



UNIVERSITY *of* WASHINGTON

CSS 343: Data Structures, Algorithms, and  
Discrete Mathematics II

# Hashing

Version 1

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# Hash Table

- Average complexity
  - Insert/delete/retrieval of binary search trees?
  - Hash table:  $O(1)$
- Worst-case complexity
  - Insert/delete/retrieval of binary search trees?
  - Hash table:  $O(n)$
- Basic idea:
  - use a **hash function** to map the items into positions in an array
  - Problem: Collision, and clustering

# Collisions



- **Collision**
  - When two keys map to the same location in the hash table
- **Flavors of Collision Resolution**
  - Closed hashing (Open Addressing)
    - Linear probing
    - Quadratic Probing
    - Double Hashing
  - Open hashing (Reconstructing the hash table)
    - Buckets
    - Separate chaining

# Closed hashing

- When a collision occurs, find another bucket for the new object
  - Probing distance function  $D(i)$  to resolve collisions
    - $i$ : number of collisions in the current attempt
  - Next bucket to exam:  $( h(x) + D(i) ) \bmod B$ 
    - $B$ : total # of buckets in a hash table
    - Generally suggested  $B$ : prime number and at least  $2 \times$  expected # of items
- $D(i)$ 
  - Linear probing:  $D(i) = i$
  - Quadratic probing:  $D(i) = i^2$
  - Double hashing:  $D(i) = i * h_2(x)$

# Linear probing

- After  $i$  collision,  $(h(x) + D(i)) \bmod B$ 
  - $(h(x) + D(i)) \bmod B$ , where  $D(i) = i$
  - That is, if  $A[i]$  is occupied
    - Try  $A[(i+1) \bmod N]$ , if still occupied
    - Try  $A[(i+2) \bmod N]$ ,  
And so on,
- Assume  $h(x) = x \% 10$ ,  $N = 10$  (hash table size)
- Example
  - 37, 98, 107, 20, 57
- Practice
- Suffer from Primary Clustering Problem
  - "Long runs of occupied slots built up (near the  $h(x)$  positions), increasing the average search time."

0	20
1	57
2	
3	
4	
5	
6	
7	37
8	98
9	107

# Quadratic probing

- After  $i$  collision,  $(h(x) + D(i)) \bmod B$
- $(h(x) + D(i)) \bmod B$ , where  $D(i) = i^2$ 
  - That is, if  $A[i]$  is occupied
    - Try  $A[i+1^2 \bmod N]$ , if still occupied
    - Try  $A[i+2^2 \bmod N]$ ,  
And so on,
- Assume  $h(x) = x \% 10$ ,  $N = 10$  (hash table size)
- Example: 37, 18, 107, 57
- Although fewer, suffer from Secondary Clustering
  - elements that hash to the same position will probe the same alternate cells (same sequence is followed to handle collision)
  - The size of the hash table should be a prime number to avoid repeats

0	
1	107
2	
3	
4	
5	
6	57
7	37
8	18
9	

# Double hashing

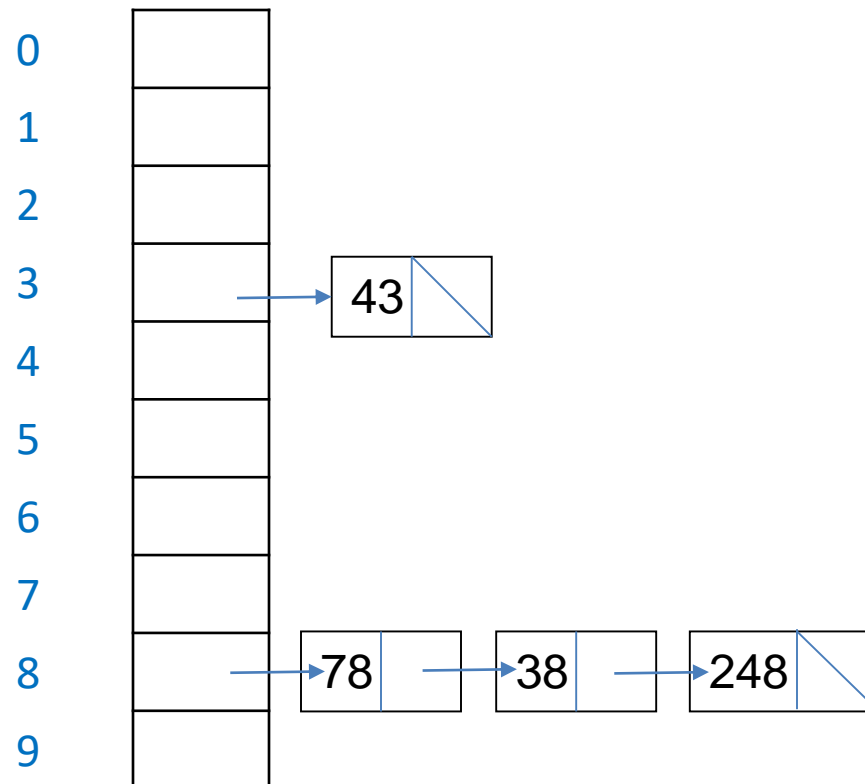
- $(h(x) + D(i)) \bmod B$ , where  $D(i) = i * h_2(x)$
- Typically,  $h_2(x) = r - (x \% r)$ ;  $r$  is a prime number smaller than  $B$
- Example
  - Assume  $h(x) = x \% 10$ ;  $h_2(x) = 7 - (x \% 7)$ , hash 89, 18, 49, 58, 50, 60, 23
- **If  $B$  is non-prime number**, possibly lead to **cycle of repeated visits**

Fail on 23

0	50
1	
2	60
3	58
4	
5	
6	49
7	
8	18
9	89

# Open hashing

- each bucket of the hash table stores a linked list of the objects that hash to that bucket
- A simple example
  - $h(x) = x \% 10$ , hash 78, 38, 43, 248
- Practice





# Define hash function

Define your own hash function:

**How to write a hash function to quickly find the following educational programs at UW Bothell in a hash table (BBUS, BEDUC, BES, BHLTH, BLS, BNURS, BPOLST, CSS)**

**Note: A good hash function is**

- 1) Easy to compute**
- 2) Will evenly distribute the possible data**

**Try:**

- 1) Use a prime number as a size of hash table**
- 2) Avoid collision**

# Define hash function

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WAY #1:

sum up ascii values of each character and use table size 17

A -> 65

B -> 66

...

Z -> 90

CSS:  $C(67) + S(83) + S(83) \Rightarrow 233 \% 17 = 12$

After finish it, you can shift the values so that the smallest number be zero

# Define hash function

Define your own hash function:

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WAY #2:

2-D array with its length as the row

- Total of 5 rows (index starts from 1)
- The second letter, ascii value % 5 (5 = column)

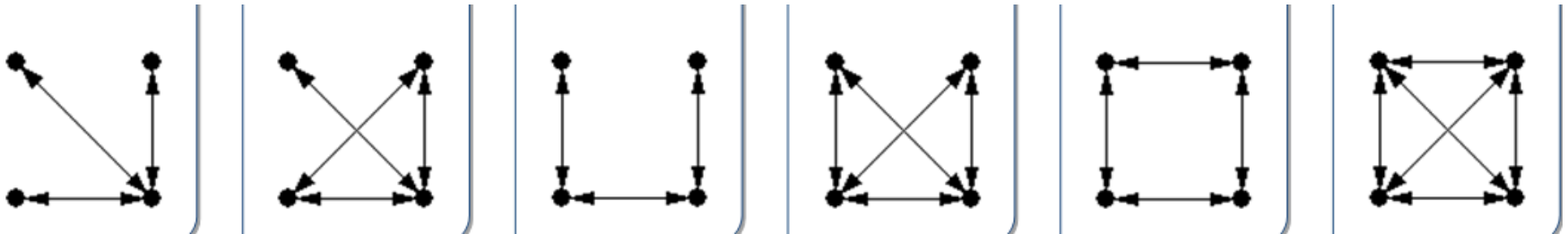
CSS

row=3, column =  $83\%5 = 3$ .

So, CSS is in array[3][3]

# Graph hash function

How to define a hash function for a graph?



Design your own hash function to store all the 6 types of graphs.

Reference: g6 format see [Files/additional materials/format.txt](#)  
on Canvas