## Lecture #2 Exercises

## Distributed Lab

July 25, 2024



**Exercise 1.** Suppose that for the given cipher with a security parameter  $\lambda$ , the adversary  $\mathcal A$  can deduce the least significant bit of the plaintext from the ciphertext. Recall that the advantage of a bit-guessing game is defined as  $\mathsf{SSAdv}[\mathcal A] = |\mathsf{Pr}[b=\hat b] - \frac12|$ , where b is the randomly chosen bit of a challenger, while  $\hat b$  is the adversary's guess. What is the maximal advantage of  $\mathcal A$  in this case?

**Hint:** The adversary can choose which messages to send to challenger to further distinguish the plaintexts.

- a) 1
- b)  $\frac{1}{2}$
- c)  $\frac{1}{4}$
- d) 0
- e) Negligible value (negl( $\lambda$ )).

**Exercise 2.** Consider the cipher  $\mathcal{E} = (E, D)$  with encryption function  $E : \mathcal{K} \times \mathcal{M} \to \mathcal{C}$  over the message space  $\mathcal{M}$ , ciphertext space  $\mathcal{C}$ , and key space  $\mathcal{K}$ . We want to define the security that, based on the cipher, the adversary  $\mathcal{A}$  cannot restore the message (security against message recovery). For that reason, we define the following game:

- 1. Challenger chooses random  $m \stackrel{R}{\leftarrow} \mathcal{M}, k \stackrel{R}{\leftarrow} \mathcal{K}$ .
- 2. Challenger computes the ciphertext  $c \leftarrow E(k, m)$  and sends to A.
- 3. Adversary outputs  $\hat{m}$ , and wins if  $\hat{m} = m$ .

We say that the cipher  $\mathcal{E}$  is secure against message recovery if the **message recovery** advantage, denoted as MRadv[ $\mathcal{A}$ ,  $\mathcal{E}$ ] is negligible. Which of the following statements is a valid interpretation of the message recovery advantage?

- a)  $\mathsf{MRadv}[\mathcal{A}, \mathcal{E}] := \left| \mathsf{Pr}[m = \hat{m}] \frac{1}{2} \right|$
- b) MRadv[ $\mathcal{A}, \mathcal{E}$ ] :=  $|\Pr[m = \hat{m}] 1|$ .
- c)  $\mathsf{MRadv}[\mathcal{A},\mathcal{E}] := \mathsf{Pr}[m = \hat{m}]$
- d)  $\mathsf{MRadv}[\mathcal{A},\mathcal{E}] := \left| \mathsf{Pr}[m = \hat{m}] \frac{1}{|\mathcal{M}|} \right|$

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**Exercise 3.** Suppose that f and g are negligible functions. Which of the following functions is not neccessarily negligible?

- a) f + g
- b)  $f \times g$
- c) f g
- d) f/g

e) 
$$h(\lambda) := \begin{cases} 1/f(\lambda) & \text{if } 0 < \lambda < 100000 \\ g(\lambda) & \text{if } \lambda \geq 100000 \end{cases}$$

**Exercise 4.** Suppose that  $f \in \mathbb{F}_p[x]$  is a d-degree polynomial with d **distinct** roots in  $\mathbb{F}_p$ . What is the probability that, when evaluating f at n random points, the polynomial will be zero at all of them?

- a) Exactly  $(d/p)^n$ .
- b) Strictly less that  $(d/p)^n$ .
- c) Exactly nd/p.
- d) Exactly d/np.