



# Field Extensions and Elliptic Curves

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## Distributed Lab

 [zkdl-camp.github.io](https://zkdl-camp.github.io)

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# Plan

## 1 Field Extensions

- A bit of intuition
- General Definition
- Polynomial Fraction Rings
- Finite Field Extensions

## 2 Algebraic Closure

- Definition

## 3 Elliptic Curve

- Definition
- Group Structure

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# Field Extensions

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$\mathbb{Q}$  vs  $\mathbb{R}$ *Question #1*

What is the difference between rational numbers  $\mathbb{Q}$  and real numbers  $\mathbb{R}$ ?

*Definition*

Rational numbers  $\mathbb{Q}$  are defined as the set  $\{\frac{n}{m} : n \in \mathbb{Z}, m \in \mathbb{N}\}$ .

*Question #2*

Why cannot we say  $m \in \mathbb{Z}$ , similarly to  $n$ ?

*Theorem*

$\sqrt{2}$  is not a rational number. Neither is  $\pi$  and  $e$ . But they are reals.

*Conclusion*

$\mathbb{R}$  is sort of “an extended version of  $\mathbb{Q}$ ”.

# What about $\mathbb{R}$ ?

## *Rhetorical Question*

Can we extend  $\mathbb{R}$ ?

Yes — just use complex numbers  $\mathbb{C}$ !

## *Definition*

Complex numbers  $\mathbb{C}$  is defined as the set of  $x + iy$  where  $i^2 = -1$ .

## *Definition*

Complex numbers  $\mathbb{C}$  are defined as the set of pairs  $(x, y) \in \mathbb{R}^2$  where addition is defined as  $(x_1, y_1) + (x_2, y_2) = (x_1 + x_2, y_1 + y_2)$ , and the multiplication is:

$$(x_1, y_1) \cdot (x_2, y_2) = (x_1x_2 - y_1y_2, x_1y_2 + x_2y_1).$$

# A bit about complex numbers

## Theorem

$(\mathbb{C}, +, \times)$  is a field.

## Example

Let us see how arithmetic is performed in  $\mathbb{C}$ .

- **Addition:**  $(2 + 3i) + (4 + 5i) = 6 + 8i$ .
- **Multiplication:**  $(1 + i)(2 + i) = 2 + 3i + i^2 = 1 + 3i$ .
- **Division:**

$$\frac{2 + i}{1 + i} = \frac{(2 + i)(1 - i)}{(1 + i)(1 - i)} = \frac{2 - i - i^2}{1 - i^2} = \frac{3 - i}{2} = \frac{3}{2} - \frac{1}{2}i$$

## Question

What is  $(1 + i) + (2 + i)$ ?  $i(1 + i)$ ?  $1/i$ ?

# Field Extension

## Conclusion + Question

$\mathbb{C}$  is sort of “an extended version of  $\mathbb{R}$ ”. Thus, we have

$$\mathbb{Q} \subset \mathbb{R} \subset \mathbb{C}, \text{ where } \mathbb{Q}, \mathbb{R}, \mathbb{C} \text{ are fields}$$

So we have two questions in mind:

- Is there any mathematical term for this?
- Can we go further?

## Definition

Let  $\mathbb{F}$  be a field. A field  $\mathbb{K}$  is called an **extension** of  $\mathbb{F}$  if  $\mathbb{F} \subset \mathbb{K}$  which we denote as  $\mathbb{K}/\mathbb{F}$ .

## Example

$\mathbb{C}/\mathbb{R}$  is a field extension. So is  $\mathbb{R}/\mathbb{Q}$ .

# $\mathbb{Q}(\sqrt{2})$ and $\mathbb{Q}(i)$

## Example

Define  $\mathbb{Q}(\sqrt{2}) := \{p + q\sqrt{2} : p, q \in \mathbb{Q}\}$ . This is a field extension of  $\mathbb{Q}$ . Arithmetic over  $\mathbb{Q}(\sqrt{2})$  looks like:

- Addition:**  $(1 + 2\sqrt{2}) + (3 + 4\sqrt{2}) = 4 + 6\sqrt{2}$ .

- Multiplication:**

$$(1 + 2\sqrt{2})(1 + \sqrt{2}) = 1 + 3\sqrt{2} + 2\sqrt{2}^2 = 5 + 3\sqrt{2}.$$

- Division:**

$$\frac{1 + 2\sqrt{2}}{1 + \sqrt{2}} = \frac{(1 + 2\sqrt{2})(\sqrt{2} - 1)}{(\sqrt{2} + 1)(\sqrt{2} - 1)}$$

## Example

Similarly,  $\mathbb{Q}(i) := \{p + qi : p, q \in \mathbb{Q}\}$  is a field extension of  $\mathbb{Q}$ .



$$\mathbb{Q}(\sqrt{2}, i)$$

### Example

Define  $\mathbb{Q}(\sqrt{2}, i) = \{\alpha + \beta\sqrt{2} : \alpha, \beta \in \mathbb{Q}(i)\}$ . Typicall element of  $\mathbb{Q}(\sqrt{2}, i)$  can be written as:

$$(a + bi) + (c + di)\sqrt{2} = a + c\sqrt{2} + b\sqrt{2}i + di\sqrt{2}$$

### Notice

Each element of  $\mathbb{Q}(\sqrt{2}, i)$  is a linear combination of  $\{1, \sqrt{2}, i, \sqrt{2}i\}$ . This is usually called a **basis**. Moreover, to denote the dimensionality of  $\mathbb{Q}(\sqrt{2}, i)$  over  $\mathbb{Q}$ , we write  $[\mathbb{Q}(\sqrt{2}, i) : \mathbb{Q}] = 4$ .

# Real Polynomials modulo $x^2 + 1$

## Definition... "Kinda"

Consider the set  $\mathcal{P}$  — a set of polynomials  $\mathbb{R}[x]$  modulo  $p(x) := x^2 + 1$ .

## Example

For example,  $1, 5 + x, 3x, 1 + 2x \in \mathcal{P}$ .

But what about  $x^2 + 2x + 4$ ? We can divide by  $x^2 + 1$ !

$$x^2 + 2x + 4 = (x^2 + 1) \cdot 1 + (2x + 3)$$

So in  $\mathcal{P}$ , we have  $x^2 + 2x + 4 = 2x + 3$ .

# Real Polynomials modulo $x^2 + 1$

## Arithmetic

Over this field, we can do arithmetic as usual.

- **Addition:**  $(1 + x) + (2 + 3x) = 3 + 4x$ .
- **Multiplication:**  $(1 + x)(2 + x) = x^2 + 3x + 2 = 3x + 1$ .
- **Inverse:**

$$\left(\frac{1+x}{2}\right)^{-1} = 1 - x$$

Indeed,

$$\frac{1+x}{2} \cdot (1-x) = \frac{1}{2} \cdot (1-x^2) = \frac{1}{2} (-(x^2+1) + 2) = 1 \text{ (in } \mathcal{P})$$

# Hold on a minute...

## Results

- $(1 + x) + (2 + 3x) = 3 + 4x$
- $(1 + x)(2 + x) = 1 + 3x$
- $\left(\frac{1+x}{2}\right)^{-1} = 1 - x$

## Same, but over $\mathbb{C}$

Let us do the same, but instead of  $X$ , use  $i$ .

- $(1 + i) + (2 + 3i) = 3 + 4i$ .
- $(1 + i)(2 + i) = 2 + 3i + i^2 = 1 + 3i$ .
- $\frac{1}{\frac{1+i}{2}} = \frac{2}{1+i} = \frac{2(1-i)}{(1+i)(1-i)} = 1 - i$ .

# Hold on a minute...



So, basically,  $\mathcal{P}$  and  $\mathbb{C}$  have the same structure! Formally, they are isomorphic:  $\mathcal{P} \cong \mathbb{C}$ .

## Question

Could we have used  $x^2 + 3$  instead of  $x^2 + 1$ ? What about  $x^2 + x + 1$ ?

Yes, any **irreducible** 2nd-degree polynomial  $p(x)$  over  $\mathbb{R}$  can be used. Typically, this is denoted as  $\boxed{\mathbb{R}[x]/(p(x))}$ .

# Isomorphisms

## Reminder

For two groups  $(\mathbb{G}, +)$  and  $(\mathbb{H}, \times)$  we defined homomorphism to be a function  $\phi : \mathbb{G} \rightarrow \mathbb{H}$  such that

$$\phi(a + b) = \phi(a) \times \phi(b)$$

However, we claim that  $\mathbb{R}/(x^2 + 1) \cong \mathbb{C}$ , which are fields!

## Definition

A **field isomorphism** is a function  $\phi : (\mathbb{F}, +, \times) \rightarrow (\mathbb{K}, \oplus, \otimes)$  such that

- $\phi(a + b) = \phi(a) \oplus \phi(b)$
- $\phi(a \times b) = \phi(a) \otimes \phi(b)$
- $\phi(1_{\mathbb{F}}) = 1_{\mathbb{K}}$

But for now,  $\cong$  means “exhibit the same structure”.

# Key Theorems

## Theorem

Let  $\mathbb{F}$  be a field and  $\mu(x)$  — irreducible polynomial over  $\mathbb{F}$  (**reduction polynomial**). Consider a set of polynomials over  $\mathbb{F}[x]$  modulo  $\mu(x) \in \mathbb{F}[x]$ , formally denoted as  $\mathbb{F}[x]/(\mu(x))$ . Then,  $\mathbb{F}[x]/(\mu(x))$  is a field.

## Theorem

Let  $\mathbb{F}$  be a field and  $\mu \in \mathbb{F}[X]$  is an irreducible polynomial of degree  $n$  and let  $\mathbb{K} := \mathbb{F}[X]/(\mu(X))$ . Let  $\theta \in \mathbb{K}$  be the root of  $\mu$  over  $\mathbb{K}$ . Then,

$$\mathbb{K} = \{c_0 + c_1\theta + \cdots + c_{n-1}\theta^{n-1} : c_0, \dots, c_{n-1} \in \mathbb{F}\}$$

# Coming back to previous examples

## Example

Again, consider  $\mathbb{Q}(\sqrt{2}) = \{q + p\sqrt{2} : p, q \in \mathbb{Q}\}$ . Then,

$$\mathbb{Q}(\sqrt{2}) \cong \mathbb{Q}[x]/(x^2 - 2)$$

## Example

Similarly,  $\mathbb{Q}(i) \cong \mathbb{Q}[x]/(x^2 + 1)$ .

## Example

And  $\mathbb{Q}(\sqrt{2}, i)$  is just a little bit more tricky. Notice that we can take

$$p(x) := (x^2 - 2)(x^2 + 1) = x^4 - x^2 - 2$$

So  $\mathbb{Q}(\sqrt{2}, i) \cong \mathbb{Q}[x]/(x^4 - x^2 - 2)$ .



# Finite Field Extension

## Definition

Recall that  $\mathbb{F}_p$  (**prime field**) is a set  $\{0, 1, \dots, p-1\}$  with arithmetic modulo  $p$ .

In many cases, we need to extend  $\mathbb{F}_p$  2, 4, 8, 12, 24 times. For this, we use the so-called **finite field extension**.

## Definition

Suppose  $p$  is prime and  $m \geq 2$ . Let  $\mu \in \mathbb{F}_p[X]$  be an irreducible polynomial of degree  $m$ . Then, elements of  $\mathbb{F}_{p^m}$  are polynomials in  $\mathbb{F}_p^{(\leq m)}[X]$  modulo  $\mu(x)$ . In other words,

$$\mathbb{F}_{p^m} = \{c_0 + c_1X + \dots + c_{m-1}X^{m-1} : c_0, \dots, c_{m-1} \in \mathbb{F}_p\},$$

where all operations are performed modulo  $\mu(X)$ .

# Examples

It would be convenient to build  $\mathbb{F}_{p^2}$  as  $\mathbb{F}_p[i]/(i^2 + 1)$ , but is it always possible? In other words, when  $X^2 = -1$  has a solution in  $\mathbb{F}_p$ ?

## Theorem

*Let  $p$  be an odd prime. Then  $X^2 + 1$  is irreducible in  $\mathbb{F}_p[X]$  if and only if  $p \equiv 3 \pmod{4}$ .*

## Example

Pick  $p = 19$ . Then  $\mathbb{F}_{361} := \mathbb{F}_{19}[i]/(i^2 + 1)$ . So typical elements are:  $1 + 3i$ ,  $10 + 15i$ ,  $18 + 18i$ ,  $5$ ,  $7i$ , ...

- **Addition:**  $(1 + 10i) + (18 + 15i) = 19 + 25i = 6i$ .
- **Multiplication:**  
 $(5 + 6i)(6 + 7i) = 30 + 71i + 42i^2 = -12 + 71i = 7 + 14i$ .

# More Examples: Binary Extension Fields

## Example

Consider the  $\mathbb{F}_{2^4}$ . Then, there are 16 elements in this set:

$$0, 1, X, X + 1,$$

$$X^2, X^2 + 1, X^2 + X, X^2 + X + 1,$$

$$X^3, X^3 + 1, X^3 + X, X^3 + X + 1,$$

$$X^3 + X^2, X^3 + X^2 + 1, X^3 + X^2 + X, X^3 + X^2 + X + 1.$$

Set  $\mu(X) := X^4 + X + 1$ . Then, operations are performed in the following manner:

- **Addition:**  $(X^3 + X^2 + 1) + (X^2 + X + 1) = X^3 + X$ .
- **Multiplication:**  $(X^3 + X^2 + 1) \cdot (X^2 + X + 1) = X^2 + 1$  since:
- **Inversion:**  $(X^3 + X^2 + 1)^{-1} = X^2$  since  
 $(X^3 + X^2 + 1) \cdot X^2 \bmod (X^4 + X + 1) = 1$ .

# More Examples: BN254

## Example

Consider the **BN254 scalar field**, used in SNARKs:

$$p = 0x30644e72e131a029 \cdots a8d3c208c16d87cfd47$$

- Then,  $\mathbb{F}_{p^2} := \mathbb{F}_p[u]/(u^2 + 1)$  since  $p \equiv 3 \pmod{4}$ .
- Define  $\xi := 9 + u \in \mathbb{F}_{p^2}$ . Then, set  $\mathbb{F}_{p^6} := \mathbb{F}_{p^2}[v]/(v^3 - \xi)$ .
- Finally, set  $\mathbb{F}_{p^{12}} := \mathbb{F}_{p^6}[w]/(w^2 - v)$ .

Equivalently, we can write:

$$\mathbb{F}_{p^{12}} := \mathbb{F}_p[w]/(w^{12} - 18w^6 + 82)$$

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# Algebraic Closure

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# Definition

## Definition

A field  $\mathbb{F}$  is called **algebraically closed** if every non-constant polynomial  $p(x) \in \mathbb{F}[X]$  has a root in  $\mathbb{F}$ .

## Example

$\mathbb{R}$  is not algebraically closed since  $X^2 + 1$  has no roots in  $\mathbb{R}$ . However,  $\mathbb{C}$  is algebraically closed, which follows from the fundamental theorem of algebra. Since  $\mathbb{C}$  is a field extension of  $\mathbb{R}$ , it is also an algebraic closure of  $\mathbb{R}$ . This is commonly denoted as  $\overline{\mathbb{R}} = \mathbb{C}$ .

## Definition

A field  $\mathbb{K}$  is called an **algebraic closure** of  $\mathbb{F}$  if  $\mathbb{K}/\mathbb{F}$  is algebraically closed. This is denoted as  $\overline{\mathbb{F}} = \mathbb{K}$ .

# Algebraic Closure for Finite Fields

Recall that we are cryptographers, not mathematicians. So we are interested in  $\overline{\mathbb{F}}_p$ . So I have two news to you:

- **Good news:**  $\overline{\mathbb{F}}_p$  exists.
- **Bad news:**  $\overline{\mathbb{F}}_p$  is infinite.

## Theorem

*No finite field  $\mathbb{F}$  is algebraically closed.*

**Proof.** Suppose  $a_1, a_2, \dots, a_n \in \mathbb{F}$  are all elements of  $\mathbb{F}$ . Consider

$$p(x) := \prod_{i=1}^n (x - a_i) + 1 = (x - a_1)(x - a_2) \cdots (x - a_n) + 1.$$

Clearly,  $p(x)$  is a non-constant polynomial and has no roots in  $\mathbb{F}$ , since for any  $a \in \mathbb{F}$ , one has  $p(a) = 1$ . ■

## So what?

But what form does the  $\overline{\mathbb{F}}_p$  have? Well, it is a union of all  $\mathbb{F}_{p^k}$  for  $k \geq 1$ . This is formally written as:

$$\overline{\mathbb{F}}_p = \bigcup_{k \in \mathbb{N}} \mathbb{F}_{p^k}$$

### Remark

But this definition is super counter-intuitive! So here how we usually interpret it. Suppose I tell you that polynomial  $q(x)$  has a root in  $\overline{\mathbb{F}}_p$ . What that means is that there exists some extension  $\mathbb{F}_{p^m}$  such that for some  $\alpha \in \mathbb{F}_{p^m}$ ,  $q(\alpha) = 0$ . We do not know how large this  $m$  is, but we know that it exists. For that reason,  $\overline{\mathbb{F}}_p$  is defined as an infinite union of all possible field extensions.



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# Elliptic Curve

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# Definition

## Definition

Suppose that  $\mathbb{K}$  is a field. An **elliptic curve**  $E$  over  $\mathbb{K}$  is defined as a set of points  $(x, y) \in \mathbb{K}^2$ :

$$y^2 = x^3 + ax + b,$$

called a **Short Weierstrass equation**, where  $a, b \in \mathbb{K}$  and  $4a^3 + 27b^2 \neq 0$ . We denote  $E/\mathbb{K}$  to denote the elliptic curve over field  $\mathbb{K}$ .

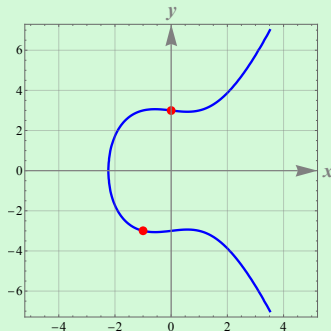
## Definition

We say that  $P = (x_P, y_P) \in \mathbb{A}^2(\mathbb{K})$  is the **affine representation** of the point on the elliptic curve  $E/\mathbb{K}$  if it satisfies the equation  $y_P^2 = x_P^3 + ax_P + b$ .

# Examples

## Example

Consider  $E/\mathbb{Q} : y^2 = x^3 - x + 9$ . Valid affine points on  $E/\mathbb{Q}$  are, for example,  $P = (0, 3)$ ,  $Q = (-1, -3) \in \mathbb{A}^2(\mathbb{Q})$ .



# More Examples

Some more examples<sup>1</sup>:

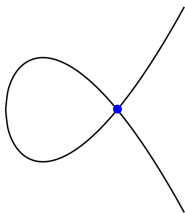


Figure 2.1:  
Singular curve  
 $y^2 = x^3 - 3x + 2$   
over  $\mathbb{R}$ .

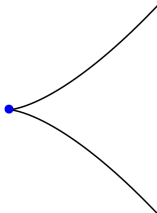


Figure 2.2:  
Singular curve  
 $y^2 = x^3$   
over  $\mathbb{R}$ .

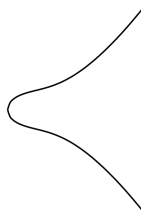


Figure 2.3:  
Smooth curve  
 $y^2 = x^3 + x + 1$   
over  $\mathbb{R}$ .

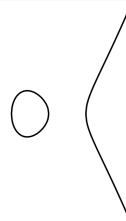


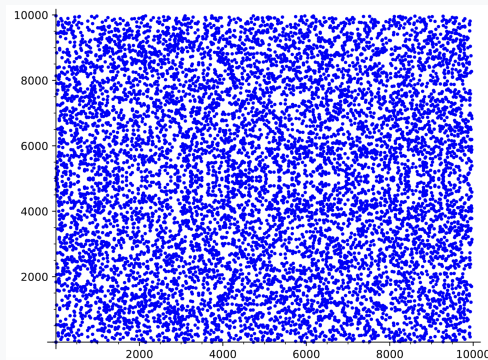
Figure 2.4:  
Smooth curve  
 $y^2 = x^3 - x$   
over  $\mathbb{R}$ .

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<sup>1</sup>Figure taken from "Pairings for Beginners"

# Real Elliptic Curves

But real elliptic curves are not that simple. Here how they look like<sup>2</sup>:



**Figure:** Curve  $E/\mathbb{F}_{9973} : y^2 = x^3 - 2x + 1$  over the finite field

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<sup>2</sup>Figure taken from “Moonmath”

# Defining a Group Structure: A Few Words

## Definition

The set of points on the curve, denoted as  $E_{a,b}(\mathbb{K})$ , is defined as:

$$E_{a,b}(\mathbb{K}) = \{(x, y) \in \mathbb{A}^2(\mathbb{K}) : y^2 = x^3 + ax + b\} \cup \{\mathcal{O}\},$$

where  $\mathcal{O}$  is the so-called **point at infinity**.

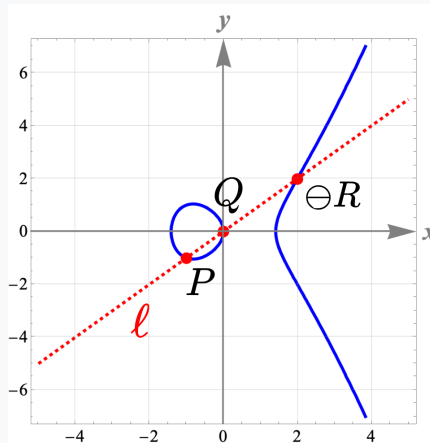
## Remark #1

If  $(x_P, y_P) \in E(\mathbb{K})$  then  $(x_P, -y_P) \in E(\mathbb{K})$ .

## Remark #2

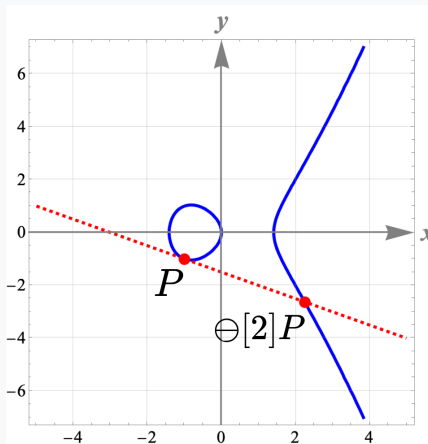
Typically,  $\mathbb{K} = \overline{\mathbb{F}}_p$ : we do not concretize over which finite field we define the elliptic curve.

# Defining a Group Structure: Chord Method



**Figure:** Chord method for adding two points

# Defining a Group Structure: Tangent Method



**Figure:** Tangent method for the point doubling



# Idea of Derivation

Line equation through  $P = (x_P, y_P)$ ,  $Q = (x_Q, y_Q)$ :

$$\ell : y = \lambda(x - x_P) + y_P, \quad \lambda = \frac{y_Q - y_P}{x_Q - x_P}$$

So all we need is to solve the system of equations:

$$\begin{cases} y^2 = x^3 + ax + b \\ y = \lambda(x - x_P) + y_P \end{cases}$$

Substituting  $y$  from the second equation to the first one, we get a cubic equation. Using Vieta's formula, one can derive

$$x_P + x_Q + x_R = \lambda^2$$

The rest is easy to finish.

# Group Law

## Definition

1. Point at infinity  $\mathcal{O}$  is an identity element.
2. If  $x_P \neq x_Q$ , use the **chord method**. Define  $\lambda := \frac{y_P - y_Q}{x_P - x_Q}$  — the slope between  $P$  and  $Q$ . Set the resultant coordinates as:

$$x_R := \lambda^2 - x_P - x_Q, \quad y_R := \lambda(x_P - x_R) - y_P.$$

3. If  $x_P = x_Q$  and  $y_P = y_Q$  (that is,  $P = Q$ ), use the **tangent method**. Define the slope of the tangent at  $P$  as  $\lambda := \frac{3x_P^2 + a}{2y_P}$  and set

$$x_R := \lambda^2 - 2x_P, \quad y_R := \lambda(x_P - x_R) - y_P.$$

4. Otherwise, define  $P \oplus Q := \mathcal{O}$ .

# One more Illustration

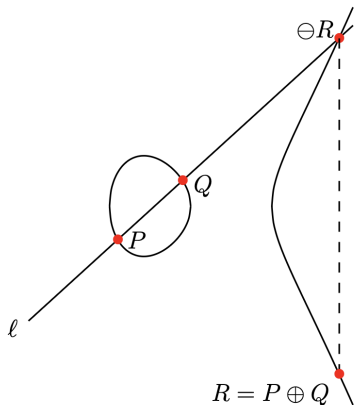


Figure 2.5: Elliptic curve addition.

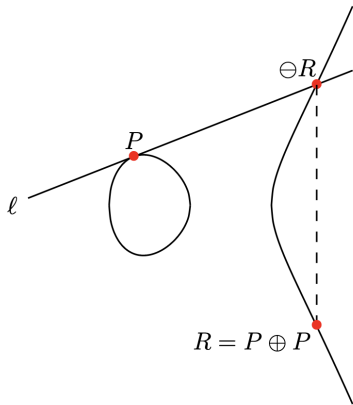


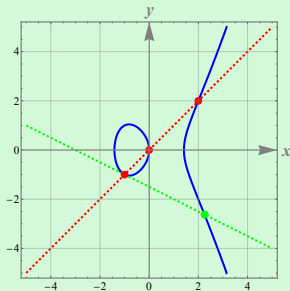
Figure 2.6: Elliptic curve doubling.

# Example

## Example

Consider  $E/\mathbb{R} : y^2 = x^3 - 2x$ .

- **Addition:**  $(-1, 1) \oplus (0, 0) = (2, -2)$ ,  $(2, 2) \oplus (-1, -1) = (0, 0)$ .
- **Doubling:**  $[2](-1, -1) = (\frac{9}{4}, -\frac{21}{8})$ .



# Hasse's Theorem

## Theorem

$(E(\overline{\mathbb{F}}), \oplus)$  forms an abelian group.

Now, let us consider the group order  $r := |E(\mathbb{F}_{p^m})|$ .

## Theorem

**Hasse's Theorem on Elliptic Curves.**  $r = p^m + 1 - t$  for some integer  $|t| \leq 2\sqrt{p^m}$ . A bit more intuitive explanation: the number of points on the curve is close to  $p^m + 1$ . The value  $t$  is called the **trace of Frobenius**.

## Remark

In fact,  $r = |E(\mathbb{F}_{p^m})|$  can be computed in  $O(\log(p^m))$ , so the number of points can be computed efficiently even for fairly large primes  $p$ .

# Discrete Logarithm

## Definition

Let  $P \in E(\overline{\mathbb{F}}_p)$  and  $\alpha \in \mathbb{Z}_r$ . Define the scalar multiplication  $[\alpha]P$  as adding  $P$  to itself  $\alpha - 1$  times (also set  $[0]P := \mathcal{O}$ ).

## Definition


Suppose  $E$  is cyclic, meaning,  $\langle G \rangle = E$  for some  $G \in E$ . The **discrete logarithm problem** on  $E$  consists in the following: suppose  $P = [\alpha]G$  for some  $\alpha \in \mathbb{Z}_r$ . Find  $\alpha$  based on  $P$ .

## Remark

If  $r$  is a product of primes  $p_1, p_2, \dots, p_k$  such that  $p_1 < p_2 < \dots < p_k$ , then the best-known algorithm to solve the discrete logarithm problem is no significantly better than  $O(\sqrt{p_1})$ .

# Thank you for your attention



 [zkdl-camp.github.io](https://zkdl-camp.github.io)

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