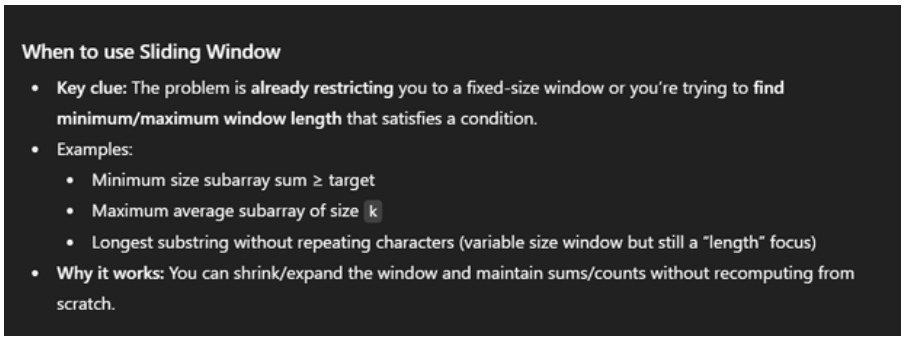
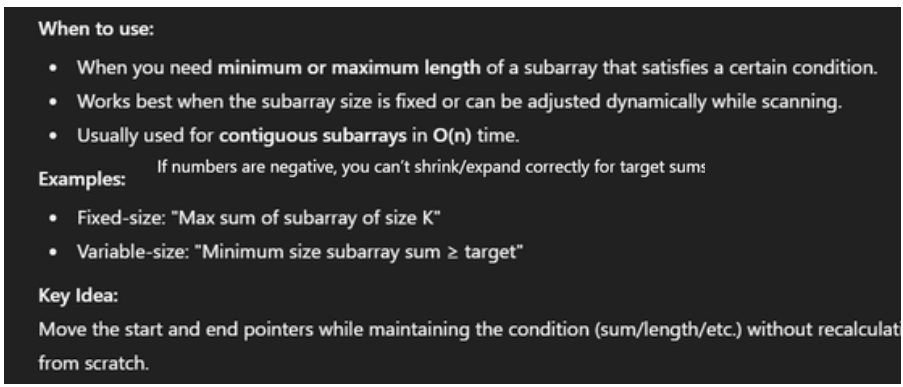
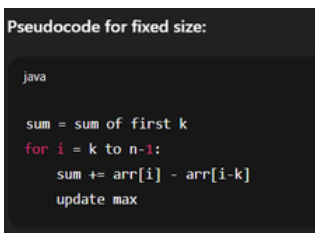
SUBARRAYS PATTERNS

1)SLIDING WINDOW



another view





SOME QUESTIONS-

1)Maximum Average Subarray I (Fixed Size)

2)Minimum Size Subarray Sum (Variable Size)

3)Longest Substring Without Repeating Characters

2) Kadane's Algorithm

2. Kadane's Algorithm

When to use:

- When you need the **maximum sum** (or max product, with modifications) of any contiguous subarray.
- Purely for **sum/product optimization, not** for length calculation.
- Works in **O(n)** time.

Examples:

- "Maximum subarray sum" (LeetCode 53)
- "Maximum product subarray" (LeetCode 152)

Key Idea:

Keep track of the best subarray ending at each position using dynamic programming (`maxEndingHere` and `maxSoFar`).

another view

When to use Kadane's Algorithm

- **Key clue:** Problem is about **maximum sum/product** of any contiguous subarray (no sum constraint, no fixed size).
- Examples:
  - Maximum subarray sum
  - Maximum product subarray
- **Why it works:** Kadane's keeps track of the best sum ending at current index and updates global max in O(n) time.

Example:

`[1, -2, 3, 4, -1, 2]` → max sum subarray is `[3, 4, -1, 2]` = `8`.

Pseudocode:

```
java

maxSoFar = arr[0];
maxEndingHere = arr[0];

for i = 1 to n-1:
    maxEndingHere = max(arr[i], maxEndingHere + arr[i])
    maxSoFar = max(maxSoFar, maxEndingHere)
```

SOME QUESTIONS-

1)Maximum Subarray

2)Maximum Sum Circular Subarray

3)Maximum Product Subarray (Kadane's variant for product)

3) Prefix Sum

3. Prefix Sum

When to use:

- When you need **sum of subarrays** frequently or need to check **sum equals target** in O(1) after preprocessing.
- Useful for problems involving **range queries** or **finding subarray sums with hashmaps** (for variable lengths).
- Works in **O(n)** preprocessing, then O(1)/O(n) query depending on approach.

Examples:

- "Count subarrays with sum = k"
- "Number of subarrays divisible by m"
- Range sum queries

Key Idea:

Store cumulative sums so that `sum(l, r) = prefix[r] - prefix[l-1]`.

Main Use:

- Any subarray sum queries quickly (can handle negative numbers).
- Especially useful for problems like **LeetCode 560 (Subarray Sum Equals K)**.
- Allows O(1) subarray sum queries after O(n) preprocessing.

Key Idea:

- Store `prefix[i] = sum of arr[0..i]`.
- Sum of subarray `(l, r)` = `prefix[r] - prefix[l-1]`.

Extra Trick for 560: Use **HashMap** to store frequency of prefix sums seen so far.
Check if `(currentPrefix - k)` exists → that means a subarray sum = k exists.

Pseudocode (560 style):

```java
map.put(0, 1) // for subarrays starting at index 0
sum = 0
count = 0
for num in arr:
    sum += num
    if map contains (sum - k):
        count += map.get(sum - k)
    map.put(sum, map.getOrDefault(sum, 0) + 1)
```

SOME QUESTIONS-

1) Subarray Sum Equals K

2) Continuous Subarray Sum

3) Range Sum Query - Immutable

Comparison of Subarray Techniques

Method	Primary Goal	When to Use	Example Problem
Kadane's Algorithm	Find maximum sum of any contiguous subarray	Max sum problems with positive and/or negative numbers	LeetCode 53: Maximum Subarray
Prefix Sum + HashMap	Find subarray(s) with exact target sum (length/count)	Target sum problems with negatives allowed	LeetCode 560: Subarray Sum Equals K
Sliding Window	Best sum/length under constraints (often positives only)	Fixed window size or variable window with positives	LeetCode 209: Minimum Size Subarray Sum
Two Pointers	Efficiently find subarrays/pairs meeting condition	Sorted arrays or constraints on difference/ratio	LeetCode 167: Two Sum II (Input Array Sorted)
Monotonic Queue / Deque	Track min/max in a sliding window efficiently	Sliding window max/min problems	LeetCode 239: Sliding Window Maximum

When to Use Which

Problem Type	Kadane's	Sliding Window	Prefix Sum
Max contiguous subarray sum (any integers)	✅	❌	❌
Fixed-size window problems	❌	✅	❌
Variable-size window, all positive numbers	❌	✅	❌
Target sum with negatives allowed (like 560)	❌	❌	✅
Many subarray sum queries	❌	❌	✅