

Cross Site Request Forgery

“Who Left Open the Cookie Jar”?

OWASP Top Ten

(2013)

A-1	Injection	Untrusted data is sent to an interpreter as part of a command or query.
A-2	Authentication and Session Management	Attacks passwords, keys, or session tokens, or exploit other implementation flaws to assume other users' identities.
A-3	Cross-site scripting	An application takes untrusted data and sends it to a web browser without proper validation or escaping
...	Various implementation problems	...expose a file, directory, or database key without access control check, ...misconfiguration, ...missing function-level access control
A-8	Cross-site request forgery	A logged-on victim's browser sends a forged HTTP request, including the victim's session cookie and other authentication information

More OWASP

- « Cross-Site Request Forgery (CSRF) is an attack that **forces an end user** to execute unwanted actions on a web application in which they're **currently authenticated**... With a little help of social engineering (such as sending a link via email or chat), an attacker may trick the users of a web application into executing actions of the attacker's choosing... »

2007: Gmail is hacked ...

- While logged into Gmail, a user visiting a malicious site would generate a request understood as originating from the victim user
- This was used to inject an email filter forwarding the victim user's email to attacker
- Allowed an attacker to gain control of davidairey.com (since the domain registrar used email based authentication ...)

Browser execution model

- Each browser window / frame
 - Uploads web content
 - Renders web content, static (HTML, subframes) or dynamic(scripts) to display the page
 - including external resources like images
 - Responds to events (see below)
- Events
 - Rendering: OnLoad
 - Timing: setTimeout(), clearTimeout()
 - Reacting to user actions: OnClick, OnMouseover

Maintaining Client State

- Web interactions are stateless by nature
 - HTTP requests sent back and forth
- How to know which browser connects?
- Methods for maintaining state:
 - Cookies: browser state
 - Sessions: server state
 - URL rewriting: browser state
 - Even more alternatives: cf. http://en.wikipedia.org/wiki/HTTP_cookie

State management: Cookies

- "Small piece of information that scripts can store on a client-side machine"
- Can be set in HTTP header

- Origin and expiration date

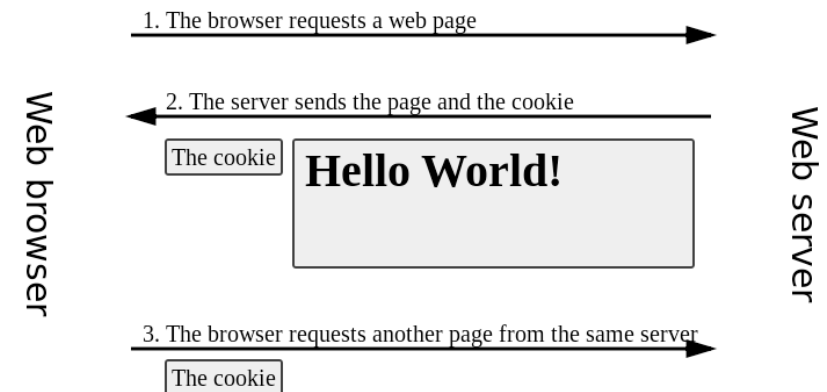
- Example:

HTTP/1.0 200 OK

Content-type: text/html

Set-Cookie: name=value

Set-Cookie: name2=value2; Expires=Wed, 09 Jun 2021 10:18:14 GMT
(content of page)



- Operation:
 - When browser connects to URL, it first checks for relevant cookie
 - If it finds a cookie for the URL, it sends the cookie info to server with the HTTP request
 - A web page can contain content from several web sites, hence several cookies can be sent during its browsing
- Long-lived: user identification (preferences, authentication, tracking ...)
 - Cookie = user ID, may be secured (integrity, confidentiality)
- Temporary: session identification
 - Cookie = random number
- "Secure" attribute instructs that cookie should only be sent over HTTPS (confidentiality to prevent man-in-the-middle attack)

HTTP Cookies Security History



- 1994: Netscape – cookies originate from and still largely based on that 4 page draft
- 1997: RFC 2109 – privacy issues, intention
- 2000: RFC 2965 – further recommendations on usage
- 2002: HttpOnly (XSS)
- 2011: RFC6265 -
- 2017-ongoing: RFC 6265bis (draft) - SameSite

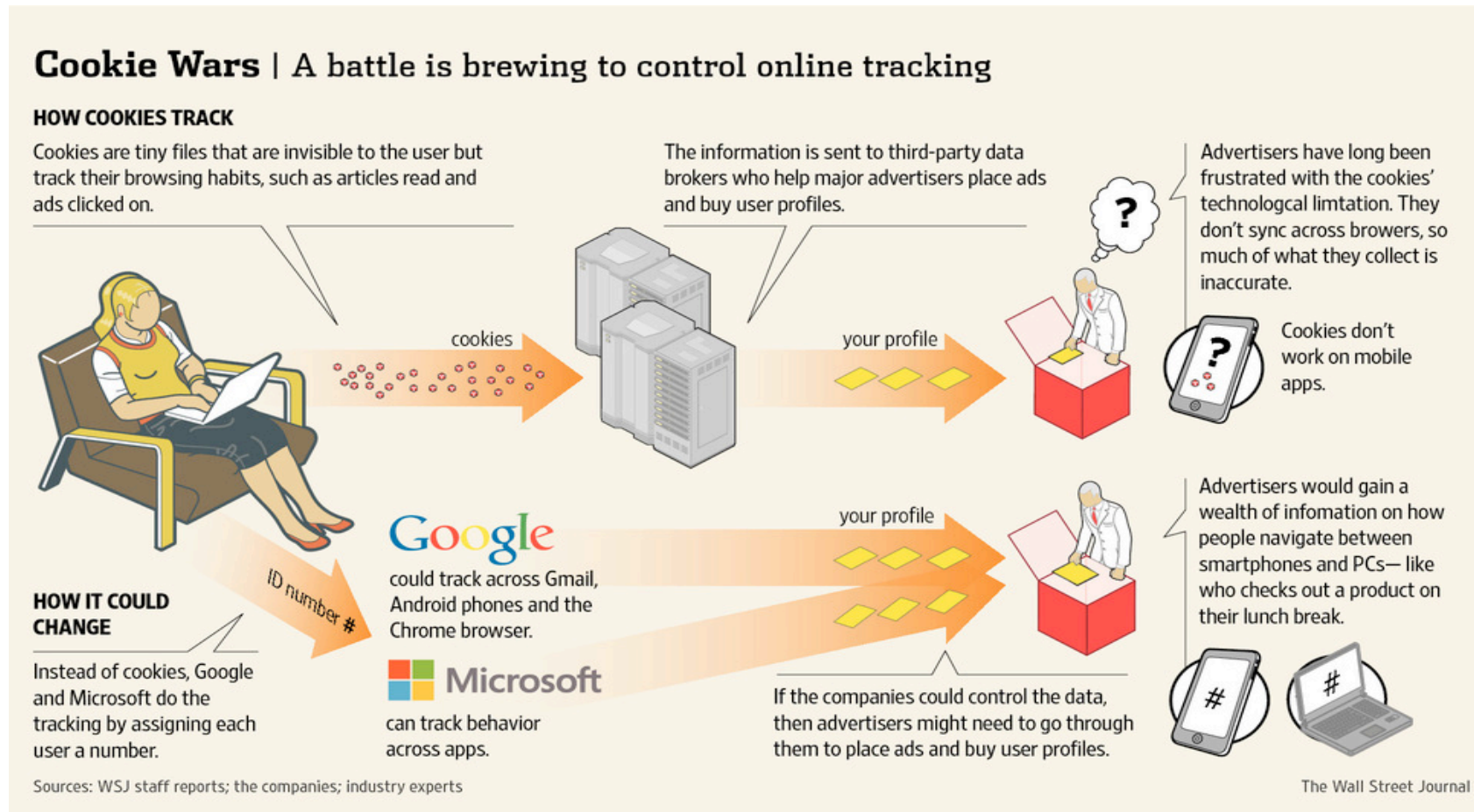
State management: Sessions

- Generally handled by a web framework
- Helps to distinguish between other simultaneous sessions
- Data storage:
 - Session stores data from ongoing transactions (workflow, shopping cart, login)
 - Information can also be removed from a session
- Operation:
 - Start session
 - Session ID is set in the browser (cookie at the beginning, or URL rewriting later on)
 - Data stored and managed on web server (costly, does not scale)
 - End session (dispose of data)
- Pros/Cons: data managed at and by server

State management: URL rewriting

- URLs modified to:
 - store parameters (RESTful approach)
E.g., `http://host:port/shopping.html;sessionid=value`
 - Force the use of a proxy: destination becomes a parameter
- Operation (example: Google)
 - Research result leads to:
`https://www.google.fr/url?q=http://fr.wikipedia.org/Cookie_(informatique)&sa=U&ei=U-9wU-27O8Gm0AWc2IGAAQ&usg=AFQjCNEItv3EUaJHvFL_fM-_7lmX9VzCLQ&sig2=Wdr5pg0cOye893nHZJO-hw&bvm=bv.66330100,d.bGQ`
 - Instead of: `http://fr.wikipedia.org/wiki/Cookie_%28informatique%29`
 - Invisible on the page (link is not displayed in plain text), only in the link bar
- Pros: cannot be suppressed by client

Big Data Wars ...



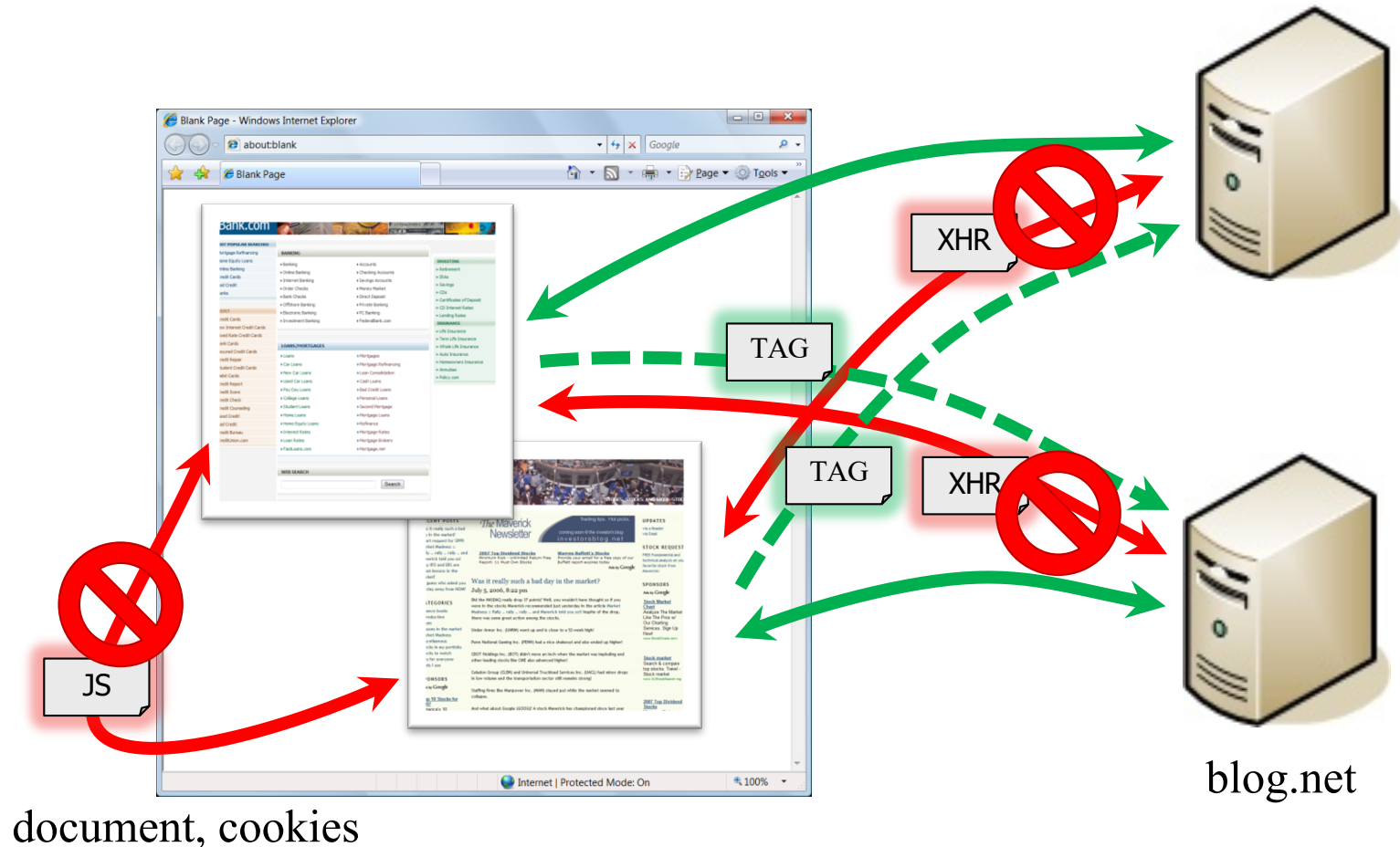
- Cookie lifetime too is a serious issue wrt. Privacy !
 - Expires / Max-Age attributes
- The application should invalidate irrelevant cookies and not rely on browser for removing them

The Browser “Same Origin” Policy (SOP)

- Every frame in a browser is associated with a domain
 - A domain is determined by the server, protocol, and port from which the frame content was downloaded
 - If a frame explicitly includes external code, this code will execute within the frame domain even though it comes from another host
- A script can only access resources (and notably cookies) associated with the same origin
 - prevents hostile script from tampering with other pages in the browser
 - prevents script from snooping on input (passwords) of other windows
- Security Problems: mostly browser bugs
 - Especially in the late 1990s – early 2000s

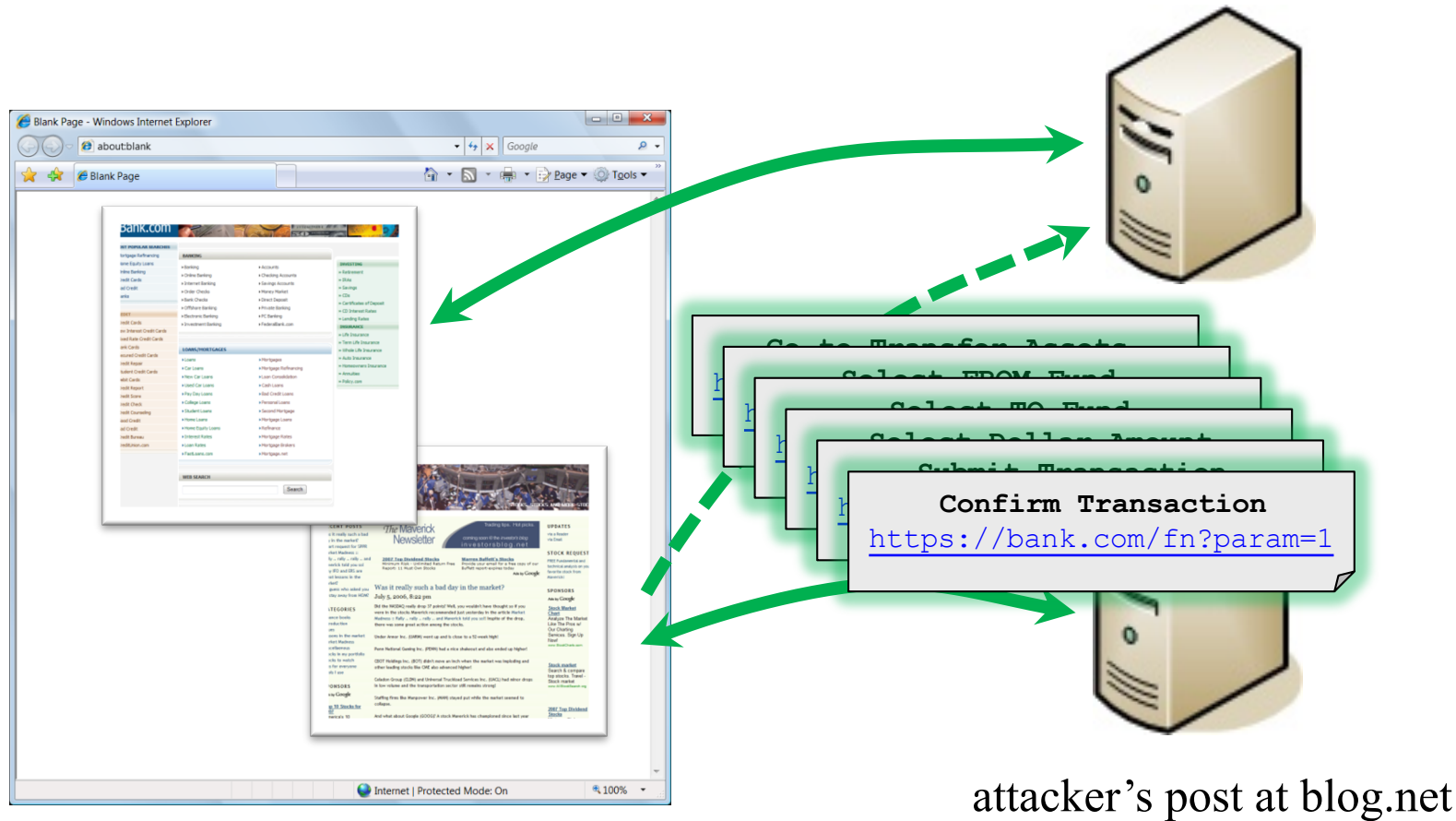
The Browser “Same Origin” Policy

bank.com



Cross-Site Request Forgery

bank.com



How Does CSRF Work?

- Hijacks inherent browser functionality and some aspects of HTTP specification
 - SOP controls and cookies
- Privilege escalation type of attack
 - “Confused deputy”: browser thinks tag/form/XHR is from same origin as destination
- Attacker performs blind attacks (cannot see server responses)
 - Unless combined with XSS ...

- Tags

```
  
<iframe src="https://bank.com/fn?param=1">  
<script src="https://bank.com/fn?param=1">
```

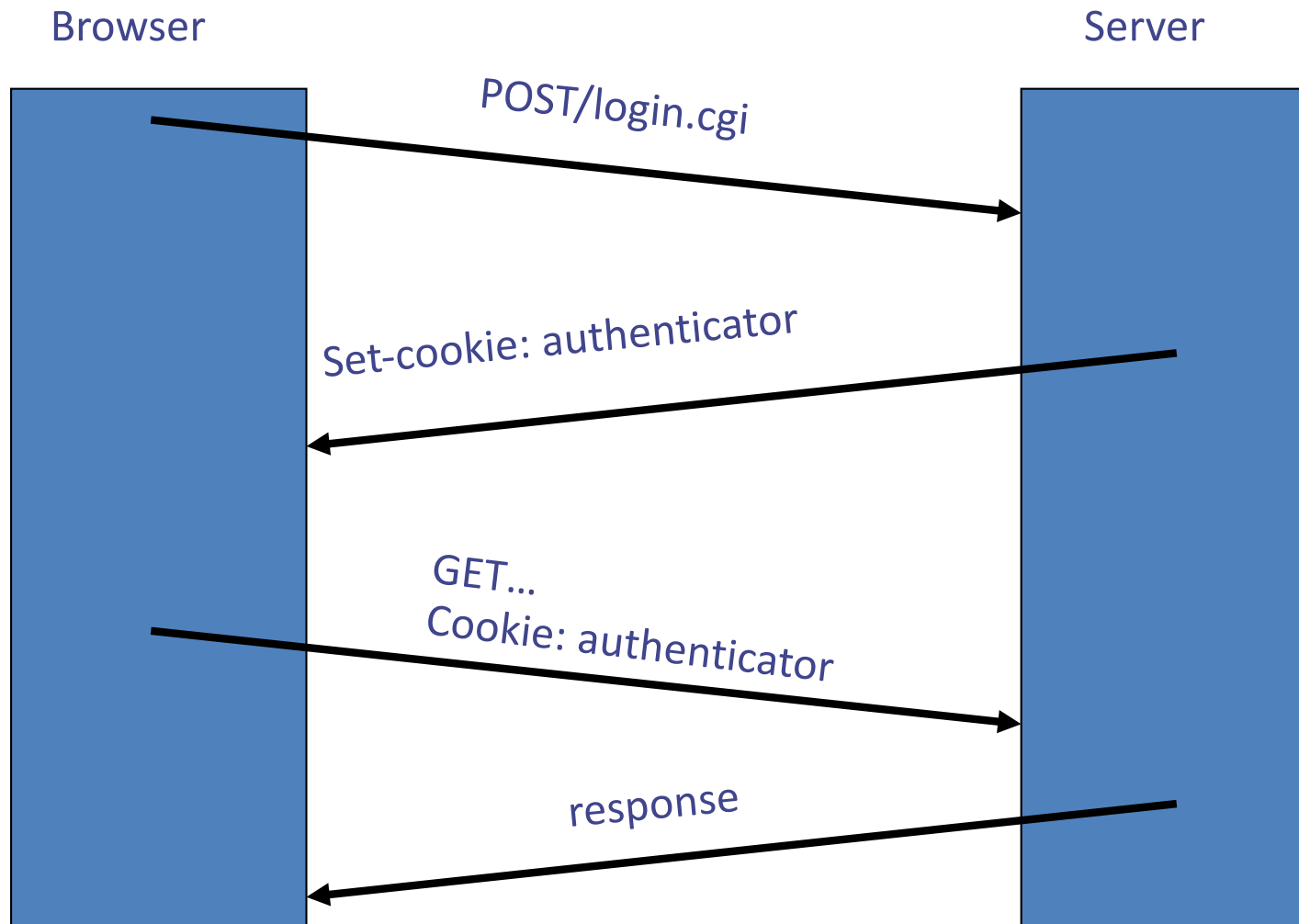
- Autoposting Forms

```
<body onload="document.forms[0].submit()">  
<form method="POST" action="https://bank.com/fn">  
  <input type="hidden" name="sp" value="8109"/>  
</form>
```

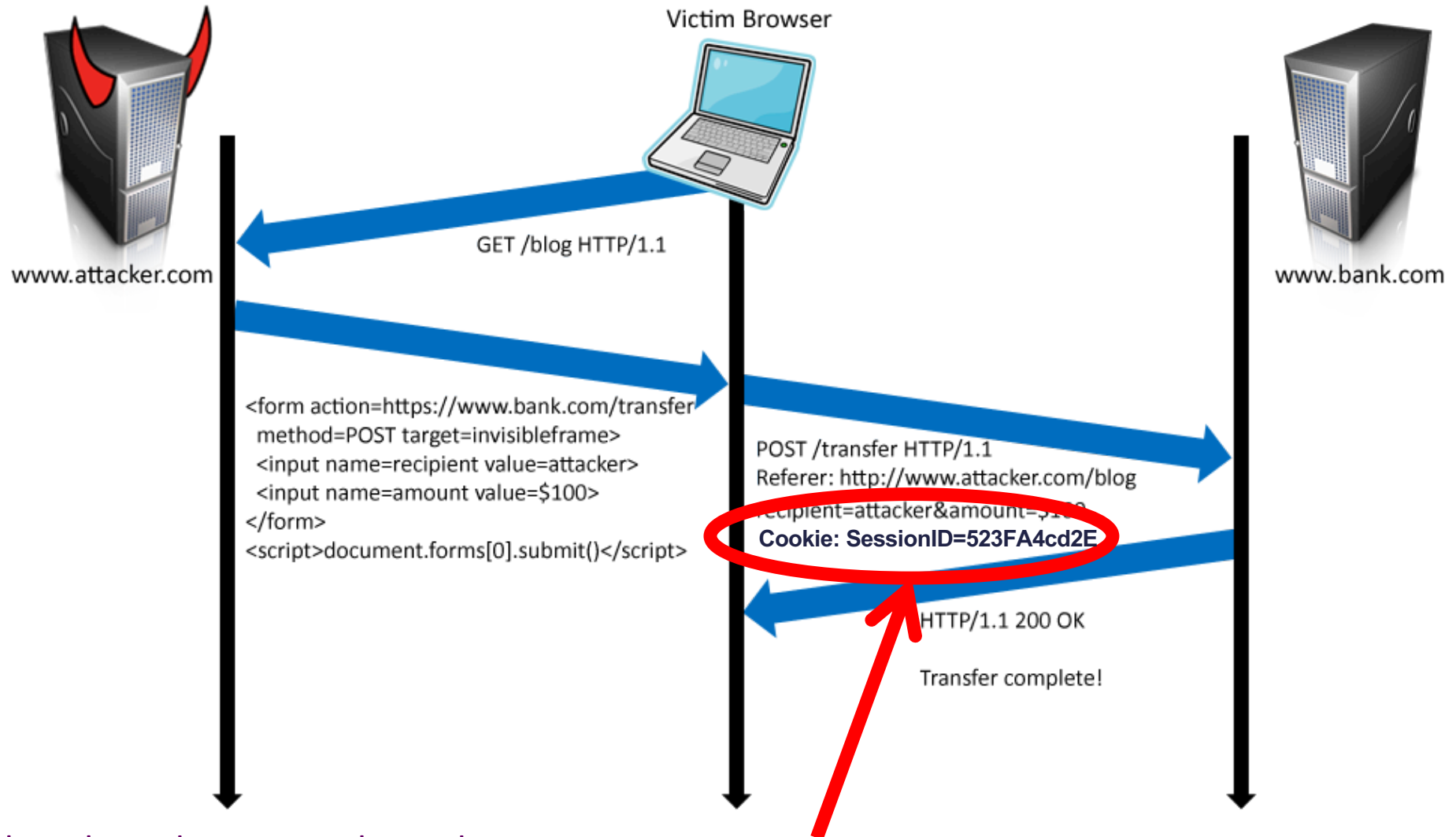
- GET requests are the most dangerous, but any request is vulnerable (POST too)
- XMLHttpRequest (AJAX)
 - ▶ Normally subject to same origin policy
 - ▶ But poorly managed CORS (Cross-Origin Resource Sharing) may relax these constraints ...
 - ▶ May be fooled by a proxy too

Authentication: session using cookies

- Browser behavior: automatically attaches cookie previously set by server



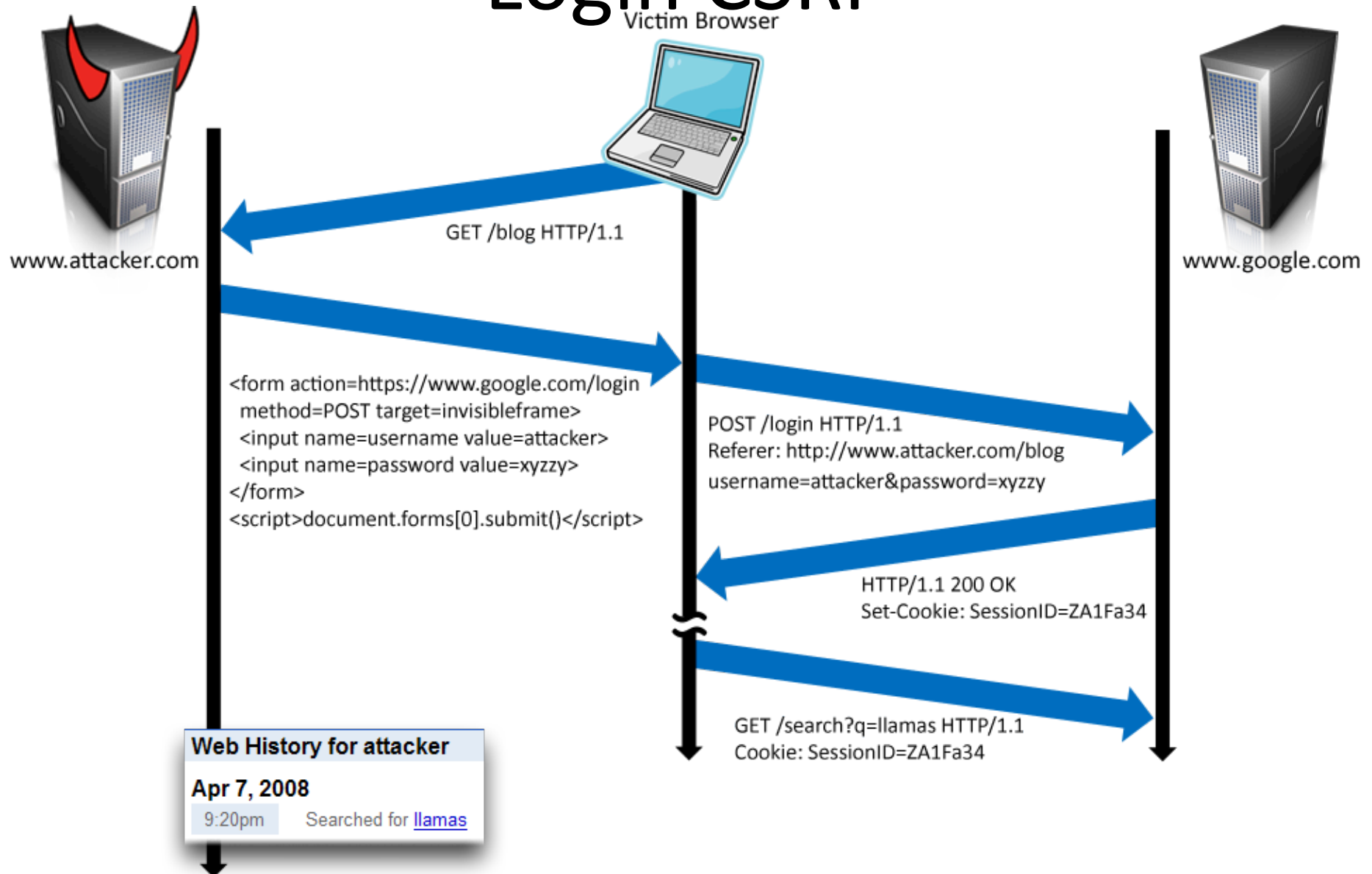
CSRF: Form post with cookie



Q: how long do you stay logged in to Gmail? Facebook?

Login CSRF

Victim Browser

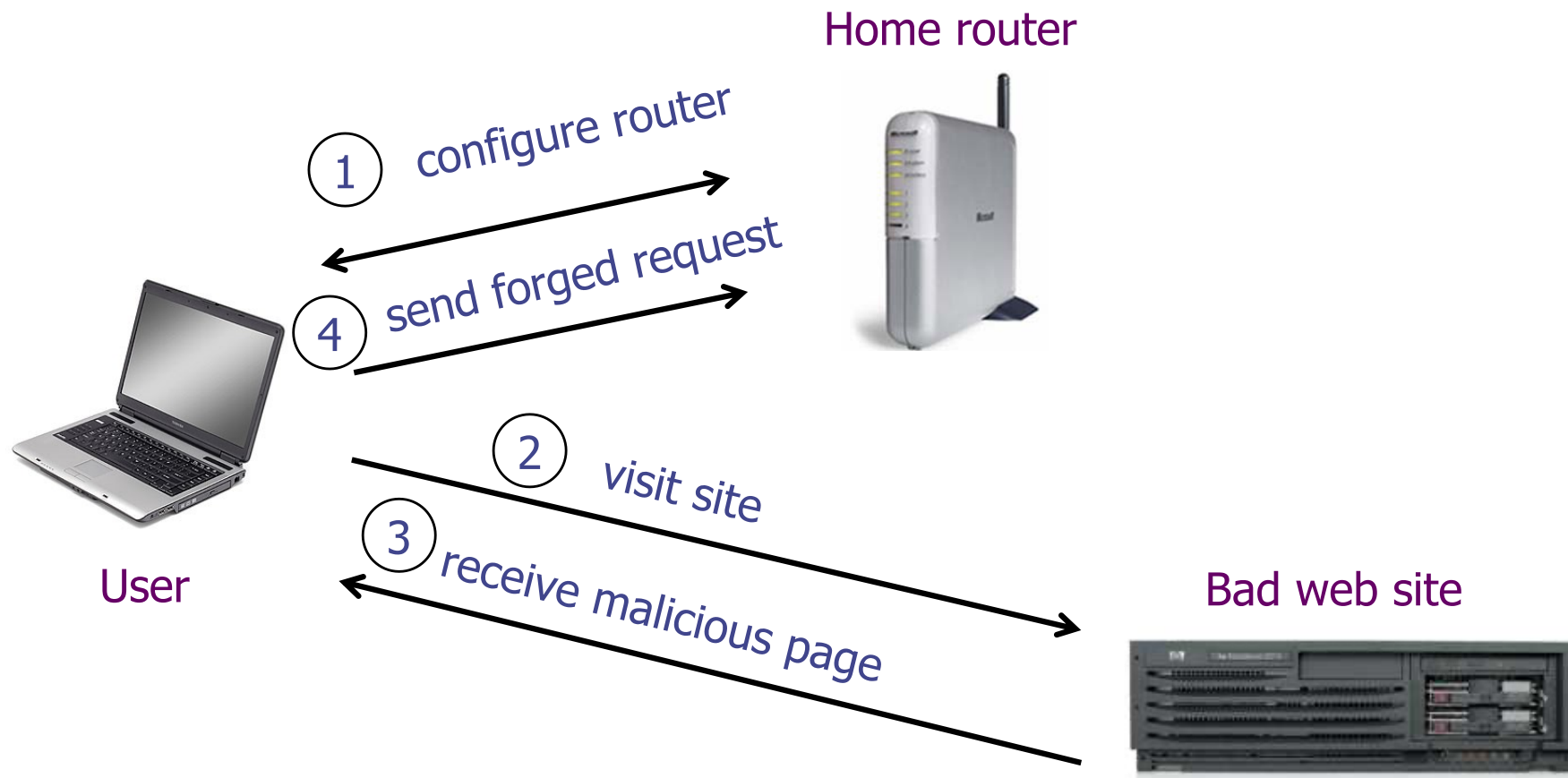


Not just Web Servers: Attacks on Home Router

[Stamm, Ramzan, Jakobsson 2006]

- Fact:
 - 50% of home users have broadband router with a default or no password
- Drive-by Pharming attack:
 - Scenario: user visits malicious site
 - Attacker script scans home network for broadband router:
 - SOP allows “send only” messages
 - Detect success using onError and likely address (e.g., 192.168.0.1):
``
 - Attacker script can login to router and change DNS server
 - Takes control of user navigation
 - Attacker can distribute malware to router
 - Attacker can block virus definition updates
 - Attacker can advertise vulnerable hosts

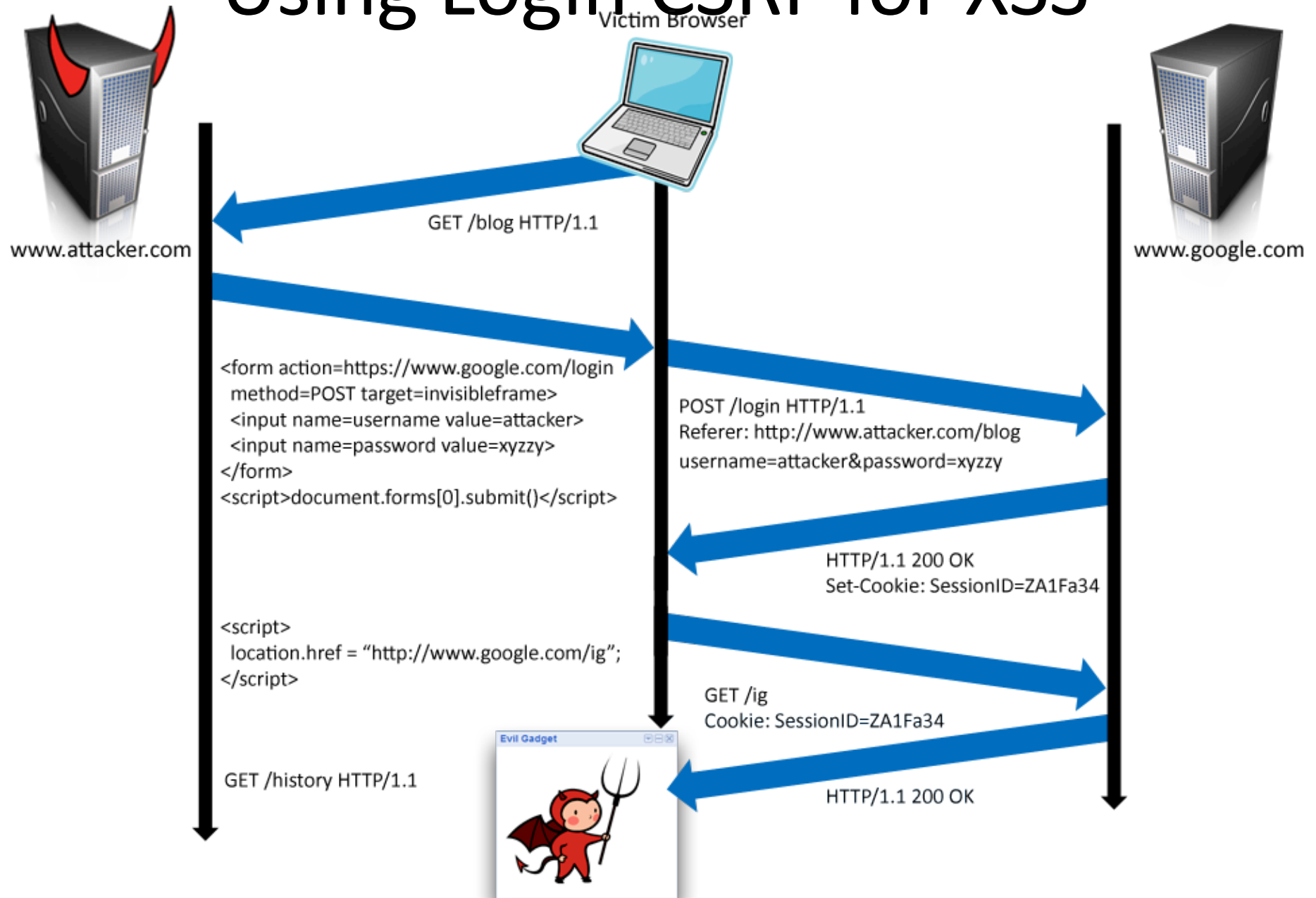
Cookieless Example: Home Router



Broader view of CSRF

- Abuse of cross-site data export feature
 - From user's browser to honest server
 - Disrupts integrity of user's session
- Why mount a CSRF attack?
 - Network connectivity
 - Read browser state
 - Write browser state
- Not just “session riding”

Using Login CSRF for XSS



The attacker's perspective

- The attacker can:
 - Control the form/XHR payload
 - Control the content type (« enctype » attribute)
 - Control the method (GET or POST)
- The attacker cannot:
 - Control other headers
 - Control cookies

CSRF Basic Defenses

- Referred Validation

```
Referer: http://www.facebook.com/home.php
```

- Persistent authentication (login/session data):
 - Client-Side Storage of session information (not effective)
 - But vulnerable to XSS attacks ...
 - ... and user manipulation of server state !
 - Server-Side session ID + Secret Token Validation

```
<input type=hidden value=23a3af01b>
```

- Custom HTTP Header: simpler approach for AJAX/XHR

```
X-Requested-By: XMLHttpRequest
```


Referer Validation Defense

- HTTP Referer header

- Referer: <http://www.gmail.com/>

OK

- Referer: <http://www.bad.com/evil.html>

KO

- Referer:

???

- Lenient Referer validation

- Doesn't block request if Referer is missing

- Strict Referer validation

- Secure, but Referer is sometimes absent...

Referer Privacy Problems

- Referer may also leak privacy-sensitive information!

`http://intranet.corp.apple.com/
projects/iphone/competitors.html`

- May be removed based on user preference in browser
- Site often cannot afford to block these users

So ... Lenient Referrer Checking?

- Other common sources of blocking:
 - Network stripping by the organization (proxy)
 - Network stripping by local machine
 - Stripped by browser for HTTPS -> HTTP transitions
 - Buggy user agents
- Insecure: attacker may strip referrer, e.g.:

```
ftp://www.attacker.com/index.html
  javascript:"<script> /* CSRF */ </script>"
  data:text/html,<script> /* CSRF */ </script>
```

Secret Token Validation

- Requests include a hard-to-guess secret
 - Unguessability replaces unforgeability
- Variations
 - Session identifier
 - Session-independent token
 - Session-dependent token
 - HMAC / MD5 / SHA-1 of session identifier for integrity protection

Secret Token Validation

slicehost

https://manage.slicehost.com/slices/new

Slices DNS Help Account

My Slices

Add a Slice

Add a Slice

Slice Size

- ☒ 256 slice \$20.00/month – 10GB HD, 100GB BW
- ☐ 512 slice \$38.00/month – 20GB HD, 200GB BW
- ☐ 1GB slice \$70.00/month – 40GB HD, 400GB BW
- ☐ 2GB slice \$130.00/month – 80GB HD, 800GB BW
- ☐ 4GB slice \$250.00/month – 160GB HD, 1600GB BW
- ☐ 8GB slice \$450.00/month – 320GB HD, 2000GB BW
- ☐ 15.5GB slice \$800.00/month – 620GB HD, 2000GB BW

System Image

Ubuntu 8.04.1 LTS (hardy)

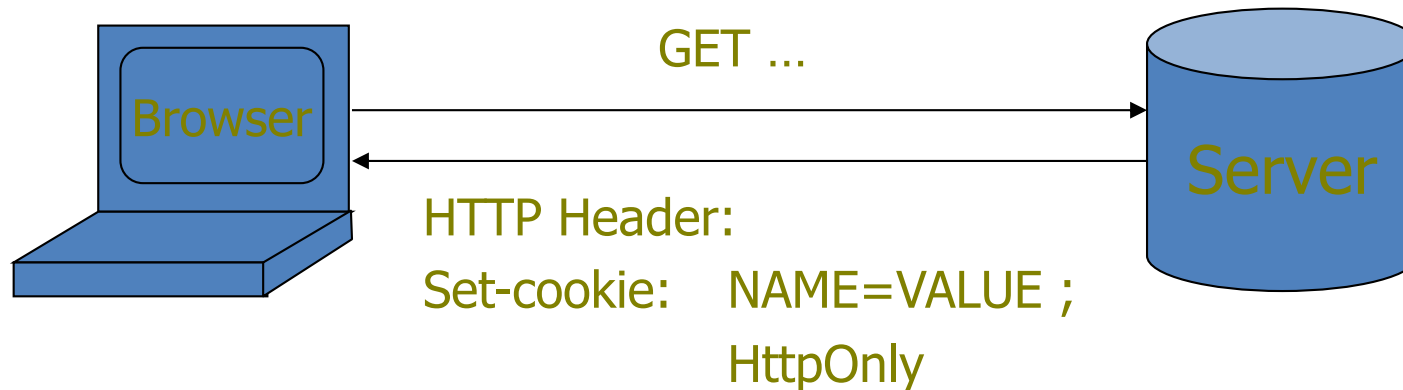
Slice Name

or [cancel](#)

NOTE: You will be charged a prorated amount based upon the number of days remaining in your

```
g:0"><input name="authenticity_token" type="hidden" value="0114d5b35744b522af8643921bd5a3d899e7fbd2" /></div>
```

XSS: HttpOnly Cookies

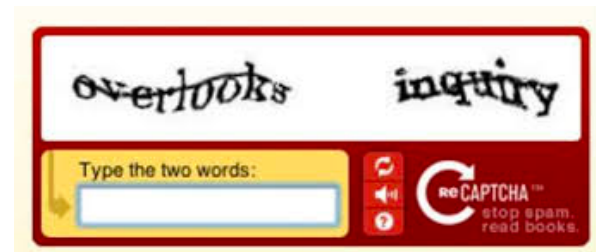


- Cookie sent over HTTP(s), but not accessible to scripts
 - cannot be read via `document.cookie`
 - Also blocks access from XMLHttpRequest headers
 - Helps prevent cookie theft via XSS

... but does not stop most other risks: typical attack is to overflow cookie repository (replace cookie with attacker value)! This is dependent on browser implementation ...

Other Mitigation Strategies

- Tokens: double-submission if maintaining CSRF token on server-side is problematic: token to be sent in header (request parameter) + cookie in body
 - Strong requirements (notably HTTPS to prevent attackers from injecting cookies, encrypted cookies)
- Additional anti-CSRF HTML elements
 - Origin header
 - SameSite cookies (draft RFC 6265bis since 2017)
 - Check https://cheatsheetseries.owasp.org/cheatsheets/Cross-Site_Request_Forgery_Prevention_Cheat_Sheet.html
- Use libraries and frameworks with built-in anti-CSRF mechanisms
 - E.g. Angular: “X-XSRF-TOKEN”
- User Interaction Based CSRF Defense before critical operation
 - Captchas: make sure a human intervenes (no automated spoofing)
 - One-time token
 - Re-authentication (Login/Password)



Defense in Depth: Origin Header

Origin: http://www.evil.com

- Alternative to Referer with fewer privacy problems
 - Sent only on POST, sends only necessary data
 - Defense against redirect-based attacks
 - Privacy
 - Identifies only principal that initiated the request (not path or query)
 - Sent only for POST requests; following hyperlink reveals nothing
 - Usability
 - Authorize subdomains and affiliate sites with simple firewall rule
- ```
SecRule REQUEST_HEADERS:Host !^www\.example\.com(:\d+)?$ deny,status:403
SecRule REQUEST_METHOD ^POST$ chain,deny,status:403
SecRule REQUEST_HEADERS:Origin !^(https?://www\.example\.com(:\d+)?)?$
```
- No need to manage secret token state
  - used with existing defenses to support legacy browsers (e.g. Referer)
- Standardization
  - Supported by W3C XHR2 and XMLHttpRequest



# Defense in Depth: SameSite Cookie

[draft-ietf-httpbis-rfc6265bis-latest](https://datatracker.ietf.org/doc/draft-ietf-httpbis-rfc6265bis-latest/) (Oct 8, 2019)

- Setting:

```
Set-Cookie: CookieName=CookieValue; SameSite=Lax;
```

```
Set-Cookie: CookieName=CookieValue; SameSite=Strict;
```

- Strict: the cookie will not be included in requests sent by third-parties (can affect browsing experience negatively)
- Lax: the cookie will be sent along with the GET request initiated by third party website, but only for top-level navigation requests (URL has to be changed in browser)
- Browsers are progressively integrating this feature

# One more thing...

- Cookie Scope:
  - based on Path attribute + Host/domain
  - Restricts usage of cookie to some application on the website
  - This is separate from SOP which is based on Host/domain+port
  - May further restrict cookie abuse

# Take-away message

- Cookie protection can be tricky, browser-specific, and is still investigated and standardized
- The prototype of a « secure » cookie ?  
**Set-Cookie: \_\_Host-SessionID=43a2;  
Path=/myapplication;Secure;HttpOnly;SameSite=Strict**
- ... that is, until the next release of RFC 6265bis...
- ... plus Tokens...
- ... and over HTTPS !
- Beware of XSS and MITM that may endanger cookie integrity (writing attack)
- .... And privacy !!!