

[🔗 Previous \(/articles/binary-tree-maximum-path-sum/\)](/articles/binary-tree-maximum-path-sum/) Next [🔗 \(/articles/smallest-string-starting-from-leaf/\)](/articles/smallest-string-starting-from-leaf/)

41. First Missing Positive [🔗 \(/problems/first-missing-positive/\)](/problems/first-missing-positive/)

Feb. 2, 2019 | 16.9K views

Average Rating: 4.42 (26 votes)

Given an unsorted integer array, find the smallest missing positive integer.

Example 1:

Input: [1,2,0]
Output: 3

Example 2:

Input: [3,4,-1,1]
Output: 2

Example 3:

Input: [7,8,9,11,12]
Output: 1

Note:

Your algorithm should run in $O(n)$ time and uses constant extra space.

Solution

Approach 1: Index as a hash key.

Data clean up
 Articles > 41. First Missing Positive ▼

First of all let's get rid of negative numbers and zeros since there is no need of them. One could get rid of all numbers larger than n as well, since the first missing positive is for sure smaller or equal to $n + 1$. The case when the first missing positive is equal to $n + 1$ will be treated separately.

number of elements is

$$n = 8$$

[1, 2, 3, 4, 5, 6, 7, 8] --> 9

[1, 2, 48, 14, 15, 16, 17, 18] --> 3

[1, 2, 3, 4, 5, 6, 7, 18] --> 8

max possible first missing number is

$$n + 1 = 9$$

What does it mean - to get rid of, if one has to keep $\mathcal{O}(N)$ time complexity and hence could not pop unwanted elements out? Let's just replace all these by 1s.

Data clean up : replace by 1s :

- negative numbers

- zeros

- numbers larger than $n = 10$

[3, 4, **-1**, **-2**, 1, 5, **16**, **0**, 2, 0] -->

[3, 4, **1**, **1**, 1, 5, **1**, **1**, 2, 1]

To ensure that the first missing positive is not 1, one has to verify the presence of 1 before proceeding to this operation.

How to solve in-place

Now there we have an array which contains only positive numbers in a range from 1 to n , and the problem is to find a first missing positive in $\mathcal{O}(N)$ time and constant space.

That would be simple, if one would be allowed to have a hash-map `positive number -> its presence` for the array.

 $\mathcal{O}(N)$ space complexity solution

with hash-map

Articles > 41. First Missing Positive ▼

`[3, 4, 1, 1, 1, 5, 1, 1, 2, 1] -->``{1: 6, 2: 1, 3: 1, 4: 1, 5: 1, 6: missing}`

Sort of "dirty workaround" solution would be to allocate a string `hash_str` with `n` zeros, and use it as a sort of hash map by changing `hash_str[i]` to `1` each time one meets number `i` in the array.

"O(1) space complexity" solution
with string

`[3, 4, 1, 1, 1, 5, 1, 1, 2, 1] -->`

number 6 is
missing
↓
"1111100000"
↑
number 5 is
present

Let's not use this solution, but just take away a pretty nice idea to use *index* as a *hash-key* for a positive number.

The final idea is to use *index in nums* as a *hash key* and *sign of the element* as a *hash value* which is presence detector.

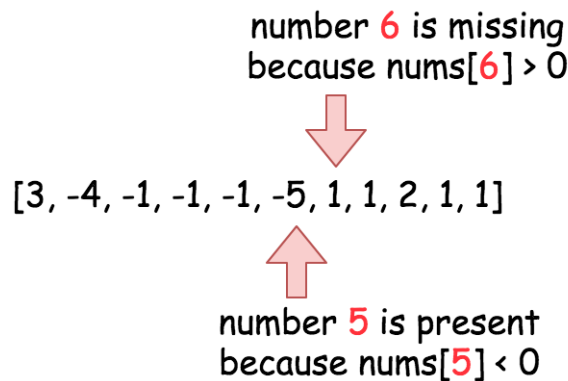
For example, negative sign of `nums[2]` element means that number `2` is present in `nums`.
The positive sign of `nums[3]` element means that number `3` is not present (missing) in `nums`.

To achieve that let's walk along the array (which after clean up contains only positive numbers),

check each element value `elem` and change the sign of element `nums[elem]` to negative to mark that number `elem` is present in `nums`. Be careful with duplicates and ensure that the sign was changed only once.

$O(1)$ space complexity solution

`[3, 4, 1, 1, 1, 5, 1, 1, 2, 1]` -->



Algorithm

Now everything is ready to write down the algorithm.

- Check if 1 is present in the array. If not, you're done and 1 is the answer.
- If `nums = [1]`, the answer is 2.
- Replace negative numbers, zeros, and numbers larger than `n` by 1s.
- Walk along the array. Change the sign of `a`-th element if you meet number `a`. Be careful with duplicates : do sign change only once. Use index 0 to save an information about presence of number `n` since index `n` is not available.
- Walk again along the array. Return the index of the first positive element.
- If `nums[0] > 0` return `n`.
- If on the previous step you didn't find the positive element in `nums`, that means that the answer is `n + 1`.

Implementation

[3, 4, -1, -2, 1, 5, 16, 0, 2, 0]

Articles 1. Check if 1 is present in nums : yes

[3, 4, 1, 1, 1, 5, 1, 1, 2, 1]

number 3 is present
make nums[3] negative

[3, -4, -1, -1, -1, -5, 1, 1, 2, 1]

number 4 is present
make nums[4] negative

[3, -4, -1, -1, -1, -5, 1, 1, 2, 1]

number 6 is the first
missing positive because
nums[6] > 0

2. Data clean up: replace negative numbers, zeros and numbers larger than n = 10 by 1s.

3. Walk along the array. Change the sign of ath element if you meet number a. Be careful with duplicates.

4. Return the index of the first positive element : return 6



4 / 4

Java

Python

Copy

```

1 class Solution:
2     def firstMissingPositive(self, nums):
3         """
4         :type nums: List[int]
5         :rtype: int
6         """
7         n = len(nums)
8
9         # Base case.
10        if 1 not in nums:
11            return 1
12
13        # nums = [1]
14        if n == 1:
15            return 2
16
17        # Replace negative numbers, zeros,
18        # and numbers larger than n by 1s.
19        # After this conversion nums will contain
20        # only positive numbers.
21        for i in range(n):
22            if nums[i] <= 0 or nums[i] > n:
23                nums[i] = 1
24
25        # Use index as a hash key and number sign as a presence detector.
26        # For example, if nums[1] is negative that means that number `1`
27        # is present in the array.

```

Complexity Analysis

- Time complexity : $O(N)$ since all we do here is four walks along the array of length N .
- Space complexity : $O(1)$ since this is a constant space solution.

Analysis written by @liaison (<https://leetcode.com/liaison/>) and @andvary (<https://leetcode.com/andvary/>)

Rate this article:

Previous (/articles/binary-tree-maximum-path-sum/)

Next (/articles/smallest-string-starting-from-leaf/)

Comments: 16

Sort By ▼



Type comment here... (Markdown is supported)

Preview

Post



(/coder_xyzyy)

coder_xyzyy (coder_xyzyy) ★ 42 ⌚ March 27, 2019 5:19 PM

Nice solution, but it's possible to do this in one pass, considering each value at most twice, and not destroying the contents of the array (though changing its order).

```
class Solution {  
public:
```

Read More

10 ^ v ⌂ Share ↩ Reply

SHOW 3 REPLIES



(/totsubo)

totsubo (totsubo) ★ 49 ⌚ June 24, 2019 5:32 PM

One could get rid of all numbers larger than n as well, since the first missing positive is for sure smaller or equal to $n + 1$

This isn't obvious and could use some more explanation.

Read More

2 ^ v ⌂ Share ↩ Reply



(/manchesterunited)

manchesterunited (manchesterunited) ★ 4 February 2, 2019 8:27 PM
 Articles > 41. First Missing Positive
 And ArrayIndexOutOfBoundsException while processing {1,3,3};

Should:
 else if (nums[a] > 0)
 nums[a] *= -1;

[Read More](#)

2 ^ v | Share | Reply

[SHOW 4 REPLIES](#)

(/manchesterunited)

manchesterunited (manchesterunited) ★ 4 February 2, 2019 7:53 PM

Should the following statement:
 If array contains only one element and it's **not** 1, the answer is 2.
 be:
 If array contains only one element and it's 1, the answer is 2.
 ?

[Read More](#)

2 ^ v | Share | Reply

[SHOW 2 REPLIES](#)

(/dloewenherz)

dloewenherz (dloewenherz) ★ 1 March 19, 2019 5:47 PM

Why would you use a hash for this when an array has O(1) lookups and would be a much cleaner solution? You could just use what you're using above as "hash value" as the array index.

```
class Solution(object):
```

[Read More](#)

1 ^ v | Share | Reply

[SHOW 1 REPLY](#)

(/kramer)

kramer (kramer) ★ 9 July 30, 2019 1:25 PM

Explained with more comments

```
int firstMissingPositive(vector& nums) {  
    int one = 0;  
    for(int i=0; i<nums.size(); i++){//check for 1s
```

[Read More](#)

0 ^ v | Share | Reply



(/fabcohen)

fabcohen (fabcohen) ★ 0 ⌚ May 10, 2019 12:58 PM
Articles > 41. First Missing Positive ▾

#include algorithm

```
int MissingInteger(vector &A)
{
    int N=A.size();
    int ret=1;
```

[Read More](#)0 ▲ ▼ | [Share](#) | [Reply](#)

(/harvey2015)

harvey2015 (harvey2015) ★ 2 ⌚ April 2, 2019 11:25 AM

two pointers

for each position i, find if i+1 exists in the array
j increases from i to len(nums), for each step it tries to find i+1 by swapping
every element is at most swapped 1 time, so o(n) time complexity

[Read More](#)0 ▲ ▼ | [Share](#) | [Reply](#)

(/michaelyta)

michaelyta (michaelyta) ★ 2 ⌚ March 10, 2019 3:01 PM

No need to have a special case where if i = n to store info at 0, you can just store info at i-1 and return i+1 when you iterate later. Will result in a bit simpler code with fewer if conditions.

```
def firstMissingPositive(self, nums: List[int]) -> int:
```

[Read More](#)0 ▲ ▼ | [Share](#) | [Reply](#)

(/hailcaesar)

hailcaesar (hailcaesar) ★ 4 ⌚ February 20, 2019 7:57 PM

My answer is not being accepted due to "memory limit" being exceeded. However, it does use O(1) space IMO. What do you guys think?

An int is O(1) in my view..

[Read More](#)0 ▲ ▼ | [Share](#) | [Reply](#)**SHOW 1 REPLY**

< 1 2 >