

Digital Logic Design

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Today's Topic

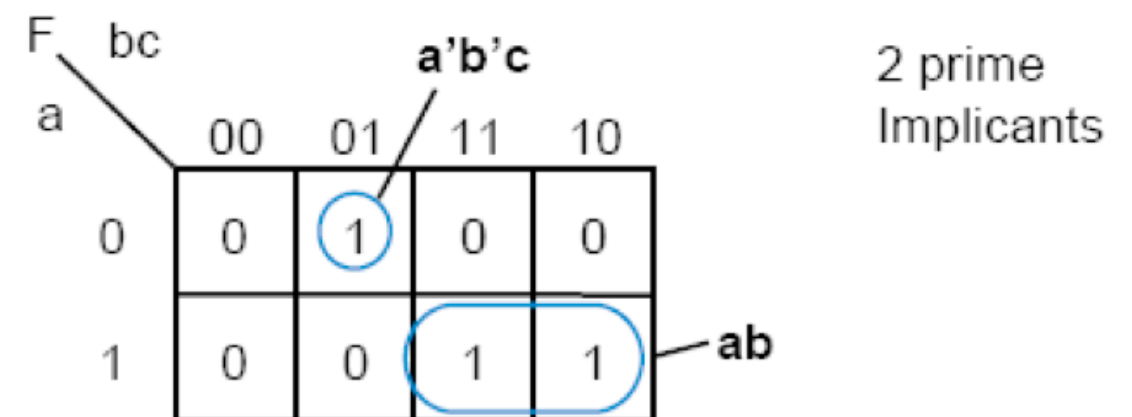
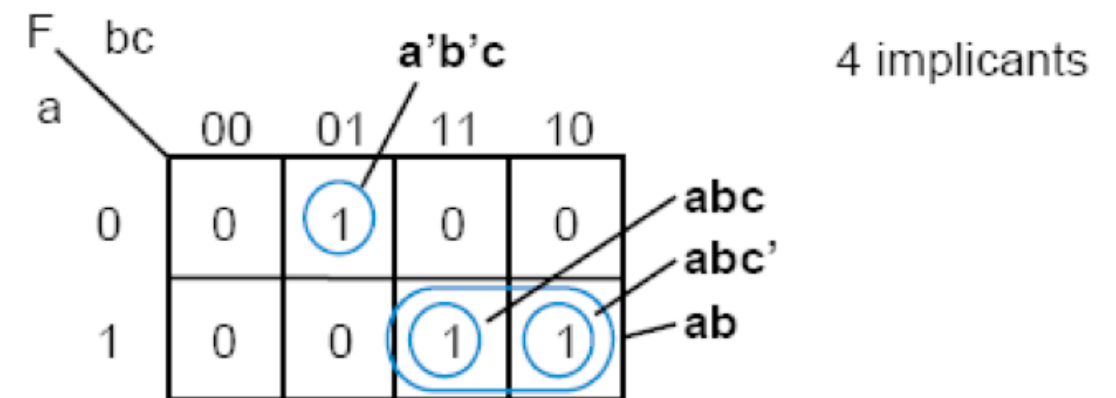
카르노맵 최적화를 체계적으로 수행하기

Terminology

- Literal
 - 변수가 나타나는 형태
- Implicant
 - 리터럴의 곱의 형태
 - 카르노맵에서는 써클로 표현
- Prime implicant
 - 더 이상 확장할 수 없는 Implicant
 - 더 이상 확장할 수 없는 써클

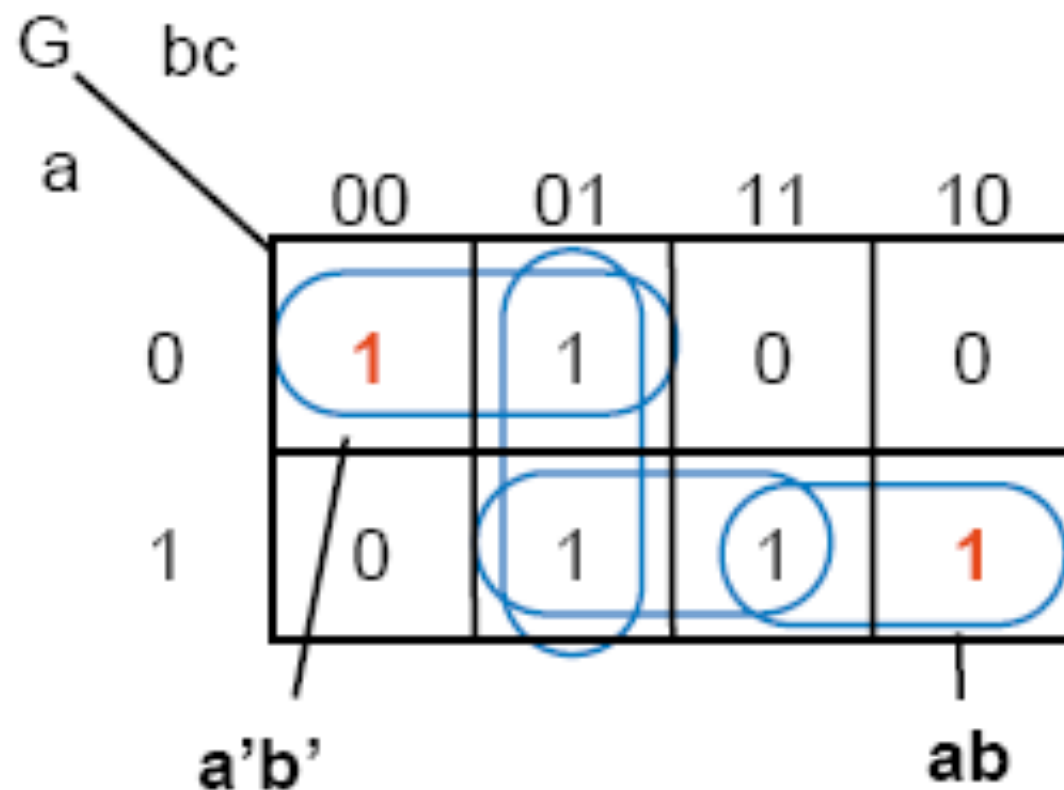
$$F = a'b'c + abc + abc'$$

6 Literals
(a, a', b, b', c, c')



Terminology

- Essential prime implicant
 - 해당 함수를 만족시키기 위해 반드시 필요한 Prime Implicant
 - 모든 EPI (Essential prime implicant)는 꼭 포함해야 함
 - Non-EPI는 포함될 수도 있고 포함되지 않을 수도 있음

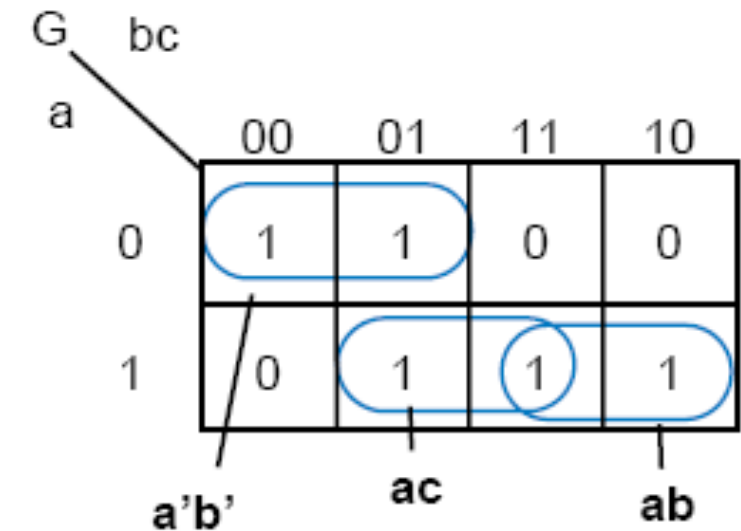


4 prime
implicants

2 essential
prime
implicants

Terminology

- Cover
 - $F = 1$ 을 만족시키는 모든 써클의 집합
 - 해당 함수를 구현할 수 있는 방법
- Cost
 - Number of gates + number of inputs
 - Ignore complement (NOT)

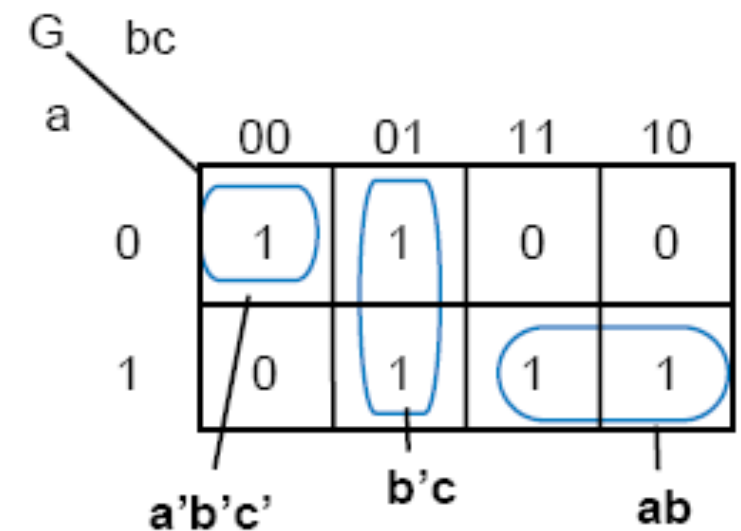


Cover 1: $G = a'b' + ac + ab$

$$G = a'b'c' + b'c + ab$$

1 3-input AND gates
2 2-input AND gates
1 3-input OR gates

$$\text{cost} = 4 + 10 = 14$$

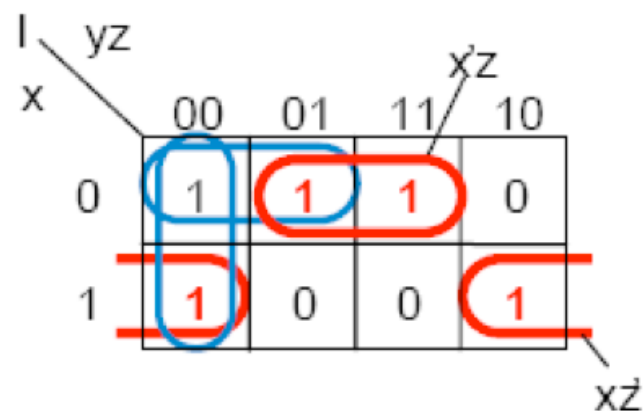
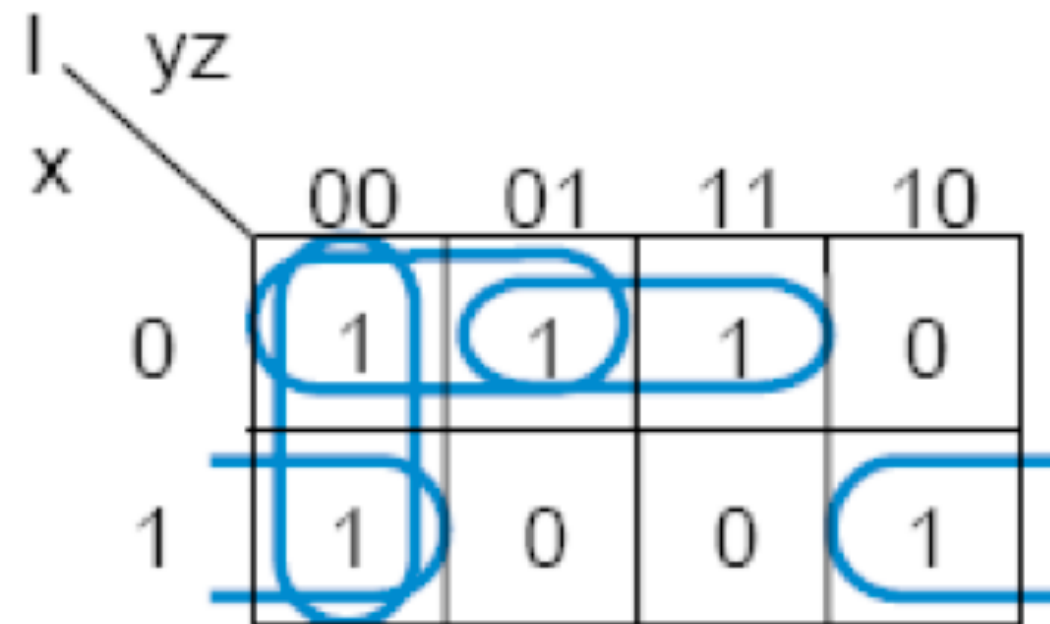


Cover 2: $G = a'b'c' + b'c + ab$

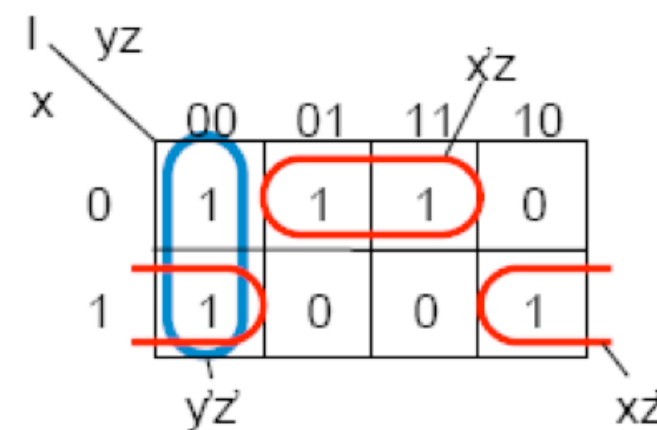
K-Map 최적화 절차

- (1) 모든 prime implicant를 생성
- (2) EPI를 모두 찾는다
- (3) EPI만으로 $F=1$ 을 만족?
 - Yes - were done!
 - No – Non-EPI 중에서 선택하여 Complete Cover를 만들 수 있는 PI들을 찾아냄

K-Map Optimization Example (1)

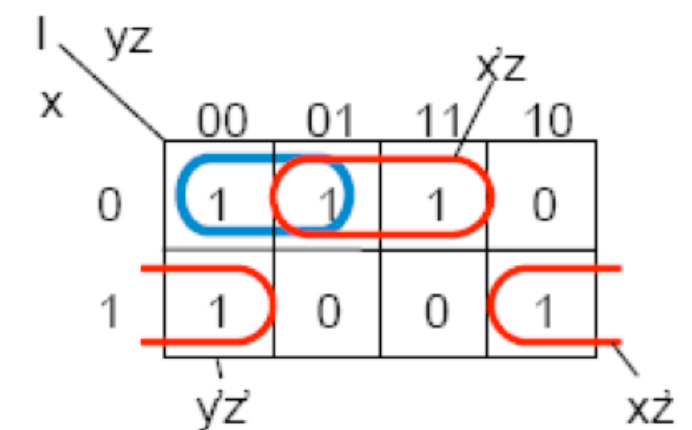


Cover:
e.p.i : $x'z, xz'$



Cover 1:
e.p.i : $x'z, xz', y'z'$

cost = 13



Cover 2:
e.p.i : $x'z, xz', x'y'$

cost = 13

K-Map Optimization Example (2)

ab \ cd	00	01	11	10
00	1	1	1	1
01	0	0	1	0
11	0	0	1	1
10	0	0	0	1

ab \ cd	00	01	11	10
00	1	1	1	1
01	0	0	1	0
11	0	0	1	1
10	0	0	0	1

Cover:
e.p.i : $a'b'$

K-Map Optimization Example (2)

ab \ cd				
	00	01	11	10
00	1	1	1	1
01	0	0	1	0
11	0	0	1	1
10	0	0	0	1

ab \ cd				
	00	01	11	10
00	1	1	1	1
01	0	0	1	0
11	0	0	1	1
10	0	0	0	1

Cover:
e.p.i : $a'b'$

K-Map Optimization Example (2)

		cd			
ab		00	01	11	10
00		1	1	1	1
01		0	0	1	0
11		0	0	1	1
10		0	0	0	1

Cover:
e.p.i : $a'b'$

		cd			
ab		00	01	11	10
00		1	1	1	1
01		0	0	1	0
11		0	0	1	1
10		0	0	0	1

Cover 1 (cost = 20)
 $F = a'b' + a'cd + abc + b'cd'$

		cd			
ab		00	01	11	10
00		1	1	1	1
01		0	0	1	0
11		0	0	1	1
10		0	0	0	1

Cover 2 (cost = 15)
 $F = a'b' + bcd + acd'$

K-Map for SOP vs. POS

- K-maps 표현 방법은 sum-of-products과 직접적으로 연관
- Product-of-sums 표현과 K-map의 관계는?

ab \ cd	cd			
	00	01	11	10
00	1	1	1	1
01	0	0	0	0
11	0	0	1	1
10	0	0	0	0

Sum-of-products
 $F = a'b' + abc$

ab \ cd	cd			
	00	01	11	10
00	1	1	1	1
01	0	0	0	0
11	0	0	1	1
10	0	0	0	0

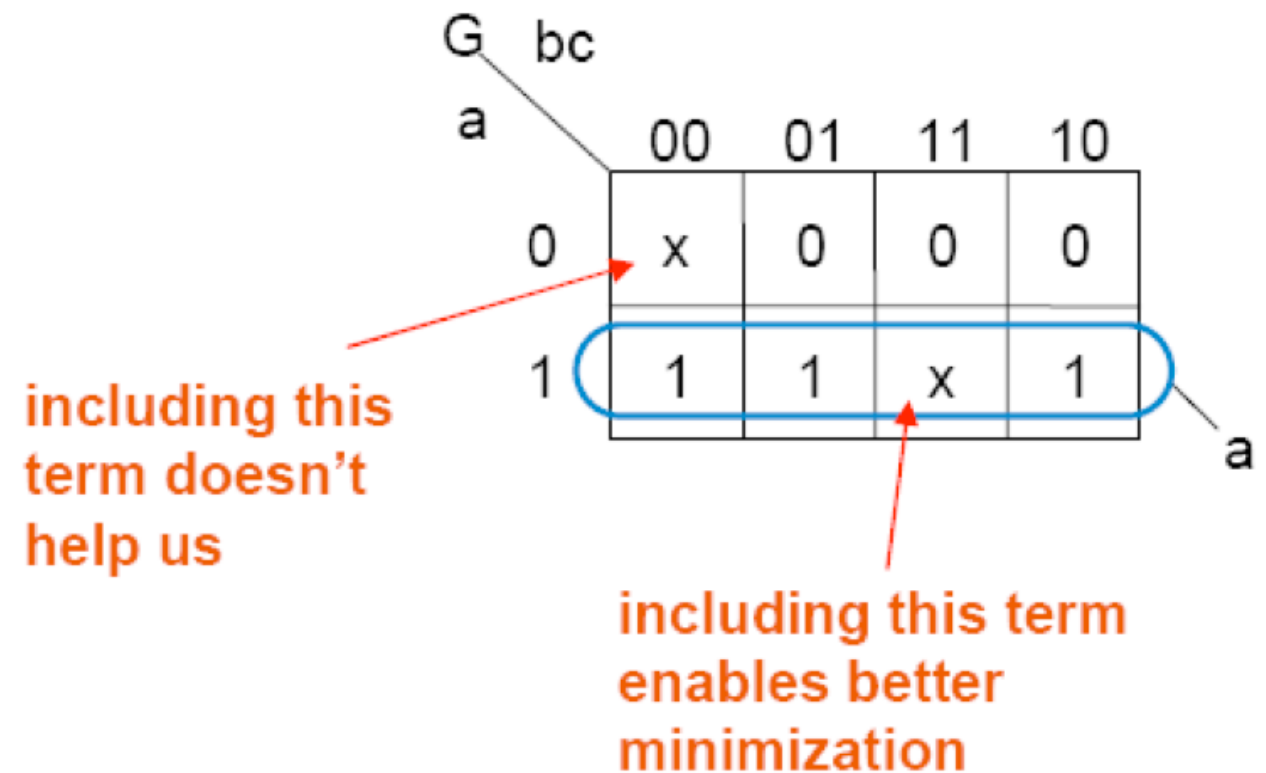
Product-of-sums
 $F = (a+b')(b'+c)(a'+b)$

Don't Care

- Don't Care 입력
 - 해당 입력에 대한 출력은 시스템 구현에 전혀 영향이 없음
 - i.e. input condition can never occur
 - K-map에서 'X'로 표시
 - 최적화에 활용

$abc = 000$ and $abc = 111$ are unused inputs

a	b	c	Z
0	0	0	x
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	x



Don't Care Input 활용하기

- Don't care 출력은 최적화에 유리한 방향으로 '0' 또는 '1'로 가정한다.
 - X 를 포함하는 경우는 최적화에 꼭 필요한 경우에만
 - 다른 X는 절대 포함하지 않는다

	yz	00	01	11	10
x					
0		X	0	0	0
1		1	X	0	0

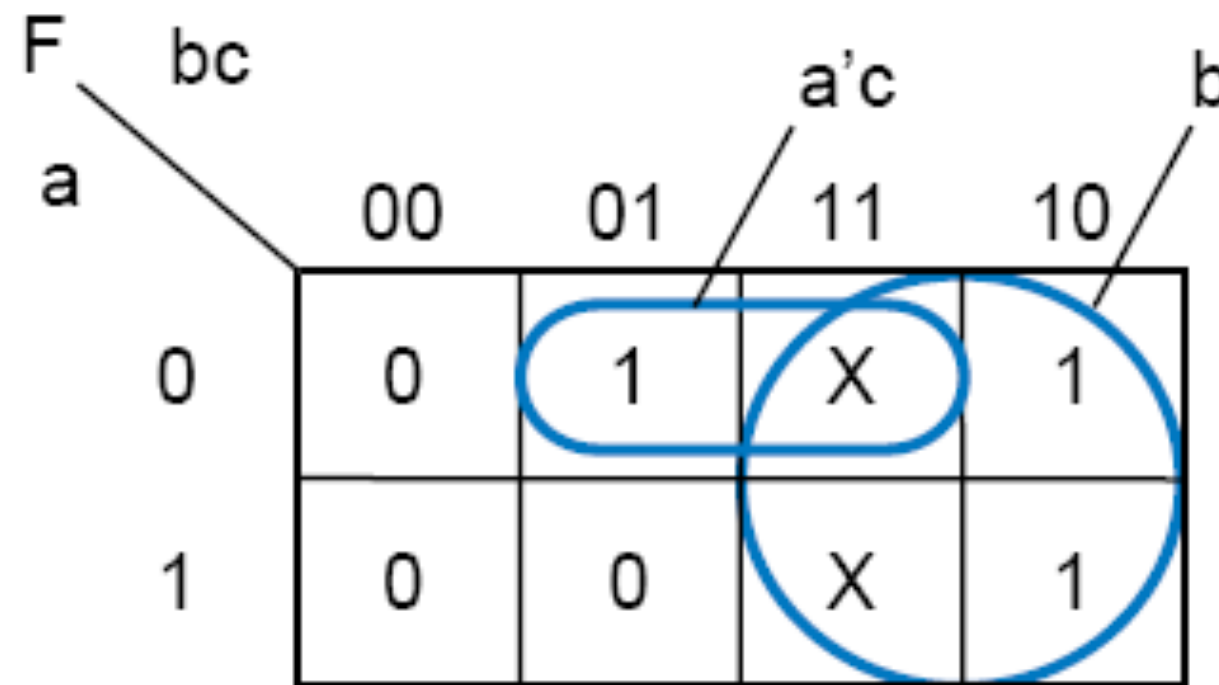
Good use of don't cares

	yz	00	01	11	10
x					
0		X	0	0	0
1		1	X	0	0

Unnecessary use of don't cares; results in extra term

Example – Using Don't Care

- Minimize:
 - $F = a'bc' + abc' + a'b'c$
 - Given don't cares: $a'bc$, abc



$$F = a'c + b$$

Using Don't Cares

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>Y</i>
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X

		<i>Y</i>			
<i>CD</i>	<i>AB</i>	00	01	11	10
00					
01					
11					
10					

Using Don't Cares

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>Y</i>
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X

<i>Y</i>		<i>AB</i>			
<i>CD</i>		00	01	11	10
00		1	0	X	1
01		0	X	X	1
11		1	1	X	X
10		1	1	X	X

Using Don't Cares

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>Y</i>
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	X
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X

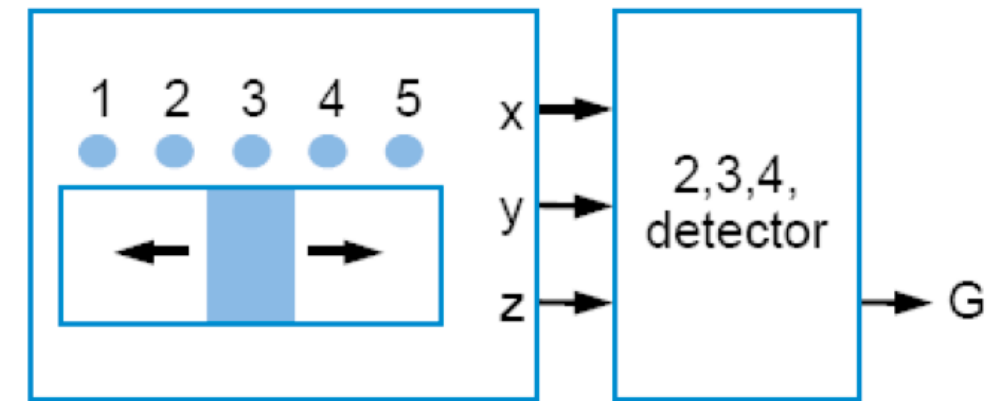
Y		AB			
CD		00	01	11	10
00		1	0	X	1
01		0	X	X	1
11		1	1	X	X
10		1	1	X	X

$$Y = A + \overline{B}\overline{D} + C$$

Example – Using Don't Care (2)

- Switch with 5 positions

- position 1, $xyz = 001$
- position 2, $xyz = 010$
- position 3, $xyz = 011$
- position 4, $xyz = 100$
- position 5, $xyz = 101$



- Want circuit that

- Outputs 1 when switch is in position 2, 3, or 4
- Outputs 0 when switch is in position 1 or 5
- Note that we never use some input combinations
 - 000, 110, 111

$$F = x'y + xy'z'$$

		yz			
x	G	00	01	11	10
0		0	0	1	1
1		1	0	0	0

Annotations: A blue oval groups the 1s in the x=0 row (yz=11 and 10), labeled $x'y$. Another blue oval groups the 1s in the x=1 row (yz=00 and 10), labeled $xy'z'$.

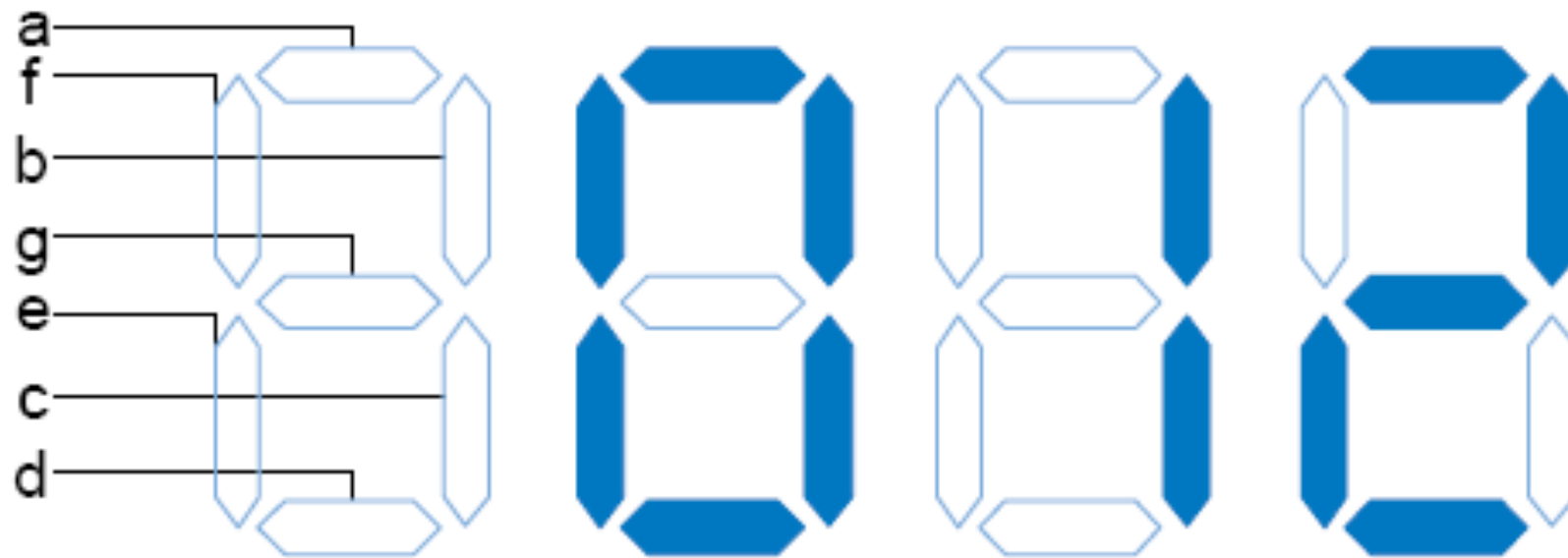
With don't cares:
 $F = y + z'$

		yz			
x	G	00	01	11	10
0		X	0	1	1
1		1	0	X	X

Annotations: Blue circles group the 1s in the y=1 column (x=0 and x=1), labeled y. Another blue circle groups the 1s in the z=0 column (x=0 and x=1), labeled z'. The cells with 'X' represent don't care conditions.

Multiple Output Circuits

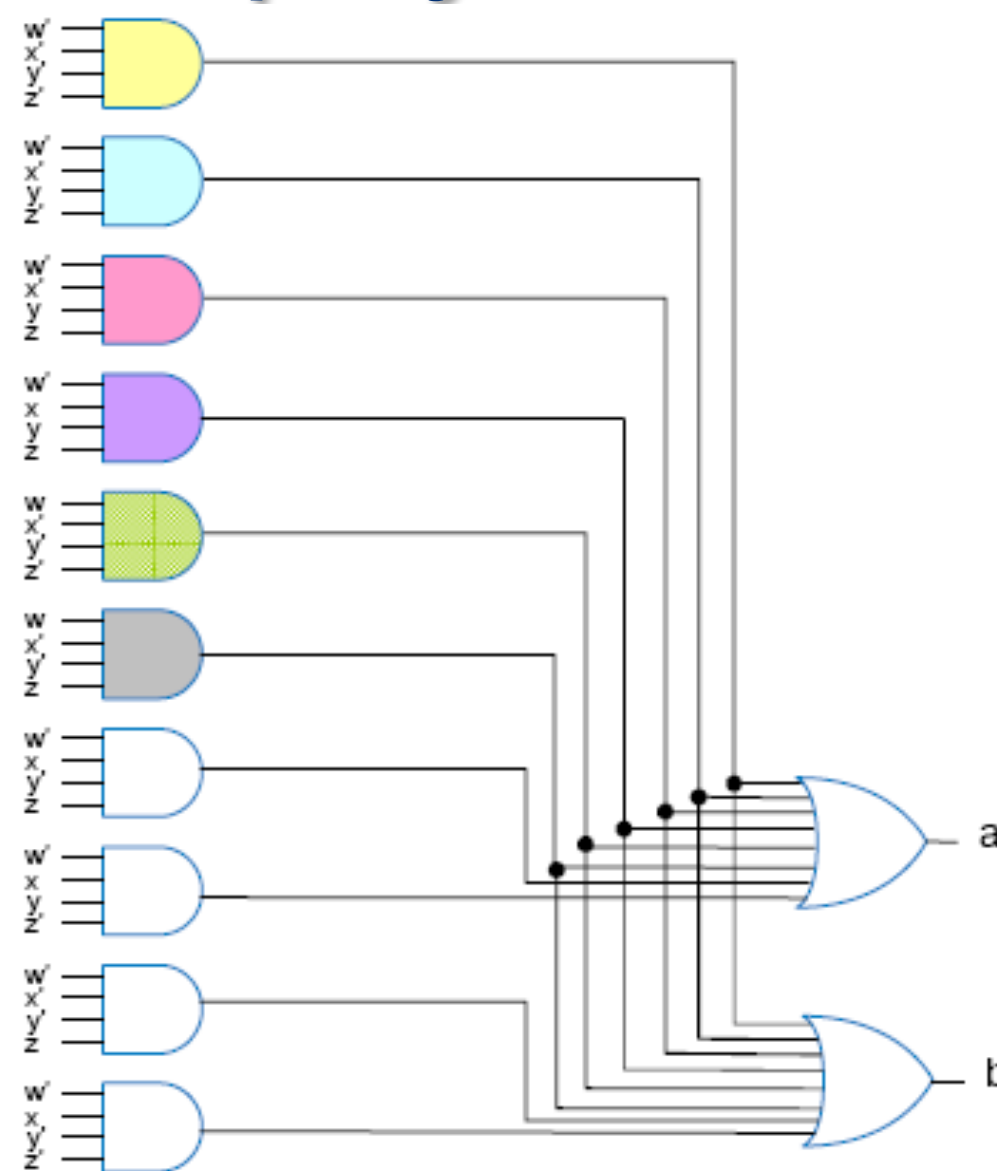
- 여러 출력이 필요한 회로 (사실상 대부분의 회로)
 - Seven-segment display
- MO (Multiple-Output) 회로를 최적화하는 기본 원리는?
 - 게이트 수를 줄여라!!



Seven Segment Display

w	x	y	z	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	0	0	0	0	0	0	0
1	0	1	1	0	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0	0	0
1	1	0	1	0	0	0	0	0	0	0
1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0	0

00000001
00000001



$$a = w'x'y'z' + w'x'yz' + w'x'yz + w'xy'z + w'xyz' + w'xyz + wx'y'z' + wx'y'z$$

$$b = w'x'y'z' + w'x'y'z + w'x'yz' + w'x'yz + w'xy'z' + w'xyz + wx'y'z' + wx'y'z$$

Multiple Output Circuits 사례 (2)

- 다중 출력 K-maps 에서 비용 최적화
 - 로컬 최적화 vs. 글로벌 최적화
 - 공유 가능한 회로를 찾아라

f1

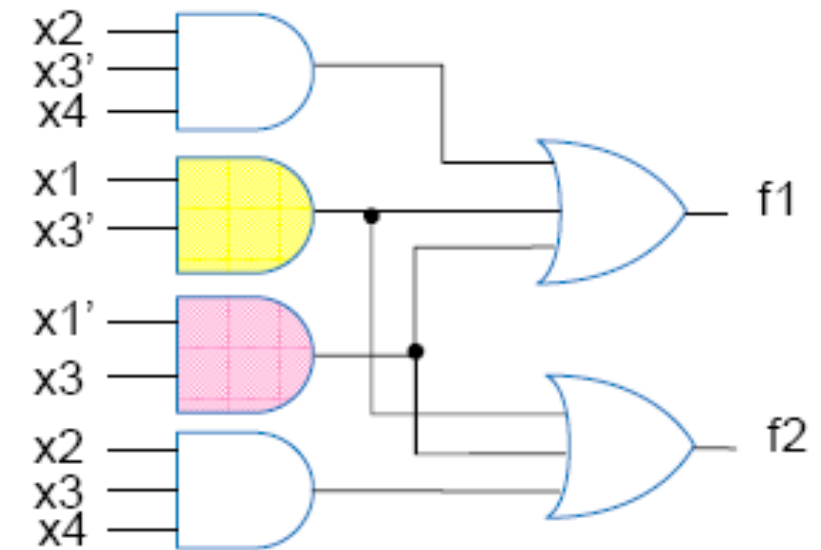
x1x2 \ x3x4	00	01	11	10
00	0	0	1	1
01	0	1	1	1
11	1	1	0	0
10	1	1	0	0

$$f1 = x1x3' + x1'x3 + x2x3'x4$$

f2

x1x2 \ x3x4	00	01	11	10
00	0	0	1	1
01	0	0	1	1
11	1	1	1	0
10	1	1	0	0

$$f2 = x1x3' + x1'x3' + x2x3x4$$



Multiple Output Circuits 사례 (3)

f3

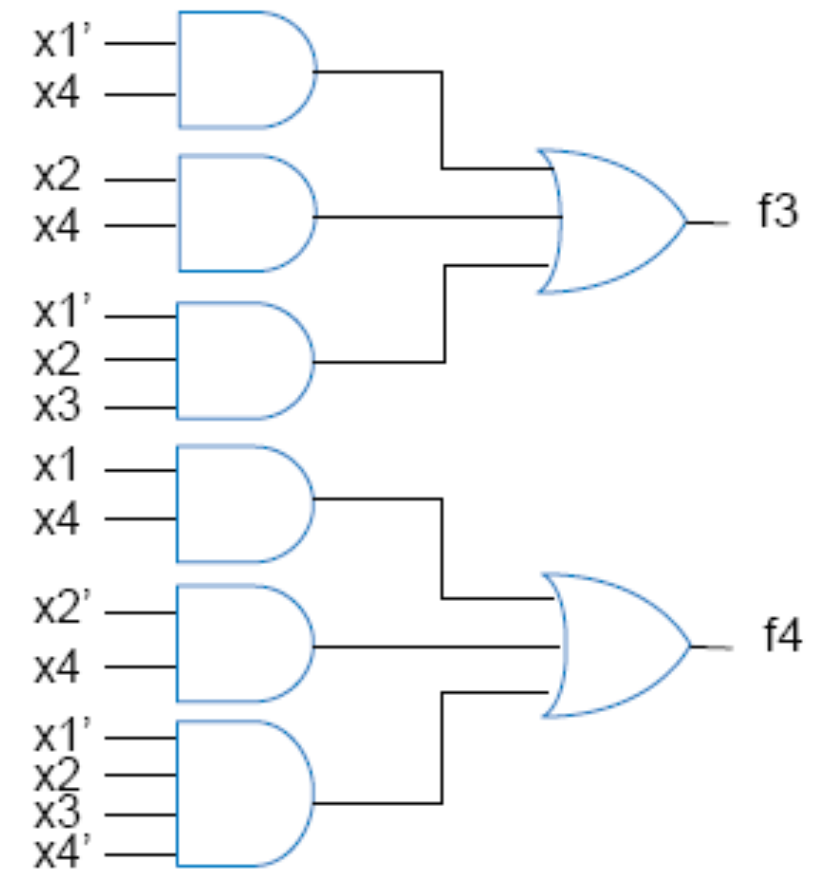
		x3x4			
x1x2		00	01	11	10
00		0	0	0	0
01		1	1	1	0
11		1	1	1	0
10		0	1	0	0

$$f3 = x1'x4 + x2x4 + x1'x2x3$$

f4

		x3x4			
x1x2		00	01	11	10
00		0	0	0	0
01		1	0	1	1
11		1	0	1	1
10		0	1	0	0

$$f4 = x1x4 + x2'x4 + x1'x2x3x4'$$



nothing to combine

Multiple Output Circuits 사례 (3)

- 어떻게든 공유 가능한 회로를 찾아내기

f3

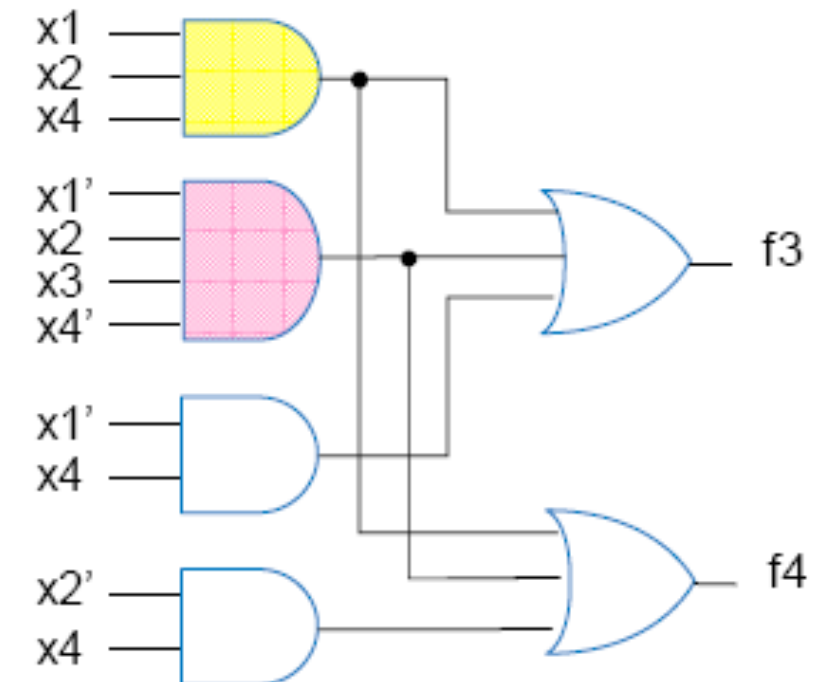
		x3x4			
x1x2		00	01	11	10
00		0	0	0	0
01		1	1	1	0
11		1	1	1	0
10		0	1	0	0

$$f3b = x1x2x4 + x1'x2x3x4' + x1'x4$$

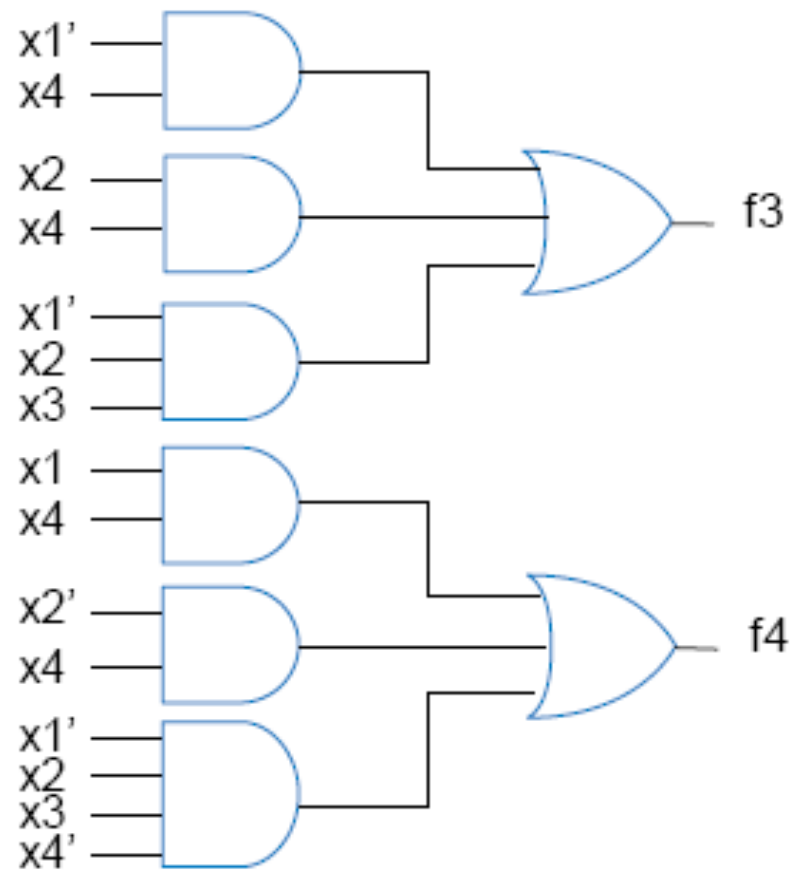
f4

		x3x4			
x1x2		00	01	11	10
00		0	0	0	0
01		1	0	1	1
11		1	0	1	1
10		0	1	0	0

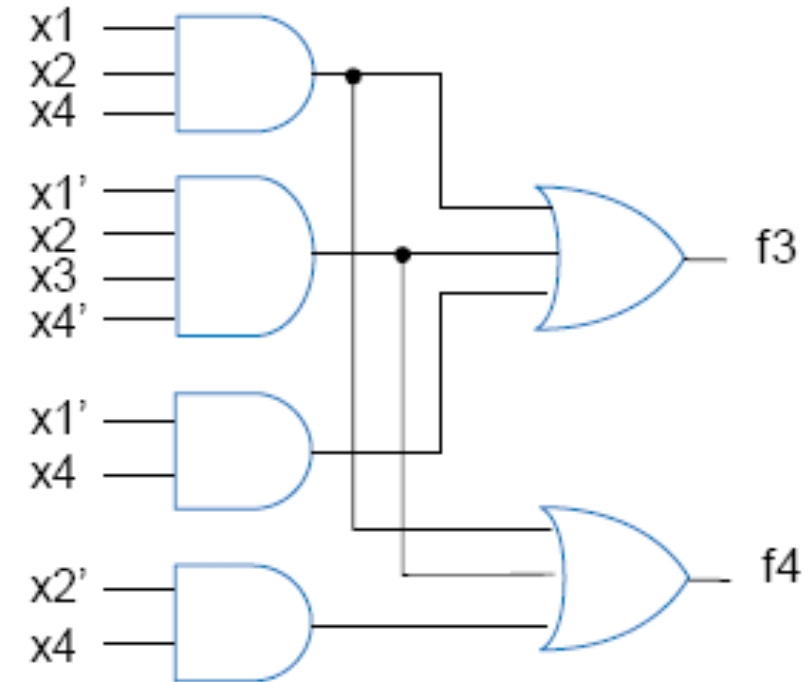
$$f4b = x1x2x4 + x1'x2x3x4' + x2'x4$$



Multiple Output Circuits Solutions 비교



$$\text{cost} = 8 + 21 = 29$$



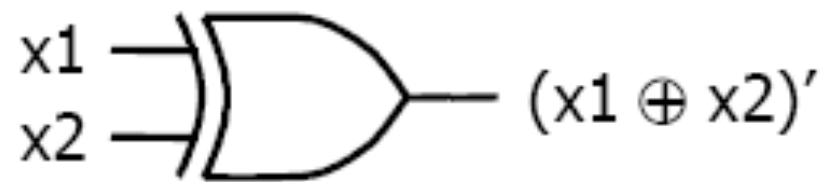
$$\text{cost} = 6 + 17 = 23$$

누가 이런 최적화를 수행해야 하나?

소프트웨어 도구!!

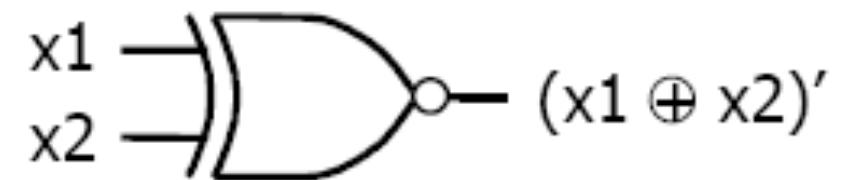
추가로 알아야 하는 논리 게이트

- XOR gate – exclusive OR
 - 두 입력이 다를 때 True
 - 입력이 3개 이상일 때에는?
- XNOR gate – complementing XOR gate
 - 입력이 3개 이상일 때에는?



XOR

a	b	F
0	0	0
0	1	1
1	0	1
1	1	0



XNOR

a	b	F
0	0	1
0	1	0
1	0	0
1	1	1