```
Consider
```

```
Input Array:
arr = [3, 5, 2, 7, 4]
Prefix Sum Array:
prefix = [3, 8, 10, 17, 21]
```

#### **Explanation:**

```
prefix[0]=arr[0]=3
```

- prefix[1]=prefix[0]+arr[1]=3+5=8
- prefix[2]=prefix[1]+arr[2]=8+2=10
- prefix[3]=prefix[2]+arr[3]=10+7=17
- prefix[4]=prefix[3]+arr[4]=17+4=21

```
#include <iostream>
#include <vector>
using namespace std;
vector<int> computePrefixSum(const vector<int>& arr) {
  int n = arr.size();
  vector<int> prefix(n);
  prefix[0] = arr[0];
  for (int i = 1; i < n; ++i) {
     prefix[i] = prefix[i - 1] + arr[i];
  }
  return prefix;
}
int main() {
  vector<int> arr = \{3, 5, 2, 7, 4\};
  vector<int> prefix = computePrefixSum(arr);
  cout << "Prefix Sum: ";</pre>
  for (int sum : prefix) {
    cout << sum << " ";
  }
  return 0;
}
Java
import java.util.Arrays;
public class PrefixSum {
```

```
public static int[] computePrefixSum(int[] arr) {
    int n = arr.length;
    int[] prefix = new int[n];
    prefix[0] = arr[0];

    for (int i = 1; i < n; i++) {
        prefix[i] = prefix[i - 1] + arr[i];
    }

    return prefix;
}

public static void main(String[] args) {
    int[] arr = {3, 5, 2, 7, 4};
    int[] prefix = computePrefixSum(arr);

    System.out.println("Prefix Sum: " + Arrays.toString(prefix));
    }
}</pre>
```

## **Applications**

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#### 1. Efficient Range Sum Queries:

- o To compute the sum of elements from index I to r, use: sum[I to r]=prefix[r]-prefix[I-1]
- o If I=0 then sum[I to r]=prefix[r].

# 2. Finding Subarray Sums:

 Quickly calculate the sum of any subarray in O(1) time after constructing the prefix sum array.

# 3. Sliding Window Optimization:

o Often combined with the sliding window technique for certain optimizations.