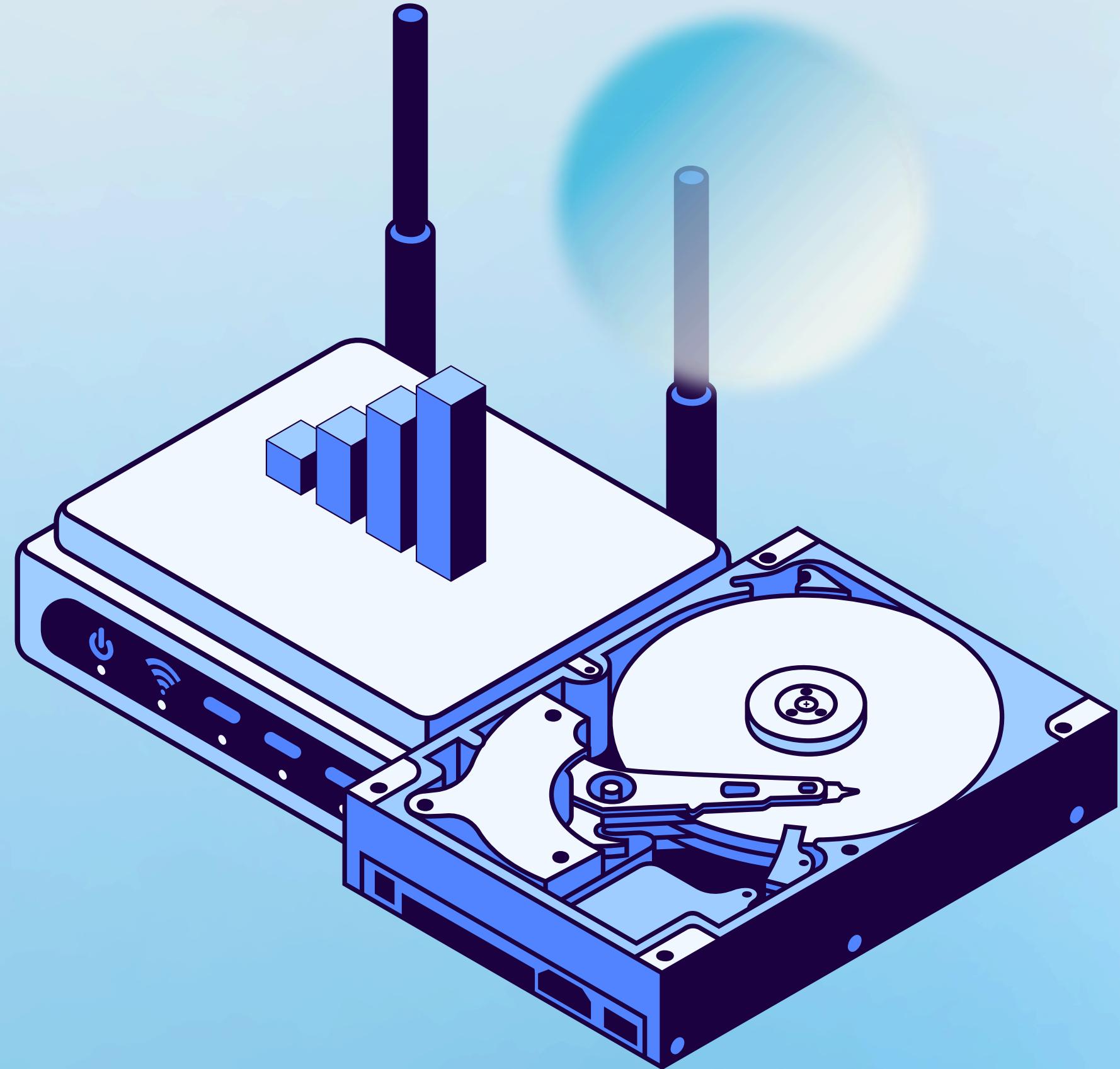
■ Night Traffic (bottom left):

A direct and minimal-cost path from Downtown Cairo to the hospital (1.5 km) reflects the effectiveness of A* under low-traffic scenarios.



These Scenarios cases validate that A* effectively balances accuracy, optimality, and responsiveness :-

- The heuristic function (Euclidean distance) guided the search effectively.
- The algorithm scaled efficiently with the size of the network.
- It supported dynamic edge weights, enabling real-time adjustments based on traffic density and road conditions.





Comparison with Alternatives

Dijkstra:

$O(m \log n)$, no heuristic, more nodes explored, it lacks heuristic guidance, it explores more nodes, increasing execution time in large city graphs.

Bellman-Ford:

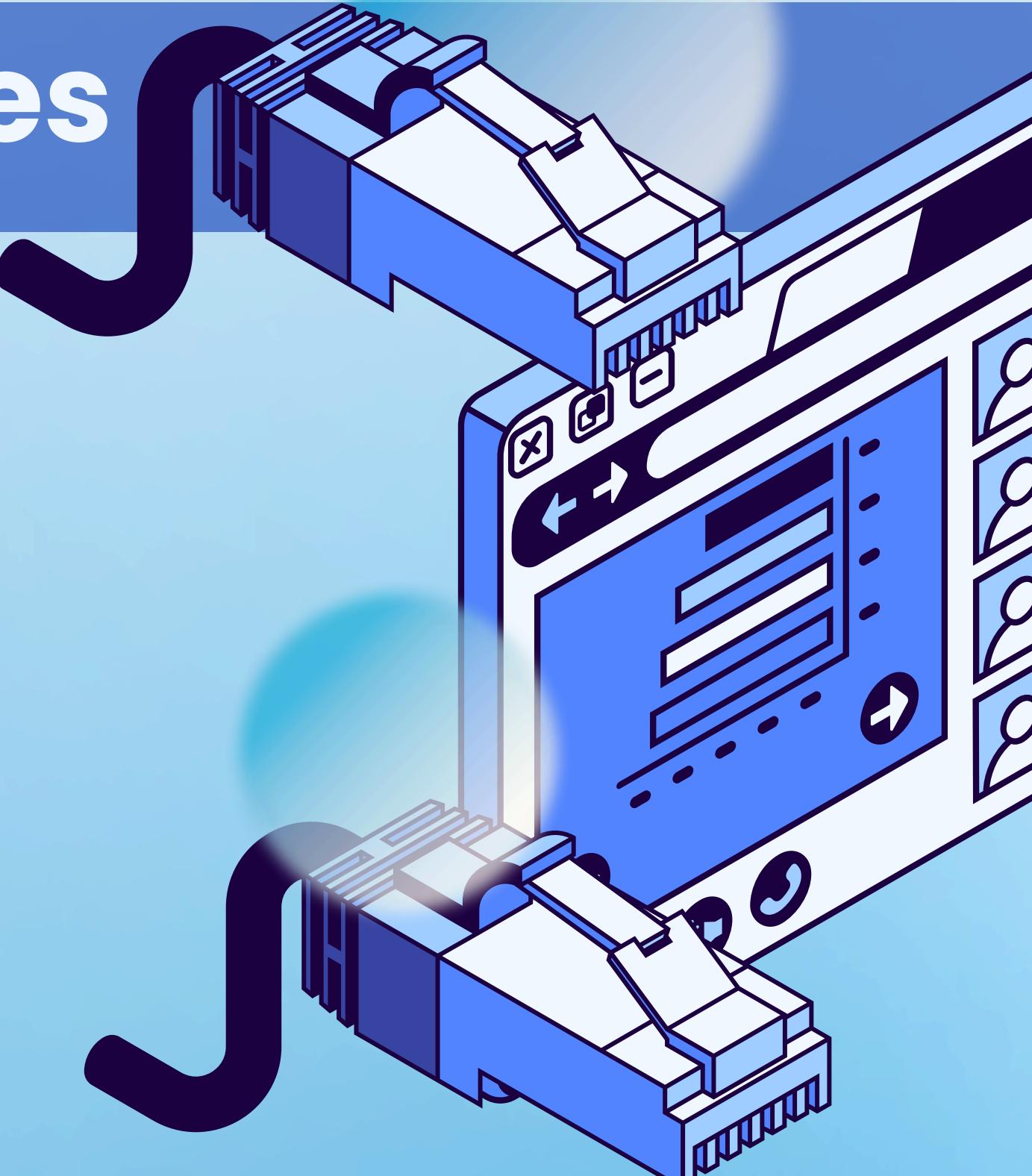
$O(n m)$, too slow, Handles negative weights and Impractical for large graphs and unnecessary here, as edge weights (distances) are positive.

Breadth-First Search:

$O(n + m)$, unweighted graphs only, not suitable for real-world traffic systems where edge weights matter, It cannot adapt to variable traffic conditions.

A*:

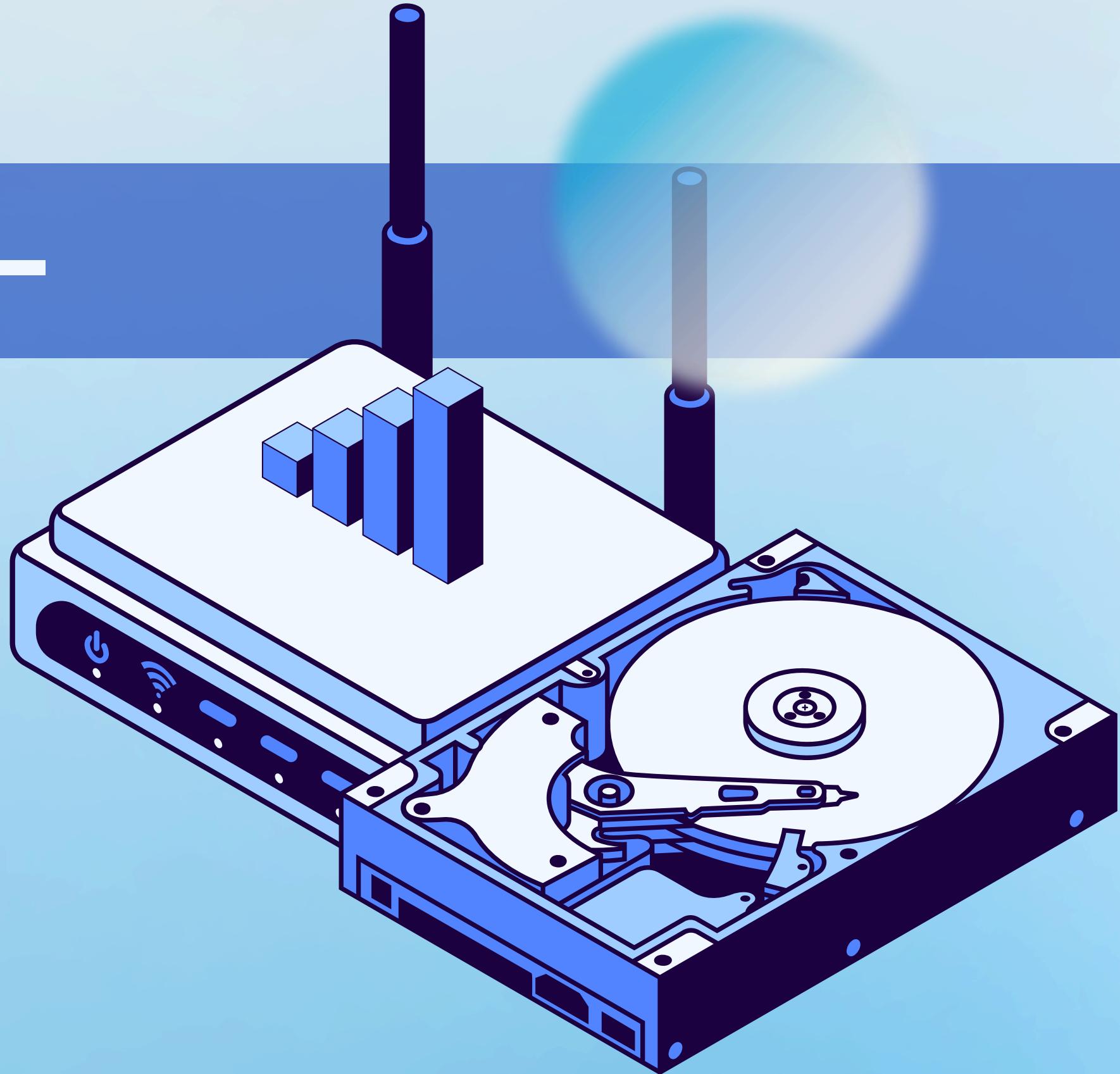
$O(m + n \log n)$, optimal for weighted, goal-directed search, balance of optimality and efficiency, leveraging the Euclidean heuristic to reduce the search space.

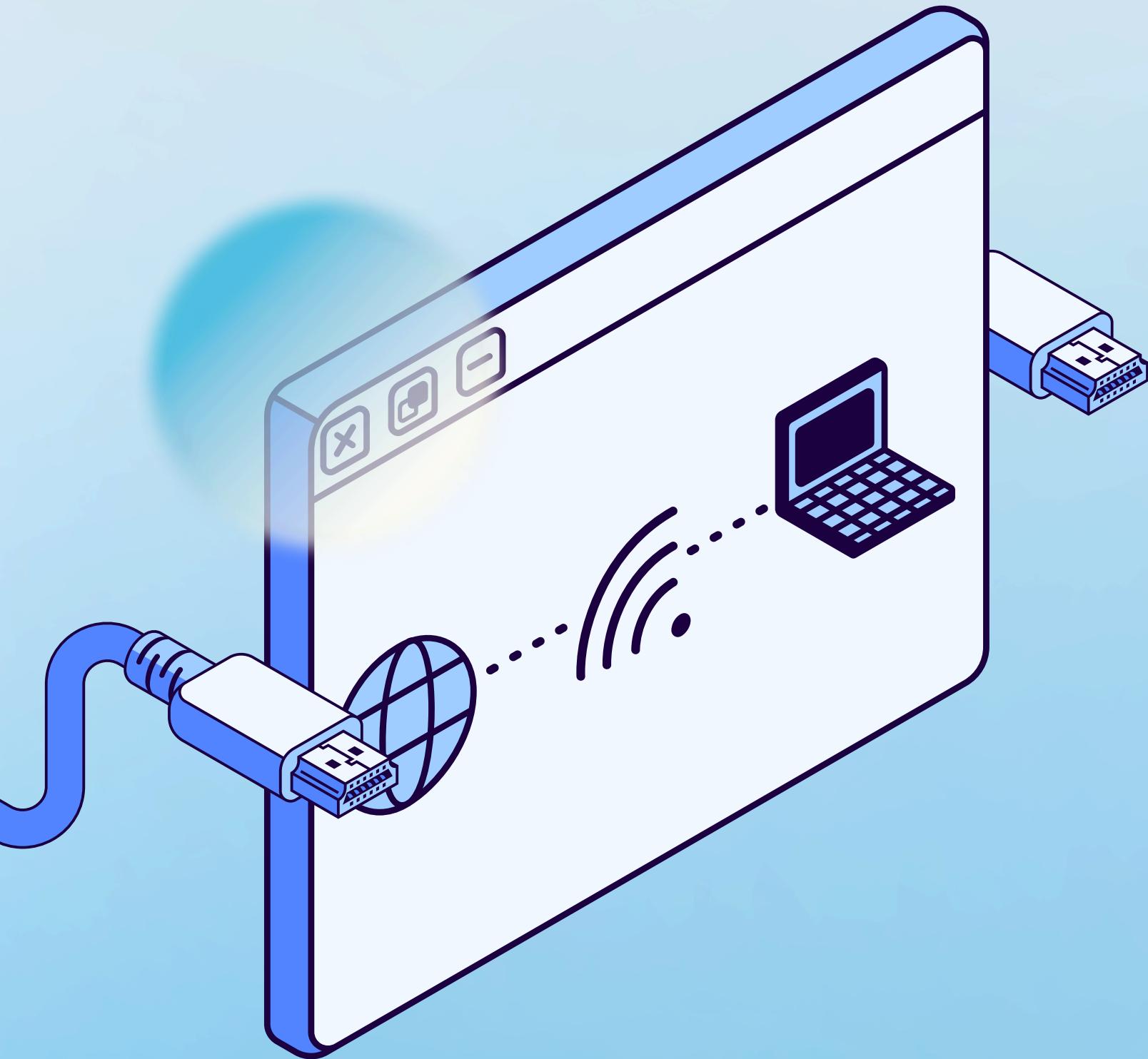




Lessons Learned :-

- Given the need for fast, accurate, and adaptable routing in emergency situations within smart cities, A* offers the best balance. Its ability to incorporate traffic data through cost functions and heuristics makes it superior for real-time urban navigation, particularly in dense areas like Greater Cairo.





Conclusion :-

- A* outperforms alternatives like Dijkstra's and BFS in emergency routing due to its goal-directed approach and ability to integrate real-time traffic data through heuristics. This makes it faster and more suitable for dynamic, realworld urban environments like Greater Cairo.



Thank You!

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