

## Assessment 2

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### 1 Problem 10.38

(a)

$aRb \Rightarrow L(a) \subset L(b)$ : If  $aRb$ , by transitivity every  $x \in A$  means that  $xRa$  and also  $xRb$ . Consequently,  $x \in L(b)$ , which means  $L(a) \subseteq L(b)$ . Also,  $b \in L(b)$  but  $b \notin L(a)$ , which means  $L(a) \subset L(b)$ .

$L(a) \subset L(b) \Rightarrow aRb$ : If  $L(a) \subset L(b)$ ,  $a \in L(a)$  and also  $a \in L(b)$ , which means  $aRb$ .

In conclusion, the statement(10.16) holds.

(b) PROOF:

If  $L(a) = L(b)$ , then for every  $x \in L(a)$  also  $x \in L(b)$ . When  $x = a$ , there is  $a \in L(a)$  also  $a \in L(b)$ . If  $a \in L(b)$  and  $a \neq b$ , there must be  $aRb$ . For the same reason, when  $x = b$  and  $a \neq b$ , there must be  $bRa$ . But  $R$  is asymmetric, so it is a contradiction. So there must be  $a = b$ .

(c) Let  $A = \{1,2,3\}$   $aRb ::= a$  is a subset of  $b$ . Let  $a = \{1,2\}$  and  $b = \{2,3\}$ , then  $L(a) = L(b) = \{\{1,2,3\}\}$ . But  $a \neq b$ , which means the conclusion of part(b) would not hold in this case.

### 2 Problem 14.37

PROOF: We know that  $f = O(g)$  and  $g = O(f)$ , which means:

$$\limsup_{x \rightarrow \infty} \frac{f(x)}{g(x)} < \infty, \quad \limsup_{x \rightarrow \infty} \frac{g(x)}{f(x)} < \infty$$

Thus,  $\frac{\ln g}{\ln f} = \frac{\ln g - \ln f + \ln f}{\ln f} = \frac{\ln \frac{f}{g}}{\ln f} + 1$ . Because  $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} < \infty$ , we can get:  $\ln \frac{f}{g} < \infty$ . And  $\lim_{x \rightarrow \infty} f(x) = 0$ . So  $\frac{\ln \frac{f}{g}}{\ln f} \rightarrow 0$ . Therefore,  $\frac{\ln \frac{f}{g}}{\ln f} + 1 = 1$ .

### 3 Problem 15.16

(a) We know that there is an order of the four groups but there is no order of the members in a specific group. For a specific sequence of four groups like  $\{\{A, B, C\}, \{D, E, F\}, \{G, H, I\}, \{J, K, L\}\}$ , consider the first group  $\{A, B, C\}$ , the following 6 lists of students can all map to it. (“—” is just to make it easier to read.)

ABC—DEFGHIJKL, ACB—DEFGHIJKL, BAC—DEFGHIJKL,  
BCA—DEFGHIJKL, CAB—DEFGHIJKL, CBA—DEFGHIJKL.

Thus,  $k = 6^4 = 1296$ .

(b) For a specific sequence of 4 groups:  $(\{\text{Group 1}, \text{Group 2}, \text{Group 3}, \text{Group 4}\})$ , if there is no order between them, it will be  $4!$  sequences map one group assignment.

Thus,  $j = 4! = 24$ .

(c) There are  $12!$  different lists of 12 students. And there is a  $k - to - 1$  mapping from the lists of students to the sequences of groups. Also, there is a  $j - to - 1$  mapping from the sequences of groups to the group assignments. Thus, assume there are  $t$  group assignments:

$$t = \frac{12!}{k \cdot j} = \frac{12!}{1296 \times 24} = 15400$$

(d)

Different lists of  $3n$  students:  $3n!$ .

There is a  $6^n - to - 1$  mapping from the lists of students to the sequences of groups.

There is a  $n! - to - 1$  mapping from the sequences of groups to the non-order  $n$  groups of 3.

Thus, assume there are  $m$  ways:

$$m = \frac{3n!}{6^n \cdot n!}$$

### 4 Problem 15.47

All the integers in  $\{1, 2, 3, \dots, 2n\}$  can be written uniquely in the form  $2^a b$ , where  $a \leq 0, b \leq 1$ , and  $b$  is not a multiple of 2. Now apply the pigeonhole principle:

**Pigeons:**  $2n$  numbers.

**Pigeonholes:** each value of  $b$  is a hole. There are  $n$  holes.

If there are  $n + 1$  numbers selected, there must be at least two numbers

share the same value of  $b$  and their  $a$  value is different, which means their quotient is a power of 2.

## 5 Problem 15.67

**Algebraic manipulation:**

$$\begin{aligned}
 \binom{n}{r} \binom{r}{k} &= \frac{n!}{(n-r)!r!} \frac{r!}{(r-k)!k!} \\
 &= \frac{n!}{(n-r)!(r-k)!k!} \\
 &= \frac{n!}{k!(n-k)!} \frac{(n-k)!}{(r-k)!(n-r)!} \\
 &= \frac{n!}{k!(n-k)!} \frac{(n-k)!}{(r-k)![(n-k)-(r-k)]!} \\
 &= \binom{n}{k} \binom{n-k}{r-k}
 \end{aligned}$$

**Combinatorial argument:** Suppose we are forming a  $r$ -person ACM team with  $k$  people from the ACM team to join the competition.

**Method 1:** First choose the ACM team in  $\binom{n}{r}$  ways. Then choose  $\binom{r}{k}$  students to join the competition. So there are totally  $\binom{n}{r} \binom{r}{k}$  ways.

**Method 2:** First choose the  $k$  students to join the competition in  $\binom{n}{k}$  ways. Then choose the other  $(r-k)$  ACM team numbers from the remaining pool of  $(n-k)$  people. So there are  $\binom{n}{k} \binom{n-k}{r-k}$  ways.

Both two methods solve the same question, so there must be:

$$\binom{n}{r} \binom{r}{k} = \binom{n}{k} \binom{n-k}{r-k}$$

## 6 Problem 15.11

(a)  $a^k$

(b)  $a^k - a$

(c) Suppose  $z$  is  $a_1 a_2 a_3 \dots a_k$ . Then  $u$  will be

$$a_{n+1} a_{n+2} a_{n+3} \dots a_k a_1 a_2 a_3 \dots a_n$$

and  $v$  will be like

$$a_{m+n+1} a_{m+n+2} a_{m+n+3} \dots a_k a_1 a_2 a_3 \dots a_n a_{n+1} a_{n+2} a_{n+3} \dots a_{m+n}.$$

Therefore, we can find that  $v$  is the length- $(n - m)$  rotation of  $z$ .

(d) PROOF: First, to prove “ $\approx$ ” is reflexive.  $v \approx v$  holds when  $n = 1$ . Second, to prove “ $\approx$ ” is symmetric. We can suppose  $\text{length}|v| = \text{length}|z| = k$ . If  $v$  is a length- $n$  rotation of  $z$ , then according to (c)  $z$  is a length- $k - n$  rotation of  $v$ , which means “ $\approx$ ” is symmetric. Third, to prove it is transitive, we can use the conclusion in (c). Thus, “ $\approx$ ” is an equivalence relation.

(e) PROOF: We can prove by induction on  $|xy|$ . Suppose  $x$  and  $y$  are not empty strings. (If they are, when  $|xy| = 0$ ,  $x, y, u$  are all empty strings. The theorem holds. When  $|xy| = 1$ , one of  $x$  and  $y$  is empty. The theorem also holds.) Let's start at  $|xy| = 2$ , then  $x = y = u$ . If  $|xy| > 2$  and  $|x| = |y|$ , when  $x = y = u$ , the theorem holds. If  $|x| < |y|$ , then there must be a  $w$  that  $xw = y = wx$ , which means  $xw = wx$ . Since  $|xw| < |xy|$ , we end up with the same problem for smaller words ( $xw = wx$  is just like  $xy = yx$ ), so we can conclude by induction. According to the induction hypothesis there exists  $u$  such that  $x, w \in u^*$ . Since  $y = wx$ ,  $y \in u^*$ .

(f) If  $p$  is prime and  $z$  is a length- $p$  string containing at least two different characters. Then after doing length-1 rotation for  $p$  times, we can get  $p$  different strings. (There is no two same strings because  $p$  is a prime number.) So  $z$  is equivalent under  $\approx$  to exactly  $p$  strings (counting itself).

(g) We can get  $a^p$  strings for a length- $p$  string from  $a$ -character alphabet.  $a^p - a$  have at least 2 different characters. At the same time, the  $\approx$  is a  $p$ -to-1 relation in all the  $a^p - a$  strings, which means  $a^p - a$  must be divisible by  $p$ .

## 7 Problem 16.64.

(a) Lets combine student B and C as student K, so there is  $(7 \times 7! - 6 \times 6!) \times 2 = 61920$ .

(b) Suppose all the lineups satisfy Rule I are in set A. Also, B is for Rule II and C is for Rule III.

$$|A|: 10! - 9! = 3265920$$

$$|B|: 9! \times 2 = 725760$$

$$|C|: 9! = 362880$$

$$|A \cap B|: (9! - 8!) \times 2 = 645120$$

$$|B \cap C|: (7 \times 7!) \times 2 = 70560$$

$$|A \cap C|: 9! - 8! = 322560$$

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |A \cap C| + |A \cap B \cap C| = 3378240$$

```

    print();
    return;
}
for(i = offset; i < N; i++)
{
    swap(i, offset);
    perm(offset + 1);
    swap(i, offset);
}

void print()

if (a[9] != 'a' || abs(indexi(a,'b')-indexi(a,'c')) == 1 || a[1] == 'd'){
    j++;
}

3378240
Program ended with exit code: 0

```

Figure 1: Using Xcode can get the same answer.

## 8 Problem 15.51.

(a)PROOF: We can use contradiction to prove. The length of the circle is  $c=1$ , so we only need to prove  $k\sqrt{2} - m\sqrt{2} \neq nc$  ( $n = 0, 1, 2, \dots$ ) unless  $m = k$ . If  $m \neq k$ , then  $(k - m)\sqrt{2}$  is irrational. (This is because  $(k - m)$  is rational and  $\sqrt{2}$  is irrational.) But  $nc$  is rational. It is a contradiction. Only when  $k = m$ , then  $k - m = 0$  and  $k\sqrt{2} - m\sqrt{2} = nc$  ( $n = 0$ ).

(b)We can prove by contradiction. Suppose any two of the  $n$  points are more than  $1/n$  distance from each other. Then the clockwise distance from some point A to this point itself is more than 1. This is a contradiction. So there have to be two that are at most distance  $1/n$  from each other.

(c)Let  $\{an\}$  denote  $an - \lfloor an \rfloor$ .  $S = \{an : a = \sqrt{2}, n = 0, 1, 2, 3, \dots\}$ . We need to prove there are at least one element of S is in  $[0, 1/n]$ . If we have proved this, then there must be at least one element in  $[(k - 1)/n, k/n]$ . Now we prove there are at least one element in S is in  $[0, 1/n]$ . First, we divide the  $[0, 1]$  into  $n$  intervals. For set  $T = \{\{at\}\}, (t = 1, 2, 3, \dots, n + 1)$ , according to the pigeonhole principle, there must be two elements of T are in the same interval. Suppose the two elements are  $\{ap\}, \{aq\}$ , then let  $j = |p - q|$ . There must be  $0 < \{aj\} < 1/n$ . So there are at least one element of S is in  $[0, 1/n]$ . In conclusion, the marked points are dense on the circle.