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Shellcode: Using the Exception Directory to find GetProcAddress

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Introduction

Let's say you want the location of the GetProcAddress API in memory, but you can't use the Import Address Table (IAT) or the Export Address Table (EAT). What other ways can you do it?. Perhaps there are many ways, but let me suggest one that's relatively simple to implement and only involves searching for immediate values in the code section. When

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GetProcAddress or GetProcAddressForCaller cannot locate the address of a function in a dynamic library, they will return the error code **STATUS_ORDINAL_NOT_FOUND**. If we search in kernelbase.dll for this immediate value, we should land somewhere in the address range of these API. From there, we locate the entry point.

Method 1 (32-bit)

Search the code section (.text) of each Dynamic-link Library (DLL) for the immediate value 0xC0000138. If we find it, reverse the direction of search until we find the prolog bytes. For stdcall convention, prolog bytes normally begin with push ebp and mov ebp, esp. If the prolog contains mov edi, edi we can safely skip that because it's only used for hot-patching systems after XP SP2. The following pseudo-code attempts to describe this idea.

```
func GetGPA
  set addr = 0

foreach (DLL in PEB) and addr is 0
  for pos = start(DLL.text) to end(DLL.text) - 4
   if pos[0] equal to STATUS_ORDINAL_NOT_FOUND
     while (pos[0] not equal to prolog (push ebp, mov ebp, esp))
     set pos = pos - 1
     set addr = pos
        break
   end if
     set pos = pos + 1
   end for
end for
```

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```
set GetGPA = addr
end func
```

The following code in C demonstrates the idea.

```
LPVOID GetGPA(VOID) {
    PPEB
                          peb;
    PPEB LDR DATA
                          ldr:
    PLDR DATA TABLE ENTRY dte;
    LPV0ID
                          addr=NULL;
    BYTE
                          С;
    PIMAGE DOS HEADER
                          dos:
    PIMAGE_NT_HEADERS
                          nt;
    PIMAGE SECTION HEADER sh;
    DWORD
                          i, j, h;
    PBYTE
                          CS;
    peb = (PPEB) readfsdword(0x30);
    ldr = (PPEB LDR DATA)peb->Ldr;
    // for each DLL loaded
    for (dte=(PLDR DATA TABLE ENTRY)ldr->InLoadOrderModuleList.Flink;
         dte->DllBase != NULL && addr == NULL:
         dte=(PLDR DATA TABLE ENTRY)dte->InLoadOrderLinks.Flink)
      // is this kernel32.dll or kernelbase.dll?
      for (h=i=0; i<dte->BaseDllName.Length/2; i++) {
        c = dte->BaseDllName.Buffer[i];
```

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```
h += (c \mid 0 \times 20);
    h = ROTR32(h, 13);
 if (h != 0x22901A8D) continue;
 dos = (PIMAGE DOS HEADER)dte->DllBase;
 nt = RVA2VA(PIMAGE NT HEADERS, dte->DllBase, dos->e lfanew);
  sh = (PIMAGE SECTION HEADER)((LPBYTE)&nt->OptionalHeader +
         nt->FileHeader.SizeOfOptionalHeader);
  for (i=0; i<nt->FileHeader.NumberOfSections && addr == NULL; i+
    if (sh[i].Characteristics & IMAGE SCN MEM EXECUTE) {
      cs = RVA2VA (PBYTE, dte->DllBase, sh[i].VirtualAddress);
      for(j=0; j < sh[i].Misc.VirtualSize - 4 && addr == NULL; j++)
        // is this STATUS ORDINAL NOT FOUND?
        if(*(DWORD*)\&cs[j] == 0xC0000138) {
          while(--j) {
            // is this the prolog?
            if(cs[j]) == 0x55 \&\&
               cs[i+1] == 0 \times 8B \&\&
               cs[i+2] == 0xEC) {
              addr = \&cs[j];
              break;
return addr;
```

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This approach should work fine on 32-bit legacy systems, but not 64-bit systems.

Method 2 (64-bit)

The first method doesn't work for x64 builds because of compiler optimizations and different calling convention. stdcall is replaced with Microsoft fastcall, and chunking can break up a function over a wider address range. For 64-bit, both problems can be solved parsing the Exception Directory (.pdata section), which is an array of IMAGE_RUNTIME_FUNCTION_ENTRY structures. When an exception occurs, the dispatcher will enumerate this array until it finds the primary function associated with the address of exception, and will use the unwind information to try fix up the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack. You can find more information about x64 exception handling heterotype-page-12 in the stack.

```
typedef struct _IMAGE_RUNTIME_FUNCTION_ENTRY {
    ULONG BeginAddress;
    ULONG EndAddress;
    ULONG UnwindInfoAddress;
} _IMAGE_RUNTIME_FUNCTION_ENTRY, *_PIMAGE_RUNTIME_FUNCTION_ENTRY;
```

The following pseudo-code attempts to describe the idea...

```
func GetGPA
  set addr = 0
```

```
foreach (DLL in PEB) and addr is 0
    foreach runtime in DLL.DataDirectory[Exception] and addr is 0
      set baddr = runtime.BeginAddress
      set start = runtime.BeginAddress + DLL.DllBase
      set end = runtime.EndAddress + DLL.DllBase
      for start to end and addr is 0
       if start[0] == near conditional jump
          set rva = (*(DWORD*)(start + 2) + 6 + start) - DLL.DllBas
         foreach runtime in DLL.DataDirectory[Exception] and addr
            if rva == runtime.BeginAddress
              set start2 = runtime.BeginAddress + DLL.DllBase
              set end2 = runtime.EndAddress + DLL.DllBase
              for start2 to end2
                if start2[0] == STATUS ORDINAL NOT FOUND
                  addr = baddr + DLL.DllBase
                 break
                end if
              end for
            end if
         end foreach
       end if
      end for
    end foreach
 end foreach
  set GetGPA = addr
end func
```

The following code has been tested on 64-bit builds of Windows 7 and Windows 10. GetGPA returned the address of GetProcAddress in both tests.

```
LPVOID GetGPA(VOID) {
    PPEB
                                   peb;
                                  ldr;
    PPEB LDR DATA
    PLDR_DATA_TABLE_ENTRY
                                  dte;
    LPV0ID
                                  addr=NULL;
    BYTE
                                  C;
    PIMAGE_DOS_HEADER
                                  dos;
    PIMAGE NT HEADERS
                                  nt;
    PIMAGE DATA DIRECTORY
                                  dir;
    PIMAGE RUNTIME FUNCTION ENTRY rf;
    DWORD
                                  i, j, h, rva, ba;
    PBYTE
                                  s1, e1, s2, e2;
    PUNWIND_INFO
                                  ui;
    peb = (PPEB) readgsqword(0x60);
    ldr = (PPEB LDR DATA)peb->Ldr;
    for (dte=(PLDR DATA TABLE ENTRY)ldr->InLoadOrderModuleList.Flink;
         dte->DllBase != NULL && addr == NULL;
         dte=(PLDR DATA TABLE ENTRY)dte->InLoadOrderLinks.Flink)
      // is this kernelbase.dll?
      for (h=0, i=0; i<dte->BaseDllName.Length/2; i++) {
        c = (BYTE)dte->BaseDllName.Buffer[i];
        h += (c \mid 0x20);
        h = ROTR32(h, 13);
      // if not, skip it
      if (h != 0x22901A8D) continue;
      dos = (PIMAGE DOS HEADER)dte->DllBase;
```

```
nt = RVA2VA(PIMAGE NT HEADERS, dte->DllBase, dos->e lfanew);
dir = (PIMAGE DATA DIRECTORY)nt->OptionalHeader.DataDirectory;
rva = dir[IMAGE DIRECTORY ENTRY EXCEPTION].VirtualAddress;
rf = (PIMAGE RUNTIME FUNCTION ENTRY) RVA2VA(ULONG PTR, dte->D1
// foreach runtime function and address not found
for(i=0; rf[i].BeginAddress != 0 && addr == NULL; i++) {
 ba = rf[i].BeginAddress;
 // we will search the code between BeginAddress and EndAddres
 s1 = (PBYTE)RVA2VA(ULONG PTR, dte->DllBase, rf[i].BeginAddres
  e1 = (PBYTE)RVA2VA(ULONG PTR, dte->DllBase, rf[i].EndAddress)
 // if chained unwind information is specified in the next ent
 ui = (PUNWIND INFO)RVA2VA(ULONG PTR, dte->DllBase, rf[i+1].Ur
  if(ui->Flags & UNW FLAG CHAININFO) {
   // find the last entry in the chain
   for(;;) {
     i++;
      e1 = (PBYTE)RVA2VA(ULONG PTR, dte->DllBase, rf[i].EndAddi
     ui = (PUNWIND INFO)RVA2VA(ULONG PTR, dte->DllBase, rf[i].
     if(!(ui->Flags & UNW FLAG CHAININFO)) break;
 // for this address range minus the length of a near condition
 while(s1 < (e1 - 6)) {
   // is the next instruction a near conditional jump?
   if(s1[0] == 0x0F \&\& s1[1] >= 0x80 \&\& s1[1] <= 0x8F) {
     // calculate the relative virtual address of jump
      rva = (DWORD)(((*(DWORD*)(s1 + 2)) + 6 + s1) - (PBYTE)dte
     // try find the rva in exception list
     for(j=0; rf[j].BeginAddress != 0 && addr == NULL; j++) {
```

```
if(rf[j].BeginAddress == rva) {
                   s2 = (PBYTE)RVA2VA(ULONG_PTR, dte->DllBase, rf[j].Beg
                   e2 = (PBYTE)RVA2VA(ULONG_PTR, dte->DllBase, rf[j].Enc
                   // try find the error code in this address range
                   while(s2 < (e2 - 4)) {
                     // if this is STATUS_ORDINAL_NOT_FOUND
                     if(*(DWORD*)s2 == 0xC0000138) {
                       // calculate the virtual address of primary funct
                       addr = (PBYTE)RVA2VA(ULONG PTR, dte->DllBase, ba)
                       break;
                     s2++;
            s1++;
      return addr;
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