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# **Advanced Programming**

1. Introduction to Functional Programing







- Motivation & Course Goals
- Functional Programming Primer + Scala Intro
- Course Organization
- Traits
- Algebraic Data Types
- Variance of Type Parameters
- Fold functions
- Primary Constructors
- In the next episode ...



# **Apache Spark**

#### A Motivating Example



- Open-source cluster-computing framework aimed at Big Data processing
- Compute queries on large amount of data in **distributed storage**
- Simple interface: like local data structures
- Powerful semantics: distribution and parallelization
- Originated in 2009 at UC Berkeley, run by **Apache Foundation**
- 600 devs from 200 companies contribute to spark
- Implemented in Scala, interfaced to Java, Python, and R
- Key reason of popularity of Scala for Big Data
- Much faster than Hadoop's MapReduce due to heavy use of in-memory operations
- In **high demand**: some of you will land a Spark/Scala job before we finish:)
- Some Users: NBCUniversal, Netflix, Uber, Capital One, Baidu, Salesforce.com, ...

## **Count Word Frequency in a File**

```
1 object StorageApp {
    def main(args: Array[String]) {
      val conf = new SparkConf()
        .setAppName ("SimpleApp")
        .setMaster ("local[6]") // use 6 cores
      val sc = new SparkContext (conf)
      val lines = sc
        .textFile("/home/wasowski/opt/spark/README.md", 2)
        cache
12
      val wordCounts = lines
13
        .flatMap (line => line.split(" ") )
14
        .map (word => (word, 1))
        .reduceByKey (_ + _)
17
      println (wordCounts.collect.map (_.toString).mkString)
18
      sc.stop()
19
21 }
```

- Resilient distributed dataset (RDD)
- lines is an RDD of strings
- **Distributed** fault-tolerant processing
- L14 split each line into words, merge into RDD of words
- L15 RDD of words to RDD of pairs
- L16 merge pairs with same word, summing counters (map-reduce)
- We use collection operations
- Transformations (flatMap,map) build representation of computation.
- Transformations are lazy.
- Actions (reduceByKey) are eager: execute (force) representations
- Pure program, no vars&side effects

(C) Andrzei Wasowski, IT University of Conenhagen

cache only works if you have referential transparency

# Things go wrong with side-effects

Side effects are officially banned in this course!

```
_{1} var counter = 0
2 var rdd = sc.parallelize(data)
3 // Wrong: Don't do this!!
4 rdd.foreach(x => counter += x)
```



- Line 2: we parallelize a computation
- Line 4: we sum values from an RDD incrementing a counter
- This cannot be done in a distributed way!!!
- Each node gets a closure containing the counter, the closure is sent to nodes.
- Each node increments a different copy of the counter!
- This is why we use functions like map and reduceByKey instead of variables

## Count Word Frequency in a Real-Time Data Stream

```
1 object StreamingApp {
    def main (args: Array[String]) {
      val sparkConf = new SparkConf()
        .setAppName("StreamingApp")
        .setMaster ("local[6]")
      // Sample every second
      val ssc = new StreamingContext (sparkConf, Seconds(1))
      val lines = ssc.socketTextStream("localhost", 9999,
              StorageLevel.MEMORY_AND_DISK_SER)
11
      val wordCounts = lines
12
        .flatMap (line => line.split(" "))
        .map(word => (word, 1))
14
        .reduceBvKev ( + )
16
17
      wordCounts.print // Nothing gets printed!
                       // The computation starts here
18
      ssc.awaitTermination
20
21 }
```

- L12-L15: identical algorithm
- Because the DStream interface is the same as RDD's
- RDDs and Streams are monads
- We will understanding this style of API really deeply
- reduceByKey needs a commutative associative operator (a monoid)
- + is a monoid on integers
- L12-L17 builds a representation of computation
- L18: streaming starts, before this nothing happens
- L17 printing every 1s until killed
- L17 behaves as if in a loop!

#### Introduction to Scala I

- A rich modern OO programming language with a functional part; eager by default, statically typed
- Compiles to JVM, compatible with Java on byte code level
- Designed by prof. Martin Odersky at EPFL in Lausanne
- First official release in 2004
- This is not a course about Scala, we use Scala to learn concepts that apply to other languages (F#, Haskell, Ocaml, Java, Python, Ruby, etc)
- Today's goals: (1) [recall] basics of functional programming and (2) teach Scala syntax and concepts
- Easy for those that have seen functional programming and Scala (focus primarily on harder exercises and on mapping your knowledge from other languages to Scala)
- Hard for those that are new to functional programming and Scala. Focus on the easiest exercises first, and really work hard the first 5-6 weeks.

#### **Basics of Scala**

A singleton class and its only instance

**object** creates a name space; used to build modules. Access the namespace with navigation: MyModule.abs(42)

```
1 object MvModule {
2
    def abs(n: Int): Int = i_i f_i (n < 0) -n else n
    private def formatAbs(x: Int) =
      s"The absolute value of $x is ${abs (x)}:
    val magic : Int = 42
    var result :Option[Int] = None
    def main(args: Array[String]): Unit = { <-----</pre>
      assert (magic - 84 == magic. -(84))
      println (formatAbs (magic-100))
15 }
```

```
def Defines a function (I.3)
```

A body **expression** (statements secondary in Scala)

Use braces if more expressions needed.

A named value declaration (final, immutable). Use this a lot.

A variable declaration. Avoid if possible.

Instantiation of a generic type

None is a singleton "constructor". Construct case classes without  $\ensuremath{\text{\textbf{new}}}$ 

```
Operators are functions, can be overloaded:
minus is Int.-(Int) :Int
Unary methods can be used infix: MyModule abs -42 legal
```

Every value is an object

Line 6 shows an interpolated character string

#### **Pure Functions**

**Def. Referentially transparent** expression (*e*)

Expression e is RT iff replacing e by its value in programs does not change their semantics

(Java) append an element to a list

a.add(5) // non RT

(Scala) append to an immutable list

val b = Cons(5,a) // RT

value void; substitution is pointless; the meaning is in the references reachable from a (change over time for the same a)

The value is a list b, identical to a, modulo the added head element

#### **Def. Pure** function (f)

Iff every expression f(x) is referentially transparent for all referentially transparent expressions x. Otherwise **impure** or **effectful**.

In practice: A function is pure if it does not have side effects (writes/reads variables, files or other streams, modifies data structures in place, sets object fields, throws exceptions, halts with errors, draws on screen)

Pure code shows dependencies in interface, good for mocking, testable

# Referential Transparency Poll

Which of the following computations are referentially transparent?

- $1 \quad a = a + 42$
- 2 a ==b + 42
- 3 a[x] == 42
- println("42")
- throw DivideByZero()
- 6 f(f(x)) if f is pure
- 7 z = z + f(f(x)) if f is pure

# **Loops and Recursion**

#### An imperative factorial

```
def factorial (n :Int) :Int = {
  var result = 1
  for (i <- 2 to n)
  result *= i
  return result</pre>
```

Loops compute with effects; cannot be used in pure code

#### Tail recursive, pure factorial

```
def factorial (n :Int) = {
  def f (n :Int, r :Int) :Int =
    if (n<=1) r
    else f (n-1, n*r)
    f (n,1)</pre>
```

Call tails are automatically compiled to loops with O(1) space overhead

#### A pure recursive factorial

```
def factorial (n :Int) :Int =
if (n<=1) 1
else n * factorial (n-1)</pre>
```

call not in tail positior

#### Example execution

```
3 \times (4 \times (factorial(3)))

5 \times (4 \times (3 \times (factorial(2))))

5 \times (4 \times (3 \times (2 \times (factorial(1)))))

5 \times (4 \times (3 \times (2 \times 1)))
```

$$\rightsquigarrow 5 * (4 * (3 * (2 * 1)))$$
  
 $\rightsquigarrow 5 * (4 * (3 * 2))$ 

$$\rightsquigarrow 5*(4*6)$$

$$\rightsquigarrow 5*24$$

Uses O(n) stack space; Technically exponential

(for this example)!

#### Def. Call in tail position

The caller immediately returns the value of the call

#### Function Values

- In functional programing functions are values
- Functions can be **passed to other functions**, composed, etc.
- Functions operating on function values are higher order (HOFs)

```
1 def map (a :List[Int]) (f :Int => Int) :List[Int] =
   a match { case Nil => Nil
            case h::tail => f(h)::map (tail) (f) }
```

```
A functional (pure) example
_{1} val mixed = List(-1, 2, -3, 4)
2map (mixed) (abs _)
1 map (mixed) ((factorial _) compose (abs _))
```

```
alternatively type it explicitly:
     (abs :Int => Int)
```

#### An imperative (impure) example

```
1 \text{ val mixed} = \text{Array} (-1, 2, -3, 4)
2 for (i <- 0 until mixed.length)</pre>
   mixed(i) = abs (mixed(i))
```

```
1 \text{ val mixed1} = \text{Array (-1, 2, -3, 4)}
2 for (i <- 0 until mixed1.length)</pre>
3 mixed1(i) = factorial(abs(mixed1(i)))
```

## Poll: How is your recursion?

```
1 def f (a :List[Int]) :Int = a match {
2   case Nil => 0
3   case h::t => h + f(t)
4 }
```

What is the result of f (List(42,-1,1,-1,1,-1)?

## **Parametric Polymorphism**

#### Monomorphic functions operate on fixed types:

```
A monomorphic map in Scala

def map (a :List[Int]) (f :Int => Int) :List[Int] =

a match { case Nil => Nil

case h::tail => f(h)::map (tail) (f) }
```

There is nothing specific here regarding Int.

```
A polymorphic map in Scala

def map[A,B] (a :List[A]) (f :A => B) :List[B] =
  a match { case Nil => Nil
      case h::tail => f(h)::map (tail) (f) }
```

An example of use (type parameters are inferred):

```
| map[Int,String] (mixed_list) { _.toString } compose
| (factorial _) compose (abs _))
```

- A polymorphic function operates on values of (m)any types (some restriction possible in Scala)
- A polymorphic type defines a parameterized family of types
- Don't confuse with OO-polymorphism roughly equal to "dynamic method dispatch" (dependent on the inheritance hierarchy)

# **HOFs in Scala Standard Library**

Methods of class List[A], operate on this list, type A is bound in the class

```
map[B](f: A =>B): List[B]
Translates this list of As into a list of Bs using f to convert the values
filter(p: A =>Boolean): List[A]
Compute a sublist of this by selecting the elements satisfying the predicate p
flatMap[B](f: A =>List[B]): List[B]
                                                             *type slightly simplified
Builds a new list by applying f to elements of this, concatenating results.
take(n: Int): List[A]
Selects first n elements.
takeWhile(p: A =>Boolean): List[A]
Takes longest prefix of elements that satisfy a predicate.
forall(p: A =>Boolean): Boolean
Tests whether a predicate holds for all elements of this sequence.
exists(p: A =>Boolean): Boolean
```

Tests whether a predicate holds for some of the elements of this sequence.

More at http://www.scala-lang.org/api/current/index.html#scala.collection.immutable.List

# **Anonymous Functions**

#### Literals

```
val 1 = List(1, -2, 3)
val a = Array(-1, 2, -3)
```

#### Function Literals (Anonymous Functions)

#### We need the same for functions

```
val negative =(x :Int) =>x < 0
negative (-42) → true
```

#### Use to create functions in place:

```
1.filter ((x:Int) =>x < 0) \rightsquigarrow ?
a.filter ((x:Int) =>x > 0) \rightsquigarrow ?
```

#### Alternative concise syntax

```
(abs _) \rightsquigarrow (x :Int) =>MyModule.abs x
```

Scala distinguishes functions and methods.

We used this syntax before to turn a method into a function (like above).

#### Currying and partial application

```
val add2 = (x : Int, v : Int) => x + v
val add =(x :Int) =>(y :Int) =>x + y
```

.....a curried function

What is the type of add? What is the value of add (2) (3) \sim ?

Curried functions can be partially applied: val incr =add (1)

a partial application

Type of incr? Value of incr  $(7) \rightsquigarrow ?$ 

Methods can also be curried: def add (x:Int) (y:Int) :Int =x + y

### **Course Organization**

- Our **website** is LearnIT + a git-repo
- **Reading:** read prescribed book chapters and papers **before class**. Without reading you may not be able to solve exercises.
- Lectures: ca. 10 weeks.

  Summarize the main points, but may skip details needed in exercises
- Exercises: ca. 11 weeks, same days as lectures implement small well defined tasks on the topic of the day's lecture
- Mini Projects: ca. 3 weeks, 2 mini-projects Programming tasks to deepen some some topics. Each mini project gets its own time, when there is no lecture, and no exercises.
- 2–3 person groups (2 is better than 3)
- Communicate in class on Thursday, and daily on the LearnIT forum. Andrzej is very good on handling LearnIT, and very bad in handling email. Use email only for sensitive matters.

### **Traits: Rich or Fat Interfaces**

```
1 // A class with a final property 'name' and
2 // a constructor. You can still add
3 // more members like in Java in braces.
4 abstract class Animal (val name :String)
6 // concrete methods
7 trait HasLegs {
  def run () :String = "after you!"
   def jump () :String = "hop!"
10 }
11 // abstract method
12 trait Audible { def makeNoise () :String }
13 // field
14 trait Registered { var id :Int = 0 }
15
16 // multiple traits mixed in
17 class Frog (name:String) extends
   Animal(name) with HasLegs with Audible {
  def makeNoise () :String = "croak!"
20 } // Frog concrete, so provide makeNoise
```

		class		abstr. class		trait
mult. inheritance	1	_	ī	-	1	+
data	1	+	T	+	1	+
concr. methods	ī	+	ī	+	T	+
abstr. methods	1	_	T	+	Ţ	+
constr. params.	1	+	Τ	+	1	-

# Algebraic Data Types (ADTs)

#### Def. Algebraic Data Type

A type generated by one or several constructors, each of which may contain zero or more arguments.

Sets generated by constructors are **summed**, each constructor is a **product** of its arguments; thus **algebraic**.

```
Example: immutable lists
1 sealed trait List[+A] .....
                                                                   Nothing: subtype of any type
2 case object Nil extends List[Nothing] .....
3 case class Cons[+A](head :A, tail :List[A]) extends List[A]
                                                                  companion object of List[+A]
  Example: operations on lists
1 object List {
   def sum(ints :List[Intl) :Int =
                                                      pattern matching uses case class constructors
     ints match { case Nil => 0
                  case Cons(x,xs) => x + sum(xs) }
   def apply[A](as..:A*):..List[A] = overload function application for the object
     if (as.isEmpty) Nil
     else Cons(as.head, apply(as.tail: *))
                                                                             variadic function
8 }
```

# Mentimeter: Dynamic Virtual Dispatch

```
1 class Printable
                                    { void hello() { print ("printable "); }}
2 class Triangle extends Printable{ void hello() { print ("triangle ");}}
3 class Square extends Printable { void hello() { print ("square "); }}
4 . . .
5 \text{ Square } x = \text{new Square ()}
6 Printable v = new Triangle ()
7 x.hello ():
8 ((Printable)x).hello ();
9 y.hello ();
10 ((Printable)y).hello ();
```

- printable printable printable
- square printable triangle printable
- square printable printable printable
- square square triangle triangle
- square square printable printable
- The program will crash, or fail to type check

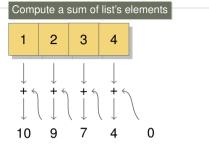
### Variance of Type Parameters

- Write A <: B to say that A is a **subtype of** B (values of A fit where Bs are expected)
- **Example**: if class A extends a class B then A <: B. Same for traits.
- Assume a generic type T[B]: B is a **covariant** parameter of T if for each A <: B we have that T[A] <: T[B] So we can use T[A] values where T[B]s are expected
- In Scala write T[+B] to specify that B is a covariant type parameter.
- Covariance common in pure programs. Scala lists are covariant (List[+B]).
- A is a contravariant parameter of T if whenever A <: B we have that T[B] <: T[A]</p>
- Contravariance is needed if A is a return type, and in some impure situations. In Scala, write TI-A1 to specify contravariance
- Invariance means that there is no automatic subtypes of generic type T: Invariance is default in Scala (when you omit the -/+)
- Recall that Java and C# generics also support variance of type parameters.
- Java has covariant Arrays (unsafe), Scala's arrays are invariant.

# **Quiz: Variance of Type Parameters**

```
1 abstract class A
3 abstract class B extends A
5 // Will the following code type check if T is
6 // (a) invariant,
7 // (b) covariant,
8 // (c) contravariant ?
10 val T[A] = new T[B]
```

# [Right]Folding: Functional Loops



What characterizes similar computations?

- An input list 1 = List(1,2,3,4)
- $\blacksquare$  An initial value z = 0
- A binary operation f :Int => Int = \_ +
- An iteration algorithm (folding)

```
1 \det foldRight[A,B] (f : (A,B) \Rightarrow B) (z :B) (l :List[A]) :B =
   1 match {
      case Cons(x,xs) \Rightarrow f(x, foldRight(f)(z)(xs))
   case Nil => z
6 \text{ val } 11 = \text{List } (1,2,3,4,5,6)
7 \text{ val sum} = \text{foldRight[Int,Int] (\_+\_) (0) (11)}
8 val product = foldRight[Int,Int] (_*_) (1) (11)
9 def map[A,B] (f :A=>B) (l: List[A])=
   foldRight[A.List[B]] ((x, z) => Cons(f(x),z)) (Nil) (1)
```

Many HOFs can be implemented as special cases of folding

## **The Primary Constructor**

```
class Person (val name: String, val age: Int) {
println ("Just constructed a person")
def description = s"$name is $age years old"
}
```

```
1 class Person {
    private String name;
    private int age:
    public String name() { return name; }
    public int age() { return age; }
    public Person(String name, int age) {
      this.name = name:
      this.age = age:
      System.out.println("Just constructed a person");
12
    public String description ()
13
    { return name + "is " + age + " years old"; }
15 }
```

- Parameters become fields
- 'val' parameters become values, 'var' become variables
- If no parameter list, primary constructor takes none
- Constructor initializes fields and executes top-level statements of the class
- Like for all functions, parameters can take default values, reducing the need for overloading
- Note: primary constructors are used with case classes

### Scala: Summary

- Basics (objects, modules, functions, expressions, values, variables, operator overloading, infix methods, interpolated strings.)
- Pure functions (referential transparency, side effects)
- Loops and recursion (tail recursion)
- Functions as values (higher-order functions)
- Parametric polymorphism (monomorphic functions, dynamic and static dispatch)
- Standard HOFs in Scala's library
- **Anonymous functions** (currying, partial function application)
- Traits (fat interfaces, multiple inheritance, mixins)
- Algebraic Data Types (pattern matching, case classes)
- Variance of type parameters (covariance, contravariance, invariance)
- Folding
- **Primary constructors** (default parameter values)

## In the next episode ...

- Basics of functional design: exceptions vs values, partial functions, the Option data type, exception oriented API of Option, for comprehensions, Either
- We will experience the first monadic computation (but refrain from defining monads yet)
- The reading should be relatively easy, so you should really try it!