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# CSE 373

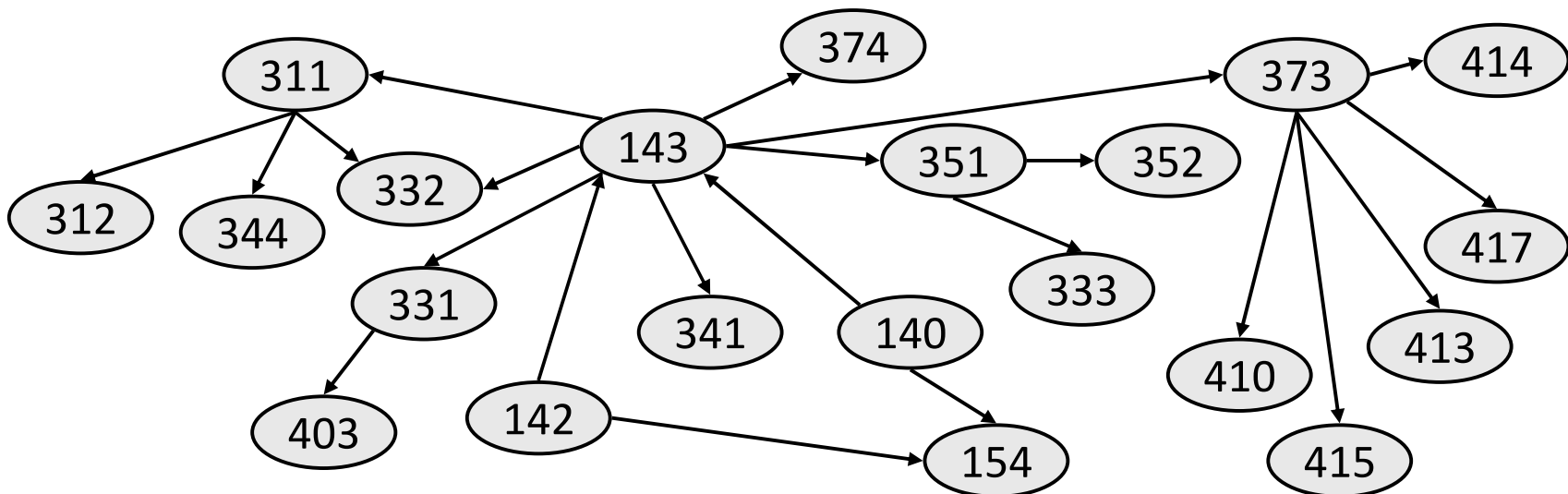
Graphs 4: Topological Sort  
reading: Weiss Ch. 9

slides created by Marty Stepp  
<http://www.cs.washington.edu/373/>

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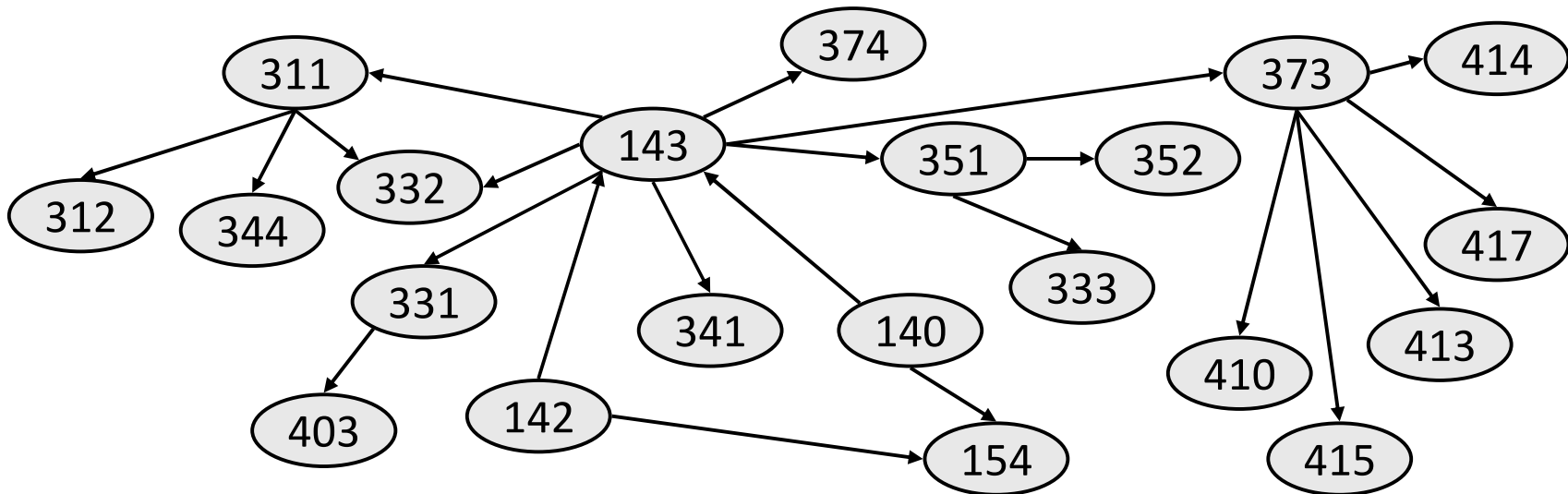
# Ordering a graph

- Suppose we have a directed acyclic graph (DAG) of courses, and we want to find an order in which the courses can be taken.
  - Must take all prereqs before you can take a given course. Example:
    - [142, 143, 140, 154, 341, 374, 331, 403, 311, 332, 344, 312, 351, 333, 352, 373, 414, 410, 417, 413, 415]
    - There might be more than one allowable ordering.
  - How can we find a valid ordering of the vertices?



# Topological Sort

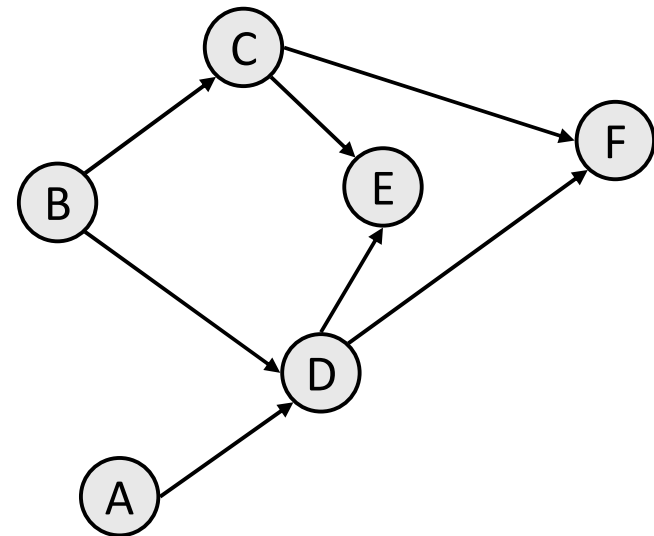
- **topological sort:** Given a digraph  $G = (V, E)$ , a total ordering of  $G$ 's vertices such that for every edge  $(v, w)$  in  $E$ , vertex  $v$  precedes  $w$  in the ordering. Examples:
  - determining the order to recalculate updated cells in a spreadsheet
  - finding an order to recompile files that have dependencies
    - (any problem of finding an order to perform tasks with dependencies)



# Topo sort example

- How many valid topological sort orderings can you find for the vertices in the graph below?

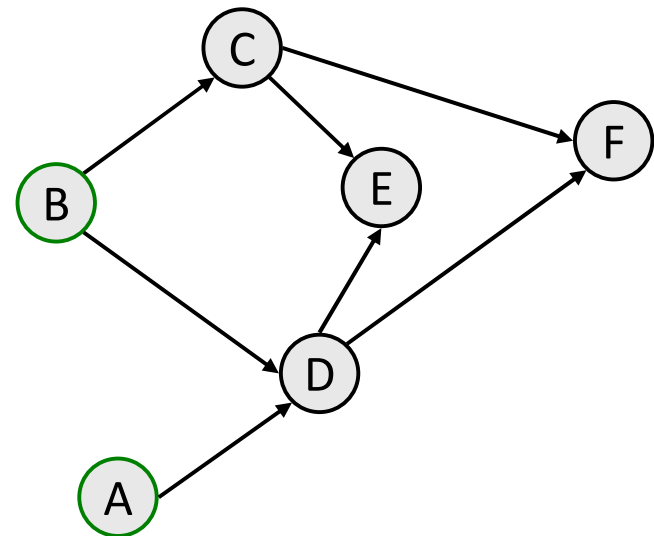
- [A, B, C, D, E, F], [A, B, C, D, F, E],
- [A, B, D, C, E, F], [A, B, D, C, F, E],
- [B, A, C, D, E, F], [B, A, C, D, F, E],
- [B, A, D, C, E, F], [B, A, D, C, F, E],
- [B, C, A, D, E, F], [B, C, A, D, F, E],
- ...



- What if there were a new vertex G unconnected to the others?

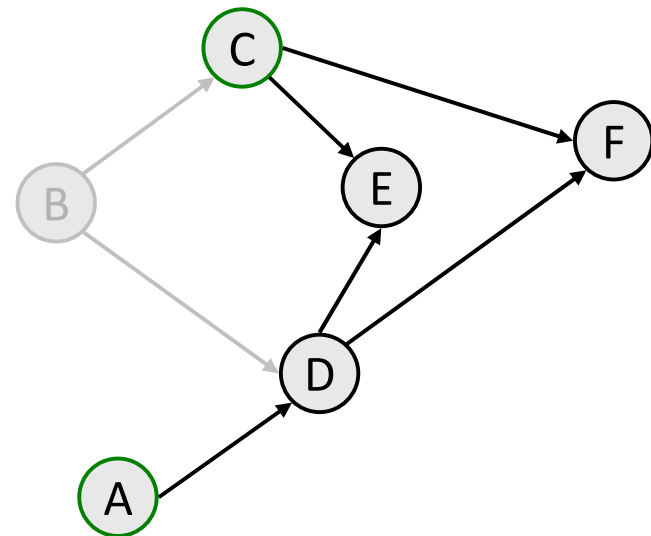
# Topo sort: Algorithm 1

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .



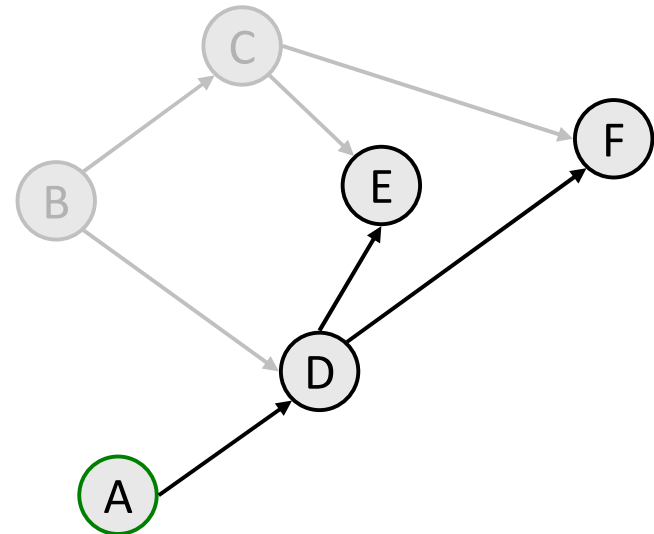
# Topo sort example

- function `topologicalSort()`:
  - `ordering := { }`.
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - `ordering += v`.
  - `ordering = { B }`



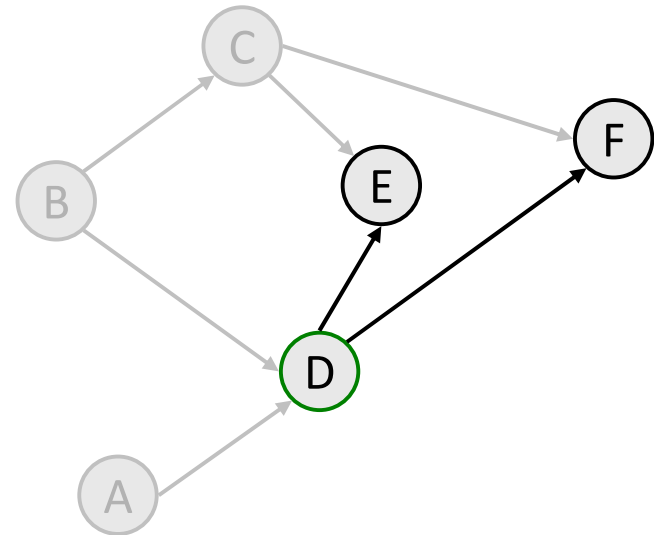
# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C \}$



# Topo sort example

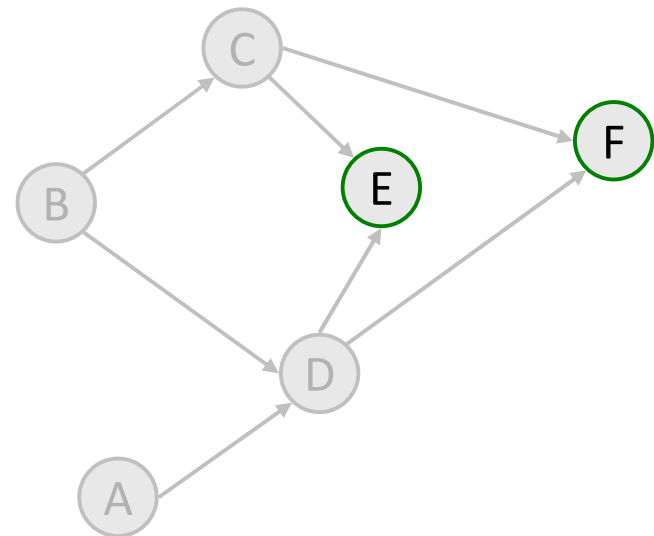
- function `topologicalSort()`:
  - `ordering := { }.`
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - `ordering += v .`
  - `ordering = { B, C, A }`





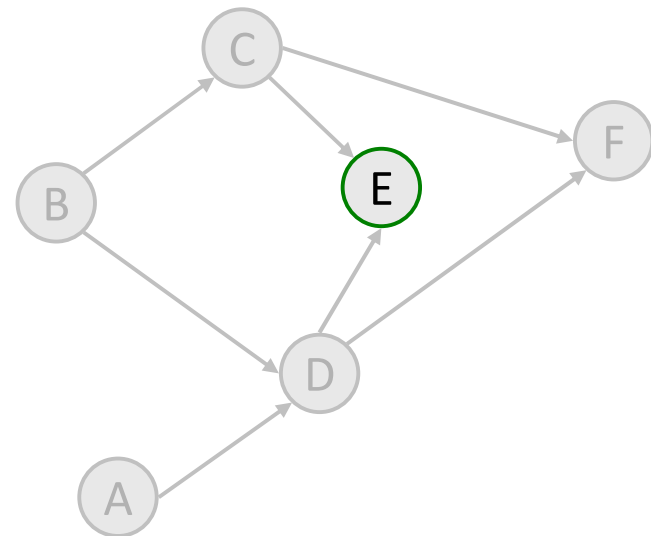
# Topo sort example

- function topologicalSort():
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C, A, D \}$



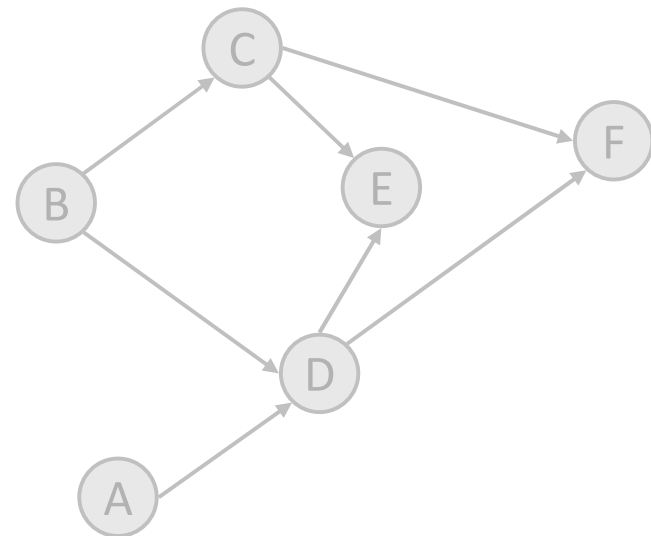
# Topo sort example

- function `topologicalSort()`:
  - `ordering := { }`.
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - `ordering += v`.
  - `ordering = { B, C, A, D, F }`



# Topo sort example

- function `topologicalSort()`:
  - `ordering := { }`.
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - `ordering += v`.
  - `ordering = { B, C, A, D, F, E }`



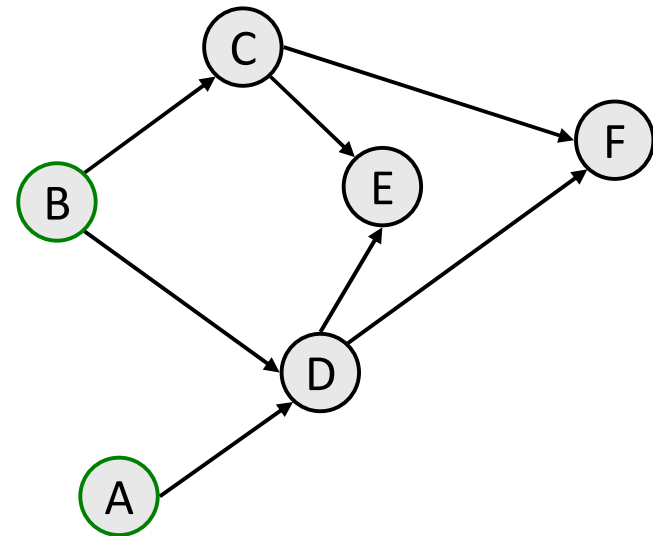
# Revised algorithm

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- We don't want to literally delete vertices and edges from the graph while trying to topological sort it; so let's revise the algorithm:
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
  - $queue := \{\text{all vertices with in-degree} = 0\}$ .
  - $ordering := \{ \}$ .
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue.
    - $ordering += v$ .
    - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .
  - If all vertices are processed, success.  
Otherwise, there is a cycle.

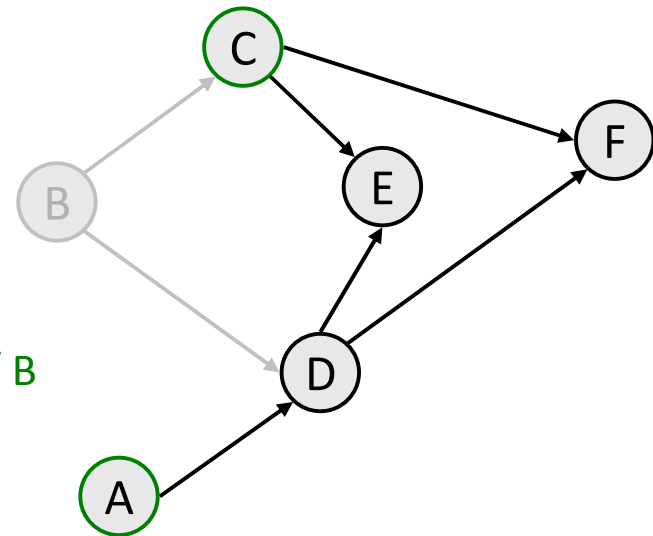
# Topo sort example 2

- function topologicalSort():
  - *map* := {each vertex  $\rightarrow$  its in-degree}.
  - *queue* := {all vertices with in-degree = 0}.
  - *ordering* := { }.
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue.
    - *ordering* +=  $v$ .
    - Decrease the in-degree of all  $v$ 's neighbors by 1 in the *map*.
    - *queue* += {any neighbors whose in-degree is now 0}.
- *map* := { A=0, B=0, C=1, D=2, E=2, F=2 }
- *queue* := { B, A }
- *ordering* := { }



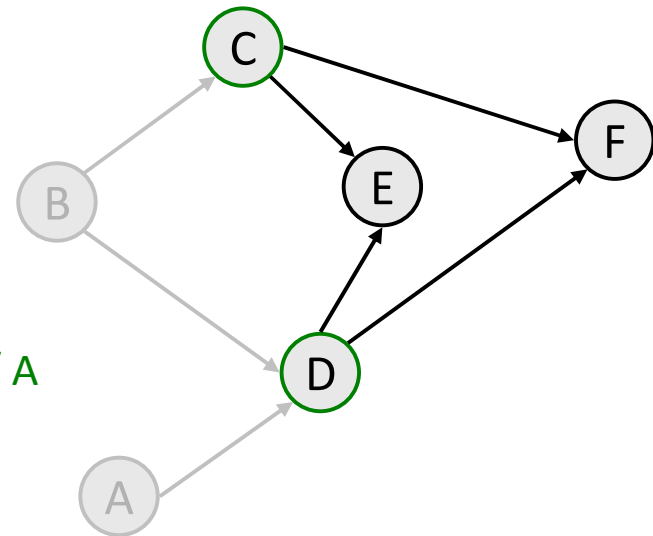
# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // B
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // C, D neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, \mathbf{C=0}, \mathbf{D=1}, E=2, F=2 \}$
- $queue \quad \quad \quad := \{ A, \mathbf{C} \}$
- $ordering \quad \quad := \{ \mathbf{B} \}$



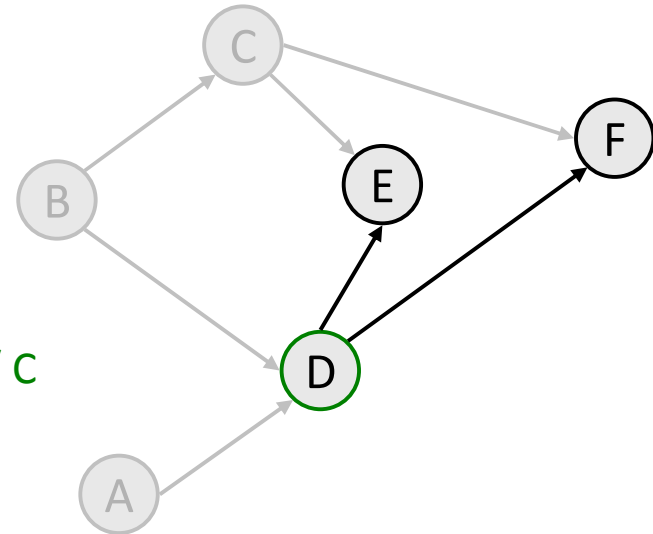
# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // A
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // D  
neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, C=0, \mathbf{D=0}, E=2, F=2 \}$
- $queue \quad \quad \quad := \{ C, \mathbf{D} \}$
- $ordering \quad \quad := \{ B, \mathbf{A} \}$



# Topo sort example 2

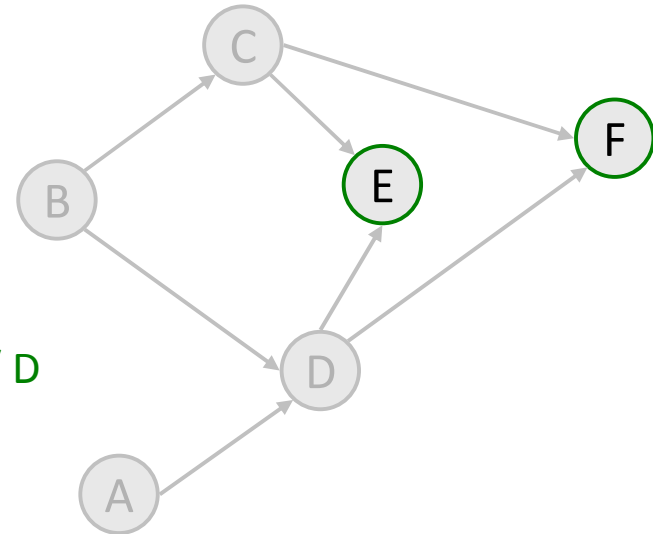
- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. *// C*
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's *// E, F* neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, C=0, D=0, E=1, F=1 \}$
- $queue \quad \quad \quad := \{ D \}$
- $ordering \quad \quad := \{ B, A, \mathbf{C} \}$





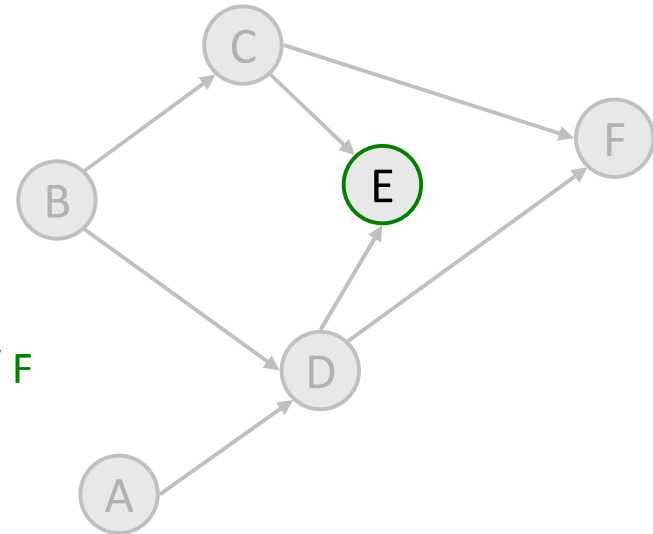
# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // D
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // F, E  
neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue \quad \quad \quad := \{ \mathbf{F}, \mathbf{E} \}$
- $ordering \quad \quad := \{ B, A, C, \mathbf{D} \}$



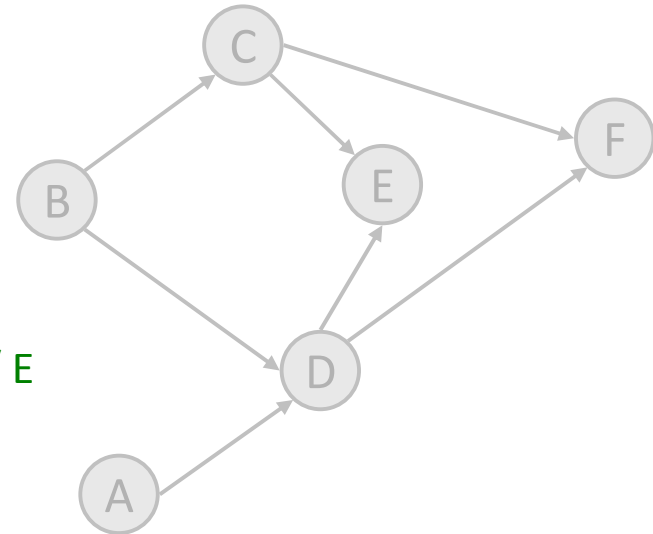
# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // F
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's // none  
neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue \quad \quad \quad := \{ E \}$
- $ordering \quad \quad := \{ B, A, C, D, F \}$



# Topo sort example 2

- function topologicalSort():
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$
  - $queue := \{\text{all vertices with in-degree} = 0\}.$
  - $ordering := \{ \}.$
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. *// E*
    - $ordering += v.$
    - Decrease the in-degree of all  $v$ 's *// none* neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
- $map \quad \quad \quad := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue \quad \quad \quad := \{ \}$
- $ordering \quad \quad := \{ B, A, C, D, F, E \}$



# Topo sort runtime

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- What is the runtime of our topological sort algorithm?
  - (with an "adjacency map" graph internal representation)
  - function topologicalSort():
    - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}.$  //  $O(V)$
    - $queue := \{\text{all vertices with in-degree} = 0\}.$
    - $ordering := \{ \}.$
    - Repeat until queue is empty: //  $O(V)$ 
      - Dequeue the first vertex  $v$  from the queue. //  $O(1)$
      - $ordering += v.$  //  $O(1)$
      - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ . //  $O(E)$  for all passes
      - $queue += \{\text{any neighbors whose in-degree is now } 0\}.$
  - Overall:  **$O(V + E)$**  ; essentially  $O(V)$  time on a sparse graph (fast!)