

# Blog Simple.



## Linux kernel exploit cheatsheet

In Security

Tags exploit, hacking, kernel exploit, linux, linux exploit, linux kernel, security

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## Drive programming

Books: [linux device driver](#)

<https://sysplay.github.io/books/LinuxDrivers/>

Dynamically assigning device numbers

```
1 int alloc_chrdev_region(dev_t *dev, unsigned int firstminor, unsigned int count, char *name);
```

Dev is the outgoing parameter, which is the dynamically obtained device number.

Firstminor specifies the first minor.

Count and name are the same as the register\_chrdev\_region parameter definition.

<https://www.oreilly.com/library/view/linux-device-drivers/0596000081/ch03s02.html>

<http://nanxiao.me/linux-kernel-note-20-device-major-minor-number>

<https://sysplay.github.io/books/LinuxDrivers>

<https://www.kernel.org/doc/Documentation/admin-guide/devices.txt>

Statically initialize character devices:

```
1 struct cdev my_cdev;
```

```
1 cdev_init(&my_cdev, &fops);
```

```
1 my_cdev.owner = THIS_MODULE;
```

[Linux character device driver cdev\\_init\(\) series](#)

Class related api

`class_create, class_register`

## Extract the rootfs in cpio format

Rootfs.cpio is packaged first cpio and then gzipped

Decompression must first change rootfs.cpio to gz suffix and then decompress, otherwise it will report an error.

## exploit tech & tricks

### tty\_struct spray

```
man 4 ptmx
```

Can understand:

```
1 When a process opens /dev/ptmx, it gets a file descriptor for a pseudoterminal master (PTM), and
```

That is to create a new one `tty_struct`, and this structure `tty_struct` has a field `const struct tty_operations* ops` we can rewrite it to an address containing a pointer to a malicious function, such as when there is no `SMAP` time can directly point to the user space, thus controlling the execution flow.

So in the case of uaf, you can spray a lot `tty_struct` to occupy the chunk we released before, then uaf will `ops` change to its own address.

Reference practice : [simple kernel exploit challenge](#)

## uaf use struct cred

And `tty_struct` spary similar, but also accounted for the pit, if `fork` a child process can just make the child process `cred` structure into the position we want, then we can directly overwrite `cred` the contents in order to put right.

Reference practice : [simple kernel exploit challenge](#)

## stack pivot

If the pointer is just called `call rax`, then you can change the pointer to `xchg eax, esp` a gadget, then mmap a memory in the user space `eax` (ie `xchg eax, esp` the address of the gadget and the `0xffffffff` bit and value), and then write the ropchain to the memory, ie It can be `stack pivot`

Eg:

```
1 159. unsigned long lower_address = xchgeaxesp & 0xFFFFFFFF;
2 160. unsigned long base = lower_address & ~0xfff;
```

```

3 161. printf("[+] Base address is %lx\n", base);
4 162. if (mmap(base, 0x30000, 7, MAP_PRIVATE | MAP_ANONYMOUS, -1, 0) != base) {
5 163.     perror("mmap");
6 164.     exit(1);
7 165. }
8 166.
9 167. unsigned long rop_chain[] = {
10      .....
11 180. };
12 181. memcpy((void*)lower_address, rop_chain, sizeof(rop_chain));

```

## Mitigation

### *checksec*

- Kaslr can `cat /proc/cmdline` be viewed by, if the option is open with kaslr, the kernel is not enabled by default.
- Smep can `cat /proc/cpuinfo` check the flags by smep

### *Bypass*

- SMEP smep is turned off by clearing the 20th bit of CR4, usually by the following gadget:

```

1 POP RDI ; RET // Place 000000000000006f0 in RDI MOV CR4 , RDI ;! RET // SMEP disbled # 64- und

```

Reference : [linux-kernel-x86-64-bypass-smep-kaslr-kptr\\_restric](#)

## something

- Rsp in the kernel is 8-byte aligned
- `tty_struct` The magic may be:

```
1 #define TTY_STRUCT_MAGIC 0x5402
2 #define TTY_MAGIC 0x5401
```

## magic number

- `mmap` Generally used when adding `O_NOCTTY`, because:

```
1 The flag O_NOCTTY can tell UNIX that this program will not become the "control terminal" on this
2 If you don't do this, all the input, such as the Ctrl+C abort signal coming from the keyboard, wi
```

- I don't know the size of the Linux kernel structure can compile a module, the module source code is used `sizeof` and then the compiler optimizes the reason, it will directly encode the size, and then `objdump -d` look at the assembly to know the size, but also pay attention to the options.
- Defining kernel functions generally takes the following form (after the function address)

```
1 typedef int __attribute__((regparm(3)))(*commit_creds_func)(unsigned long cred);
2 typedef unsigned long __attribute__((regparm(3)))(*prepare_kernel_cred_func)(unsigned long cred);
```

```
1 commit_creds_func commit_creds = (commit_creds_func)0xffffffff810a1420;
2 prepare_kernel_cred_func prepare_kernel_cred = (prepare_kernel_cred_func) 0xffffffff810a1810 ;
```

The use `regparm` is because the kernel function calling convention is different from the user mode. 32 bits use three registers to pass the first three parameters.

Refer to [gcc function attribute](#).

```
1 Your are probably thinking normal calling convention (arguments on the stack). Modern Linux kernel
```

Reference : [Function parameter passing in a Linux kernel interrupt handler \(from asm to C\)](#)

- Save the state of the user state:

```
1 unsigned long user_cs, user_ss, user_eflags;
```

```
1 void save_stats() {  
2     asm(  
3         "movq %%cs, %0\n"  
4         "movq %%ss, %1\n"  
5         "pushfq\n"  
6         "popq %2\n"  
7         : "=r"(user_cs), "=r"(user_ss), "=r"(user_eflags)  
8         :  
9         : "memory"  
10        );  
11    }
```

## Kernel Debug

### start up

Qemu can directly add parameters `-gdb tcp::23333`, but pay attention to gdb connection server error, so use the `set architecture i386:x86-64` specified framework and then `target remote :23333` connect.

Reference: <https://stackoverflow.com/questions/8662468/remote-g-packet-reply-is-too-long>



## kallsyms

Most of the time we are debugging drivers, and kallsyms has all the symbols in the kernel, including the driver module, so you can view the breakpoints of the module under kallsyms.

## Making cpio, initramfs file system

Need two tools, one is to enter the kernel source code compilation:

```
make -C /usr/src/linux/usr/ gen_init_cpio
```

one is under the kernel source directory script

```
chmod +x usr/gen_init_cpio scripts/gen_initramfs_list.sh
```

Then create the initramfs file system with the following command:

```
1 gen_initramfs_list.sh initrd/ &gt; filelist
2 gen_init_cpio filelist &gt; initrd.img
3 gzip initrd.img
4 mv initrd.img initrd-`uname -r`.img
```

## About cred structure

The first five fields of the default compiled cred are:

```
1 struct cred {  
2     unsigned long usage;  
3     unsigned int uid;  
4     unsigned int gid;  
5     unsigned int south;  
6     unsigned int sgid;  
7     unsigned int euid;  
8     ...  
9 };
```

The test should clear all these fields to get root privileges. The usage field is related to the bit. Under 32 bits `unsigned int`, under 64 bits, the default cred structure is 0xa8 bytes.

## Return user mode

The 64-bit uses the `swapgs` sum `iretq`, the former exchanges the data of gs and MSR, and the latter pops up the data from the stack in the following order:

```
1 the next RIP  
2 user land CS  
3 user land EFLAGS  
4 user land RSP  
5 user land SS
```

## Single file source kernel module compilation

Basic programming of kernel modules and writing of Makefiles

## Get kernel compilation options

Sometimes we want to get the kernel compile time options, such as to get the size of a structure (in this case, there is not enough source code), you can get it by:

At runtime:

```
1 #Current kernel config:
```

```
1 cat /boot/config-`uname -r`
```

```
1 #Other installed kernels:
```

```
1 ls /boot/config-*
```

```
1 #The following three are possible
```

```
1 /proc/config.gz
2 /boot/config
3 /boot/config-$(uname -r)
```

When you have a mirror: this time need to use the kernel source under `scripts/extract-ikconfig` script to get the mirror in the config, but it also requires kernel enabled at compile time `CONFIG_IKCONFIG_PROC` option.

## Extract vmlinux from the compressed image

It `scripts/extract-vmlinux` can be easily extracted using the kernel source : `./extract-vmlinux bzImage > vmlinx`

## linux kernel

### stack

Various stacks in Linux: process stack thread stack kernel stack interrupt stack

### tty/pty/ptmx

Under Linux `tty` / `pty` / `pts` / `ptmx` Detailed `ptmx` and analysis terminal when the relationship is created `pts`

## Call convention 64 bit

32-bit can refer to the above

```
1 A.2 AMD64 Linux Kernel Conventions
```

```
1 1. User-level applications use as integer registers for passing the sequence
2 %rdi, %rsi, %rdx, %rcx, %r8 and %r9. The kernel interface uses %rdi,
3
```

```
4 %rsi, %rdx, %r10, %r8 and %r9.
```

```
1 2. A system-call is done via the syscall instruction. The kernel destroys  
2 registers %rcx and %r11.
```

```
1 3. The number of the syscall has to be passed in register %rax.  
2 4. System-calls are limited to six arguments, no argument is passed directly on  
3 the stack.
```

```
1 5. Returning from the syscall, register %rax contains the result of the  
2 system-call. A value in the range between -4095 and -1 indicates an error,  
3 it is -errno.
```

```
1 6. Only values of class INTEGER or class MEMORY are passed to the kernel.
```

## source

free-electrons

## I don't understand for a while

It found that watching someone else rewrite wp `tty_struct` of `ops` even apply for a piece of memory pointer after the write pointer `ops` points to `proc_fops`:

```
1 struct tty_operations* fake_tty_operations = (struct tty_operations*) calloc(1, sizeof(struct tty  
2 void *fake_file_operations = calloc(1, 0x1000);
```

```
1 fake_tty_operations->proc_fops = fake_file_operations;  
2 fake_tty_operations->ioctl = (int (*)( ))xchg_eax_esp_ret;
```

I don't understand the intention for the time being...

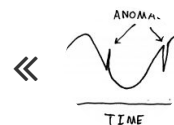
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Writing kernel exploits



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## About Me



My name is Nguyen Anh Tai. I am an independent security researcher, bug hunter and leader a security team. Security Researcher at CMC INFOSEC. I developed the every system for fun :D. My aim is to become an expert in security and xxx!

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