### Lecturer

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### **Main References**

Wen Shen,
 An Introduction to Numerical Computation

- Suli,
   Lecture Notes on Finite Element Methods for Partial Differential Equations
- Liu & Quek,
   The Finite Element Method (2nd Edition)

### **Lecture Plan**

- Computer arithmetic
- Interpolation
- Differentiation & Integration
- Solving non-linear equations, Least Squares
- Linear systems
- ODE IVP
- Midterm exam.

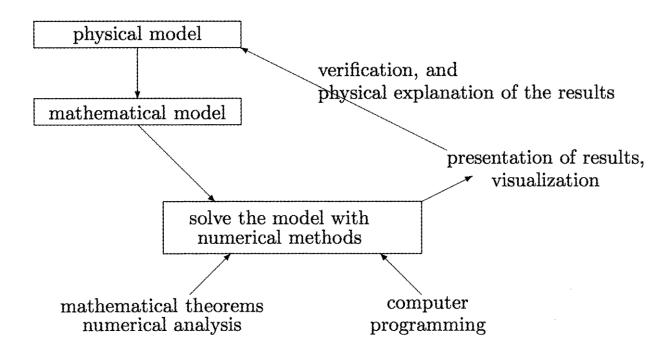
- ODE BVP
- PDE general
- FDM
- FEM
- BEM/SEM
- Final exam.
- LBM

# Computer Arithmetic

## IPCST Seoul National University

### **Numerical Methods**

- They are algorithms that compute approximate solutions to a number of problems for which exact solutions are not available.
  - Wen Shen



### **Representation of Numbers**



- Bases used in human history
  - 10
  - 2 (computer)
  - 8
  - 16 (ancient China)
  - 20 (ancient France)
  - 60 (Babylon)
- In base  $\beta$ , integer part fractional part

$$\left(\overline{a_n a_{n-1} \cdots a_1 a_0} \cdot \overline{b_1 b_2 b_3 \cdots}\right)_{\beta}$$

$$= a_n \beta^n + a_{n-1} \beta^{n-1} + \cdots + a_1 \beta + a_0 \qquad \text{(integer part)}$$

$$+ b_1 \beta^{-1} + b_2 \beta^{-2} + b_3 \beta^{-3} + \cdots \qquad \text{(fractonal part)}$$





$$x = \pm r \times \beta^e$$
,  $1 \le r < \beta$   
 $\beta = 2$  in computer arithmetic

Single-precision (32 bits)

$$-126 \le e \le 127$$

- About 10<sup>-38</sup>~10<sup>38</sup>, 7 or 8 decimal digits
   (sign: 1 bit, exponent: 8 bits, mantissa: 23 bits)
- Double-precision (64 bits)

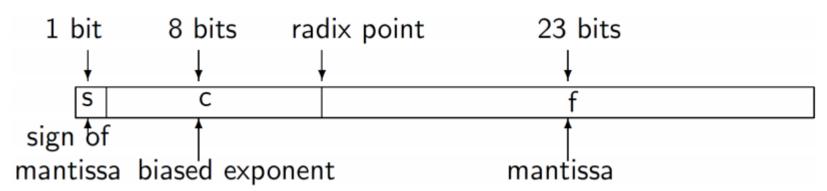
$$-1022 \le e \le 1023$$

About 10<sup>-308</sup>~10<sup>308</sup>, 15 or 16 decimal digits
 (sign: 1 bit, exponent: 11 bits, mantissa: 52 bits)





- Information to be stored:
  - The sign
  - The exponent *e*
  - The value  $r (= 1 + f = 1 + n/2^b)$ 
    - b = 23 for single precision,  $0 \le n < 2^b$  (integer)



32-bit single precision bit layout (Figure from Wen Shen)



- Let's say  $-m \le e \le M$
- Maximum absolute value =  $(2 2^{-b}) \times 2^{M}$
- Minimum absolute value =  $2^{-m}$
- Inf: infinity (e = M + 1 & f = 0)
  - Cf.) NaN: not a number (e = M + 1 & f > 0)
- 0: zero (e = -m 1)
- (Arithmetic) overflow:  $|x| > (2 2^{-b}) \times 2^{M} (\rightarrow Inf)$
- (Arithmetic) underflow:  $|x| < 2^{-m} (\rightarrow 0)$





- The floating-point representation fl(x)  $fl(x) = x \cdot (1 + \delta)$
- Relative error  $\delta = \frac{fl(x) x}{x}$
- Absolute error =  $fl(x) x = \delta \cdot x$  $|\delta| \le \varepsilon$ 
  - In case of round-off,  $|\delta| \leq \varepsilon/2$
- Machine epsilon  $\varepsilon = 2^{-b}$



$$\alpha \leq \gamma < \beta$$

$$\alpha = \left(1 + \frac{n}{2^b}\right) \times 2^e$$

$$\beta = \left(1 + \frac{n+1}{2^b}\right) \times 2^e$$



$$x = \alpha \text{ or } \beta$$



- Error propagation
  - Errors can be accumulated through multiple arithmetic operations.
  - Example 1.4 in Wen Shen

$$\begin{aligned} \mathsf{fl}(x) &= x(1+\delta_x), & \mathsf{fl}(y) &= y(1+\delta_y) \\ \mathsf{fl}(z) &= & \mathsf{fl}\left(\mathsf{fl}(x)+\mathsf{fl}(y)\right) \\ &= & \left(x(1+\delta_x)+y(1+\delta_y)\right)(1+\delta_z) \\ &= & \left(x+y\right)+x\cdot\left(\delta_x+\delta_z\right)+y\cdot\left(\delta_y+\delta_z\right)+\left(x\delta_x\delta_z+y\delta_y\delta_z\right) \\ &\approx & \left(x+y\right)+x\cdot\left(\delta_x+\delta_z\right)+y\cdot\left(\delta_y+\delta_z\right) \\ \mathsf{relative\ error} &= & \frac{\mathsf{fl}(z)-(x+y)}{x+y} = \underbrace{\frac{x\delta_x+y\delta_y}{x+y}}_{\mathsf{propagated\ err}} + \underbrace{\delta_z}_{\mathsf{propagated\ err}} \end{aligned}$$



- Error propagation
  - Errors can be accumulated through multiple arithmetic operations.
  - Example 1.4 in Wen Shen

$$\text{fl}(x) = x(1+\delta_x), \qquad \text{fl}(y) = y(1+\delta_y)$$
 
$$\text{fl}(z) = (x+y) + x \cdot (\delta_x + \delta_z) + y \cdot (\delta_y + \delta_z) + (x\delta_x\delta_z + y\delta_y\delta_z)$$
 
$$\approx (x+y) + x \cdot (\delta_x + \delta_z) + y \cdot (\delta_y + \delta_z)$$
 absolute error 
$$= \text{fl}(z) - (x+y) = x \cdot (\delta_x + \delta_z) + y \cdot (\delta_y + \delta_z)$$
 
$$= \underbrace{x \cdot \delta_x}_{\text{abs. err.}} + \underbrace{y \cdot \delta_y}_{\text{youth }} + \underbrace{(x+y) \cdot \delta_z}_{\text{round off err}}$$
 abs. err. abs. err. round off err 
$$\underbrace{\text{for } x}_{\text{propagated error}}$$



 Loss of significant digits typically happens in subtraction of two very close numbers

1.2345678

-1.2344444

0.0001234

4 significant digits disappear.

## **Loss of Significance**



- Example 1.5 in Wen Shen
  - Find roots of  $x^2 40x + 2 = 0$ . Use 4 significant digits.

$$r_{1,2} = \frac{1}{2a} \left( -b \pm \sqrt{b^2 - 4ac} \right)$$

$$x_{1,2} = 20 \pm \sqrt{398} \approx 20.00 \pm 19.95$$

$$x_1 \approx 20 + 19.95 = 39.95$$

$$x_2 \approx 20 - 19.95 = 0.05$$
 Loss of 3 significant digits

Change the algorithm to get x<sub>2</sub>

$$x_2 = \frac{c}{ax_1} = \frac{2}{1 \cdot 39.95} \approx 0.05006$$

## Loss of Significance

Example 1.6 in Wen Shen

$$f(x) = \frac{1}{\sqrt{x^2 + 2x} - x - 1}$$

- Numerical computation of this function may lead to loss of significance. Find a remedy.
  - For large x,  $\sqrt{x^2 + 2x}$  and x + 1 are very close to each other.  $\rightarrow$  loss of significant digits in subtraction
  - Change the form of the function

$$f(x) = \frac{\sqrt{x^2 + 2x} + x + 1}{\left(\sqrt{x^2 + 2x} - x - 1\right)\left(\sqrt{x^2 + 2x} + x + 1\right)}$$
$$= \frac{\sqrt{x^2 + 2x} + x + 1}{x^2 + 2x - (x + 1)^2} = -\left(\sqrt{x^2 + 2x} + x + 1\right).$$

## **Choice of Language**

- MATLAB
- Python (with NumPy & MatPlotLib)

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FortranC (+ any graph tools)
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### Investigation

 Find the floating-point representation applied in your programming language

### References

Wen Shen,
 An Introduction to Numerical Computation

- C. Moler,
   Numerical Computing with MATLAB
- Wikipedia