

**Data Structures and Algorithms
(ES221)**

Counting and Radix Sort

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Counting Sort



- Suppose that there are N integers to be sorted, $A[1 \cdots N]$, such that the range of the integers is from 1 to M
- Start by defining an array $B[1 \cdots M]$ and initializing all its elements to 0. $\rightarrow (O(M))$
- Scan through $A[i]$, for $i = 1, \cdots, N$, and copy each $A[i]$ into $B[A[i]]$. $\rightarrow (O(N))$
- Scan through $B[i]$, for $i = 1, \cdots, M$, and retrieve all the non-zero elements of $B[i]$. $\rightarrow (O(M))$
- Complexity: $O(M + N)$
 - If $M \sim N \rightarrow$ Complexity: $O(N)$
 - If $M \sim N^2 \rightarrow$ Complexity: $O(N^2)$
- Example
 - Sort 5 7 9 2 4 6 1

Counting Sort



A

5 (0)	7 (1)	9 (2)	2 (3)	4 (4)	6 (5)	1 (6)
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$$B[A[i]] = A[i]$$

B

(0)	(1)	(2)	(3)	(4)	5 (5)	(6)	(7)	(8)	(9)
-----	-----	-----	-----	-----	-------	-----	-----	-----	-----

$$B[A[0]] = A[0]$$

$$B[5] = 5$$

(0)	(1)	(2)	(3)	(4)	5 (5)	(6)	7 (7)	(8)	(9)
-----	-----	-----	-----	-----	-------	-----	-------	-----	-----

$$B[7] = 7$$

(0)	(1)	(2)	(3)	(4)	5 (5)	(6)	7 (7)	(8)	9 (9)
-----	-----	-----	-----	-----	-------	-----	-------	-----	-------

$$B[9] = 9$$

(0)	(1)	2 (2)	(3)	(4)	5 (5)	(6)	7 (7)	(8)	9 (9)
-----	-----	-------	-----	-----	-------	-----	-------	-----	-------

$$B[2] = 2$$

(0)	(1)	2 (2)	(3)	4 (4)	5 (5)	(6)	7 (7)	(8)	9 (9)
-----	-----	-------	-----	-------	-------	-----	-------	-----	-------

$$B[4] = 4$$

(0)	(1)	2 (2)	(3)	4 (4)	5 (5)	6 (6)	7 (7)	(8)	9 (9)
-----	-----	-------	-----	-------	-------	-------	-------	-----	-------

$$B[6] = 6$$

(0)	1 (1)	2 (2)	(3)	4 (4)	5 (5)	6 (6)	7 (7)	(8)	9 (9)
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$$B[1] = 1$$

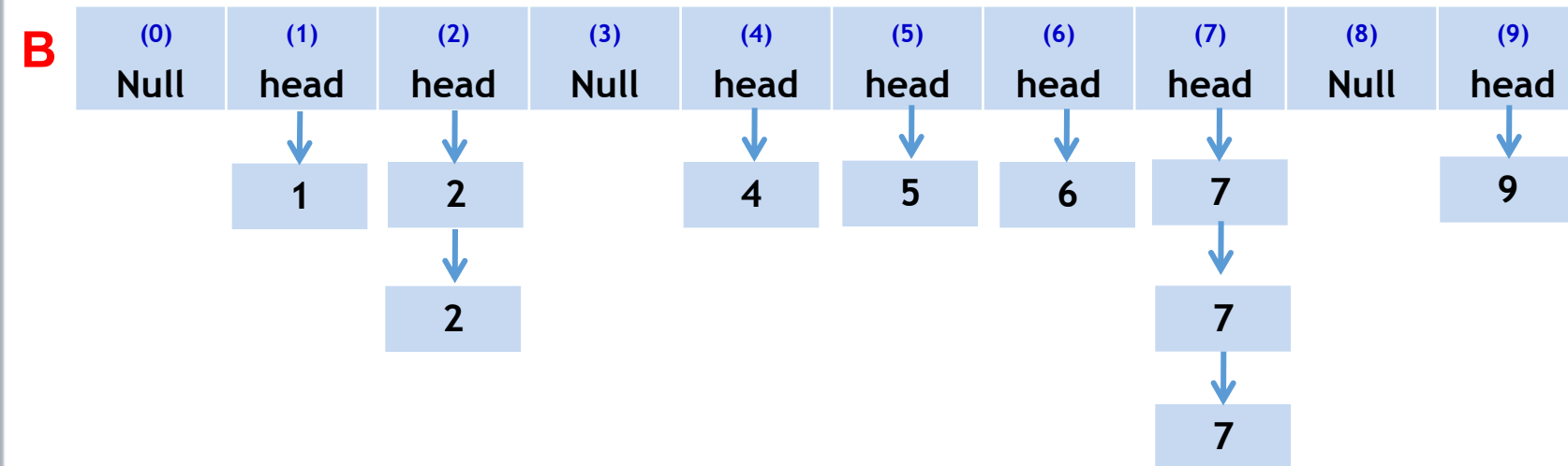
Final Output: 1 2 4 5 6 7 9

Counting Sort



- How to cope up with duplicates?
- B becomes a an array of pointers, (rather than an array of integers). Each element has a
 - head pointer (points to the head of a linked list)
 - tail pointer (points to the tail of the linked list)
- $A[j]$ is added at the tail of the list $B[A[j]]$
- Finally, the array B is sequentially traversed and each nonempty list retrieved.
- Complexity: $O(M + N)$

Counting Sort – Duplicate keys



Final Output: 1 2 2 4 5 6 7 7 7 9

Radix Sort



- Counting sort becomes expensive if **M** is large compared to **N**
- Radix sort may be built upon counting sort
- **Basic Idea:** every integer can be represented by at most k digits
 - $d_1d_2\cdots d_k$ where d_i are digits in base r



Radix Sort



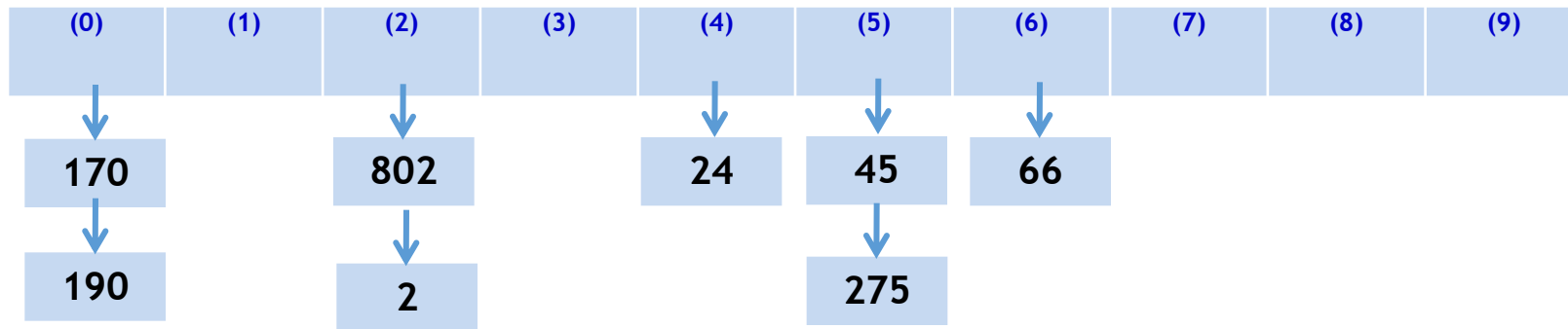
- Algorithm
 - Take the least significant digit (or group of bits) of each key.
 - Group the keys based on that digit, but otherwise keep the original order of keys
 - Repeat the grouping process with each more significant digit.
 - The sort in step 2 is usually done using **bucket sort** or **counting sort**, which are efficient in this case since there are usually only a small number of digits.

Radix Sort



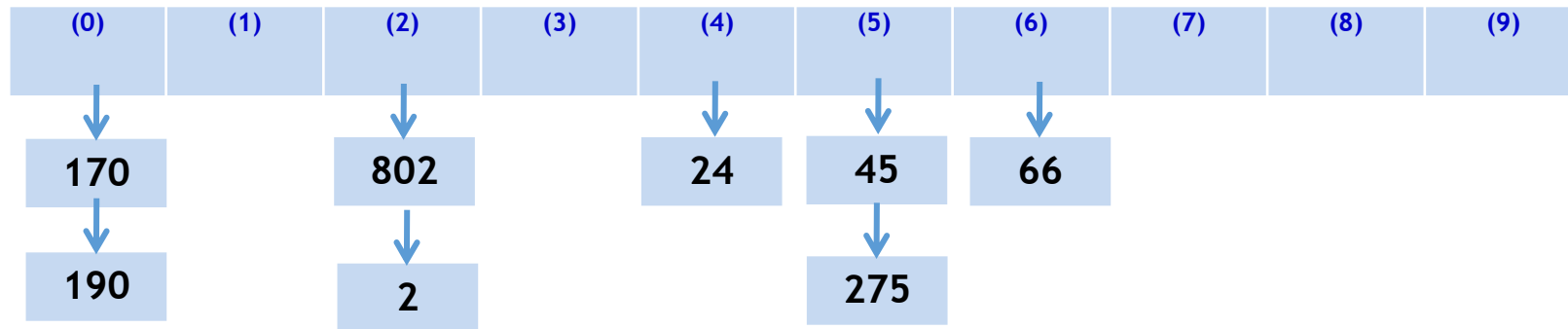
- Example
- Sort: 170, 45, 275, 190, 802, 24, 2, 66

- **First pass**

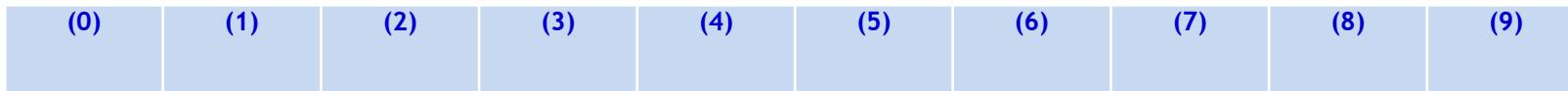


Radix Sort

- **First pass**

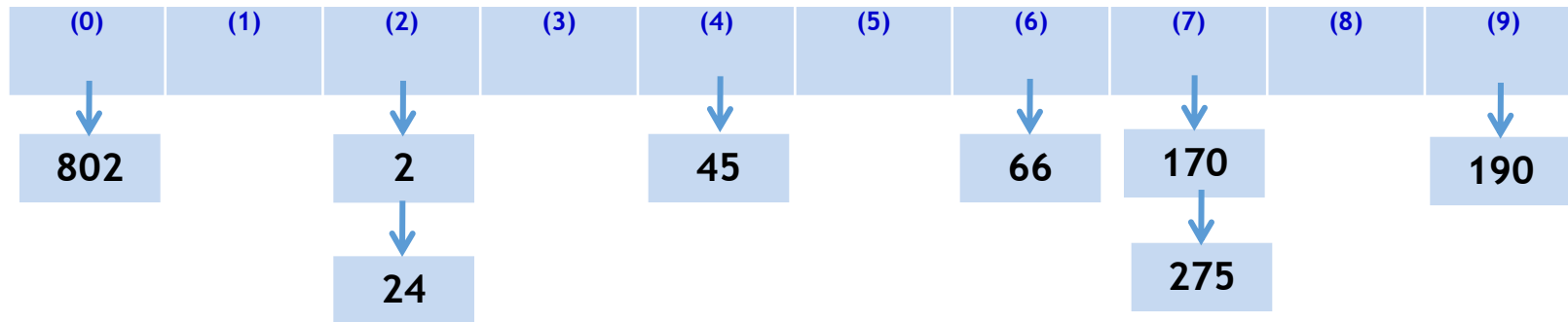


- **Second pass**

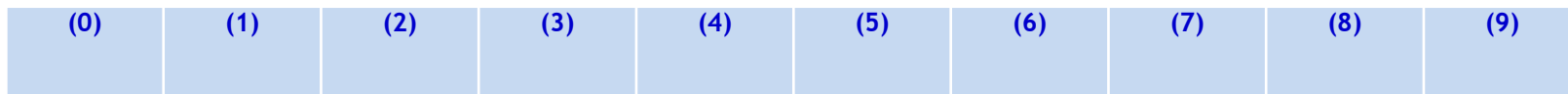


Radix Sort

- **Second pass**



- **Third pass**

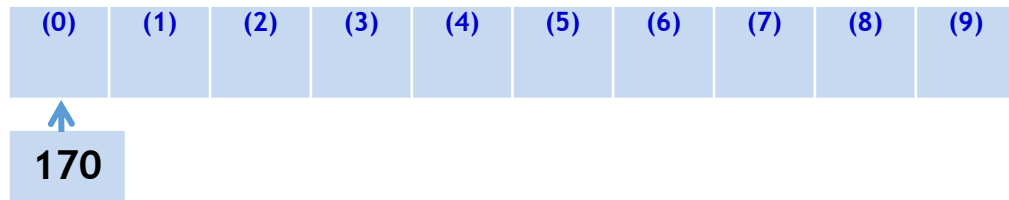


Final Output: 2 24 45 66 170 190 275 802

Radix Sort

170, 45, 275, 190, 802, 24, 2, 66

- Algorithm RadixSort(A, N, d)



for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

$D = 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10)$

enqueue($A[i]$, $Q[t]$)

$j = 0$

for $p = 0$ to N

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$

$i = 0$

$t = (A[0] \bmod 10) \div (1)$
 $= 170 \bmod 10 = 0$

enqueue(170, $Q[0]$)

Radix Sort

170, 45, 275, 190, 802, 24, 2, 66

- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

$D = 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10)$

enqueue($A[i]$, $Q[t]$)

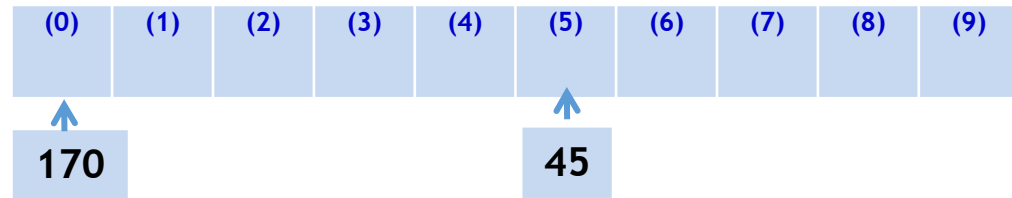
$j = 0$

for $p = 0$ to N

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



$i = 1$

$t = (A[1] \bmod 10) \div (1)$
 $= 45 \bmod 10 = 5$

enqueue(45, $Q[5]$)

Radix Sort

170, 45, 275, 190, 802, 24, 2, 66

- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10)$

enqueue($A[i]$, $Q[t]$)

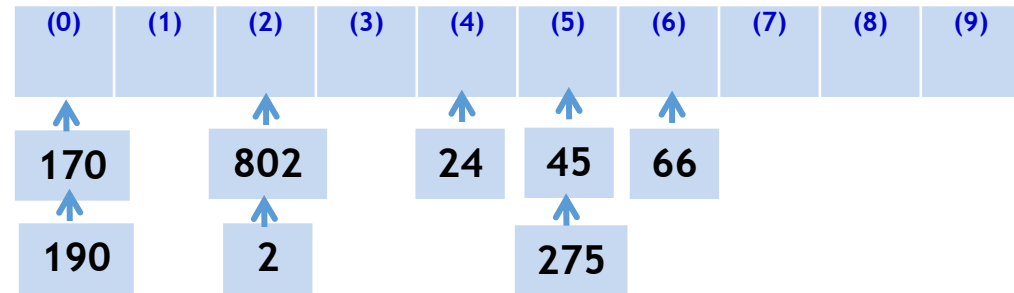
$j = 0$

for $p = 0$ to 9

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



$D = 10$

$i = 2, 4, 5, 6, 7, 8, 9$

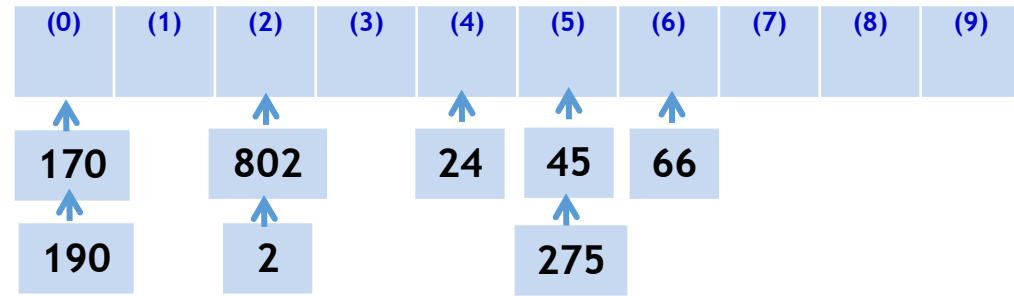
$t = (A[2] \bmod 10) \div (1)$

$= 275 \bmod 10 = 5$

enqueue(245, $Q[5]$)

Radix Sort

170, 45, 275, 190, 802, 24, 2, 66



D = 10

- Algorithm RadixSort(A, N, d)

for p = 0 to 9

Q[p] = Null

D = 1

for k = 1 to d

D = D * 10

for i = 0 to N

t = (A[i] mod D) div (D/10)

enqueue(A[i], Q[t])

j = 0

for p = 0 to N

do while Q[p] is not empty

A[j] = dequeue(Q[p])

j = j + 1

A = { 170, 190, 802, 2, 24, 45, 275, 66 }

Radix Sort

$A = \{ 170, 190, 802, 2, 24, 45, 275, 66 \}$



- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \text{ div } (D/10)$

enqueue($A[i]$, $Q[t]$)

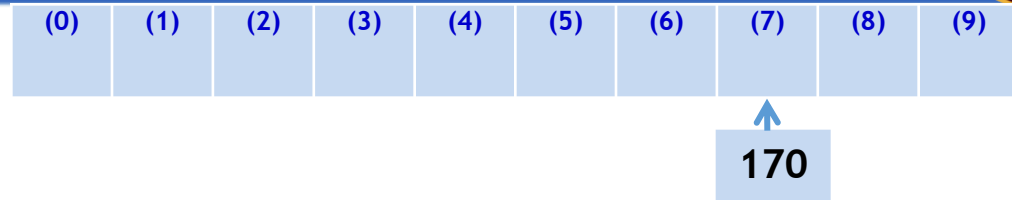
$j = 0$

for $p = 0$ to N

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



$D = 100$

$i = 0$

$t = (A[0] \bmod 100) \text{ div } (10)$

$= (170 \bmod 100) \text{ div } (10) = 70 / 10 = 7$

enqueue(170, $Q[7]$)

Radix Sort

$A = \{ 170, 190, 802, 2, 24, 45, 275, 66 \}$



- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10)$

enqueue($A[i]$, $Q[t]$)

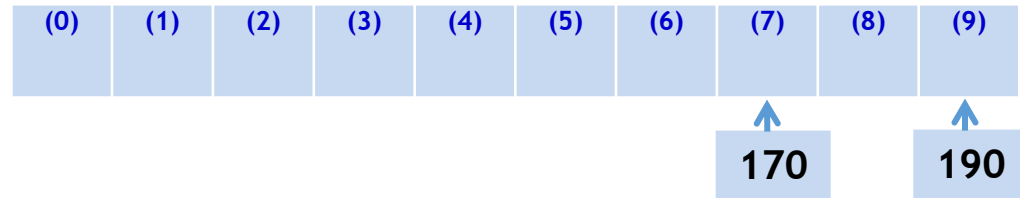
$j = 0$

for $p = 0$ to N

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



$D = 100$

$i = 1$

$t = (A[1] \bmod 100) \div (100/10)$

$= (190 \bmod 100) \div (100/10) = 90/10 = 9$

enqueue(190, $Q[9]$)

Radix Sort

$A = \{ 170, 190, 802, 2, 24, 45, 275, 66 \}$



- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10)$

enqueue($A[i], Q[t]$)

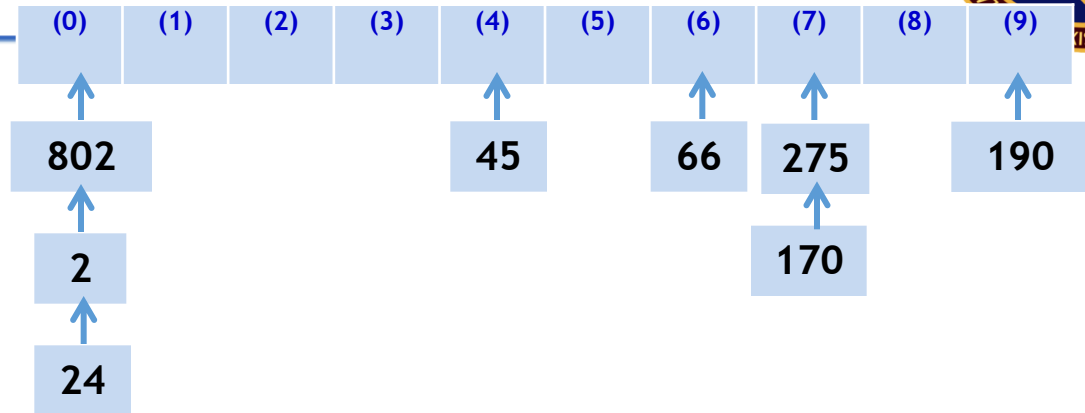
$j = 0$

for $p = 0$ to N

do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



$D = 100$

$i = 2, 3, 4, 5, 6, 7, 8, 9$

$A = \{ 802, 2, 24, 45, 66, 275, 170, 190 \}$

Radix Sort

$A = \{802, 2, 24, 45, 66, 275, 170, 190\}$



- Algorithm RadixSort(A, N, d)

for $p = 0$ to 9

$Q[p] = \text{Null}$

$D = 1$

for $k = 1$ to d

$D = D * 10$

 for $i = 0$ to N

$t = (A[i] \bmod D) \div (D/10) = (802 \bmod 1000) \div (1000/10) = 802/100 = 8$

 enqueue($A[i], Q[t]$)

$j = 0$

 for $p = 0$ to 9

 do while $Q[p]$ is not empty

$A[j] = \text{dequeue}(Q[p])$

$j = j + 1$



802

$D = 1000$

$i = 0$

$t = (A[0] \bmod 1000) \div (1000/10)$

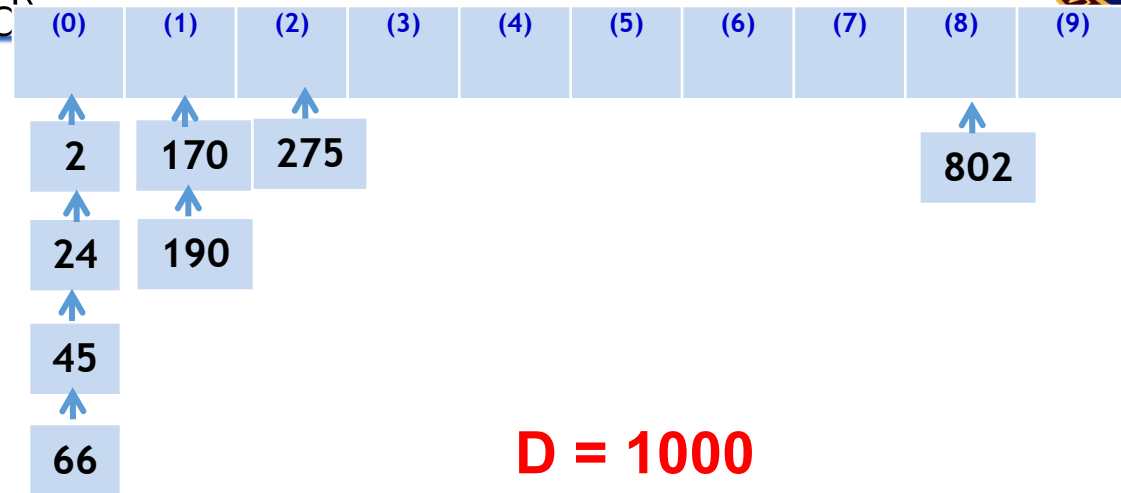
$= (802 \bmod 1000) \div (1000/10) = 802/100 = 8$

enqueue(802, $Q[8]$)

Radix Sort



A = {802, 2, 24, 45, 66, 275, 170, 190}



D = 1000

i = 1, 2, 3, 4, 5, 6, 7, 8, 9

Final Output

A = {2, 24, 45, 66, 170, 190, 275, 802}

Questions?

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