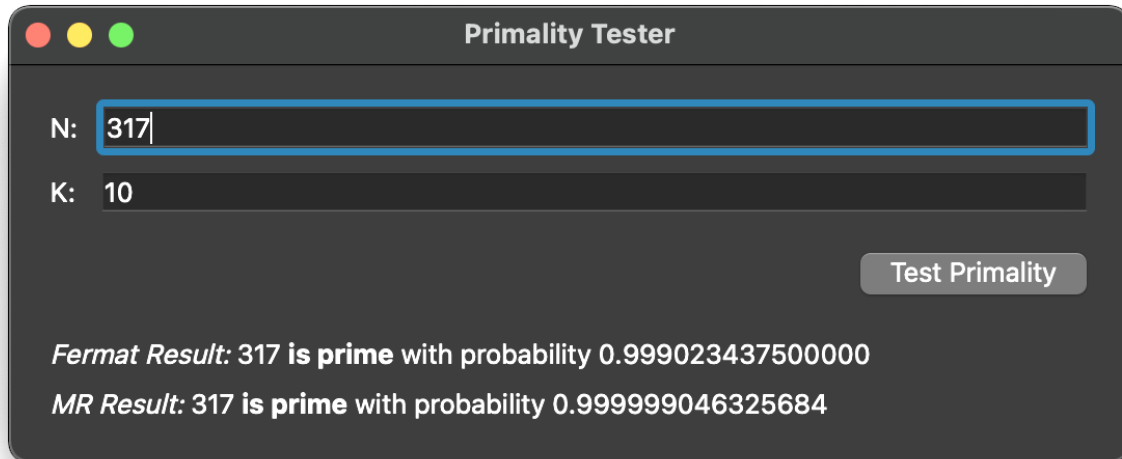


Project 1 - Fermat's Primality Tester

Example of Program Being Run



Primality Tester

N: 317

K: 10

Test Primality

Fermat Result: 317 is prime with probability 0.999023437500000

MR Result: 317 is prime with probability 0.999999046325684

Mod_Exp Function

```
def mod_exp(x, y, N):  
    # TOTAL TIME COMPLEXITY  $O(n^3)$   
    # TOTAL SPACE COMPLEXITY  $O(5n)$   
    if y == 0:  
        return 1  
    z = mod_exp(x, y // 2, N)  
    if y % 2 == 0:  
        return (z ** 2) % N  
    else:  
        return (x * (z ** 2)) % N
```

SPACE COMPLEXITY $O(3n)$ {Int} x y N

SPACE COMPLEXITY $O(1n)$ {Int} Return value

SPACE COMPLEXITY $O(1n)$ {Int} Z

SPACE COMPLEXITY $O(1n)$ {Int} Return value

SPACE COMPLEXITY $O(2n^2)$

SPACE COMPLEXITY $O(1n)$ {Int} Return value

Probability Function

```
def f_probability(k):  
    # TOTAL TIME COMPLEXITY  $O(2n^2 * 2^k)$   
    # TOTAL SPACE COMPLEXITY  $O(1)$   
    return 1 - (1 / (2 ** k))  
  
def m_probability(k):  
    # TOTAL TIME COMPLEXITY  $O(2n * 2^k)$   
    # TOTAL SPACE COMPLEXITY  $O(1)$   
    return 1 - (4 ** -k)
```

EXPONENTIAL, Div, Sub $O(2n^2 * 2^k)$ SPACE COMPLEXITY $O(1)$ {Int} Return Value

EXPONENTIAL, Sub $O(2n * 2^k)$ SPACE COMPLEXITY $O(1)$ {Int} Return Value

Fermat's

```
def fermat(N, k):  
    # TOTAL TIME COMPLEXITY  $O(n^4 + 1) \rightarrow O(n^4)$   
    # TOTAL SPACE COMPLEXITY  $O(6n + 4)$   
    M = N - 1  
    for i in range(0, k):  
        y = random.randint(1, M)  
        if mod_exp(y, M, N) != 1:  
            return 'composite'  
    return 'prime'
```

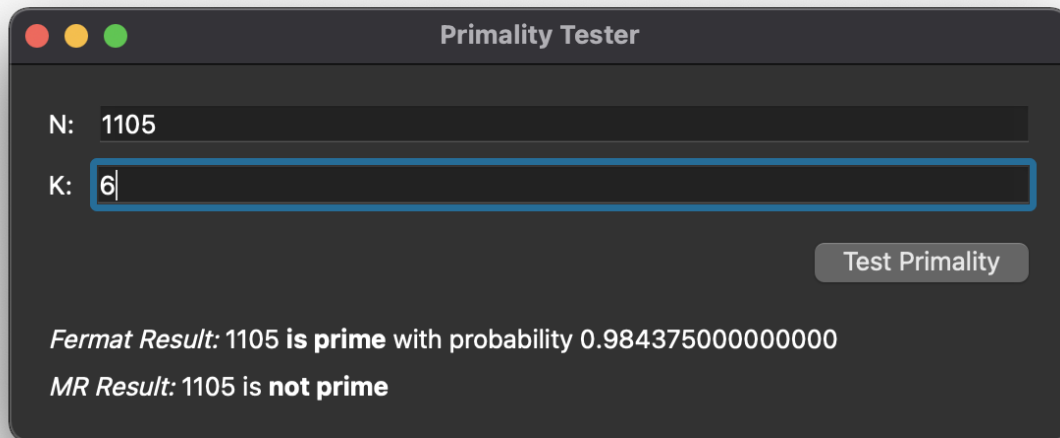
SPACE COMPLEXITY $O(2)$ {Int} k N
SPACE COMPLEXITY $O(1)$ {Int} M
SPACE COMPLEXITY $O(1n)$ {int} y
SPACE COMPLEXITY $O(5n)$
SPACE COMPLEXITY $O(1)$ {String} Return Value
SPACE COMPLEXITY $O(1)$ {String} Return Value

Miller-Rabin

```
def miller_rabin(N, k):  
    # TOTAL TIME COMPLEXITY  $O(2n^4 + n^2 + 2N + 2) \rightarrow O(n^4)$   
    # TOTAL SPACE COMPLEXITY  $O(6n + 4)$   
    squares = 0  
    M = N - 1  
    # Determine how many times the MR algorithm should loop by squaring N-1 through a bit shift until it is odd  
    # TIME COMPLEXITY  $O(n+2)$   
    # SPACE COMPLEXITY N/A  
    while M % 2 == 0:  
        M >>= 1  
        squares += 1  
    # MR Algorithm  
    # TIME COMPLEXITY  $O(2n^3 * n)$   
    # SPACE COMPLEXITY  $O(5n+2)$   
    def test_evaluation(X):  
        if mod_exp(X, M, N) == 1:  
            return False  
        for j in range(0, squares):  
            if mod_exp(X, 2 ** j * M, N) == N - 1:  
                return False  
        return True  
    # Accuracy Loop  
    # TIME COMPLEXITY  $O(2n^4 * n^2 + n)$   
    # SPACE COMPLEXITY  $O(6n + 4)$   
    for i in range(0, k):  
        a = random.randint(2, N)  
        if test_evaluation(a):  
            return 'composite'  
    return 'prime'
```

SPACE COMPLEXITY $O(2)$ {Int} k N
SPACE COMPLEXITY $O(1)$ {Int} squares
SPACE COMPLEXITY $O(1)$ {Int} M
SPACE COMPLEXITY $O(5n)$
SPACE COMPLEXITY $O(1)$ {Bool} Return
SPACE COMPLEXITY $O(1)$ {Int} j
SPACE COMPLEXITY $O(5n)$
SPACE COMPLEXITY $O(1)$ {Bool} Return
SPACE COMPLEXITY $O(1)$ {Bool} Return
SPACE COMPLEXITY $O(1)$ {Int} i
SPACE COMPLEXITY $O(1)$ {Int} a
SPACE COMPLEXITY $O(5n + 2)$
SPACE COMPLEXITY $O(1)$ {Str} Return
SPACE COMPLEXITY $O(1)$ {Str} Return

2.) d.) One of the things that I tried in order to find a number was exploit the weakness in the Fermat result by pitting it against a pseudoprime/Carmichael number, although with a high enough K it is unlikely that one of these numbers completely gets past the program it's still possible to introduce a discrepancy in the results.



3.) Although it was described in the code quite explicitly the overall time complexity for both the Fermat and the Miller-Rabin algorithms is $O(n^4)$ [After selecting the most impactful part of the polynomial -- $O(n^4+1)$ and $O(2n^4+n^2+2n+2)$ respectively] This is because both rely quite heavily on the Mod_Exp [$O(n^3)$] function existing within a loop to guarantee an accurate result. Logically the two probability functions are both computed in finite time $O(1)$, because we are calculating with a very fine granularity we get $O(2n^2 * 2^k)$ for the fermat probability and $O(2n * 2^k)$ for the MR function. Both could be calculated to the exact same $O(n)$ however by swapping the division in the fermat calculation for the negative exponent in the MR calculation we can bring the time down to a linear $O(n)$ time.

Space Complexity is largely inconsequential in a program of this size but it is mostly built on the complexity of the mod_exp function as it is recursive and leaves a lot of memory on the stack as it drills down further into the recursive element.

4.) The probability of the fermat problem is calculated with the following formula

$$1 - 1/2^k$$

The probability of the MR problem is calculated by the following formula

$$1 - 4^{-k}$$

Both provide very small numbers which can be subtracted from 1 to determine the chance that we have determined a correct number. Because of its exponential nature the MR is more reliable and faster to determine than the Fermat while $k < \sim 50$ but it becomes negligible after that.