并行与分布式作业

Homework-5 第5次作业

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一、问题描述

- 1. Consider a sparse matrix stored in the compressed row format (you may find a description of this format on the web or any suitable text on sparse linear algebra). Write an OpenMP program for computing the product of this matrix with a vector. Download sample matrices from the Matrix Market (http://math.nist.gov/MatrixMarket/) and test the performance of your implementation as a function of matrix size and number of threads.
- 2. Implement a producer-consumer framework in OpenMP using sections to create a single producer task and a single consumer task. Ensure appropriate synchronization using locks. Test your program for a varying number of producers and consumers.
- 3. 利用MPI通信程序测试本地进程以及远程进程之间的通信时延和带宽。

二、解决方案

Problem1:用OpenMP做稀疏矩阵向量乘法

矩阵形式

采用**MatrixMarket matrix coordinate real symmetric**形式,即用记录矩阵中每一个非零项的值和 坐标,形式如下:

```
%%MatrixMarket matrix coordinate real general
% This ASCII file represents a sparse MxN matrix with L
% nonzeros in the following Matrix Market format:
% +-----+
% |%%MatrixMarket matrix coordinate real general | <--- header line
% |%
                                <--+
% |% comments
                                  |-- 0 or more comment
lines
% |%
% | M N L
                                 | <--- rows, columns, entries</pre>
% | I1 J1 A(I1, J1)
                                 <--+
% | I2 J2 A(I2, J2)
% | I3 J3 A(I3, J3)
                                    |-- L lines
                                    % | IL JL A(IL, JL)
                                 <--+
% Indices are 1-based, i.e. A(1,1) is the first element.
```

其中,

```
      val[i]:第i个非零项的值

      I[i]:第i个非零项的行坐标

      J[i]:第i个非零项的列坐标
```

并行方案

朴素的实现:

```
for(int i = 0; i < M; i++){
   for(int j = 0; j < N; j ++){
      ans[i] += matrix[i][j] * vector[j];
   }
}</pre>
```

由于直接记录了每一个非零项的值和坐标,所以上面的朴素实现变为:

```
for(int i = 0; i < nz; i++){//nz为非零项的个数
ans[I[i]] += val[i] * vector[J[i]];
}
```

使用guided的调度,并行执行for循环。进行计时。

```
clock_t start, end;
start = clock();
#pragma omp parallel for num_threads(threadNum) schedule(guided)
for (int i = 0; i < nz; i++) {
    ans[I[i]] += val[i] * vector[J[i]];
}
end = clock();
printf("time=%f\n", ((double)end - start) / (CLOCKS_PER_SEC / 1000));</pre>
```

Problem2: 生产者消费者模型

设计buffer的结构如下:

```
struct bufferStruct {
   int buf[SIZE];
   int head, tail, count;
} buffer;
```

buffer是一个环形缓冲区,head是下一个要读取的数据的位置,tail是下一个要写入的数据的位置,count表示当前的环形缓冲区内有多少个数据。

buffer的基本操作如下:

```
void init() {
   buffer.head = buffer.tail = buffer.count = 0;
}
```

```
int isfull() {
    return (buffer.count == SIZE);
}
int isempty() {
   return (buffer.count == 0);
}
int put(int tag) {//向环形缓冲区写入数据
   if (isfull()) return -1;
#pragma omp critical
   {
        buffer.buf[buffer.tail] = tag;
       buffer.tail = (buffer.tail + 1) % SIZE;
       buffer.count++;
   return 1;
}
int get() {//从环形缓冲区读取数据
   if (isempty()) return -1;
   int tag;
#pragma omp critical
   {
       tag = buffer.buf[buffer.head];
        buffer.head = (buffer.head + 1) % SIZE;
       buffer.count--;
   return tag;
}
```

设计生产者和消费者,生产者把[st, ed)的数据写入到buffer,消费者依次读取cnt个数据。

单一生产者和单一消费者:

多个生产者和多个消费者(示例):

Problem3: 利用MPI通信程序测试本地进程以及远程进程之间的通信时延和带宽

搭建MPI环境

由于只有一台机器,所以使用了虚拟机来作为另一个节点,搭建了一个只有两个节点的mpi集群。搭建方法参考了下面这个链接:

https://wenku.baidu.com/view/85be4649a6c30c2259019ee9.html

测试通信时延

创建两个进程,分别跑在本地节点跟虚拟机节点上。进程0开始计时,发送一个字节的消息给进程1,进程1接受消息,并且把该一个字节的消息返回给进程0,进程0接收到返回的消息后停止计时。由于消息只有一个字节,所以拷贝到缓存的时间可以忽略不计,中间经过的时间就是消息从进程0发送到进程1,再从进程1返回进程0的时间,也就是往返时间。则单程的时间为往返时间的一半。

由于只执行一次的时间误差可能很大,所以执行1000次,取平均值以减少误差。

进程0的代码如下:

```
if (rank == 0) {
        dest = 1;
        source = 1;
        for (n = 1; n \le reps; n++) \{
            T1 = MPI_Wtime();
            rc = MPI_Send(&msg, 1, MPI_BYTE, dest, tag, MPI_COMM_WORLD);
            if (rc != MPI_SUCCESS) {
                printf("Send error in task 0!\n");
                MPI_Abort(MPI_COMM_WORLD, rc);
                exit(1);
            }
            rc = MPI_Recv(&msg, 1, MPI_BYTE, source, tag, MPI_COMM_WORLD,
&status);
            if (rc != MPI_SUCCESS) {
                printf("Receive error in task 0!\n");
                MPI_Abort(MPI_COMM_WORLD, rc);
                exit(1);
            }
            T2 = MPI_Wtime();
            deltaT = T2 - T1;
            sumT += deltaT;
        }
        avgT = (sumT * 1000000) / reps;
        printf("\nAvg round trip time = %d microseconds\n", avgT);
        printf("Avg one way latency = %d microseconds\n", avgT / 2);
```

}

进程1的代码如下:

```
else if (rank == 1) {
        dest = 0;
        source = 0;
        for (n = 1; n \le reps; n++) \{
            rc = MPI_Recv(&msg, 1, MPI_BYTE, source, tag, MPI_COMM_WORLD,
&status);
            if (rc != MPI_SUCCESS) {
                printf("Receive error in task 1!\n");
                MPI_Abort(MPI_COMM_WORLD, rc);
                exit(1);
            }
            rc = MPI_Send(&msg, 1, MPI_BYTE, dest, tag, MPI_COMM_WORLD);
            if (rc != MPI_SUCCESS) {
                printf("Send error in task 1!\n");
                MPI_Abort(MPI_COMM_WORLD, rc);
                exit(1);
            }
        }
    }
```

测试带宽

进程0发送和接收n个double类型的数据给进程1,进程1也发送和接收n个double类型的数据给进程0。通过对进程0的发送和接收操作进行计时,从而计算带宽,带宽就是 $\frac{2\times n\times sizeof(double)\times 10^{-6}}{t}$ MB/s。

使用MPI_Isend和MPI_Irecv进行测试,由于这两个函数是non-blocking的,需要用MPI_Waitall来等待两个操作执行完毕。

为了减少误差,对于每个n测试10次,取时间的最小值作为使用的时间。

```
for (int k = 0; k < NUMBER_OF_TESTS; k++) {</pre>
    if (rank == 0) {
         MPI_Sendrecv(MPI_BOTTOM, 0, MPI_INT, 1, 14,
                         MPI_BOTTOM, 0, MPI_INT, 1, 14, MPI_COMM_WORLD, &status);
         t1 = MPI_Wtime();
         MPI_Isend(sbuf, n, MPI_DOUBLE, 1, k, MPI_COMM_WORLD, &r[0]);
         MPI_Irecv(rbuf, n, MPI_DOUBLE, 1, k, MPI_COMM_WORLD, &r[1]);
         MPI_Waitall(2, r, statuses);
         t2 = MPI_Wtime() - t1;
         if (t2 < tmin) tmin = t2;
    } else if (rank == 1) {
         MPI_Sendrecv(MPI_BOTTOM, 0, MPI_INT, 0, 14,
                         MPI_BOTTOM, 0, MPI_INT, 0, 14, MPI_COMM_WORLD, &status);
         \label{eq:mpi_recv} \texttt{MPI\_Irecv}(\texttt{rbuf}, \ \mathsf{n}, \ \texttt{MPI\_DOUBLE}, \ \mathsf{0}, \ \mathsf{k}, \ \texttt{MPI\_COMM\_WORLD}, \ \&\texttt{r}[\mathsf{0}]);
         MPI_Isend(sbuf, n, MPI_DOUBLE, 0, k, MPI_COMM_WORLD, &r[1]);
         MPI_Waitall(2, r, statuses);
    }
}
```

三、实验结果

Problem1: 用OpenMP做稀疏矩阵向量乘法

(1) 使用4个线程时,所用时间随矩阵规模的变化:

```
en-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 494_bus.mtx 4
WMatrixMarket matrix coordinate real symmetric
494 494 1080
time=0.322000
%%MatrixMarket matrix coordinate real symmetric
662 662 1568
time=0.333000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 685_bus.mtx 4
%MatrixMarket matrix coordinate real symmetric
685 685 1967
time=0.320000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 4
 %MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=0.306000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main bcsstk16.mtx 4
%MatrixMarket matrix coordinate real symmetric
4884 4884 147631
time=0.307000
```

可以看到,一开始的矩阵规模为494*494,且有1080个非零项,后者同理。time为多线程计算矩阵向量乘法所用的毫秒数。

随着矩阵规模和非零项的同时增加,使用四个线程时,所用的时间近乎不变,甚至随着矩阵规模的增加 而减少,原因是矩阵此时的矩阵规模比较小,随着矩阵规模的增大,线程的开销的占比逐渐减少,多线 程带来的加速比更大,所以总时间反而随着任务的增加而减少了。

(2) 所用时间随线程数量的变化:

```
wen-G3-3574:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 6
%%MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=0.609000
gwen@gwen-63-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 7
9%MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=0.417000
%%MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=0.159000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 9
%MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=0.251000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 10
 %MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=7.981000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$./main 1138_bus.mtx 11
 MatrixMarket matrix coordinate real symmetric
1138 1138 2596
time=3.443000
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem1$ ./main 1138_bus.mtx 12
%MatrixMarket matrix coordinate real symmetric
time=80.232000
```

可以看到,随着线程数量的增加,一开始所用的时间逐渐减少,在使用8个线程时到达最优,随后线程数量继续增加时,所用时间也增加了。原因是超过8个线程之后,线程的开销(创建线程和调度线程)的占比越来越大,当线程数量到达一定数目以后,线程的开销已经不可承受,使总时间大幅增加。

Problem2: 生产者消费者模型

单一生产者和单一消费者:

```
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem2$ ./main
Producing 0 by thread 0
Consuming 0 by thread 1
Producing 1 by thread 0
Consuming 1 by thread
Producing 2 by thread 0
Consuming 2 by thread 1
Producing 3 by thread 0
Consuming 3 by thread 1
Producing 4 by thread 0
Consuming 4 by thread 1
Producing 5 by thread 0
Consuming 5 by thread 1
Producing 6 by thread 0
Consuming 6 by thread 1
Producing 7 by thread 0
Consuming 7 by thread 1
Producing 8 by thread 0
Consuming 8 by thread 1
Producing 9 by thread 0
Consuming 9 by thread 1
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem25
```

可以看到,thread 0为生产者线程,thread 1为消费者线程,生产者和消费者的运行是正确的,没有出现数据竞争问题。

多个生产者和多个消费者:

三个生产者、三个消费者: 生产者分别生产[0, 10), [10, 20), [20, 30)的数据,每个消费者消费10个数据。

```
gwen@gwen-G3-3579:/media/gwen/DATA/A_GW/Courses/Parallel_and_distributed_computing/hw5/Problem2$ ./main
Buffer is empty
Producing 10 by thread 1
Producing 0 by thread 4
Consuming 20 by thread 2
Consuming 0 by thread 5
                                                                                 Ţ
Producing 20 by thread 3
Consuming 10 by thread 0
Producing 11 by thread 1
Producing 1 by thread 4
Consuming 11 by thread 2
Consuming 1 by thread 5
Producing 21 by thread 3
Consuming 21 by thread 0
Producing 2 by thread 4
Producing 12 by thread 1
Producing 22 by thread 3
Consuming 2 by thread 2
Consuming 12 by thread 5
Consuming 22 by thread 0
Producing 3 by thread 4
Producing 13 by thread 1
Producing 23 by thread 3
Consuming 3 by thread 2
```

由于输出太长,这里只截图一部分。一共分配了6个线程,分别对应3个生产者和3个消费者,为了让程序的性能更优,在不影响消费者读取数据的正确性的前提下,没有限制输出为临界区,所以输出顺序有一定的混乱(因为消费者获取了数据之后,并不是马上输出数据,中间有一定的时间,而这里就会有线程之间的争用,这对于消费者的正确性没有影响,只是会影响现在这样把结果输出调试的顺序)。可以看到消费者和生产者的运行是正确的。

Problem3: 利用MPI通信程序测试本地进程以及远程进程之间的通信时延和带宽

通信带宽和时延的测试结果如下:

```
gwen@gwen-G3-3579:~/cluster$ ./run.sh
Avg round trip time = 34 microseconds
Avg one way latency = 17 microseconds
       time (sec)
                    Rate (MB/sec)
1000
      0.000039
                    409.200390
2000
      0.000042
                    766.958446
4000
      0.000067
                    948.535180
      0.000234
8000
                    547.827461
16000 0.000348
                     736.448439
32000
      0.000813
                    629.575974
64000 0.001490
                    687.194767
128000 0.003082
                    664.598421
256000 0.005013
                     816.999676
512000 0.010493
                      780.725707
1024000 0.020680
                      792.264944
gwen@gwen-G3-3579:~/cluster$
```

四、遇到的问题及解决方法

这次实验中遇到的最大的问题就是mpi集群的搭建,花了几乎一整天的时间。我用的是虚拟机,按照网上的教程搭建好后,一旦遇到跨节点的通信就会报错,为此花费了很长时间,我用的是mpich,尝试了把mpich换成openmpi、更新版本等操作,都没有找到错误的原因是什么,最后发现了如下一个问题解答:

If you have a machinefile with hostnames instead of ip-addresses and have the machines connected locally then you should have a nameserver running locally as well or else change the entries in your machine file to ip-address instead of hostnames. Having just /etc/hosts will not solve the issue.

我在etc/hosts中配置了两个节点的hostname、ip,在machinefile中使用的是hostname,由于是虚拟机所以也符合connected locally的条件,所以把machinefile中的hostname全部改成ip地址即可!