

MAST 20018 - Exam Solutions 2017

Q1 a)

 $x_1 :=$ # of brownies eaten $x_2 :=$ # of scoops of choc ice cream eaten $x_3 :=$ # of bottles of cola drunk daily $x_4 :=$ # of pieces of cheesecake eaten

$$\min Z = 50x_1 + 20x_2 + 30x_3 + 80x_4 \quad \checkmark$$

s.t

$$400x_1 + 200x_2 + 150x_3 + 500x_4 \geq 500 \quad \checkmark$$

$$3x_1 + 2x_2 \geq 6 \quad \checkmark$$

$$2x_1 + 2x_2 + 4x_3 + 6x_4 \geq 10 \quad \checkmark$$

$$2x_1 + 4x_2 + x_3 + 5x_4 \geq 8 \quad \checkmark$$

$$x_i \geq 0 \quad \forall i \quad \checkmark$$

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b) i) Activity levels are not completely continuous. \checkmark

ii) The objective function contribution from each activity is not completely proportional. i.e., can you buy 70% of a brownie for 10c? \checkmark

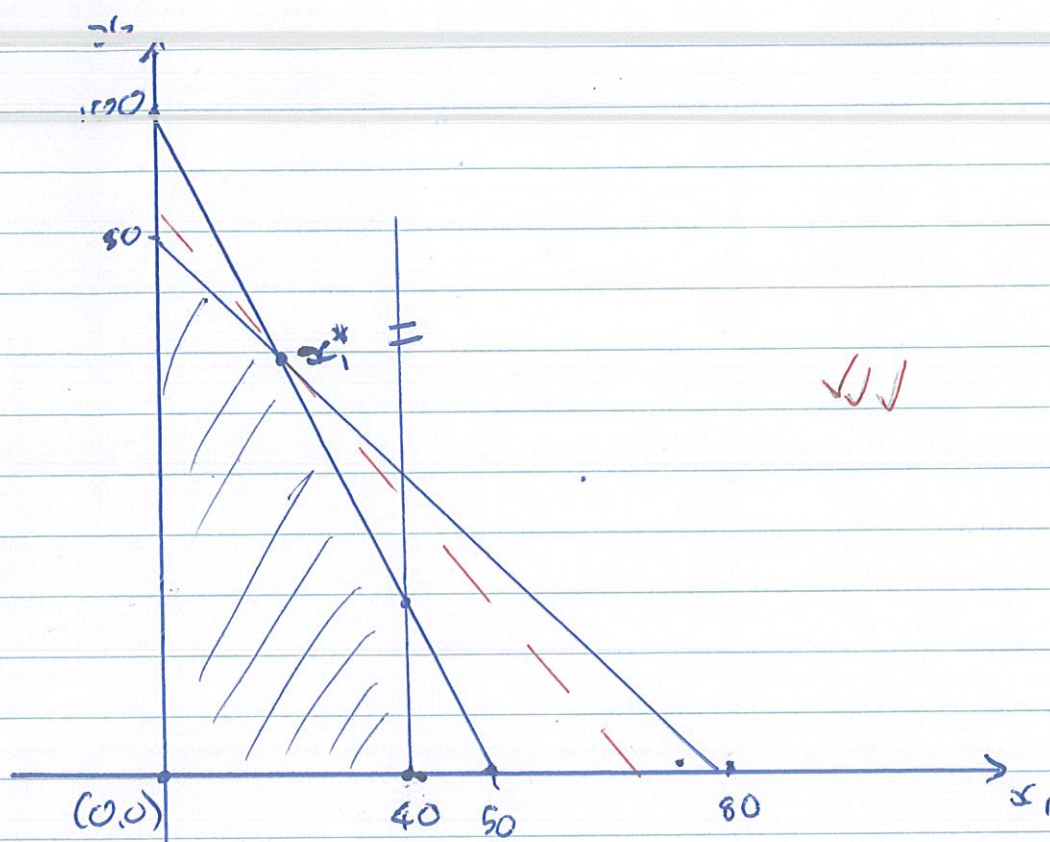
iii) Activities are not 100% independent (they may interfere through digestion, for eg) \checkmark

iv) There would be natural variation in eg the calories in each item. The exact nutritional values would have some uncertainty. \checkmark

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Q2).

a)



b).

If $k=0$ then $Z=0$ and all feasible points are optimal. ✓

If $k \neq 0$ then

$$x_2 = -\frac{3}{2}x_1 + \frac{Z}{2k} \quad \checkmark$$

If $k > 0$ then we wish to minimize the y-intercept.
 $\therefore (0,0)$ is optimal ✓. $Z=0$ ✓

If $k < 0$ then we wish to maximise the y-intercept.
 \therefore move the gradient to the point x_1^*
 which is the intersection of

$$2x_1 + x_2 = 100 \quad \text{and} \quad x_1 + x_2 = 80.$$

$$\therefore x_1 = 20 \quad \checkmark \quad \text{and} \quad x_2 = 60. \quad \checkmark$$

$$z = 3k(20) + 2k(60)$$

$$= 60k + 120k = 180k. \quad \checkmark$$

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Q3)

a) $\max -z = -2x_1 - 3x_2$ (or min)

$$\frac{1}{2}x_1 + \frac{1}{4}x_2 + s_1 = 4 \quad \checkmark$$

$$x_1 + 3x_2 - t_1 + y_1 = 36 \quad \checkmark$$

$$x_1 + x_2 + y_2 = 10 \quad \checkmark$$

$$x_1, x_2, s_1, t_1, y_1, y_2 \geq 0. \quad \checkmark \quad [4]$$

b) min $w = y_1 + y_2$

BV	x_1	x_2	s_1	t_1	y_1	y_2	RHS
s_1	$\frac{1}{2}$	$\frac{1}{4}$	1	0	0	0	4
y_1	1	3	0	-1	1	0	36
y_2	1	(1)	0	0	0	1	10
w	0	0	0	0	-1	-1	0
$\rightarrow w$	2	4	0	-1	0	0	46

BV	x_1	x_2	s_1	t_1	y_1	y_2	RHS
s_1	$\frac{1}{4}$	0	1	0	0	$-\frac{1}{4}$	$\frac{3}{2}$
y_1	-2	0	0	-1	1	-3	6
x_2	1	1	0	0	0	1	10
w	-2	0	0	-1	0	-4	6

c) The LP is infeasible since the phase 1 table has not eliminated the artificial variables. \leq

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Q4).

Find optimal solution:

$$\max z = x_1 + 4x_2$$

s.t

$$x_1 + 2x_2 + s_1 = 6$$

$$2x_1 + x_2 + s_2 = 8$$

$$x_1, x_2, s_1, s_2 \geq 0.$$

\therefore

$$A_B = \begin{bmatrix} 2 & 0 \\ 1 & 1 \end{bmatrix} \quad \checkmark$$

$$A_B | I = \left[\begin{array}{cc|cc} 2 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{array} \right]$$

using Gauss-Jordan.

$$A_B^{-1} = \begin{bmatrix} \frac{1}{2} & 0 \\ -\frac{1}{2} & 1 \end{bmatrix} \quad \checkmark$$

$$\therefore A_B^{-1} \cdot b = \begin{bmatrix} \frac{1}{2} & 0 \\ -\frac{1}{2} & 1 \end{bmatrix} \begin{bmatrix} 6 \\ 8 \end{bmatrix} = \begin{bmatrix} 3 \\ 5 \end{bmatrix} \quad \text{RHS.}$$

The column for x_1 in the optimal tableau

$$\text{is } A_B^{-1} \cdot \begin{bmatrix} 1 \\ 2 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} \\ \frac{3}{2} \end{bmatrix} \quad //$$

The column for s_1 in the optimal tableau is

$$A_B^{-1} \cdot \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} \\ -\frac{1}{2} \end{bmatrix} \quad //$$

$$C_B = [4, 0]$$

$$\therefore C_B \cdot A_B^{-1} = [4, 0] \begin{bmatrix} \frac{1}{2} & 0 \\ -\frac{1}{2} & 1 \end{bmatrix} = [2, 0]$$

\therefore Coeff of x_1 in Z -row is

$$[2, 0] \begin{bmatrix} 1 \\ 2 \end{bmatrix} - c_1 = 1$$

Coeff of s_1 in Z -row is

$$[2, 0] \begin{bmatrix} 1 \\ 0 \end{bmatrix} - 0 = 2$$

$$Z \text{ at the RHS is } C_B \cdot A_B^{-1} \cdot b = [2, 0] \begin{bmatrix} 6 \\ 8 \end{bmatrix}$$

$$= 12$$

\therefore

BV	x_1	x_2	s_1	s_2	RHS
x_2	$1/2$	1	$1/2$	0	3
s_2	$3/2$	0	$-1/2$	1	5
Z	1	0	2	0	12

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Q5.

check feasibility in primal

$$2 \cdot \frac{5}{26} + 5/2 + 3 \cdot \frac{27}{26}$$

$$= \frac{10}{26} + \frac{65}{26} + \frac{81}{26} = \frac{156}{26} = 6 \quad \therefore s_1 = 0$$

$$5 \cdot \frac{5}{26} + 4 \cdot 5/2 + \frac{27}{26} = 12 \quad \therefore s_2 = 0$$

$$2 \cdot \left(\frac{5}{2}\right) = 5 \quad \therefore s_3 = 0$$

Dual $\min 6y_1 + 12y_2 + 5y_3$

$$2y_1 + 5y_2 \geq 9$$

$$y_1 + 4y_2 + 2y_3 \geq 14$$

$$3y_1 + 4y_2 \geq 7$$

$$y_1, y_2, y_3 \geq 0$$

$$x_1 > 0, x_2 > 0, x_3 > 0.$$

$$\therefore t_1 = 0 \quad t_2 = 0 \quad t_3 = 0.$$

$$\therefore \begin{aligned} 2y_1 + 5y_2 &= 9 \\ 3y_1 + y_2 &= 7 \end{aligned}$$

$$-13y_1 = -26$$

$$y_1 = 2$$

$$y_2 = 1$$

$$y_1 + 4y_2 + 2y_3 = 14$$

$$\therefore 2 + 4 + 2y_3 = 14$$

$$y_3 = \frac{5}{2} = 2.5$$

$$w = 6y_1 + 12y_2 + 5y_3$$

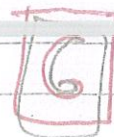
$$= 6(2) + 12(1) + 5 \cdot \frac{5}{2}$$

$$= 12 + 12 + \frac{25}{2} = 36\frac{1}{2}$$

$$z = 9\left(\frac{5}{26}\right) + 14\left(\frac{5}{12}\right) + 7\left(\frac{2}{26}\right)$$

$$= 44$$

$\therefore x^* = \left(\frac{5}{26}, \frac{5}{12}, \frac{2}{26}\right)$ is optimal.



Q6.)

BV	x_1	x_2	x_3	x_4	RHS
x_2	1	1	1	0	4
x_4	1	0	-1	1	1
Z	$1-3\alpha$	4α	4	0	16 ✓

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=>

BV	x_1	x_2	x_3	x_4	RHS
x_2	1	1	1	0	4 ✓
x_4	①	0	-1	1	1
Z	$1-7\alpha$	0	$4-4\alpha$	0	$16-16\alpha$

which is optimal iff $1 - 7\alpha \geq 0$ $\alpha \leq 1/7$

and $4 - 4\alpha \geq 0$ $\alpha \leq 1$

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(i) iff $\alpha \leq 1/7$. $\therefore \alpha \in [0, 1/7]$

(ii)

BV	x_1	x_2	x_3	x_4	RHS
x_2	0	1	2	-1	3
x_1	1	0	-1	1	1
Z	0	0	$5 - 13\alpha$	$7\alpha - 1$	$15 - 9\alpha$

[2]

which is optimal iff $5 - 13\alpha \geq 0$ $\alpha \leq 5/13$

$7\alpha - 1 \geq 0$ $\alpha \geq 1/7$

$\therefore \alpha \in [1/7, 5/13]$

$\Rightarrow \alpha \in [1/7, 5/13]$

[2]

(iii) If $\alpha = 1/7$ then $x_2 = 4$ and $x_4 = 1$ is optimal, with $x_1 = x_3 = 0$

i.e., the point $(0, 4)$ is optimal in the original program

Also, $x_2 = 3$, $x_1 = 1$ and $x_3 = x_4 = 0$ is optimal. i.e. $(1, 3)$ is optimal in the original program.

\therefore all points $\lambda(0, 4) + (1 - \lambda)(1, 3)$ are optimal, where $\lambda \in [0, 1]$.

$= (1 - \lambda, \lambda + 3)$

[2]