Semantic Analysis
Using
Semantic Actions,
Semantic Action Stack
&
Operator Stack

#### Semantic Analysis

Semantic Actions are added to the grammar to indicate where a Semantic Routine is to be executed.

Each grammar symbol terminal or non-terminal my have zero or more associated Semantic Actions.

Semantic Routines are placed in a parser to help perform Semantic Analysis.

Semantic Routines communicate via Semantic Action Records.

Semantic Action Records (SAR) contain information needed by Semantic Routines to perform Semantic Analysis.

The Semantic Action Stack (SAS) is used to store and pass Semantic Records between Semantic Routines.

#### Operator Precedence

Infix expressions must be converted to Postfix expressions to correctly evaluate.

#### Postfix expression

No Parenthesis

Enforce Operator Precedence

Shuntyard Algorithm used to convert Infix to Postfix

$$J = I + K * E;$$

$$JIKE*+=$$

Hint: Semantic Analysis (2 Pass) is only done after the first pass has already loaded the Symbol table for the entire program.

Hint: Focus on how nested or complex operations are always converted to simpler generally binary operations.

$$x+y+z+g \equiv t1=x+y$$

$$t2=t1+z$$

$$t3=t2+g$$

$$x.y.z.g \equiv t1=x.y$$

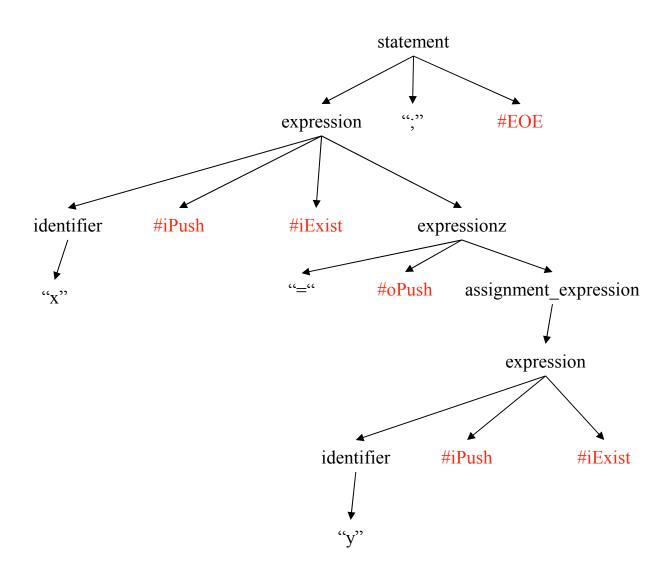
$$t2=t1.z$$

t3=t2.q

Hint: Be sure you understand how #rExists and #iExists work before adding lots of complexity to your Compiler. A common mistake is to write naïve versions of these critical semantic functions and then try to compensate by adding complexity all over the Compiler.

X = y

### Simple Assignment Statement



$$x = y$$

 $\mathbf{X}$ 

- #iPush

Operator Stack

$$x = y$$

Note: The Semantic Routine for #iExist creates an id\_sar for x, however most example will simply show this as:

x - #iExist

x id\_sar - #iExist

Semantic Action Stack

Operator Stack

Operator Stack

x id\_sar - #iExist

$$X = V$$

Operator Stack

y - #iPush

x id sar - #iExist

$$x = y$$
;

Operator Stack

y id sar - #iExist

x id sar - #iExist

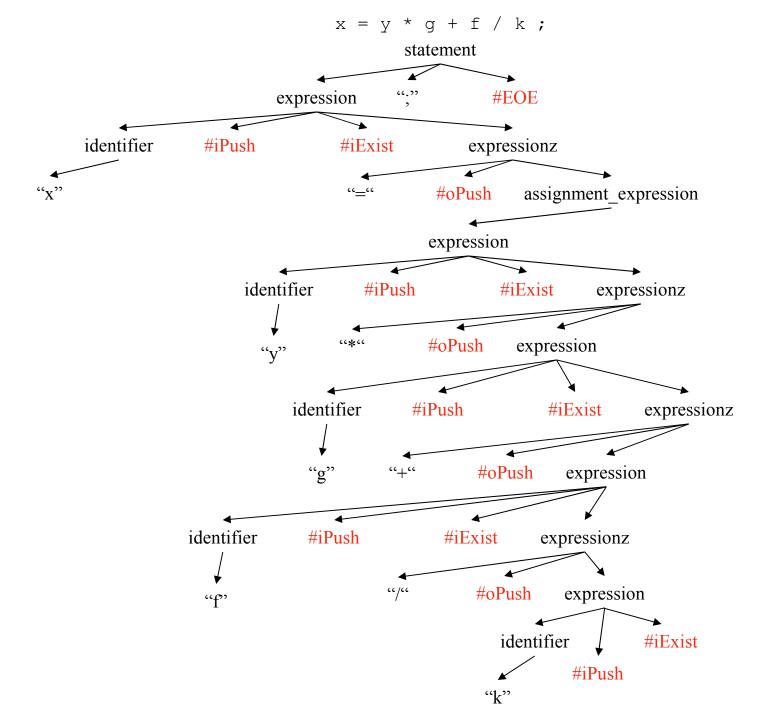
$$x = y$$
parse

#### **-** #EOE

Operator Stack

x = y \* g + f / k ;

## Assignment Statement with infix to Postfix Conversion



$$\frac{x}{y} = y * g + f / k ;$$

X

- #iPush

Operator Stack

$$x = y * g + f / k ;$$

 $\mathbf{X}$ 

- #iExist

Operator Stack

$$x = y * g + f / k$$
;

Operator Stack

- #iExist

X

$$x = \underbrace{v} * g + f / k ;$$
parse

Operator Stack

y - #iPush

x - #iExist

$$x = y * g + f / k ;$$

Operator Stack

y - #iExist

x - #iExist

$$x = y * g + f / k$$
;

- \* #oPush
- = -#oPush

Operator Stack

y - #iExist

x - #iExist

$$x = y * q + f / k$$
;

Operator Stack

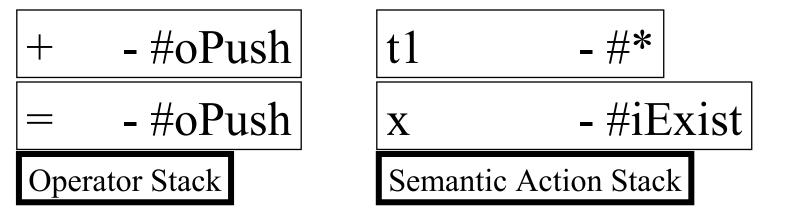
g - #iPush

y - #iExist

x - #iExist

$$x = y * g + f / k ;$$

```
    g - #iExist
    * - #oPush
    g - #iExist
    y - #iExist
    x - #iExist
    Operator Stack
    Semantic Action Stack
```



$$x = y * g + f / k$$
;

+ - #oPush

= -#oPush

Operator Stack

f - #iPush

t1 - #\*

x - #iPush

$$x = y * g + f / k ;$$

Operator Stack

f -#iExist

t1 - #\*

x - #iExist

$$x = y * g + f k$$
;

/	- #oPush	$\int \int f$	•	- #iE	xist
+	- #oPush		1	- #*	
= -#oPush		X	-	- #iExist	
Operator Stack			Semantic Action Stack		

$$x = y * g + f / k$$
;

	k	- #iPush	
/ - #oPush	f	- #iExist	
+ - #oPush	t1	<b>-</b> #*	
= -#oPush	X	- #iExist	
Operator Stack	Semantic	Semantic Action Stack	

$$x = y * g + f / k$$
;

```
k -#iExist

/ -#oPush

+ -#oPush

t1 -#*

= -#oPush

Operator Stack

k -#iExist

f -#iExist

-#iExist

x -#iExist
```

$$x = y * g + f / k$$

$$- \#EOE$$

$$pop / \#/$$

$$k = SAS.pop()$$

$$f = SAS.pop()$$

$$Test f / k \text{ is valid}$$

$$t2 = f / k$$

$$SAS.push(t2)$$

- #oPush

- #oPush

Operator Stack

t2 - #/

**-** #\*

- #iExist X

$$x = y * g + f / k$$

$$- #EOE$$

$$pop + #+$$

$$t2 = SAS pop()$$

t3 **-** #+ - #oPush - #iExist X Semantic Action Stack

Operator Stack

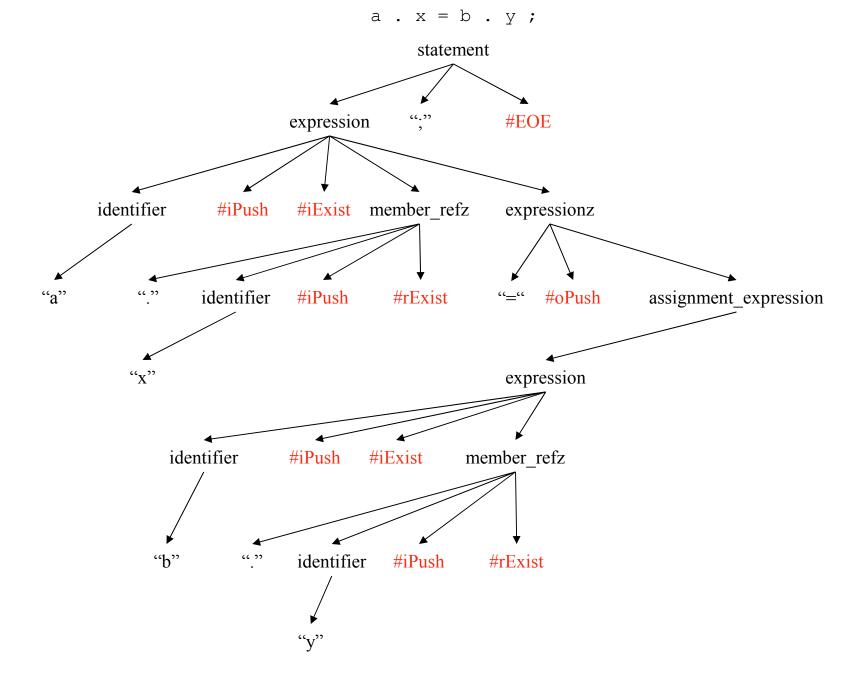
$$x = y * g + f / k$$

$$\# F \cap F$$
parse

#### **-** #EOE

<u>a</u>. x = b. y;

# Simple Assignment Statement with references to a Class Member



$$\begin{array}{c} a \\ \\ \text{parse} \end{array} \quad x = b \quad y \quad ;$$

a

- #iPush

Operator Stack

$$a \cdot x = b \cdot y ;$$

a - #iExist

Semantic Action Stack

Operator Stack

$$a = b \cdot y;$$

- #iExist

a

$$a \cdot x = b \cdot y ;$$

x - #iPush

a - #iExist

Semantic Action Stack

Operator Stack

Note: #rExist creates a ref sar for a.x by checking that the class (type) of a has an element named x. Also note there is only one SAR on the Semantic Action Stack after the #rExist routine has executed.

Ivalue (location) vs. rvalue (contents)

a.x ref sar - #rExist

Operator Stack

$$a \cdot x = b \cdot y$$

Operator Stack

a.x ref\_sar

- #rExist

a. 
$$x = b$$
.  $y$ ;

Operator Stack

b - #iPush

a.x ref\_sar - #rExist

a . 
$$x = b$$
 .  $y$ ;

Operator Stack

b - #iExit

a.x ref\_sar - #rExist

a 
$$x = b$$
 y;

Operator Stack

b - #iExit

a.x ref\_sar - #rExist

a. 
$$x = b$$
.  $y$ 

y - #iPush

b - #iExit

a.x ref\_sar - #rExist

Semantic Action Stack

= -#oPush

Operator Stack

a . 
$$x = b$$
 .  $y$ 

Note: #rExist creates a ref\_sar for b.y by checking that the class (type) of b has an element named y. Also note there are only two SAR on the Semantic Action Stack after the #rExist routine has executed.

= -#oPush

Operator Stack

b.y ref sar - #rExist

a.x ref sar - #rExist

a . 
$$x = b$$
 .  $y$  :

Note: The ref\_sar for a.x and b.y, ultimately will relate to a collection of code that will compute a location in memory (e.g., Base Address + Offset). The computed address should be stored in the Symbol Table like any other identifier. Thus there will need to be a symbol id (e.g., t9 and t10) associated with each of the ref\_sar's which is not shown in the example. The symbol id, can be created and added to the Symbol Table when the ref\_sar is created or as I prefer to think of it when the ref\_sar is popped from the Semantic Action Stack.

= - #oPush
Operator Stack

b.y ref\_sar - #rExist

a.x ref\_sar - #rExist

Semantic Action Stack

$$a \cdot x = b \cdot y$$
parse

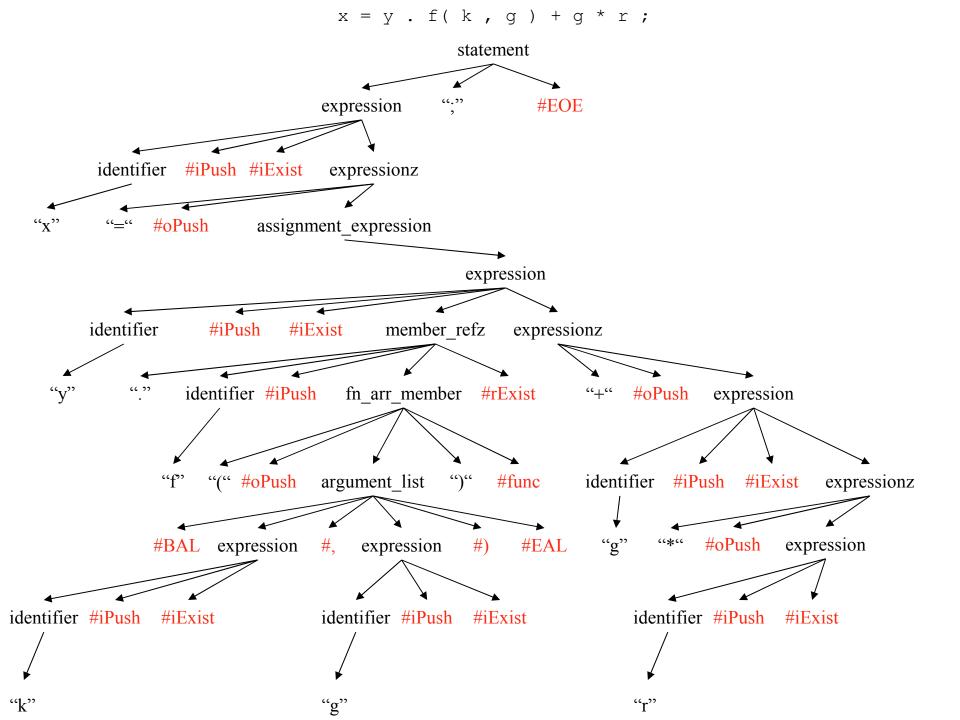
Test a.x = b.y is valid

Note: Just as when x = y, the semantic action #= only has to deals with two SAR's (each ref\_sar should have an entry in the Symbol Table), one for the LHS and one for the RHS. The assignment statement is not more complex, it simply must deals with the references before processing the assignment statement.

Operator Stack

 $x = y \cdot f(k, g) + g * r;$ 

## Assignment Statement of member function with infix to postfix conversion



$$x = y$$
. f ( k , g ) + g \* r;

X

- #iPush

Operator Stack

$$x = y \cdot f (k, g) + g * r;$$

x - #iExist

$$x = y$$
. f (k, g) + g \* r;

= - #oPush
Operator Stack

x - #iExist

$$x = y$$
. f ( k , g ) + g \* r ;

= -#oPush x
Operator Stack Sema

y - #iPush x - #iExist

$$x = y$$
 f ( k , g ) + g \* r ;

- #iExist - #iExist - #oPush X Semantic Action Stack

Operator Stack

$$x = y f (k, g) + g * r;$$

y -#iExist

= -#oPush x -#iExist

Operator Stack Semantic Action Stack

$$x = y \cdot f (k, g) + g * r;$$

f - #iPush
y - #iExist
x - #iExist
Semantic Action Stack

= - #oPush
Operator Stack

$$x = y$$
 . f (k, g) + g \* r;

f - #iPush
y - #iExist
x - #iExist
Semantic Action Stack

= -#oPush
Operator Stack

$$x = y$$
 .  $f(x) + g * r;$ 

$$x = y$$
 . f (k, g) + g \* r;

- #iPush bal sar - #BAL - #iPush - #iExist - #oPush - #oPush - #iExist X Semantic Action Stack Operator Stack

k

$$x = y \cdot f (k \cdot g) + g * r;$$

```
- #iExist
                    k
                    bal sar
                             - #BAL
                               - #iPush
     - #oPush
                               - #iExist
     - #oPush
                               - #iExist
                    X
Operator Stack
                    Semantic Action Stack
```

$$x = y$$
 . f  $(k, q) + g * r;$ 

- #iPush

- #iExist bal sar - #BAL - #iPush - #oPush - #iExist - #oPush - #iExist X Operator Stack Semantic Action Stack

$$x = y \cdot f \cdot (k \cdot g) + g \cdot r ;$$

- #iExist - #iExist bal sar - #BAL - #iPush - #iExist - #iExist X Semantic Action Stack

- #oPush- #oPush

Operator Stack

$$x = y$$
 . f ( k , g) + g \* r ;

pop (
No infix to postfix
conversion necessary
on expression g

al\_sar - #EAL

f - #iPush

y - #iExist

x - #iExist

Semantic Action Stack

- #oPushOperator Stack

```
x = y . f ( k , g ) + g * r ;

-#EAL
g = SAS.pop()
k = SAS.pop()
bal_sar == SAS.pop()
Add g & k to al sar
```

SAS.push(al sar)

al\_sar - #EAL

f - #iPush

y - #iExist

x - #iExist

Semantic Action Stack

- #oPushOperator Stack

```
x = y . f ( k , g) + g * r ;
```

## - #func

al\_sar = SAS.pop() f = SAS.pop() Add f & al\_sar to func\_sar

SAS.push(func sar)

Note: #func <u>Can't determine if the function exist</u> it simply combines the argument list with the function name to form the function sar.

func sar - #func

parse

y - #iExist

x - #iExist

Semantic Action Stack

= -#oPush

Operator Stack

```
x = y \cdot f (k, g) + g * r;
```

## - #rExist

top = SAS.pop()

next = SAS.pop()

Test that top is a valid member of next (next.top)

Create ref\_sar for top and next

SAS.push(ref\_sar)

Note: #rExist <u>will determine if the function exist</u> because it is part of a reference [e.g, y.f(k, g)]. <u>#iExist would be used</u> if the function is not part of reference (see later slides).

= -#oPush

Operator Stack

ref sar - #rExist

- #iExist

$$x = y \cdot f (k, g) + g * r;$$

Note: When evaluating function calls from semantic analysis a ref sar created via a #rExists or id sar created via an #iExist will be on the Semantic Action Stack after the function has been determined to exist. While the SAR on the SA Stack is sufficient at this time, it may help to understand that a temporary variable must be associated with the SAR if the function returns a value. The temporary variable will not be used until Intermediate and Target Code Generation. Hint: Just let the ref sar and id sar have a temporary variable attribute, if you want.

= -#oPush x
Operator Stack Semantic A

ref\_sar - #rExist

x - #iExist

Semantic Action Stack

$$x = y$$
. f (k, g) + g \* r;

- + #oPush
- = -#oPush

Operator Stack

ref\_sar - #rExist

x - #iExist

$$x = y$$
 . f ( k , g ) +  $q$  \* r ;

+ - #oPush

= -#oPush

Operator Stack

g - #iPush

ref sar - #rExist

x - #iExist

$$x = y \cdot f (k, g) + g * r;$$

+ - #oPush

= -#oPush

Operator Stack

g - #iExist

ref\_sar - #rExist

x - #iExist

$$x = y$$
 . f ( k , g ) + g \* r ;

\* - #oPush

+ - #oPush

ref\_sar - #rExist

- #oPush

x - #iExist

Operator Stack

Semantic Action Stack

$$x = y$$
 . f ( k , g ) + g \* r;

r - #iPush

\* - #oPush

+ - #oPush

= - #oPush

Operator Stack

ref\_sar - #rExist

x - #iExist

Semantic Action Stack

$$x = y \cdot f (k, g) + g * r :$$

```
r - #iExist

* - #oPush

+ - #oPush

= - #oPush

Operator Stack

ref_sar - #rExist

x - #iExist

Semantic Action Stack

Semantic Action Stack
```

$$x = y$$
 . f ( k , g ) + g \* r ;

# - #EOE

```
pop *
       #*
       r = SAS.pop()
       g = SAS.pop()
       Test g * r is valid
       t4 = g * r
       SAS.push(t4)
```

t4 - #EOE

- #oPush

- #oPush

Operator Stack

- #iExist X

- #rExist

Semantic Action Stack

ref sar

$$x = y$$
 . f ( k , g ) + g \* r ;

# - #EOE

```
pop + #+
     t4 = SAS.pop()
     ref_sar = SAS.pop()
     Test ref_sar + t4 is valid
     t6 = ref_sar + t4
     SAS.push(t6)
```

Note: An entry should be created in the Symbol Table for the y.f(k,g) ref sar.

t6 - #EOE

x - #iExist

Semantic Action Stack

= -#oPush

$$x = y$$
 . f ( k , g ) + g \* r ;

x . f () . g () . y = r;

Assignment Statement to a member element returned by a function

$$x$$
 . f ( ) . g ( ) .  $y = r$ ;

X

- #iPush

Operator Stack

$$x = f() \cdot g() \cdot y = r;$$

x - #iExist

$$x = f() \cdot g() \cdot y = r;$$

 $\mathbf{X}$ 

- #iExist

Operator Stack

$$x \cdot f(x) \cdot g(x) \cdot y = r;$$

f - #iPush

x - #iExist

Semantic Action Stack

$$x \cdot f() \cdot g() \cdot y = r;$$

f - #iPush x - #iExist

Semantic Action Stack

$$x \cdot f(x) \cdot g(x) \cdot y = r;$$

bal\_sar - #BAL f - #iPush x - #iExist

Operator Stack Semantic Action Stack

- #oPush

al\_sar - #EAL

f - #iPush

x - #iExist

Semantic Action Stack

$$x \cdot f() \cdot g() \cdot y = r;$$

$$x \cdot f \cdot y \cdot g \cdot y = r ;$$

# - #func

al\_sar = SAS.pop()
f = SAS.pop()
Add f & al\_sar to func\_sar
SAS.push(func\_sar)

func sar - #func

x - #iExist

Semantic Action Stack

x . f () g () . y = r .

#### - #rExist

func\_sar = SAS.pop()

x = SAS.pop()

Test that func\_sar (function f) is a valid member of the class of x

Create ref\_sar for the result

SAS.push(ref\_sar)

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f () = g () \cdot y = r ;$$

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f () \cdot g () \cdot y = r ;$$

g - #iPush

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f () \cdot g () \cdot y = r ;$$

g - #iPushref\_sar - #rExist

Semantic Action Stack

$$x \cdot f () \cdot g () \cdot y = r ;$$

bal\_sar - #BAL

g - #iPush

ref\_sar - #rExist

Semantic Action Stack

al\_sar - #EAL

g - #iPush

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f() \cdot g() \cdot y = r$$

#### - #EAL

bal\_sar == SAS.pop()
Create empty al\_sar

SAS.push(al\_sar)

al\_sar - #EAL

g - #iPush

ref\_sar - #rExist

Semantic Action Stack

 $x \cdot f () \cdot g () \cdot y = r ;$ 

parse

# - #func

al\_sar = SAS.pop()
g = SAS.pop()
Add g & al\_sar to func\_sar
SAS.push(func\_sar)

func sar - #func

ref sar - #rExist

Semantic Action Stack

 $x \cdot f () \cdot g () \cdot y = r ;$ 

# - #rExist

 $func\_sar = SAS.pop()$ 

ref\_sar = SAS.pop(), the x.f() ref\_sar

Test that the func\_sar (function g) is a valid member of the ref\_sar (the return type of function f)

Create ref\_sar for the result

SAS.push(ref\_sar)

Note: A a Symbol Table entry is needed for the x.f() ref\_sar result (e.g., t1).

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f () \cdot g () = r ;$$

ref\_sar - #rExist

$$x \cdot f() \cdot g() \cdot v = r;$$

y - #iPush

ref\_sar - #rExist

Semantic Action Stack

 $x \cdot f() \cdot g() \cdot y = r;$ 

# - #rExist

y = SAS.pop() ref sar = SAS.pop(), t1.g() ref sar

Test that variable y is a valid member of the ref\_sar (the return type of function g)

Create ref\_sar for top and next

SAS.push(ref\_sar)

Note: The result of t1.g() will be placed in the Symbol Table at t2.

ref\_sar - #rExist

Semantic Action Stack

$$x \cdot f() \cdot g() \cdot y = r;$$

= -#oPush

Operator Stack

ref\_sar - #rExist

$$x \cdot f () \cdot g () \cdot y = r;$$

- #oPush

Operator Stack

r - #iPush

ref\_sar - #rExist

$$x \cdot f () \cdot g () \cdot y = r \cdot ;$$

= -#oPush

Operator Stack

r - #iExist

ref\_sar - #rExist

x . f () . g () . y = r

**-** #EOE

pop = #=
 r = SAS.pop()
 ref\_sar = SAS.pop()

Test ref\_sar (variable y which is a public member of the return type of function g ) = r is valid

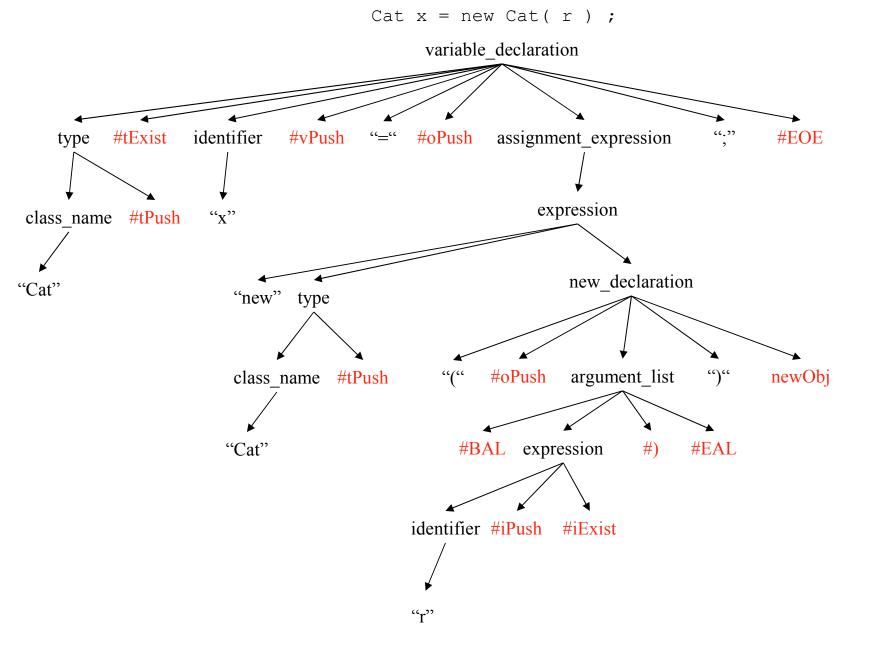
parse

Note: The ref\_sar result is placed in the Symbol Table at t3.

Operator Stack

Cat x = new Cat(r);

Creating a instance of a Class



Cat

- #tPush

Operator Stack

$$Cat x = new Cat (r);$$

- #tExists

Cat = SAS.pop()

Test that Cat is a valid Class

Operator Stack

Cat 
$$x = \text{new Cat}(r)$$
;

Operator Stack

Cat 
$$x = new Cat(r)$$
;

X

- #vPush

Operator Stack

Operator Stack

X

- #vPush

Cat 
$$x = \underline{new}$$
 Cat  $(r)$ ;

Operator Stack

X

- #vPush

Cat 
$$x = new Cat(r)$$
;

Operator Stack

Cat - #tPush

x - #vPush

( - #oPush

= -#oPush

Operator Stack

bal sar - #BAL

Cat - #tPush

x - #vPush

Cat 
$$x = new Cat(r)$$
;

r - #iPush

bal\_sar - #BAL

Cat - #tPush

x - #vPush

Semantic Action Stack

( - #oPush

= -#oPush

Operator Stack

$$Cat x = new Cat (r)$$

```
| r - #iExist |
| bal_sar - #BAL |
| ( - #oPush | Cat - #tPush |
| = - #oPush | x - #vPush |
| Operator Stack | Semantic Action Stack |
```

```
Cat x = new Cat (
         - #)
   pop (
        No infix to postfix
        conversion necessary
        on expression r
                                 - #iExist
                     bal sar
                                 - #BAL
                     Cat
                                 - #tPush
     - #oPush
                                 - #vPush
                     X
Operator Stack
                     Semantic Action Stack
```

```
Cat x = new Cat (r)
```

- #EAL

r = SAS.pop()
bal\_sar == SAS.pop()
Create al\_sar for r
SAS.push(al\_sar)

al sar - #EAL

Cat - #tPush

x - #vPush

Semantic Action Stack

= -#oPush

Operator Stack

```
Cat x = new Cat (
                        parse
```

## - #newObj

al sar = SAS.pop()

Cat type sar == SAS.pop()

Test that class Cat has a constructor that takes the arguments of al sar

Create a new sar for the constructor of Cat with arguments al sar

SAS.push(new sar)

new sar - #newObj

- #oPush

Operator Stack

- #vPush X

Cat x = new Cat (r)



## **-** #EOE

Test that  $x = new_sar is valid$ 

Note: As new sar

Operator Stack

Cat x[] = new Cat[r];

Creating an instance of an Array

Cat

- #tPush

Operator Stack

Cat x = new Cat r;
-#tExists

Cat = SAS.pop()

Test that Cat is a valid Class

Cat - #tPush

Operator Stack

Operator Stack

Operator Stack

Cat x[] = new Cat[r];

Note: While we show the vPush as x[], your stored information might be more like x, @:Cat. Indicating variable x is an array of type Cat.

x[] - #vPush

Semantic Action Stack

Operator Stack

= - #oPush
Operator Stack

x[] - #vPush

Cat 
$$x[] = \underline{new} Cat[r]$$
;

= -#oPush x[]
Operator Stack Semantic

Semantic Action Stack

- #vPush

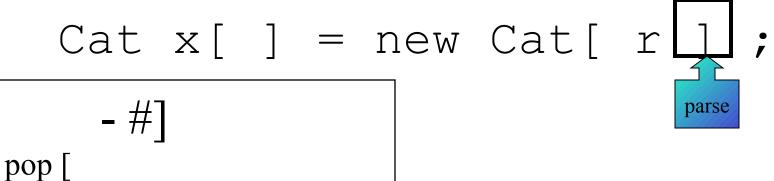
Operator Stack

Cat - #tPush

x[] - #vPush

Cat 
$$x[] = new Cat[\underline{r}];$$

Cat 
$$x[] = new Cat[r];$$



No infix to postfix conversion necessary on expression r

r - #iExist
 Cat - #tPush
 x[] - #vPush
 Semantic Action Stack

- #oPushOperator Stack

Cat x[] = new Cat[r]

- #new[]

r = SAS.pop() this is an expression Test that r is an integer

Cat - #tPush

x[] - #vPush

Semantic Action Stack

= -#oPush

Operator Stack

```
Cat x[] = new Cat[r];
```

- #new[]

r = SAS.pop();

Test that r is an integer

Cat type\_sar = SAS.pop()

Test that an array of type Cat can be can be created Create a new\_sar for an array of type Cat with r elements SAS.push(new sar)

Note: A new\_sar is used when constructing a new object or new array. This is similar to using a ref\_sar when referencing a member attribute or member function.

= - #oPush
Operator Stack

new\_sar

- #new[]

parse

- #vPush

Cat x[] = new Cat[r]



## **-** #EOE

f (r \* 3, g < k);

## Function call with infix to postfix conversion of arguments

f

- #iPush

Operator Stack

- #oPush Operator Stack

Semantic Action Stack

- #iPush

bal\_sar - #BAL

( - #oPush

Operator Stack

Semantic Action Stack

r - #iPush

bal sar - #BAL

f - #iPush

Semantic Action Stack

( - #oPush

Operator Stack

r - #iExist

bal\_sar - #BAL

- #iPush

Semantic Action Stack

- #oPushOperator Stack

r -#iExist

\* -#oPush

bal\_sar -#BAL

( -#oPush

Operator Stack

Semantic Action Stack

```
lg < k);
              - #,
                            parse
       pop *
              3 = SAS.pop()
              r = SAS.pop()
              Test that r * 3 is valid
              t7 = r * 3
              SAS.push(t7)
                       bal sar
     - #oPush
                                     - #iPush
                        Semantic Action Stack
Operator Stack
```

g - #iPusht7 - #,

- #BAL

f - #iPush
Semantic Action Stack

bal sar

( - #oPush Operator Stack

g - #iExist
t7 - #,
bal\_sar - #BAL
f - #iPush
Semantic Action Stack

( - #oPush Operator Stack

g -#iExist

t7 -#,

< -#oPush

( -#oPush

Operator Stack

g -#iExist

t7 -#,

bal\_sar -#BAL

-#iPush

Semantic Action Stack

k - #iPush - #iExist t7 bal sar - #BAL - #iPush Semantic Action Stack

< -#oPush

```
k
           - #iExist
           - #iExist
           -#,
t7
bal sar
           - #BAL
           - #iPush
Semantic Action Stack
```

< - #oPush

( - #oPush

```
f (r * 3, q < k)
     -#)
pop < #<
     k = SAS.pop()
     g = SAS.pop()
     Test that g < k is valid
     t8 = g < k
      SAS.push(t8)
                     t8
                                 - #`
                                 -#,
                     bal sar
                                 - #BAL
     - #oPush
                                 - #iPush
                      Semantic Action Stack
Operator Stack
```

```
f (r * 3, g < k);

-#)
pop(
```

```
f (r * 3, g < k)

-#EAL

t8 = SAS.pop()

t7 = SAS.pop()

bal sar == SAS.pop()
```

Create al\_sar for t7 & t8 SAS.push(al\_sar)

al\_sar - #EAL f - #iPush

Operator Stack

# - #func

func\_sar

- #func

Operator Stack

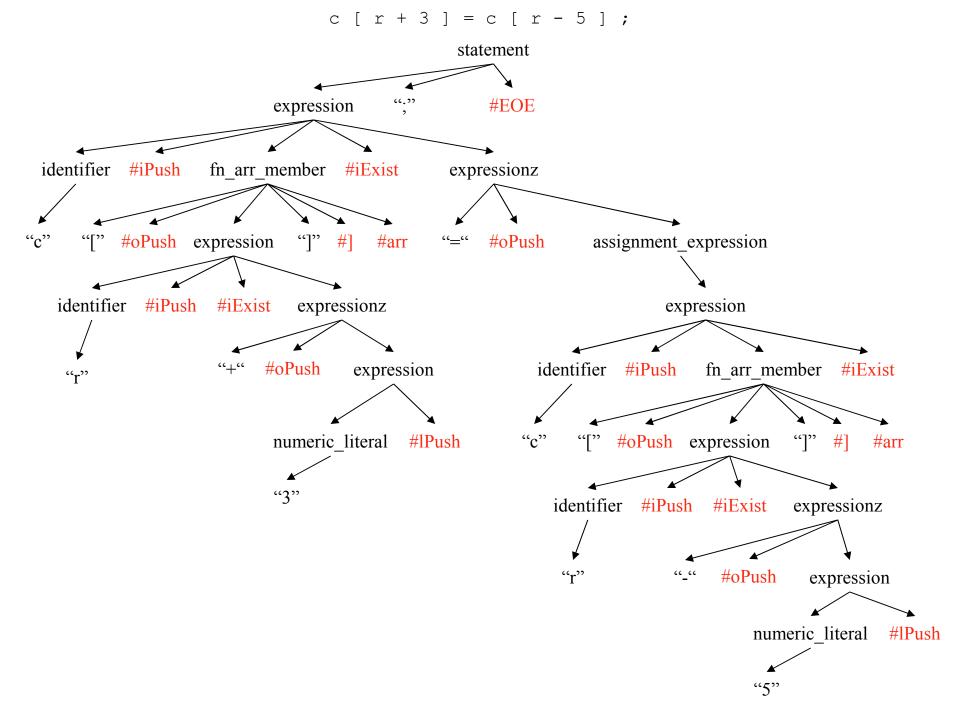
$$f (r * 3, g < k)$$
;

f(t7,t8) id\_sar - #iExist

f (r \* 3, g < k)
-#EOE

c [r + 3] = c [r - 5];

Array to Array assignment with infix to postfix conversion of index expressions



$$[r + 3] = c[r - 5];$$

C

- #iPush

Operator Stack

$$c [r + 3] = c [r - 5];$$

- #oPush

Operator Stack

C

- #iPush

$$c [r + 3] = c [r - 5];$$

r - #iPushc - #iPushSemantic Action Stack

Operator Stack

- #oPush

$$c [r + 3] = c [r - 5];$$

$$c [r + 3] = c [r - 5];$$

- + #oPush
  - #oPush

Operator Stack

r - #iExist

c - #iPush

$$c [r + 3] = c [r - 5];$$

+ - #oPush

- #oPush

Operator Stack

- #lPush

r - #iExist

c - #iPush

c [ r + 3] = c [

-#]

pop + #+

$$3 = SAS.pop()$$
 $r = SAS.pop()$ 

Test that r + 3 is valid

 $t23 = r + 3$ 
 $SAS.push(t23)$ 

**-** #]

$$c [r + 3] = c [r - 5];$$
pop[

t23 -#]
c -#iPush
Semantic Action Stack

c [r + 3] = c [r - 5];

- #arr

t23 = SAS.pop()

c = SAS.pop()

Test that t23 is an integer

parse

Create an arr\_sar for the array c with index t23

SAS.push(arr\_sar)

arr\_sar

- #arr

Operator Stack

$$c [r + 3] = c [r - 5];$$

### - #iExist

arr\_sar = SAS.pop()

Test that an arry c exist in the current scope Create an id\_sar for the array c with index t23 SAS.push(c[t23])

Note: How the id\_sar must be able to work with a simple variable (e.g., r), a function call (e.g., f(t7,t8)) or an array reference (e.g., c[t23]).

Think of a sar as a base class, with an id\_sar as a derived class and the variable, function and array sars as classes derived from id\_sar.

c[t23] id sar - #iExist

Semantic Action Stack

$$c [r+3] = c [r-5];$$

Operator Stack

c[t23] id\_sar - #iExist

$$c [r + 3] = c [r - 5];$$

c - #iPush

c[t23] id\_sar - #iExist

Semantic Action Stack

= -#oPush

$$c [r + 3] = c [r - 5];$$

$$c [r + 3] = c [r - 5];$$

- #oPush

= -#oPush

Operator Stack

r - #iPush

c - #iPush

c[t23] id\_sar - #iExist

$$c [r + 3] = c [r - 5];$$

r - #iExist
c - #iPush

c[t23] id\_sar - #iExist

Semantic Action Stack

Operator Stack

- #oPush

$$c [r + 3] = c [r - 5];$$

- - #oPush
- #oPush
- = -#oPush

Operator Stack

r - #iExist

c - #iPush

c[t23] id\_sar - #iExist

$$c [r+3] = c [r-5];$$

- #oPush

= -#oPush

Operator Stack

5 - #lPush

r - #iExist

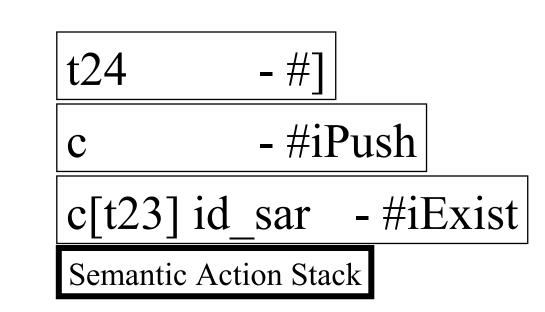
c - #iPush

c[t23] id sar - #iExist

- #oPush

- #oPush

Operator Stack



parse

$$c [r + 3] = c [r - 5];$$
pop[

t24 - #]
c - #iPush

c[t23] id\_sar - #iExist

Semantic Action Stack

- #oPush

c [ r + 3 ] = c [ r - 5];
$$- \# arr$$

$$t24 = SAS pop()$$

$$t24 = SAS.pop()$$
  
 $c = SAS.pop()$ 

Test that t24 is an integer

Create an arr\_sar for the array c with index t24

SAS.push(arr\_sar)

arr\_sar - #arr

c[t23] id\_sar - #iExist

Semantic Action Stack

= -#oPush

$$c[r+3]=c[r-5]$$
;

### - #iExist

arr\_sar = SAS.pop()

Test that an arry c exist in the current scope Create an id\_sar for the array c with index t24 SAS.push(c[t24])

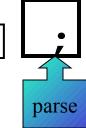
= -#oPush

Operator Stack

c[t24] - #iExist

c[t23] id sar - #iExist

c [r + 3] = c [r - 5]



#### - #EOE

## Multi-Threaded Semantic Actions

- •Implementing Multi-Threaded features is not part of CS4490 only parsing the statements is require.
- •DO NOT Implement these features until you have completed your Compiler!

Spawn a Thread

d

- #iPush

Operator Stack

d

- #iExists

Operator Stack

go - #iPush

d - #iExist

Semantic Action Stack

bal\_sar - #BAL

go - #iPush

( - #oPush d - #iExist

Operator Stack Semantic Action Stack

i - #iPush
bal\_sar - #BAL
go - #iPush
d - #iExist
Semantic Action Stack

( - #oPush Operator Stack

al sar - #EAL

go - #iPush

d - #iExist

Semantic Action Stack

func\_sar - #func

d - #iExist

Semantic Action Stack

- #iPush

ref\_sar - #rExist

Semantic Action Stack

# - #spawn

sar = SAS.pop()

sar must be an integer variable defined in the local scope (just like #iExists)

ref\_sar == SAS.pop()

Test that the ref\_sar is a valid function call. Note: The variable will be used to set the status of the spawned thread (-1 failed, > 0 for the thread spawned successfully). It is possible for the variable to be passed as a parameter to the function allowing the function to know it's thread id (cool).

Operator Stack