A GRADUATE COURSE IN OPTIMIZATION

A GRADUATE COURSE

IN

OPTIMIZATION

CIE6010 Notebook

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Notations and Conventions

X

Set

```
\inf X \subseteq \mathbb{R} Infimum over the set X
\mathbb{R}^{m \times n}
                  set of all m \times n real-valued matrices
\mathbb{C}^{m \times n}
                  set of all m \times n complex-valued matrices
                  ith entry of column vector \boldsymbol{x}
x_i
                  (i,j)th entry of matrix \boldsymbol{A}
a_{ij}
                  ith column of matrix A
\boldsymbol{a}_i
\boldsymbol{a}_{i}^{\mathrm{T}}
                  ith row of matrix A
                  set of all n \times n real symmetric matrices, i.e., \mathbf{A} \in \mathbb{R}^{n \times n} and a_{ij} = a_{ji}
\mathbb{S}^n
                  for all i, j
                  set of all n \times n complex Hermitian matrices, i.e., \mathbf{A} \in \mathbb{C}^{n \times n} and
\mathbb{H}^n
                  \bar{a}_{ij} = a_{ji} for all i, j
\boldsymbol{A}^{\mathrm{T}}
                  transpose of \boldsymbol{A}, i.e, \boldsymbol{B} = \boldsymbol{A}^{\mathrm{T}} means b_{ji} = a_{ij} for all i,j
                  Hermitian transpose of \boldsymbol{A}, i.e, \boldsymbol{B} = \boldsymbol{A}^{H} means b_{ji} = \bar{a}_{ij} for all i,j
A^{H}
trace(A)
                  sum of diagonal entries of square matrix A
1
                  A vector with all 1 entries
0
                  either a vector of all zeros, or a matrix of all zeros
                  a unit vector with the nonzero element at the ith entry
e_i
C(A)
                  the column space of \boldsymbol{A}
\mathcal{R}(\boldsymbol{A})
                  the row space of \boldsymbol{A}
\mathcal{N}(\boldsymbol{A})
                  the null space of \boldsymbol{A}
\operatorname{Proj}_{\mathcal{M}}(\mathbf{A}) the projection of \mathbf{A} onto the set \mathcal{M}
```

10.2. Wednesday

In this lecture we will discuss the main penalty algorithms formally.

10.2.1. Introduction to penalty algorithms

Given the equality constraint problem

$$\begin{aligned} & \min \quad f(\boldsymbol{x}) \\ & h(\boldsymbol{x}) = 0 \\ & \boldsymbol{x} \in X \subseteq \mathbb{R}^n \end{aligned} \tag{10.10}$$

The augmented Lagrangian function is given by:

$$L_{\boldsymbol{c}}(\boldsymbol{x}, \boldsymbol{\lambda}) = f(\boldsymbol{x}) + \boldsymbol{\lambda}^{\mathrm{T}} h(\boldsymbol{x}) + \frac{\boldsymbol{c}}{2} \|h(\boldsymbol{x})\|^{2}$$
(10.11)

There are two penalty algorithms formally listed below:

Quadratic Penalty (Courant, 1943).

$$m{x}^r \leftarrow rg \min_{m{x} \in X} L_{m{c}^r}(m{x}, m{\lambda}^r, m{c}^r)$$
 increase $m{c}^{r+1} > m{c}^r
ightarrow \infty$, $\|m{\lambda}^r\| < +\infty$

Augmented Lagrangian Multiplier Method.

$$\mathbf{x}^r \leftarrow \arg\min_{\mathbf{x} \in X} L_{\mathbf{c}}(\mathbf{x}, \boldsymbol{\lambda}^r)$$

 $\boldsymbol{\lambda}^{r+1} \leftarrow \boldsymbol{\lambda}^r + \mathbf{c}h(\mathbf{x})$

where c is sufficiently large in general, but not goes to infinite.

- ullet The reason why the λ converges to the optimal Lagrange multiplier will be discussed later.
- Why do we need to take sufficiently large *c*? Let's raise an example first.

■ Example 10.3 Given the problem

$$\min \quad \frac{1}{2}(-x_1^2 + x_2^2)$$
$$x_1 = 1$$

with optimal solution $\pmb{x}^*=(1,0); \lambda^*=1.$ The augmented Lagrangian function is given by:

$$L_c(\boldsymbol{x},\lambda) = \frac{1}{2}(-x_1^2 + x_2^2) + \lambda(x_1 - 1) + \frac{c}{2}(x_1 - 1)^2,$$

Applying the optimality condition, we derive

$$\mathbf{x}^*(c,\lambda) = \begin{pmatrix} \frac{c-\lambda}{c-1} \\ 0 \end{pmatrix}$$

The second order necessary condition should be

$$\nabla_{xx}^2 L_c = \begin{pmatrix} c - 1 & 0 \\ 0 & c + 1 \end{pmatrix} \succeq 0$$

Therefore, from this example we can see that if c is not large enough, the second order optimality condition will be violated. Later we will give a formal proof for this statement.

10.2.2. Convergence Analysis

Theorem 10.2 Given the sequence of optimization problem

$$\min \quad L_{\boldsymbol{c}^k}(\boldsymbol{x}, \boldsymbol{\lambda}^k), \boldsymbol{x} \in X \tag{10.12}$$

with associated local minimum \mathbf{x}^k . Suppose $\{\mathbf{c}^k\}$ is monotone increaseing to infinite, and $\{\boldsymbol{\lambda}^k\}$ is bounded, then every limit point in the sequence $\{\mathbf{x}^k\}$ is the global minimum for original problem (10.10)

Proof. 1. Firstly note that the optimal value f^* for (10.10) is the infimum of the

optimal value for (10.12), i.e.,

$$f^* = \inf_{h(\mathbf{x}) = 0, \mathbf{x} \in X} f(\mathbf{x})$$

$$= \inf_{h(\mathbf{x}) = 0, \mathbf{x} \in X} \left\{ f(\mathbf{x}) + (\boldsymbol{\lambda}^k)^{\mathrm{T}} h(\mathbf{x}) + \frac{\mathbf{c}^k}{2} ||h(\mathbf{x})||^2 \right\}$$

$$= \inf_{h(\mathbf{x}) = 0, \mathbf{x} \in X} L_{\mathbf{c}^k}(\mathbf{x}, \boldsymbol{\lambda}^k)$$

We study the lower bound for $L_{c^k}(x, \lambda^k)$. By definition,

$$L_{\mathbf{c}^k}(\mathbf{x}, \boldsymbol{\lambda}^k) \ge L_{\mathbf{c}^k}(\mathbf{x}^k, \boldsymbol{\lambda}^k) \tag{10.13}$$

Taking infimum over \boldsymbol{x} both sides, we obtain:

$$f^* \ge L_{c^k}(\boldsymbol{x}^k, \boldsymbol{\lambda}^k) \tag{10.14}$$

2. Then we show that the limit point \bar{x} is such that $f^* \geq f(\bar{x})$ and \bar{x} is feasible by using (10.14). Suppose $\{x^k, \lambda^k\} \to \{\bar{x}, \bar{\lambda}\}$, and by taking limsup both sides for (10.14), we obtain:

$$f(\bar{\boldsymbol{x}}) + (\bar{\boldsymbol{\lambda}})^{\mathrm{T}} h(\bar{\boldsymbol{x}}) + \limsup_{k \to \infty} \frac{\boldsymbol{c}^k}{2} \|h(\boldsymbol{x}^k)\|^2 \le f^*$$
 (10.15)

Since $c^k \to \infty$ and $||h(x)||^2 \ge 0$, and the LHS of (10.15) is bounded above, we derive $h(x^k) \to 0$, and in particular,

$$h(\bar{\mathbf{x}}) = 0 \tag{10.16}$$

Combining (10.15) and (10.16), we conclude that $f^* \geq f(\bar{x})$ and \bar{x} is feasible, i.e., \bar{x} is global minimum.