CS3243 AY20/21 Sem1

Lecture 5: September 9

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5.1 $\alpha - \beta$ Search

5.1.1 Algorithm

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Algorithm 1 – Pruning Strategy
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1: function MaxValue(state, \alpha, \beta)
        if Terminal(state) == True then return Utility(state)
 3:
        for each a in Actions(state) do
 4:
             result \leftarrow RESULT(state, a)
 5:
 6:
             v \leftarrow \max(v, MinValue(result, \alpha, \beta))
             if v \ge \beta then return v
 7:
             \alpha \leftarrow \max(\alpha, \mathbf{v})
 8:
         return v
 9: function MINVALUE(state, \alpha, \beta)
        if Terminal(state) == True then return Utility(state)
10:
         v \leftarrow +\infty
11:
        for each a in Actions(state) do
12:
             result \leftarrow RESULT(state, a)
13:
             v \leftarrow \min(v, MaxValue(result, \alpha, \beta))
14:
             if v \leq \beta then return v
15:
             \alpha \leftarrow \min(\alpha, \mathbf{v})
16:
         return v
```

5.1.2 Analysis

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For a game tree of depth d and branch factor b,
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Minimax: O(b^d)

\alpha - \beta: O(b^{d/2}) (best case)
```

5.2 Real-time decisions

5.2.1 Evaluation functions

Although $\alpha - \beta$ search reduces workload by pruning some branches, it still needs to reach terminals, which can be very consuming when the game tree has large depth.

To solve this problem, we modify the $\alpha - \beta$ search.

$$H-Minimax(s,d) = \begin{cases} EVAL(s), & \text{if } CUTOFF-TEST(s,d) \\ max_{a \in Actions(s)} \ H-Minimax(result(s,a),d+1), & \text{if } player(s) = MAX \\ min_{a \in Actions(s)} \ H-Minimax(result(s,a),d+1), & \text{if } player(s) = MIN \end{cases}$$

This method modifies $\alpha - \beta$ search in two ways:

- replace the utility function by a heuristic evaluation function **EVAL**, which estimates the position's utility
- replace the terminal test by a cutoff test that decides when to apply EVAL

Here we introduce one kind of evaluation function called weighted linear function

$$EVAL(s) = w_1 * f_1(s) + w_2 * f_2(s) + \dots + w_n * f_n(s) = \sum w_i * f_i(s)$$

where each w_i is a weight and each f_i is a feature of the position.

5.3 Local search

"Local search algorithms operate using a single current node and generally move only to neighbors of that node." - AIMA

5.3.1 Hill-climbing search

Algorithm 2 Hill-climbing search

- 1: function HILL-CLIMBING(state s) returns a state that is a local best
- 2: $\min Val \leftarrow Val(s)$
- $3: minState \leftarrow$
- 4: $\mathbf{for} \ \mathbf{u} \ \mathbf{in} \ \mathbf{N(s)} \ \mathbf{do}$
- 5: **if** Val(u) < minVal **then** minState = u, minVal = Val(u) **return** minState

Drawback of this approach:

This algorithm may get stuck at local maximum.

5.3.2 N-Queens Puzzle

On a N*N chess board, place N queens in such a way that none of them attack each other.

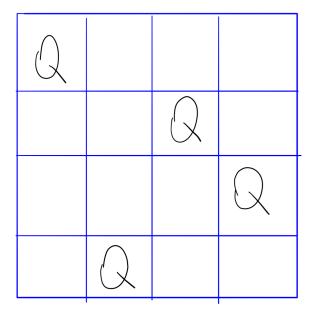


Figure 5.1: Hill-climbing search gets stuck in this non-optimal state

However, if we allow the algorithm to tolerate some regressions, then it is likely to get out of local maximum, as shown below.

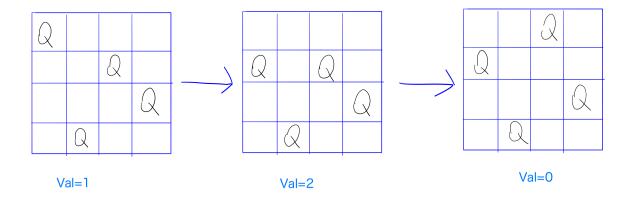


Figure 5.2: A tolerable algorithm can jump out of local maximum