

Secure Software Design and Engineering (CY-321)

Code Obfuscation

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Informally, to obfuscate a program P means to transform it into a program P' that is still executable but for which it is hard to extract information.

"Hard?" ⇒ Harder than before!



static obfuscation ⇒ obfuscated programs that remain fixed at runtime.

dynamic obfuscators ⇒ transform programs continuously at runtime, keeping them in constant flux.

Obfuscating: Expression Equivalence



$$x+0\rightarrow x$$

$$x+0\rightarrow x$$

$$(a-(-b))$$

$$a\times b$$

$$eln(a)+ln(b)$$

Obfuscating: Expression Equivalence



$$y = x + 1$$

temp =
$$x * 2 - x$$

y = temp + 1

$$y = (x + 5) - 3;$$

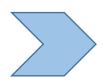


$$y = x + (5 - 3);$$

Obfuscating: Expression Equivalence



$$y = x * 42;$$





```
int compute(int x, int y) {
int result = (x * y) + (x - y);
    return result;
    }
```



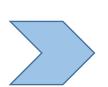
```
int multiply(int a, int b) {
    return a * b;
    }

int difference(int a, int b) {
    return a - b;
    }

int compute(int x, int y) {
    int part1 = multiply(x, y);
    int part2 = difference(x, y);
    return part1 + part2;
    }
```



```
int modexp ( int y , int x []
, int w , int n ) {
   int R , L ;
   int k = 0;
   int s = 1;
   while (k < w) {
       f (x[k],s,y,n
, & R ); // Call to `f`
       s = R * R % n ; //
Square the result
       L = R;
       k ++;
   return L;
```



```
void f ( int xk , int s , int y , int
n , int * R) {
   if ( xk == 1)
      * R = ( s* y ) % n;
   else
      * R = s;
}
```

Splitting



```
int modexp ( int y , int x []
, int w , int n ) {
   int R , L ;
   int k = 0;
   int s = 1;
   while (k < w) {
       f (x[k],s,y,n
, & R ); // Call to `f`
       s = R * R % n ; //
Square the result
       L = R;
       k ++;
   return L;
```

It takes four parameters:

- •y: the base
- •x[]: an array of binary exponent values
- •w: the number of bits in the exponent
- •n: the modulus

It initializes:

- •k = 0: index for looping through x[]
- •s = 1: stores intermediate results
- •L: stores the final computed value

Inside the while loop:

- •Calls f(x[k], s, y, n, &R), which modifies R
- •Computes s = R * R % n
- •Stores R in L
- •Increments k
- •When k == w, L holds the result and is returned.



The function f takes the following parameters:

- •xk: A single bit from an exponent array (x[]) in the original modexp function.
- •s: An intermediate computation value (used in modular exponentiation).
- •y: The base of exponentiation.
- •n: The modulus for modular exponentiation.
- •R: A pointer to store the result.

```
void f ( int xk , int s , int y , int n , int * R)
{
   if ( xk == 1)
     * R = ( s* y ) % n;
   else
     * R = s;
}
```



```
int add(int a, int b) {
    return a + b;
    }

int multiply(int a, int b) {
    return a * b;
    }
```





```
float foo[100];
void f(int a, float b) {
    foo[a] = b; // Stores `b` at
index `a` in the array
float g(float c, char d) {
    return c * (float)d; //
Multiplies `c` with `d`
(converted to float)
int main() {
                  // Calls
    f(42, 42.0);
`f()
   float v = g(6.0, 'a'); //
Calls `g()`
```



Merging

```
float foo[100];
float fg(int a, float bc, char d,
int which) {
    if (which == 1)
       foo[a] = bc; // If
`which == 1`, perform `f()` logic
(store in array)
    return bc * (float)d; //
Always return `bc * d`, mimicking
`g()`
int main() {
   fg(42, 42.0, 'b', 1); //
Equivalent to calling `f(42,
42.0)
   float v = fg(99, 6.0, 'a',
2); // Equivalent to g(6.0)
'a')`
```

Obfuscating: Copying Code



Make the program larger by cloning pieces of it:



Make the copied code look different from the original:

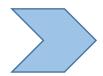


Now the attacker must examine all pairs of code blocks to see which ones are the same

Obfuscating: Copying Code



```
int compute(int x) {
    return (x * 2) + 5;
}
```



Makes it harder to recognize identical logic.

Code analysis tools may **fail to detect function similarity**.

```
int compute(int x) {
    int temp = x * 2;
    int y = temp + 5;
    return y;
}

int compute_alternative(int x) { //
Same logic but slightly changed
structure
    int y = (x * 2);
    return y + 5;
}
```

Obfuscating: Interpretation



Add a level of interpretation

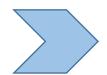
- 1 Define your own instruction set
- 2 Translate your program to this instruction set
- 3 Write an interpreter for the instruction set

Your program: 10-100x slower than before.

Obfuscating: Interpretation



```
int compute(int x) {
    return (x * 2) + 5;
}
```



```
int interpretBytecode(int *bytecode, int
x) {
    int result = 0;
    for (int i = 0; i < 3; i++) { // Loop
over bytecode instructions
        if (bytecode[i] == 1) result = x *
2; // Instruction 1 = Multiply by 2
        else if (bytecode[i] == 2) result
+= 5; // Instruction 2 = Add 5
    return result;
int main() {
    int bytecode[] = {1, 2}; // Encoded
sequence of operations
    int result =
interpretBytecode(bytecode, 10);
    printf("%d\n", result);
    return 0;
```



Opaque values: These are values that appear unpredictable to an adversary but are known at compile or runtime.

Array aliasing: This refers to using multiple references (aliases) to the same memory location in an array, making it difficult to track actual variable values.



Aliasing occurs in two pointers can refer to the same memory location

Two reference parameters can also alias each other

A reference parameter and a global variable

Two array elements indexed by different variables



Two Pointers Referring to the Same Memory Location

```
#include <stdio.h>
int main() {
   int num = 42;
   int *ptr1 = #
    int *ptr2 = ptr1; // Both ptr1 and ptr2 point
to 'num'
    *ptr2 = 99; // Changing value using ptr2
   printf("%d\n", *ptr1); // Output: 99 (ptr1 also
sees the change)
   return 0;
```



Two Reference Parameters Can Also Alias Each Other

```
#include <iostream>
void modify(int &a, int &b) {
    a = 10; // Modifies b as well if they alias
each other
    b = 20;
int main() {
    int x = 5;
    modify(x, x); // x is passed twice, creating
aliasing
    std::cout << x << std::endl; // Output: 20</pre>
(last modification applies)
    return 0;
```



A Reference Parameter and a Global Variable

```
#include <iostream>
int globalVar = 50;
void modifyGlobal(int &param) {
    param = 100; // Modifies globalVar because
param is an alias
int main() {
    modifyGlobal(globalVar);
    std::cout << globalVar << std::endl; // Output:</pre>
100
    return 0;
```



Two Array Elements Indexed by Different Variables

```
#include <stdio.h>
int main() {
    int arr[5] = \{1, 2, 3, 4, 5\};
    int *ptr1 = &arr[2]; // Points to arr[2]
    int *ptr2 = arr + 2; // Also points to arr[2]
(same memory location)
    *ptr1 = 99; // Modify arr[2]
    printf("%d\n", *ptr2); // Output: 99
    return 0;
```



Two Array Elements Indexed by Different Variables

```
#include <stdio.h>
int main() {
    int arr[5] = \{1, 2, 3, 4, 5\};
    int i = 1, j = 2; // Two different index
variables
    int *ptr1 = &arr[i]; // Points to arr[1]
    int *ptr2 = &arr[j - 1]; // Also points to
arr[1] (j - 1 = 1)
    *ptr1 = 99; // Modify arr[1] using ptr1
    printf("%d\n", *ptr2); // Output: 99 (because
ptr2 also points to arr[1])
    return 0;
```



```
#include <stdio.h>
int main() {
    int arr[3] = \{42, 99, 7\};
    int *ptr = arr; // Pointer
aliasing the array
    int x = ptr[1]; // Opaque value
from aliasing
    if (x == 99) {
        printf("Secret branch!\n");
    } else {
        printf("Normal branch\n");
```

Instead of directly accessing arr[1], we retrieve its value through a pointer alias (ptr[1]).

The value 99 (retrieved via aliasing) determines whether the conditional executes.

Reverse engineers trying to analyze the control flow statically may not recognize that x == 99 is always true.

Instead of writing if (x == 99), the program retrieves 99 dynamically from the array



```
#include <stdio.h>
int main() {
    int arr[3] = \{10, 20, 30\};
    int *alias1 = arr;
    int *alias2 = alias1 + 1; //
Points to arr[1]
    int x = *alias2; // Opaque value
from aliasing
    if (x == 20) {
        printf("Access Granted!\n");
    } else {
        printf("Access Denied!\n");
```

Obfuscating: Renaming



```
int add(int a, int b)
{
    return a + b;
}

int main() {
    int sum = add(5,
10);
    return sum;
}
```



```
int x1(int x2, int x3) {
    return x2 + x3;
}
int main() {
    int x4 = x1(5, 10);
    return x4;
}
```

Obfuscating: Code Flattening



```
void process(int x) {
    if (x > 10)

printf("High\n");
    else

printf("Low\n");
}
```



```
void process(int x) {
    int state = (x > 10) ? 1 : 0;

    switch (state) {
        case 1:
            printf("High\n");
            break;
        case 0:
            printf("Low\n");
            break;
    }
}
```

Obfuscating: Junk Code Insertion



```
int add(int a, int b)
{
    return a + b;
}
```



```
int add(int a, int b) {
   int noise = 42;
   noise += 10;
   return a + b;
}
```



Questions??

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