0.1 Relations

Before dvelving into the projective coordinates and further zero-knowledge topics, let us first discuss the concept of relations, which will be intensively used from now on. Now, what is a relation? The definition is incredibly concise.

Definition 0.1. Let \mathcal{X}, \mathcal{Y} be some sets. Then, \mathcal{R} is a **relation** if

$$\mathcal{R} \subset \mathcal{X} \times \mathcal{Y} = \{(x, y) : x \in \mathcal{X}, y \in \mathcal{Y}\}$$
 (1)

Interpretation is approximately the following: suppose we have sets \mathcal{X} and \mathcal{Y} . Then, relation \mathcal{R} gives a set of pairs (x, y), telling that $x \in \mathcal{X}$ and $y \in \mathcal{Y}$ are related.

Example. Let $\mathcal{X} = \{\text{Oleksandr}, \text{Phat}, \text{Anton}\}\$ and $\mathcal{Y} = \{\text{Backend}, \text{Frontend}, \text{Research}\}\$. Define the following relation of "person x works in field y":

$$\mathcal{R} = \{ (Oleksandr, Research), (Phat, Frontend), (Anton, Backend) \}$$
 (2)

Obviously, $\mathcal{R} \subset \mathcal{X} \times \mathcal{Y}$, so \mathcal{R} is a relation.

Remark. There are many ways to express that $(x,y) \in \mathcal{R}$. Most common are $x\mathcal{R}y$ and $x \sim y$. Also, sometimes, one might encounter relation definition as a boolean function $\mathcal{R}: \mathcal{X} \times \mathcal{Y} \to \{0,1\}$, where $\mathcal{R}(x,y)$ is 1 if (x,y) is in the relation, and 0 otherwise. Further, we will use notation $x \sim y$ to denote that $(x,y) \in \mathcal{R}$.

Example. Let E be a cyclic group of points on the Elliptic Curve of order $r \geq 2$ with a generator $\langle G \rangle = E$. Let $\mathcal{X} = \mathbb{Z}_r$ and $\mathcal{Y} = E$. Define a relation $\mathcal{R} \subset \mathcal{X} \times \mathcal{Y}$ by:

$$\mathcal{R} = \{ (\alpha, P) \in \mathbb{Z}_r \times E : [\alpha]G = P \}$$
 (3)

Essentially, such a relation is a set of secret keys α and corresponding public keys P. In this case, for example, $0\mathcal{RO}$ and $1\mathcal{RG}$ or $0\sim\mathcal{O}$ and $1\sim G$.

Remark. When we say that \sim is a relation on a set \mathcal{X} , we mean that \sim is a relation \mathcal{R} on the following Cartesian product: $\mathcal{R} \subset \mathcal{X} \times \mathcal{X}$.

Now, let us formally define the term **equivalence relation**.

Definition 0.2. Let \mathcal{X} be a set. A relation \sim on \mathcal{X} is called an **equivalence relation** if it satisfies the following properties:

- 1. **Reflexivity:** $x \sim x$ for all $x \in \mathcal{X}$.
- 2. **Symmetry:** If $x \sim y$, then $y \sim x$ for all $x, y \in \mathcal{X}$.
- 3. **Transitivity:** If $x \sim y$ and $y \sim z$, then $x \sim z$ for all $x, y, z \in \mathcal{X}$.

Example. Let \mathcal{X} be the set of all people. Define a relation \sim on \mathcal{X} by $x \sim y$ if $x, y \in \mathcal{X}$ have the same birthday. Then \sim is an equivalence relation on \mathcal{X} . Let us demonstrate that:

1. **Reflexivity:** $x \sim x$ since x has the same birthday as x.

- 2. **Symmetry:** If $x \sim y$, then $y \sim x$ since x has the same birthday as y.
- 3. **Transitivity:** If $x \sim y$ and $y \sim z$, then $x \sim z$ since x has the same birthday as y and y has the same birthday as z.

Example. Suppose $\mathcal{X} = \mathbb{Z}$ and n is some fixed integer. Let $a \sim b$ mean that $a \equiv b \pmod{n}$. It is easy to verify that \sim is an equivalence relation:

- 1. **Reflexivity:** $a \equiv a \pmod{n}$, so $a \sim a$.
- 2. **Symmetry:** If $a \equiv b \pmod{n}$, then $b \equiv a \pmod{n}$, so $b \sim a$.
- 3. **Transitivity:** If $a \equiv b \pmod{n}$ and $b \equiv c \pmod{n}$, then $a \equiv c \pmod{n}$. It is not that obvious, so we can prove it: from the first equality we have $\exists q \in \mathbb{Z} : a b = nq$. From the second, $\exists r \in \mathbb{Z} : b c = nr$. Adding both we get (a b) + (b c) = n(r + q) or, equivalently, a c = n(r + q), meaning $a \equiv c \pmod{n}$.

The example below is less obvious with a bit more difficult proof, which we will skip. Yet, it is quite curious, so here it is.

Example. Let $\mathcal G$ be the set of all possible groups. Define a relation \sim on $\mathcal G$ by $\mathbb G \sim \mathbb H$ if $\mathbb G \cong \mathbb H$ (in other words, $\mathbb G$ and $\mathbb H$ are isomorphic). Then \sim is an equivalence relation.

Now, suppose I give you a set \mathcal{X} with some equivalence relation \sim (say, $\mathcal{X} = \mathbb{Z}$ and $a \equiv b \pmod{n}$). Notice that you can find some subset $\mathcal{X}' \subset \mathcal{X}$ in which all elements are equivalent (and any other element from $\mathcal{X} \setminus \mathcal{X}'$ is not). In the case of modulo relation above, \mathcal{X}' could be the set of all integers that are congruent to 1 modulo n, so $\mathcal{X}' = \{\ldots, -n+1, 1, n+1, 2n+1, \ldots\}$. This way, we can partition the set \mathcal{X} into disjoint subsets, where all elements in each subset are equivalent. Such subsets are called **equivalence classes**. Now, let us give a formal definition.

Definition 0.3. Let \mathcal{X} be a set and \sim be an equivalence relation on \mathcal{X} . For any $x \in \mathcal{X}$, the **equivalence class** of x is the set

$$[x] = \{ y \in \mathcal{X} : x \sim y \} \tag{4}$$

The **set of all equivalence classes** is denoted by \mathcal{X}/\sim (or, if the relation \mathcal{R} is given explicitly, then \mathcal{X}/\mathcal{R}), which is read as " \mathcal{X} modulo relation \sim ".

Example. Let $\mathcal{X} = \mathbb{Z}$ and n be some fixed integer. Define \sim on \mathcal{X} by $x \sim y$ if $x \equiv y \pmod{n}$. Then the equivalence class of x is the set

$$[x] = \{ y \in \mathbb{Z} : x \equiv y \pmod{n} \}$$
 (5)

For example, $[0] = \{..., -2n, -n, 0, n, 2n, ...\}$ while $[1] = \{..., -2n + 1, -n + 1, 1, n + 1, 2n + 1, ...\}$. Note that a set $\{[0]\}$

Now, as we have said before, a set of all equivalence classes form a partition of the set \mathcal{X} . This means that any element $x \in \mathcal{X}$ belongs to exactly one equivalence class. This is a very important property, which we will use in the next section. Formally, we have the following

lemma.

Lemma 0.4. Let \mathcal{X} be a set and \sim be an equivalence relation on \mathcal{X} . Then,

- 1. For each $x \in \mathcal{X}$, $x \in [x]$ (quite obvious, follows from reflexivity).
- 2. For each $x, y \in \mathcal{X}$, $x \sim y$ if and only if [x] = [y].
- 3. For each $x, y \in \mathcal{X}$, either [x] = [y] or $[x] \cap [y] = \emptyset$.

Example. Let $n \in \mathbb{N}$ and, again, $\mathcal{X} = \mathbb{Z}$ with a "modulo n" equivalence relation \mathcal{R}_n . Define the equivalence class of x by $[x]_n = \{y \in \mathbb{Z} : x \equiv y \pmod{n}\}$. Then,

$$\mathbb{Z}/\mathcal{R}_n = \{[0]_n, [1]_n, [2]_n, \dots, [n-2]_n, [n-1]_n\}$$
(6)

forms a partition of \mathbb{Z} , that is

$$\bigcup_{i=0}^{n-1} [i]_n = \mathbb{Z},\tag{7}$$

and for all $i, j \in \{0, 1, ..., n-1\}$, if $i \neq j$, then $[i]_n \cap [j]_n = \emptyset$. Commonly, we denote the set of all equivalence classes as $\mathbb{Z}/n\mathbb{Z}$ or, as we got used to, \mathbb{Z}_n . Moreover, we can naturally define the addition as:

$$[x]_n + [y]_n = [x + y]_n \tag{8}$$

Then, the set $(\mathbb{Z}/n\mathbb{Z}, +)$ with the defined addition is a group.

The primary reason we considered equivalence relations is that we will define the projective space as a set of equivalence classes. Besides this, when defining proofs of knowledge, argument of knowledge and zero-knowledge protocols, we will use the concept of relations and equivalence relations intensively.

0.2 Elliptic Curve in Projective Coordinates

0.2.1 Projective Space

Recall that we defined the elliptic curve as

$$E(\overline{\mathbb{F}}_p) := \{ (x, y) \in \mathbb{A}^2(\overline{\mathbb{F}}_p) : y^2 = x^3 + ax + b \} \cup \{ \mathcal{O} \}$$
 (9)

The above definition is the definition of the elliptic curve in *the affine space*. However, notice that in this case we need to append a somewhat artificial point \mathcal{O} to the curve. This is done to make the curve a group since without this point it is unclear how to define addition of two, say, negative points on the curve (since the resultant vertical line does not intersect the curve at any other point). The way to unify all the points $E/\overline{\mathbb{F}}_p$ with this magical point at infinity \mathcal{O} is to use the **projective space**.

Essentially, instead of working with points in affine n-space (in our case, with two-dimensional points $\mathbb{A}^2(\mathbb{K})$), we work with lines that pass through the origin in (n+1)-dimensional space (in our case, 3-dimensional space $\mathbb{A}^3(\mathbb{K})$). We say that two points from this (n+1)-dimensional space are **equivalent** if they lie on the same line that passes through the origin (we will show the illustration a bit later).

It seems strange that we need to work with 3-dimensional space to describe 2-dimensional points, but this is the way to unify all the points on the curve. Because, in this case, the point

at infinity is represented by a set of points on the line that passes through the origin and is parallel to the y-axis. We will get to understanding how to interpret that. Moreover, by defining operations on the projective space, we can make the operations on the curve more efficient.

Now, to the formal definition.

Definition 0.5. Projective coordinate, denoted as $\mathbb{P}^2(\mathbb{K})$ (or sometimes simply \mathbb{KP}^2) is a triple of elements (X:Y:Z) from $\mathbb{A}^3(\overline{\mathbb{K}})\setminus\{0\}$ modulo the equivalence relation^a:

$$(X_1:Y_1:Z_1) \sim (X_2:Y_2:Z_2) \text{ iff } \exists \lambda \in \overline{\mathbb{K}}: (X_1:Y_1:Z_1) = (\lambda X_2:\lambda Y_2:\lambda Z_2)$$
 (10)

^aAlthough we specify the definition for n=2, the definition can be generalized to any $\mathbb{P}^n(\overline{\mathbb{K}})$.

This definition on itself might be a bit too abstract, so let us consider the concrete example for projective space $\mathbb{P}^2(\mathbb{R})$.

Example. Consider the projective space $\mathbb{P}^2(\mathbb{R})$. Then, two points $(x_1, y_1, z_1), (x_2, y_2, z_2) \in \mathbb{R}^3$ are equivalent if there exists $\lambda \in \mathbb{R}$ such that $(x_1, y_1, z_1) = (\lambda x_2, \lambda y_2, \lambda z_2)$. For example, $(1, 2, 3) \sim (2, 4, 6)$ since (1, 2, 3) = 0.5(2, 4, 6).

Example. Now, how to geometrically interpret $\mathbb{P}^2(\mathbb{R})$? Consider the Figure below.

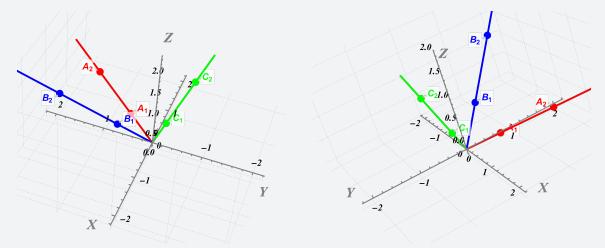


Illustration: Geometric interpretation of $\mathbb{P}^2(\mathbb{R})$, the same scene from different perspectives. The **red** line is represented by equation (2t, 3t, t), **blue** line by (-2t, 3t, 3t), and **green** line is represented by (t, -2t, 5t) for parameter $t \in \mathbb{R}$.

Here, the figure demonstrates three equivalence classes, being a set of points on the **red**, **blue**, and **green** lines (except for the origin).

The reason why geometrically the set of equivalence classes lie on the same line that passes through the origin is following: suppose we have a point $\vec{v}_0 = (x_0, y_0, z_0) \in \mathbb{R}^3$, represented as a vector. Then, the set of all points that are equivalent to (x_0, y_0, z_0) is the set of all points $(\lambda x_0, \lambda y_0, \lambda z_0) = \lambda \vec{v}_0$ for $\lambda \in \mathbb{R} \setminus \{0\}$. So \vec{v}_0 is the representative of equivalence class $[\vec{v}_0] = \{\lambda \vec{v}_0 : \lambda \in \mathbb{R}, \lambda \neq 0\}$. Now notice, that this is a parametric equation of a line that passes through the origin and the point \vec{v}_0 : notice that for $\lambda = 0$ (if we assume that expression is also defined for zero λ) we have the origin $\vec{0}$, while for $\lambda = 1$ we have the point \vec{v}_0 . Then, any other values of λ in-between [0,1] or outside define the set of points lying on the same line.

Now, projective coordinates are not that useful unless we can come back to the affine space. This is done by defining the map $\phi: \mathbb{P}^2(\overline{\mathbb{K}}) \to \mathbb{A}^2(\overline{\mathbb{K}})$ as follows: $\phi: (X:Y:Z) \mapsto (X/Z,Y/Z)$. If, in turn, we want to go from the affine space to the projective space, we can define the map $\psi: \mathbb{A}^2(\overline{\mathbb{K}}) \to \mathbb{P}^2(\overline{\mathbb{K}})$ as follows: $\psi: (x,y) \mapsto (x:y:1)$. Geometrically, map ϕ means that we take a point (X:Y:Z) and project it onto the plane Z=1.

Example. Again, consider three lines from the previous example. Now, we additionally draw a plane $\pi: z=1$ in our 3-dimensional space (see Illustration below).

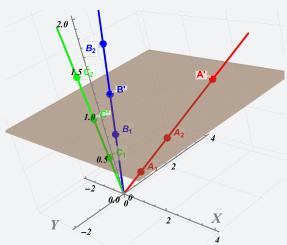


Illustration: Geometric interpretation of converting projective form to the affine form.

By using the map $(X:Y:Z)\mapsto (X/Z,Y/Z)$, all points on the line get mapped to the itersection of the line with the plane $\pi:z=1$. This way, for example, points on the **red line** ℓ_{red} get mapped to the point A'=(2,3,1), corresponding to (2,3) in affine coordinates. So, for example, point $(6,9,3)\in\ell_{red}$, lying on the same line, gets mapped to (6/3,9/3)=(2,3). Similarly, all **blue line** points get mapped to the point B'=(-2/3,1,1), while all **green line** points get mapped to the point $C'=(0.2,-0.4,1)^a$.

^aOne can verify that based on the equations provided from the previous example

0.2.2 Elliptic Curve Equation in Projective Form

Now, quite an interesting question is following: how to represent (basically, rewrite) the "affine" elliptic curve equation¹

$$E_{\mathbb{A}}(\overline{\mathbb{F}}_p): y^2 = x^3 + ax + b, \ a, b \in \overline{\mathbb{F}}_p$$
 (11)

in the projective form? Since currently, we defined the curve as the 2D curve, but now we are working in 3D space! The answer is following: recall that if $(X:Y:Z) \in \mathbb{P}^2(\overline{\mathbb{F}}_p)$ lies on the curve, so does the point (X/Z,Y/Z). The condition on the latter point to lie on $E_{\mathbb{A}}(\overline{\mathbb{F}}_p)$ is following:

$$\left(\frac{Y}{Z}\right)^2 = \left(\frac{X}{Z}\right)^3 + a \cdot \frac{X}{Z} + b \tag{12}$$

 $^{^{1}}$ Further, we will use notation $E_{\mathbb{A}}$ to represent the elliptic curve equation in the affine form, and $E_{\mathbb{P}}$ to represent the elliptic curve in the projective form.

But now multiply both sides by Z^3 to get rid of the fractions:

$$E_{\mathbb{P}}(\overline{\mathbb{F}}_p): Y^2 Z = X^3 + aXZ^2 + bZ^3 \tag{13}$$

This is an equation of the elliptic curve in **projective form**.

Now, one of the motivations to work with the projective form was to unify affine points $E_{\mathbb{A}}/\overline{\mathbb{F}}_p$ and the point at infinity \mathcal{O} , which acted as an identity element in the group $E_{\mathbb{A}}(\overline{\mathbb{F}}_p)$. So how do we encode the point at infinity in the projective form?

Well, notice the following observation: all points $(0:\lambda:0)$ always lie on the curve $E_{\mathbb{P}}(\overline{\mathbb{F}}_p)$. Moreover, the map from the projective form to the affine form is ill-defined for such points, since we would need to divide by zero. So, we can naturally make the points $(0:\lambda:0)$ to be the set of points at infinity. This way, we can define the point at infinity as $\mathcal{O} = (0:1:0)$.

Finally, let us summarize what we have observed so far.

Definition 0.6. The homogenous projective form of the elliptic curve $E_{\mathbb{P}}(\overline{\mathbb{F}}_p)$ is defined as the set of all points $(X:Y:Z) \in \mathbb{P}^2(\overline{\mathbb{F}}_p)$ in the projective space that satisfy the equation

$$E_{\mathbb{P}}(\overline{\mathbb{F}}_p): Y^2Z = X^3 + aXZ^2 + bZ^3, \quad a, b \in \overline{\mathbb{F}}_p,$$
 (14)

where the point at infinity is encoded as $\mathcal{O} = (0:1:0)$.

Example. Consider the BN254 curve $y^2 = x^3 + 3$ over reals \mathbb{R} . Its projective form is given by the equation $Y^2Z = X^3 + 3XZ^2$, which gives a surface, depicted below.

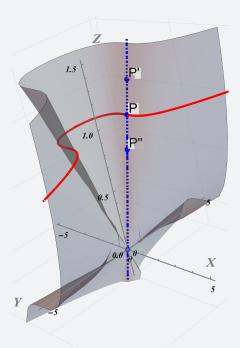


Illustration: BN254 Curve Elliptic Curve in Projective Form over \mathbb{R} . In gray is the surface, while **red** points are the points on the affine curve (lying on the plane $\pi: z = 1$).

Points $P' \approx (0:2.165:1.25)$ and $P'' \approx (0:1.3:0.75)$ in projective form both lie on the curve and get mapped to the same point $P \approx (0,1.732)$ in affine coordinates.

0.2.3 General Projective Coordinates

Hold on, but why did we use the term *homogenous*? The reason why is because we defined equivalence as follows: $(X:Y:Z)\sim (\lambda X:\lambda Y:\lambda Z)$ for some $\lambda\in\overline{\mathbb{K}}$, called **homogenous coordinates**. However, this is not the only way to define equivalence. Consider a more general form of equivalence relation:

$$(X:Y:Z) \sim (X':Y':Z') \text{ iff } \exists \lambda \in \overline{\mathbb{K}}: (X,Y,Z) = (\lambda^n X', \lambda^m Y', \lambda Z')$$
 (15)

In this case, to come back to the affine form, we need to use the map $\phi:(X:Y:Z)\mapsto (X/Z^n,Y/Z^m)$.

Example. The case n=2, m=3 is called the **Jacobian Projective Coordinates**. An Elliptic Curve equation might be then rewritten as:

$$Y^2 = X^3 + aXZ^4 + bZ^6 (16)$$

The reason why we might want to use such coordinates is that they can be more efficient in some operations, such as point addition. However, we will not delve into this topic much further.

Example. Consider the BN254 curve $y^2 = x^3 + 3$ over reals \mathbb{R} , again. Its *Jacobian projective* form is given by the equation $Y^2 = X^3 + 3XZ^4$, which gives a surface, depicted below.

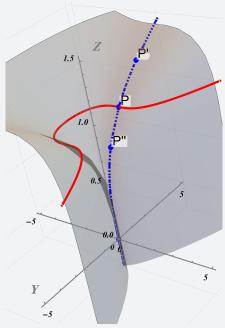


Illustration: BN254 Curve Elliptic Curve in Jacobian Projective Form over \mathbb{R} . In gray is the surface, while **red** points are the points on the affine curve (lying on the plane $\pi: z=1$).

Notice that now, under the map $(X : Y : Z) \mapsto (X/Z^2, Y/Z^3)$, points in the same equivalence class (in \mathbb{R}^3) do not lie on the same line, but rather on the same *curve*. Namely, equivalence class has a form $[(x_0, y_0, z_0)] = \{t^2x_0, t^3y_0, tz_0 : t \in \mathbb{R} \setminus \{0\}\}$.

0.2.4 Fast Addition

Let us come back to the affine case and assume that the underlying field is the prime field \mathbb{F}_p . Recall that for adding two points $P = (x_P, y_P)$ and $Q = (x_Q, y_Q)$ to get $R = (x_R, y_R) \leftarrow P \oplus Q$ one used the following formulas (there is no need to understand the derivation fully, just take it as a fact):

$$x_R \leftarrow \left(\frac{y_Q - y_P}{x_Q - x_P}\right)^2 - x_P - x_Q, \quad y_R \leftarrow \left(\frac{y_Q - y_P}{x_Q - x_P}\right) (x_P - x_R) - y_P \tag{17}$$

Denote by M the cost of multiplication, by S the cost of squaring, and by I the cost of inverse operation in \mathbb{F}_p . Note that we do not count addition/inverse costs as they are significantly lower than operations listed. Then, the cost of additing two points using above formula is 2M+1S+1I. Indeed, our computation can proceed as follows:

- 1. Calculate $t_1 \leftarrow (x_Q x_P)^{-1}$, costing 11.
- 2. Calculate $\lambda \leftarrow (y_Q y_P)t_1$, costing 1M.
- 3. Calculate $t_2 \leftarrow \lambda^2$, costing 1S.
- 4. Calculate $x_R \leftarrow t_2 x_P x_Q$, costing almost nothing.
- 5. Calculate $y_R \leftarrow \lambda(x_P x_R) y_P$, costing 1M.

Well, there are just 4 operations in total, so what can go wrong? The problem is that we need to calculate the inverse of $(x_Q - x_P)$, which is a very, very costly operation. In fact, typically $11 \gg 20M$ or even worse, the ratio might reach 80 in certain cases.

Now imagine we want to add 4 points, say $P_1 \oplus P_2 \oplus P_3 \oplus P_4$: this costs 6M + 3S + 3I. Now we have 3 inverses, which is a lot. Finally, if we are to add much larger number of points (for example, when finding the scalar product), this gets even worse.

Projective coordinates is a way to solve this problem. The idea is to represent points in the projective form (X:Y:Z), so when adding two numbers in projective form, you still get a point in a form (X:Y:Z). Then, after conducting a series of additions, you can convert the point back to the affine form.

Bug why adding two points, say, $(X_P : Y_P : Z_P)$ and $(X_Q : Y_Q : Z_Q)$, in the projective form is more efficient? We will not derive the formulas, but trust us that they have the following form:

$$X_{R} = (X_{P}Z_{Q} - X_{Q}Z_{P})(Z_{P}Z_{Q}(Y_{P}Z_{Q} - Y_{Q}Z_{P})^{2} - (X_{P}Z_{Q} - X_{Q}Z_{P})^{2}(X_{P}Z_{Q} + X_{Q}Z_{P}));$$

$$Y_{R} = Z_{P}Z_{Q}(X_{Q}Y_{P} - X_{P}Y_{Q})(X_{P}Z_{Q} - X_{Q}Z_{P})^{2} - (Y_{P}Z_{Q} - Y_{Q}Z_{P})((Y_{P}Z_{Q} - Y_{Q}Z_{P})^{2}Z_{P}Z_{Q} - (X_{P}Z_{Q} + X_{Q}Z_{P})(X_{P}Z_{Q} - X_{Q}Z_{P})^{2});$$

$$Z_{R} = Z_{P}Z_{Q}(X_{P}Z_{Q} - X_{Q}Z_{P})^{3}.$$

$$(18)$$

Do not be afraid, you do not need to understand how this formula is derived. But notice that despite the very scary look, there is no inversions involved! Moreover, this formula can be calculated in only 12M + 2S! So all in all, this is much more effective than 2M + 1S + 1I.

The only inversion which is unavoidable in the projective form is the inversion of Z since after all additions (and doublings) have been made, we need to use map $(X:Y:Z) \mapsto (X/Z,Y/Z)$ to return back to the affine form. However, this inversion is done only once at the end of the computation, so it is not that costly.

Proposition 0.7. To conclude, typically, when working with elliptic curves, one uses the following strategy:

- 1. Convert affine points to projective form using the map $(x, y) \mapsto (x, y, 1)$.
- 2. Perform all operations in the projective form, which do not involve inversions.
- 3. Convert the result back to the affine form using the map $(X : Y : Z) \mapsto (X/Z, Y/Z)$. This is illustrated in the Figure below.

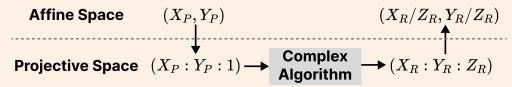


Illustration: General strategy when performing operations over Elliptic Curves.

- 0.2.5 Fast Scalar Multiplication
- 0.3 Elliptic Curve Pairing
- 0.3.1 Definition
- 0.3.2 Case Study: BLS Signature