EE2026

Digital Fundamentals

Combinational Building Blocks and Structural/Dataflow Verilog Description

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Outline

- Introduction
- Binary adders
 - Half adders, full adders, ripple adders.
- Magnitude comparators
- Decoders, BCD to 7-segment decoders
- Encoders, Multiplexers
- Demultiplexers
- Tri-state logic elements

Introduction

- There are two types of logic circuits
 - Combinational and sequential logic circuits
- Combinational logic
 - The output depends only on the current inputs
- Sequential logic
 - The output depends on both past and present inputs, which implies that there is a memory element in the sequential circuit
- Combinational building blocks that are commonly used in digital systems

Half Adders

It is a one bit binary adder with two inputs of A_i and B_i

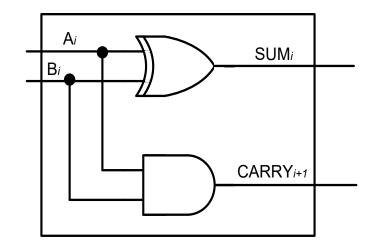
$$0 + 0 = 0$$

 $0 + 1 = 1$
 $1 + 0 = 1$
 $1 + 1 = 10$

Carry
$$\rightarrow C_{i+1}$$

$$A: A_n \dots A_{i+1} A_i \dots A_0$$

$$B: B_n \dots B_{i+1} B_i \dots B_0$$
Sum $\rightarrow S_i$



A _i	B _i	Sum _i	Carry _{i+1}
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

Carry in from i-1 bit cannot be added

$$S_i = \overline{A}_i \cdot B_i + A_i \cdot \overline{B}_i \cdot C_{i+1} = A_i \cdot B_i$$

Half Adders (cont.)

Dataflow Verilog description

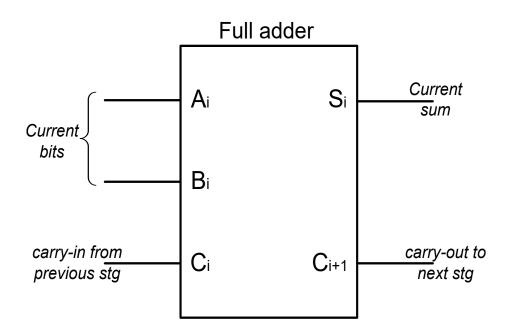
$$S_i = A_i \oplus B_i$$
$$C_{i+1} = A_i \cdot B_i$$

A _i	B _i	Sum _i	Carry _{i+1}
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

```
module ha(S,Cout,A,B);
   input A, B;
   output S, Cout; // Cout is the carry output
   assign S = A ^ B;
   assign Cout = A & B;
endmodule
```

Full Adders

 Full adders can use the carry bit from the previous stage of addition



A _i	B_i	C_i	Si	C _{i+1}
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Full Adders (cont.)

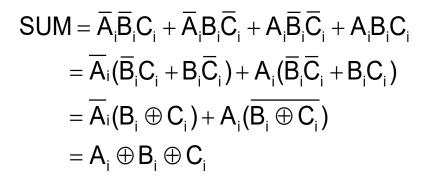
K-map for SUM

A _i B _i C _i	0	1
00	0	1
01	1	0
11	0	1
10	1	0

Note: C_{i+1} is not a MSOP, but less overall hardware is reqd. if we use this expression. It allows sharing of A_i XOR B_i between SUM_i and C_{i+1} .

K-map for CARRY

A _i B _i C _i	0	1
00	0	0
01	0	1
11	1	1
10	0	1



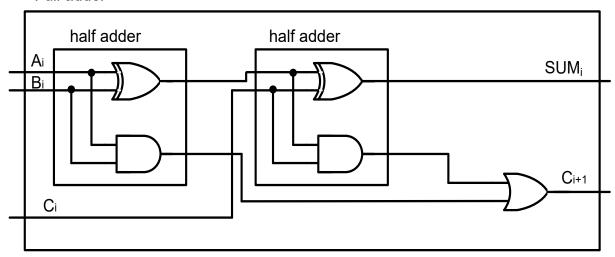
$$\begin{split} C_{i+1} &= A_i B_i + A_i \overline{B}_i C_i + \overline{A}_i B_i C_i \\ &= A_i B_i + C_i (A_i \overline{B}_i + \overline{A}_i B_i) \\ &= A_i B_i + C_i (A_i \oplus B_i) \end{split}$$

Full Adder Circuit

$$SUM = (A_i \oplus B_i) \oplus C_i$$

$$C_{i+1} = A_i B_i + C_i (A_i \oplus B_i)$$

Full adder



Note: A full adder adds 3 bits. Can also consider as first adding first two and then the result with the carry

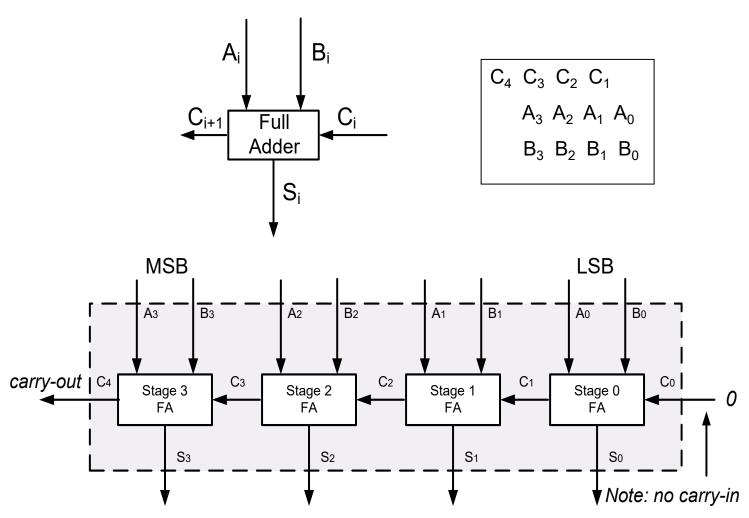
Full Adder Circuit

Verilog dataflow description

$$SUM = (A_i \oplus B_i) \oplus C_i$$
$$C_{i+1} = A_i B_i + C_i (A_i \oplus B_i)$$

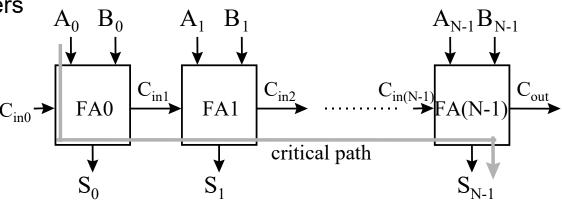
```
module fa(S,Cout,A,B,Cin);
  input A, B, Cin; // Cin is the carry input
  output S, Cout; // Cout is the carry output
  assign S = A ^ B ^ Cin;
  assign Cout = A & B | Cin & (A ^ B);
endmodule
```

Parallel Adders



Parallel Adders (cont.)

- In general, n full adders need to be used to form an n-bit adder
- Carry ripple effect
 - output of each full adder is not available until the carry-in from the previous stage is delivered
 - carry bits have to propagate from one stage to the next
 - as the carries ripple through the carry chain → also known as ripple carry adders



 This slow rippling effect is substantially reduced by using carry look ahead adders

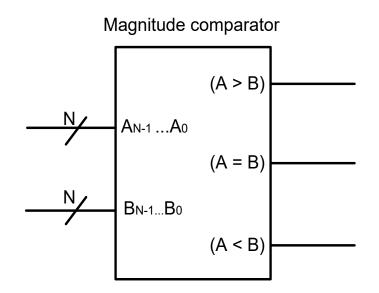
Parallel Adders (cont.)

Structural Verilog description (parameterized, arbitrary bit width)

```
module rca(S,Cout,A,B,Cin); // 4-bit ripple carry adder
    parameter N = 4; // parameterized bit width
    input [N-1:0] A, B;
    input Cin; // Cin is the adder carry input (at LSB)
    output [N-1:0] S;
    output Cout; // Cout is the adder carry output (at MSB)
    wire [N:0] C; // carry inputs of all full adders + carry output of last one
    assign C[0] = Cin;
    assign Cout = C[N];
    genvar i; // temp variable used only in generate loop
       generate for(i=0;i<N;i=i+1) begin
            fa FAinstance (.S(S[i]),.Cout(C[i+1]),.A(A[i]),.B(B[i]),.Cin(C[i]));
            end
       endgenerate
endmodule
```

Magnitude Comparator

 Outputs are functions of relative magnitudes of input binary numbers A and B



Functional block diagram

Magnitude Comparator: Truth Table

2-bit magnitude comparator

A_1	A_0	B ₁	B_{0}	(A > B)	(A = B)	(A < B)				
0	0	0	0	0	1	0				
0	0	0	1	0	0	1				
0	0	1	0	0	0	1				
0	0	1	1	0	0	1				
0	1	0	0	1	0	0				
0	1	0	1	0	1	0				
0	1	1	0	0	0	1				
0	1	1	1	0	0	1				
1	0	0	0	1	0	0				
1	0	0	1	1	0	0				
1	0	1	0	0	1	0				
1	0	1	1	0	0	1				
1	1	0	0	1	0	0				
1	1	0	1	1	0	0				
1	1	1	0	1	0	0				
1	1	1	1	0	1	0				

K-maps for A>B and A<B



A_1A_0 B_1B_0	00	01	1 1	10
00	0	1	1	1
01	0	0	1	1
11	0	0	0	0
10	0	0	(1	0

$$(A > B) = A_1 \overline{B}_1 + A_0 \overline{B}_1 \overline{B}_0 + A_1 A_0 \overline{B}_0$$

A<B

A_1A_0 B_1B_0	00	01	11	10
00	0	0	0	0
01	1	0	0	0
11	1) 1	0 (1
10)	1	0	0

$$(A < B) = \overline{A}_1 B_1 + \overline{A}_1 \overline{A}_0 B_0 + \overline{A}_0 B_1 B_0$$

K-map for A=B

A=B

A_1A_0 B_1B_0	00	01	11	10
00	1	0	0	0
01	0	1	0	0
11	0	0	1	0
10	0	0	0	1

$$(A = B) = \overline{A}_1 \overline{A}_0 \overline{B}_1 \overline{B}_0 + \overline{A}_1 A_0 \overline{B}_1 B_0 + A_1 A_0 B_1 B_0 + A_1 \overline{A}_0 B_1 \overline{B}_0$$

This can be generated indirectly using (A<B) and (A>B)



$$(A = B) = \overline{(A < B)} \cdot \overline{(A > B)}$$

Magnitude Comparator: Verilog

Dataflow Verilog description (parameterized, arbitrary bit width)

```
module magcomp(AgreaterB,AequalB,AlowerB,A,B);
    parameter N = 4;

input [N-1:0] A, B;
    output AgreaterB, AequalB, AlowerB;

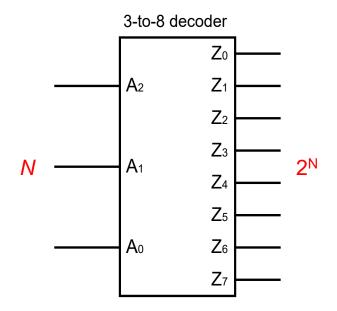
assign AgreaterB = (A > B);
    assign AequalB = (A == B);
    assign AlowerB = (A < B);

/* to reduce complexity at the cost of slightly worse performance: assign AlowerB = ~AgreaterB & ~AequalB */
endmodule</pre>
```

Decoder

- Input: N-bit input code
- A decoder activates a (single) appropriate output line among M (more than N, usually 2^N) as a function of the input

Functional block diagram

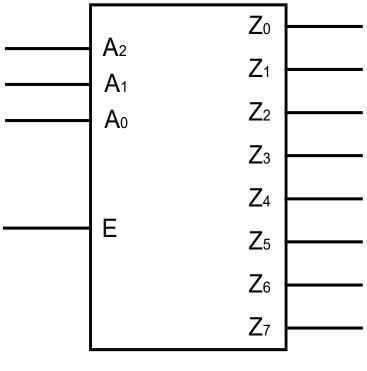


Truth Table

A ₂	A ₁	A ₀	Z ₀	Z ₁	Z ₂	Z ₃	Z ₄	Z ₅	Z ₆	Z ₇
0	0	0	1	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0	0
0	1	0	0	0	1	0	0	0	0	0
0	1	1	0	0	0	1	0	0	0	0
1	0	0	0	0	0	0	1	0	0	0
1	0	1	0	0	0	0	0	1	0	0
1	1	0	0	0	0	0	0	0	1	0
1	1	1	0	0	0	0	0	0	0	1

Example: Decoder 3-8

- A decoder has 2^N output lines for N inputs
 - named N-2^N decoder
 - output can be single- or multi-bit
- Enable signal
 - if E = 1, normal operation
 - if E= 0, disable outputs (all 0's)
 - N enables permit to combine multiple decoders (see videolecture on MSIs)



Functional block diagram

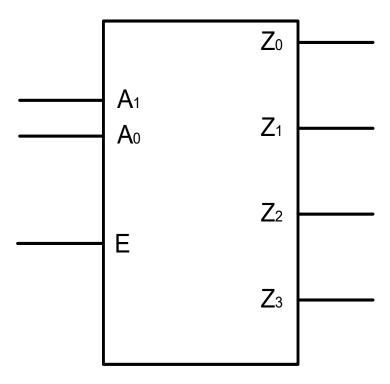
Decoder with Enable Signal

Truth Table including Enable signal (single-bit output)

Inputs							Out	outs			
Е	A ₂	A ₁	A ₀	Z ₀	Z ₁	Z ₂	Z ₃	Z ₄	Z ₅	Z ₆	Z ₇
0	Χ	Χ	Χ	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	0	0	0	0
1	0	0	1	0	1	0	0	0	0	0	0
1	0	1	0	0	0	1	0	0	0	0	0
1	0	1	1	0	0	0	1	0	0	0	0
1	1	0	0	0	0	0	0	1	0	0	0
1	1	0	1	0	0	0	0	0	1	0	0
1	1	1	0	0	0	0	0	0	0	1	0
1	1	1	1	0	0	0	0	0	0	0	1

Example: Decoder 2-4

- A 2:4 decoder has 2² output lines for N inputs
 - output can be single- or multi-bit
- Enable signal
 - if E = 1, normal operation
 - if E= 0, disable outputs (all 0's)



Functional block diagram

Decoder with Enable Signal

Truth Table including Enable signal (single-bit output)

	nput	8		Outputs				
Е	A ₁	A ₀	Z ₀	Z ₁	Z ₂	Z ₃		
0	X	Χ	0	0	0	0		
1	0	0	1	0	0	0		
1	0	1	0	1	0	0		
1	1	0	0	0	1	0		
1	1	1	0	0	0	1		

Decoder: Verilog

Dataflow Verilog description of a 2:4 decoder

```
module decoder24(Z,A,E);
    input [1:0] A;
    input E;
    output [0:3] Z;

assign Z = ((A == 2'b00) & E) ? 4'b1000 :
        ((A == 2'b01) & E) ? 4'b0100 :
        ((A == 2'b10) & E) ? 4'b0010 :
        ((A == 2'b11) & E) ? 4'b0001 :
        4'b0000;
        // 0000 is assigned if E=0
endmodule
```

Inputs			Outputs				
Е	A ₁	A ₀	Z ₀	Z ₁	Z ₂	Z ₃	
0	Χ	Χ	0	0	0	0	
1	0	0	1	0	0	0	
1	0	1	0	1	0	0	
1	1	0	0	0	1	0	
1	1	1	0	0	0	1	

Decoder: Verilog

Dataflow Verilog description of a 2:4 decoder

```
module decoder24(Z,A,E);
    input [1:0] A;
    input E;
    output [0:3] Z;
    assign Z = ((A == 2'b00) \& E) ? 4'b1000 :
                ((A == 2'b01) \& E) ? 4'b0100 :
                ((A == 2'b10) \& E) ? 4'b0010 :
                ((A == 2'b11) \& E) ? 4'b0001 :
                4'b0000; // 0000 if E=0
endmodule
```

Inputs			Outputs				
E	A ₁	A ₀	Z ₀	Z ₁	Z ₂	Z ₃	
0	Χ	Χ	0	0	0	0	
1	0	0	1	0	0	0	
1	0	1	0	1	0	0	
1	1	0	0	0	1	0	
1	1	1	0	0	0	1	

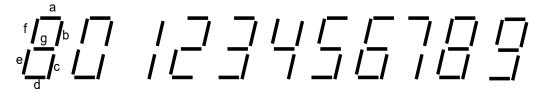
Decoder: Verilog

Parameterized dataflow Verilog description (arbitrary bit width)

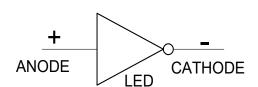
```
module decoder(Z,A,E);
    parameter M = 4; // parameterized design (sets # inputs)
    parameter N = 16; // parameterized design (sets # outputs=2^M)
    input [M-1:0] A;
    input E;
    output [0:N-1] Z;
    wire [N-1:0] zerovec = {N{1'b0}}; // replication operator N(.)
    /* other option: define constant zerovec=1'b0, it will be extended to the left
       with zeros to the correct bit width */
    assign Z = (enable) ? (1 << A) : zerovec;
    // if enable=0, output is set to to zerovec = 00...0
    // if enable=1, shift "1" A times and fill all other positions with zeros
endmodule
```

Example: BCD-to-7 Segment Decoder

 Converts a BCD number into signals required to display that number on a 7-segment display

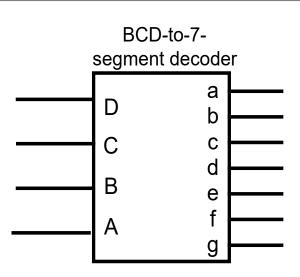


A 7-segment display. Each segment is an LED which will light when a logic T signal is applied to it

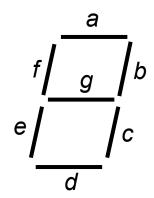


- 7-segment displays are of 2 types: common anode and common cathode
- Common anode display has all LED anodes connected and is active low, whereas the common cathode display is active high

BCD-to-7 Segment Decoder – cont.



Functional block diagram



Truth Table

D	С	В	Α	а	b	С	d	е	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	X	X	X	X	X	X	X
1	0	1	1	X	X	X	X	X	X	X
1	1	0	0	Χ	X	X	X	X	X	X
1	1	0	1	X	X	X	X	X	X	X
1	1	1	0	X	X	X	X	X	X	X
1	1	1	1	X	X	X	X	X	X	X

- Verilog: essentially, description of a truth table
 - no clever dataflow Verilog description (behavioral, see later)

Encoder

- For different input bits (usually 2^N), encoder generates a code with fewer bits (usually N bits) uniquely identifying the input
 - performs the inverse of the decoding function Functional block diagram

 C_1 **I**5 **I**7

Truth Table (an 8-3 encoder)

I ₀	I ₁	l ₂	I ₃	I ₄	I ₅	I ₆	I ₇	C ₂	C ₁	C ₀
1	0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1
0	0	1	0	0	0	0	0	0	1	0
0	0	0	1	0	0	0	0	0	1	1
0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	1	0	0	1	0	1
0	0	0	0	0	0	1	0	1	1	0
0	0	0	0	0	0	0	1	1	1	1

No clever dataflow Verilog description (behavioral, see later)

Example: Priority Encoder

- Generic encoders: error flagged if multiple input bits are 1
- Priority encoder allows multiple input bits to be 1
 - output set by the input bit with highest priority (i.e., most significant position), ignoring those with lower priority

I ₀	I ₁	l ₂	l ₃	l ₄	l ₅	l ₆	l ₇	C ₂	C ₁	C ₀
1	0	0	0	0	0	0	0	0	0	0
X	1	0	0	0	0	0	0	0	0	1
X	X	1	0	0	0	0	0	0	1	0
X	X	X	1	0	0	0	0	0	1	1
X	X	X	X	1	0	0	0	1	0	0
X	X	X	X	X	1	0	0	1	0	1
X	X	X	X	X	X	1	0	1	1	0
X	X	X	X	X	X	X	1	1	1	1

Example: Priority Encoder

Dataflow Verilog description of 4-2 priority encoder

I ₀	I ₁	l ₂	I ₃	C ₁	C ₀
1	0	0	0	0	0
X	1	0	0	0	1
X	X	1	0	1	0
Χ	Χ	Χ	1	1	1

 Use nested conditional operators, starting from MSB and progressively moving to the LSB

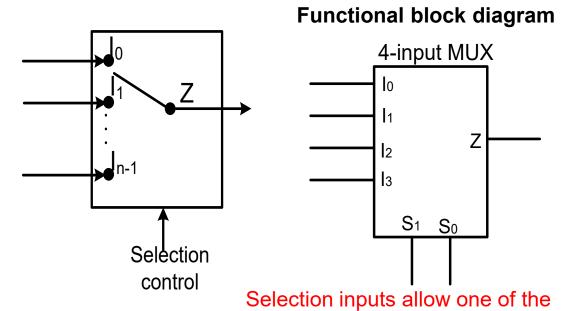
No clever parameterized dataflow Verilog description

Multiplexer

 A multiplexer (MUX) is a combinational circuit element that selects data from one of 2^N inputs and directs it to a single output, according to an N-bit selection signal

inputs to pass through to the output

- inputs/outputs can be 1 or M>1 bit wide
- examples with 1-bit inputs/output



Condensed truth table

S ₁	S ₀	Z
0	0	I ₀
0	1	l ₁
1	0	l ₂
1	1	l ₃

Actual truth table would have 2^6 rows corresponding to I_0 , I_1 , I_2 , I_3 , S_0 and S_1

Example: 4:1 MUX

 Sometimes include enable input signal

$$Z = E \cdot (\overline{S_0}\overline{S_1}I_0 + S_0\overline{S_1}I_1 + \overline{S_0}S_1I_2 + S_0S_1I_3)$$

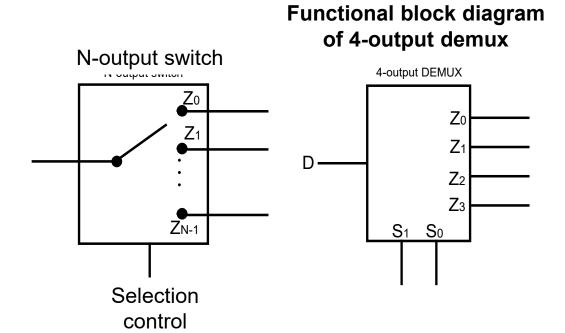
 M-bit inputs/output: use M 1-bit MUXes

E	S₁	So	l _o	I₁	l ₂		Z
0	Χ	Χ	Χ	Χ	Χ	Χ	0
1	0	0	0	Χ	Χ	Х	0
1	0	0	1	X	Χ	Χ	1
1	0	1	Χ	0	Χ	Χ	0
1	0	1	Χ	1	Χ	Χ	1
1	1	0	Х	Χ	0	Х	0
1	1	0	Χ	Χ	1	Χ	1
1	1	1	Χ	X	Χ	0	0
1	1	1	Χ	Χ	Χ	1	1

```
module mux41(Z,S,I0,I1,I2,I3,E);
    parameter M = 16; // 16-bit inputs and outputs
    input [M-1:0] I0, I1, I2, I3; // inputs
    input [1:0] S; // 2-bit selection signal
    input E; // enable
    output [M-1:0] Z;
    assign Z = E ? (S[1] ? (S[0] ? I3 : I2) : (S[0] ? I1 : I0)) : 0;
endmodule
```

Demultiplexer

- A Demultiplexer (DEMUX) connects an input signal to any of 2^N output lines, based on an N-bit selection control
 - inputs/outputs can be 1 or M>1 bit wide
 - examples with 1-bit inputs/output



D	S ₁	S ₀	Z ₀	Z ₁	Z ₂	Z_3
0	X					
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

Truth table

Example: 1:4 DEMUX

Boolean expression of output

$$Z_{0} = D \cdot \overline{S_{0}} \cdot \overline{S_{1}}$$

$$Z_{1} = D \cdot \underline{S_{0}} \cdot \overline{S_{1}}$$

$$Z_{2} = D \cdot \overline{S_{0}} \cdot S_{1}$$

$$Z_{3} = D \cdot S_{0} \cdot S_{1}$$

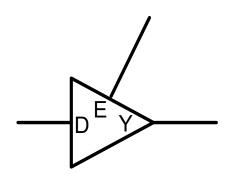
D	S ₁	S ₀	Z ₀	Z ₁	Z ₂	Z_3
0	X	Χ	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

```
module demux41(Z0,Z1,Z2,Z3,S,D);
    parameter M = 16; // 16-bit inputs and outputs
    input [M-1:0] D; // input
    input [1:0] S; // 2-bit selection signal
        output [M-1:0] Z0, Z1, Z2, Z3;
    assign Z0 = (S == 2'b00) ? D : 1'b0; // zeros extended to other bits of Z0
    assign Z1 = (S == 2'b01) ? D : 1'b0; // zeros extended to other bits of Z0
    assign Z2 = (S == 2'b10) ? D : 1'b0; // zeros extended to other bits of Z0
    assign Z3 = (S == 2'b11) ? D : 1'b0; // zeros extended to other bits of Z0
endmodule
```

Tri-State Logic Elements

- Ordinarily, a digital device has 2 states
 - tri-state devices also have high impedance state (Z)
 - floating output: the device does not force any voltage
 - voltage set by the output of some other device
 - if only one device is enabled at a time (all others in Z),
 multiple devices can drive the same node without conflicting
 - several tri-state logic gates
 - example: tri-state buffer with active-high enable

Functional block diagram



Voltage table

Е	D	Υ
1	0	0
1	1	1
0	X	Z

← Z = high impedance

Tri-State Logic Gates: Verilog

Dataflow Verilog description of various logic gates

tristate buffer with active-high enable

```
module tristatebuffer(Y,D,E);
  input D, E;
  output Y;
  assign Y = E ? D : 1'bz;
endmodule
```

tristate inverter with active-high enable

```
module tristateinv(Y,D,E);
input D, E;
output Y;
assign Y = E ? ~D : 1'bz;
endmodule
```

tristate buffer with active-low enable

```
module tristatebuffer(Y,D,E);
  input D, E;
  output Y;
  assign Y = E ? 1'bz : D;
endmodule
```

tristate inverter with active-low enable

```
module tristateoinv(Y,D,E);
input D, E;
output Y;
assign Y = E ? 1'bz : ~D;
endmodule
```

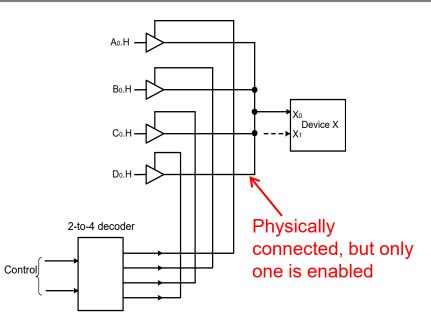
Can you write the dataflow Verilog description of tristate NAND2 and MUX4:1?

MUXes Based on Tri-State Elements

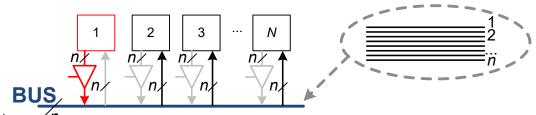
 Tri-state gates with common output implement MUXes

When Control = 00, tri-state device for A_0 is enabled, others are disabled. Hence A_0 is connected to X_0 , etc.

Control signals select which input goes to X ⇒ effectively it behaves like a MUX



- Useful to connect several resources to same bus
 - avoids expensive point-to-point interconnection
 - the enabled resource drives the bus (others in Z may receive)



Summary

- Introduction to combinational building blocks and their structural/dataflow Verilog description
- Binary adders
 - half adders, full adders, ripple carry adders
- Magnitude comparators
- Decoders, BCD-to-7-segment decoders
- Encoders, Priority encoders
- Multiplexers
- Demultiplexers
- Tri-state logic elements
- Behavioral Verilog description style is also possible: see next lectures...

Suggestions for Self-Improvement

- In addition to the lecture/tutorials/lab sessions on Verilog, you may want to read chapter 4 of the textbook (see IVLE Workbin)
 - description of logic functions

