

MEAM 520

Lecture 13: Graph Representations

Practical

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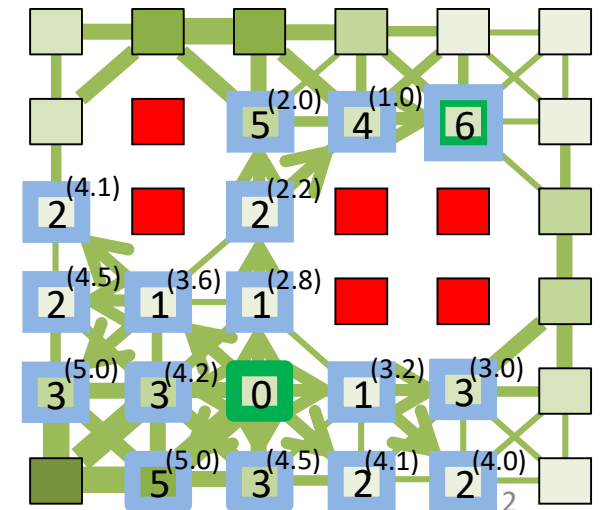
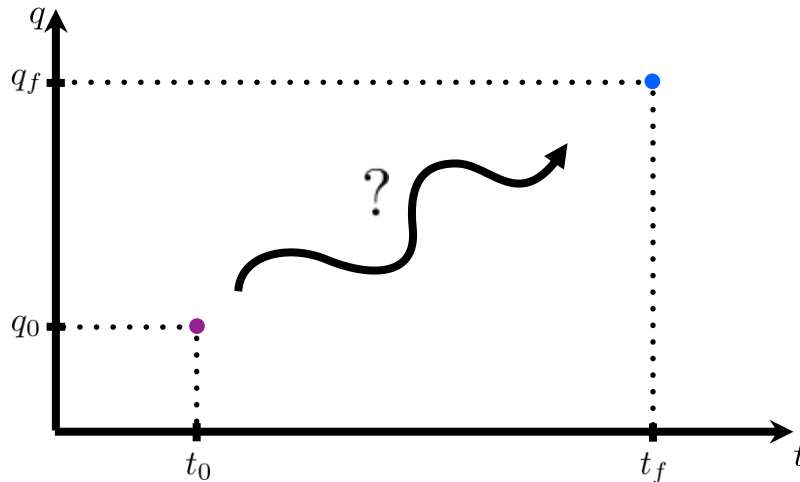
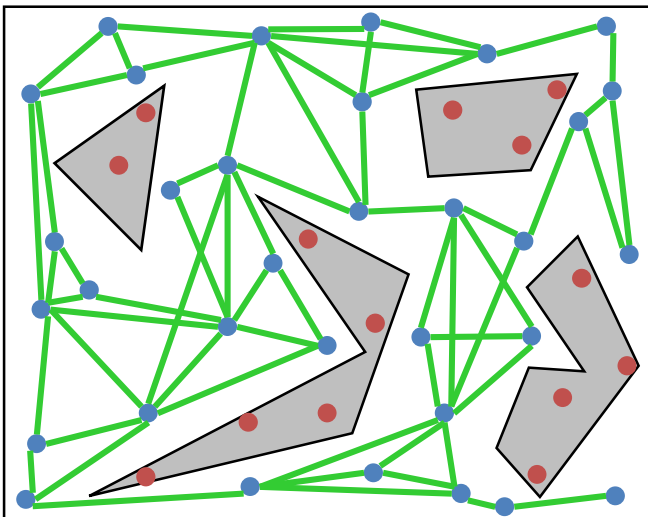
Mechanical Engineering & Applied Mechanics

University of Pennsylvania

Last Time: Trajectory Planning

Planning strategy:

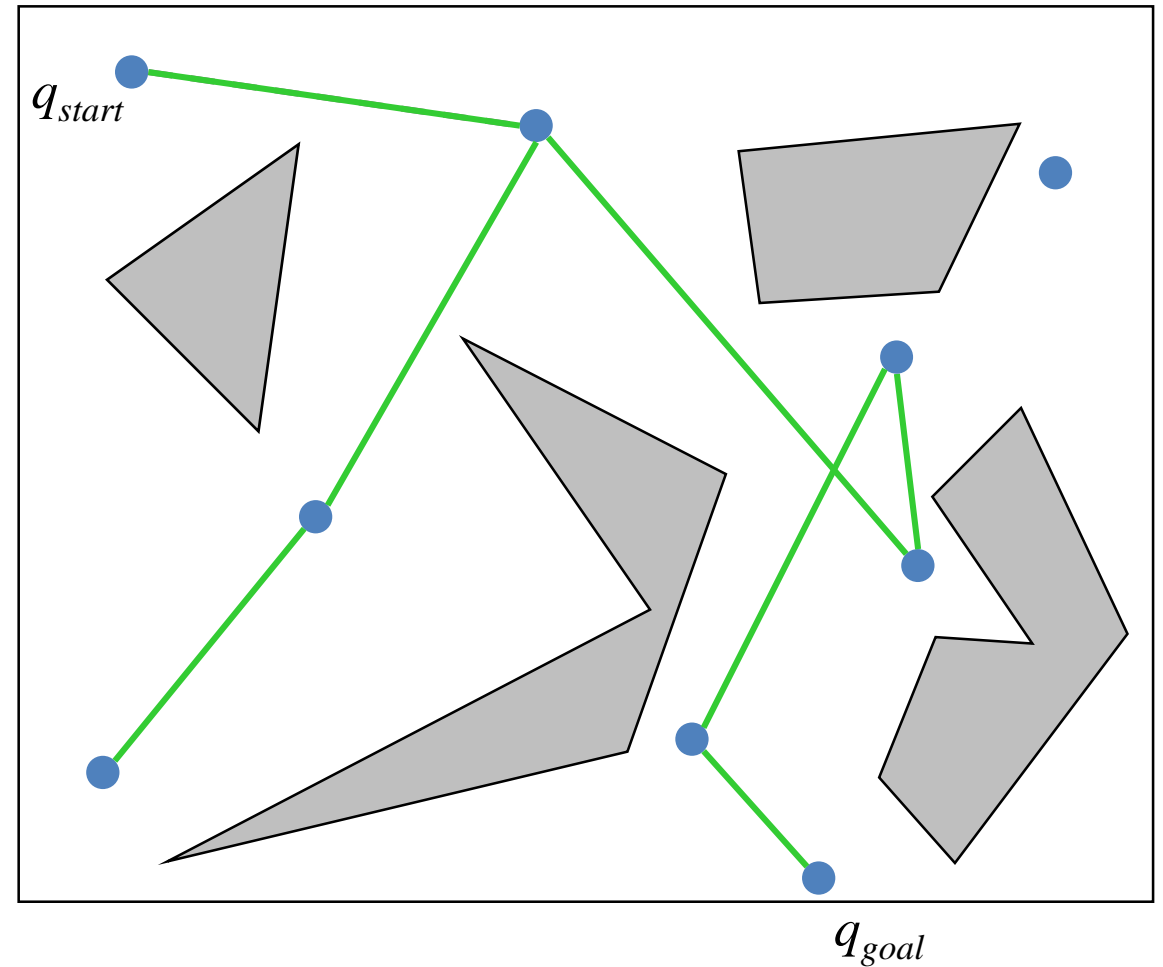
1. Convert your free C-space into a graph/roadmap
2. Find a path from q_{start} to a node q_a that is in the roadmap
3. Find a path from q_{goal} to a node q_b that is in the roadmap
4. Search the roadmap for a path from q_a to q_b



Last Time: Rapidly-exploring Random Trees (RRTs)

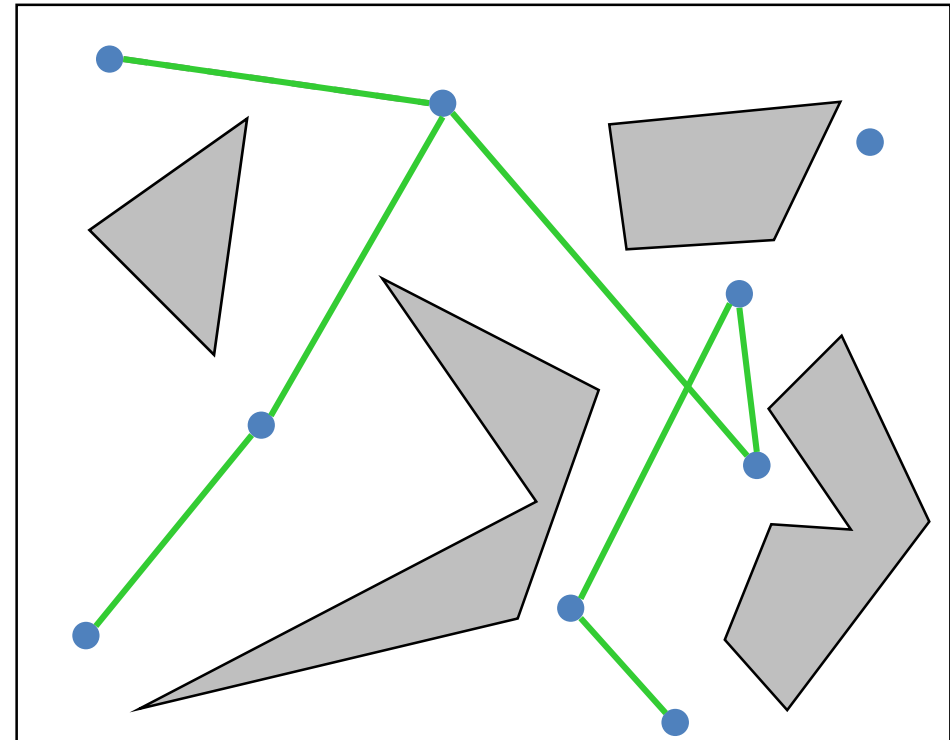
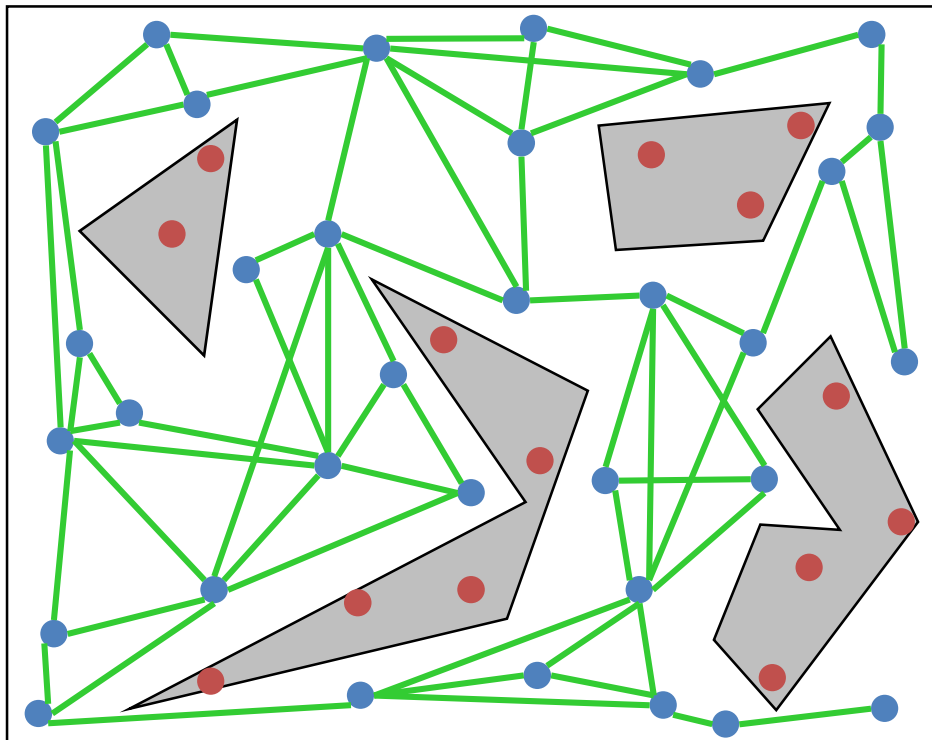
Combine graph construction and path search

Bias the search by changing
your sampling distribution or
connection strategy

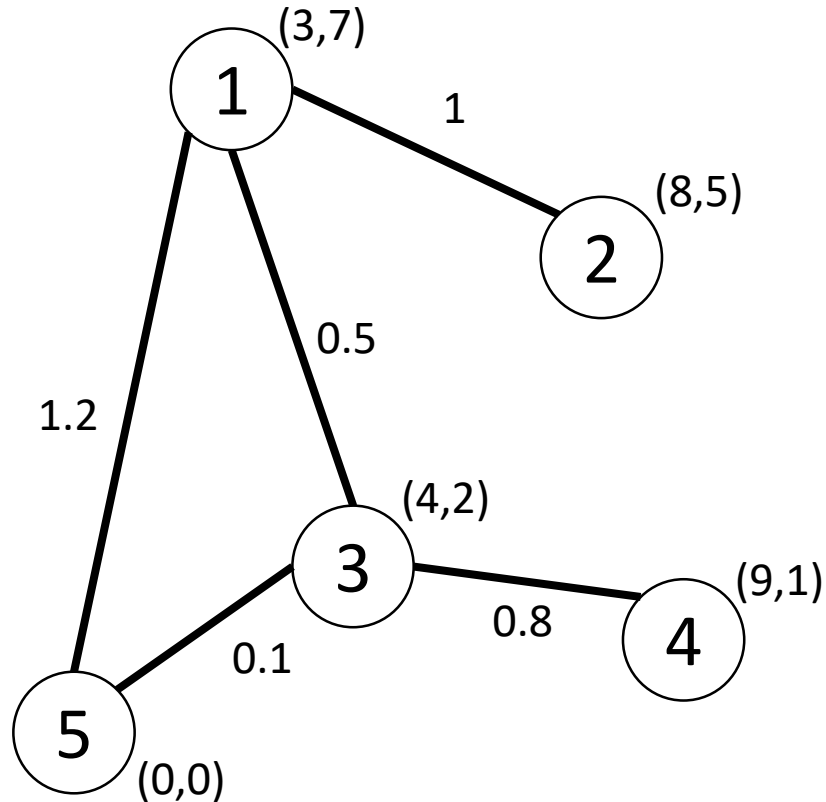


These algorithms are fine to describe pictorially.

But how do we actually represent this graph structure?



Graph Components



$$V = \{v_1, v_2, v_3, v_4, v_5\}$$

The edge $\{v_i, v_j\}$ connects vertices v_i and v_j

$$E = \{ \{v_1, v_2\}, \{v_1, v_3\}, \{v_1, v_5\}, \{v_3, v_4\}, \{v_3, v_5\} \}$$

In code, we can write:

$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0 \end{bmatrix}$$

Each row is a coordinate location

$$E = \begin{bmatrix} 1, & 2, & 1.0; \\ 1, & 3, & 0.5; \\ 1, & 5, & 1.2; \\ 3, & 4, & 0.8; \\ 3, & 5, & 0.1 \end{bmatrix}$$

vertices cost

Example Use: Dijkstra

V = [3, 7;
8, 5;
4, 2;
9, 1;
0, 0]

E = [1, 2, 1.0;
1, 3, 0.5;
1, 5, 1.2;
3, 4, 0.8;
3, 5, 0.1]

Say we are currently expanding vertex i

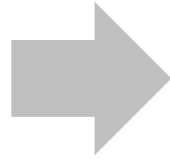
This step can be slow

1. Search first 2 columns of E for i
2. Add the other vertex to the neighbor list and the cost from the 3rd column

Note: if you have a cost metric, you may have to go back into the V matrix to calculate the cost.

Adjacency Matrix

$E = \begin{bmatrix} 1, & 2, & 1.0; \\ 1, & 3, & 0.5; \\ 1, & 5, & 1.2; \\ 3, & 4, & 0.8; \\ 3, & 5, & 0.1 \end{bmatrix}$



$E = \begin{bmatrix} 0, & 1, & 1, & 0, & 1; \\ 1, & 0, & 0, & 0, & 0; \\ 1, & 0, & 0, & 1, & 1; \\ 0, & 0, & 1, & 0, & 0; \\ 1, & 0, & 1, & 0, & 0 \end{bmatrix}$

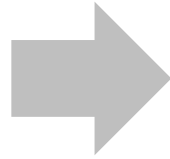
Entry (i, j) is 1 if there is an edge
0 if there is no edge

Looking up neighbors is really fast!

It is possible to also store costs

Adjacency Matrix

E = [1, 2, 1.0;
1, 3, 0.5;
1, 5, 1.2;
3, 4, 0.8;
3, 5, 0.1]



E = [Inf, 1.0, 0.5, Inf, 1.2;
1.0, Inf, Inf, Inf, Inf;
0.5, Inf, Inf, 0.8, 0.1;
Inf, Inf, 0.8, Inf, Inf;
1.2, Inf, 0.1, Inf, Inf]

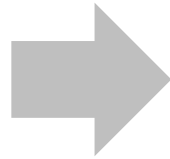
Entry (i, j) is cost if there is an edge
Inf if there is no edge

Looking up neighbors is really fast!

It is possible to also store costs

Adjacency Matrix

E = [1, 2, 1.0;
1, 3, 0.5;
1, 5, 1.2;
3, 4, 0.8;
3, 5, 0.1]



E = [Inf, 1.0, 0.5, Inf, 1.2;
1.0, Inf, Inf, Inf, Inf;
0.5, Inf, Inf, 0.8, 0.1;
Inf, Inf, 0.8, Inf, Inf;
1.2, Inf, 0.1, Inf, Inf]

Entry (i, j) is cost if there is an edge
Inf if there is no edge

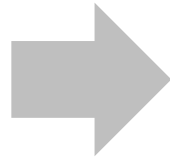
Looking up neighbors is really fast!

It is possible to also store costs

BUT, these extra 0/Inf entries take a lot of space

Adjacency List: Store edges by vertex

$E = [1, 2, 1.0;$
1, 3, 0.5;
1, 5, 1.2;
3, 4, 0.8;
3, 5, 0.1]



$E = \{[2, 1.0; 3, 0.5; 5, 1.2],$
[1, 1.0],
[1, 0.5; 4, 0.8; 5, 0.1],
[3, 0.8],
[1, 1.2; 3, 0.1]\}

Entry i contains all of the edges incident on i

Only take as much storage space as needed

It is possible to also store costs

Looking up neighbors is fast (though not as fast as Adjacency Matrices)

Adjacency Matrices vs Adjacency Lists

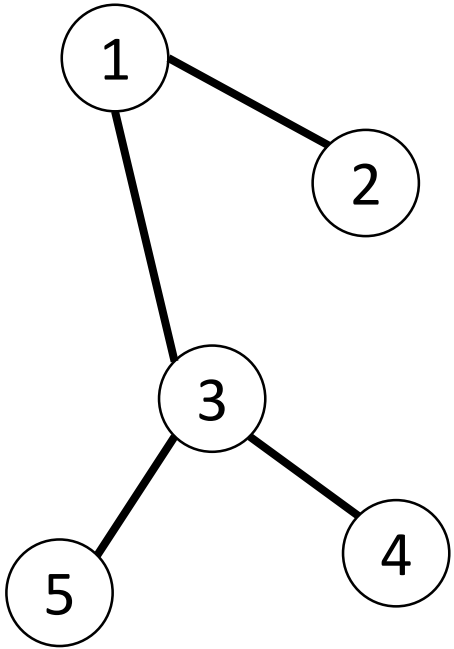
$E = \begin{bmatrix} \text{Inf} & 1.0 & 0.5 & \text{Inf} & 1.2 \\ 1.0 & \text{Inf} & \text{Inf} & \text{Inf} & \text{Inf} \\ 0.5 & \text{Inf} & \text{Inf} & 0.8 & 0.1 \\ \text{Inf} & \text{Inf} & 0.8 & \text{Inf} & \text{Inf} \\ 1.2 & \text{Inf} & 0.1 & \text{Inf} & \text{Inf} \end{bmatrix}$

- Use this when you have a dense graph (lots of edges)
- Use this when neighbor lookup has to be VERY fast

$E = \{[2, 1.0; 3, 0.5; 5, 1.2], [1, 1.0], [1, 0.5; 4, 0.8; 5, 0.1], [3, 0.8], [1, 1.2; 3, 0.1]\}$

- Use this when you have a sparse graph (not many edges)
- Use this when adding new vertices has to be very fast

Example Use: RRT



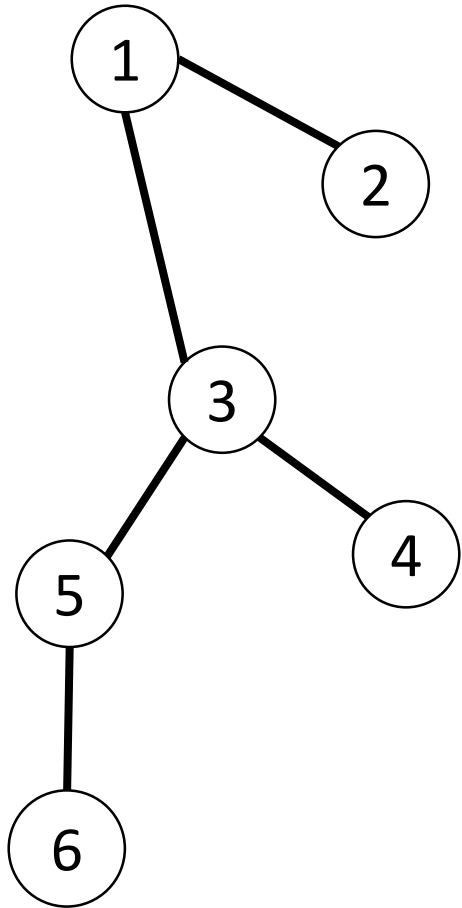
$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0 \end{bmatrix}$$

$$E = \begin{bmatrix} 1, & 2; \\ 1, & 3; \\ 3, & 4; \\ 3, & 5 \end{bmatrix}$$

$$E = \begin{bmatrix} 0, 1, 1, 0, 0; \\ 1, 0, 0, 0, 0; \\ 1, 0, 0, 1, 1; \\ 0, 0, 1, 0, 0; \\ 0, 0, 1, 0, 0 \end{bmatrix};$$

$$E = \{ [2, 3]; \\ [1]; \\ [1, 4, 5]; \\ [3]; \\ [3] \}$$

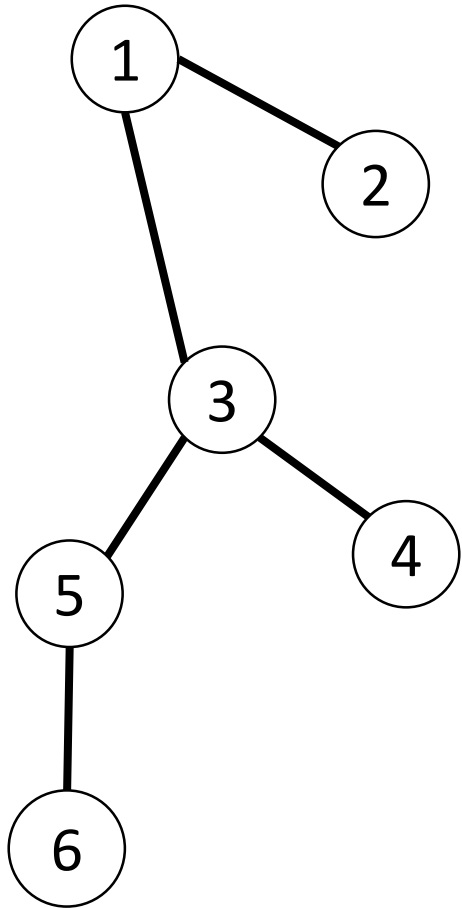
Example Use: RRT


$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0; \\ 0, & -2 \end{bmatrix}$$
$$E = \begin{bmatrix} 1, & 2; \\ 1, & 3; \\ 3, & 4; \\ 3, & 5 \end{bmatrix}$$
$$E = \begin{bmatrix} 0, 1, 1, 0, 0; \\ 1, 0, 0, 0, 0; \\ 1, 0, 0, 1, 1; \\ 0, 0, 1, 0, 0; \\ 0, 0, 1, 0, 0 \end{bmatrix};$$
$$E = \{ [2, 3]; \\ [1]; \\ [1, 4, 5]; \\ [3]; \\ [3] \}$$

Needed operations:

1. Add a new vertex

Example Use: RRT

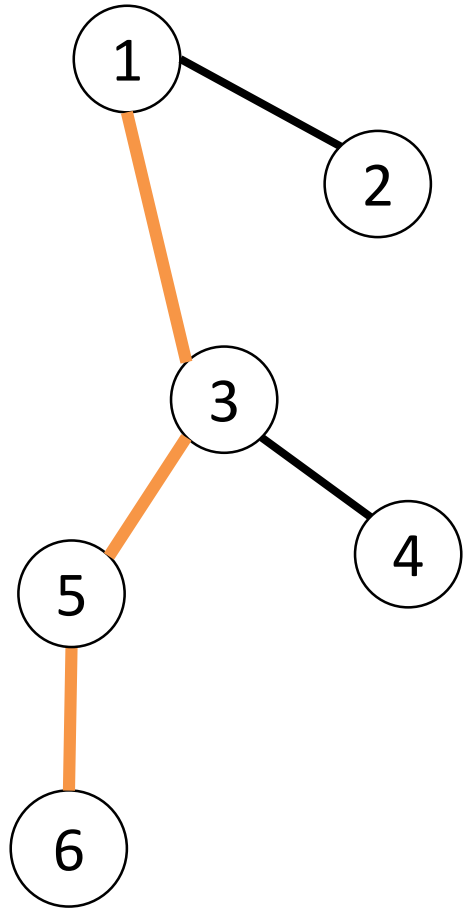

$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0; \\ 0, & -2 \end{bmatrix}$$
$$E = \begin{bmatrix} 1, & 2; \\ 1, & 3; \\ 3, & 4; \\ 3, & 5; \\ 5, & 6 \end{bmatrix}$$

Needed operations:

1. Add a new vertex
2. Add a new edge

$$E = \begin{bmatrix} 0, & 1, & 1, & 0, & 0, & 0; \\ 1, & 0, & 0, & 0, & 0, & 0; \\ 1, & 0, & 0, & 1, & 1, & 0; \\ 0, & 0, & 1, & 0, & 0, & 0; \\ 0, & 0, & 1, & 0, & 0, & 1; \\ 0, & 0, & 0, & 0, & 1, & 0 \end{bmatrix};$$
$$E = \{ [2, 3]; \\ [1]; \\ [1, 4, 5]; \\ [3]; \\ [3, 6]; \\ [5] \}$$

Example Use: RRT



$$V = \begin{bmatrix} 3, & 7; \\ & 8, & 5; \\ & 4, & 2; \\ & 9, & 1; \\ & 0, & 0; \\ & 0, & -2 \end{bmatrix}$$

$$E = \begin{bmatrix} 1, & 2; \\ 1, & \leftarrow 3; \\ 3, & \nearrow 4; \\ 3, & \leftarrow 5; \\ 5, & \nearrow 6 \end{bmatrix}$$

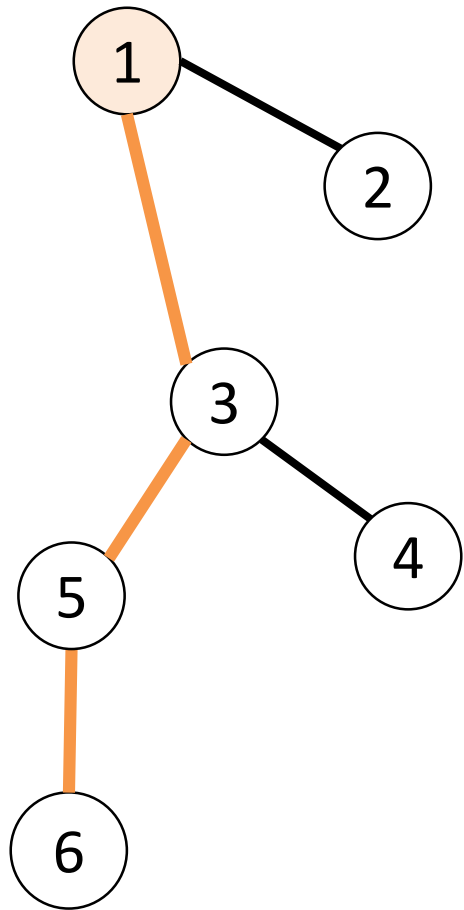
Needed operations:

1. Add a new vertex
2. Add a new edge
3. Backtrack the path

$$E = \begin{bmatrix} 0, 1, 1, 0, 0, 0; \\ 1, 0, 0, 0, 0, 0; \\ 1, 0, 0, 1, \cancel{1}, 0; \\ 0, 0, 1, \nearrow 0, 0, 0; \\ 0, 0, 1, 0, 0, \cancel{1}; \\ 0, 0, 0, 0, 1, 0 \end{bmatrix};$$

$$E = \{ [2, 3]; \\ [1]; \\ [1, 4, \cancel{5}]; \\ [3]; \\ [3, \cancel{6}]; \\ [5] \}$$

Example Use: RRT



$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0; \\ 0, & -2 \end{bmatrix}$$

$$E = \begin{bmatrix} 1, & 2; \\ 1, & \leftarrow 3; \\ 3, & \nearrow 4; \\ 3, & \leftarrow 5; \\ 5, & \nearrow 6 \end{bmatrix}$$

Needed operations:

1. Add a new vertex
2. Add a new edge
3. Backtrack the path

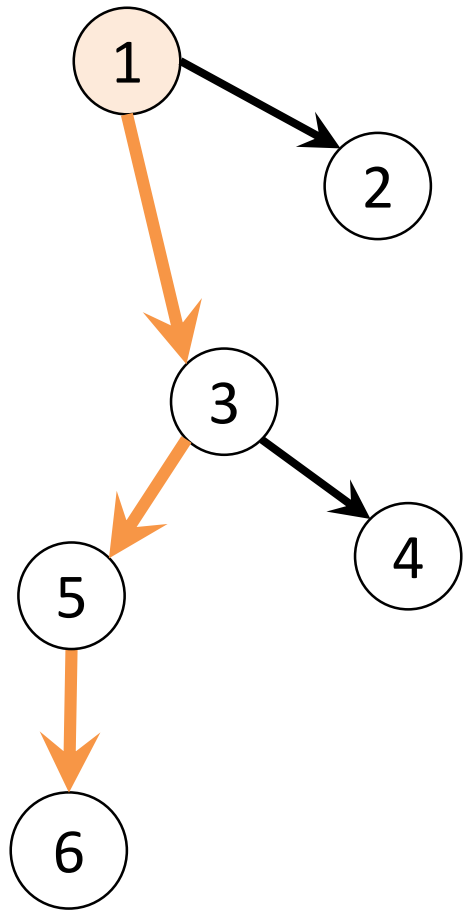
Trees are special because there is a ROOT node

We are trying to find a path back to the root

Add some directionality to the representation!

$$E = \begin{bmatrix} 0, 1, 1, 0, 0, 0; \\ 1, 0, 0, 0, 0, 0; \\ 1, 0, 0, 1, \cancel{1}, 0; \\ 0, 0, 1, \nearrow 0, 0, 0; \\ 0, 0, 1, 0, 0, \cancel{1}; \\ 0, 0, 0, 0, 1, 0 \end{bmatrix}; \quad E = \{ [2, 3]; [1]; [1, 4, \cancel{5}]; [3]; [3, 6]; [5] \}$$

Example Use: RRT



$$V = \begin{bmatrix} 3, & 7; \\ 8, & 5; \\ 4, & 2; \\ 9, & 1; \\ 0, & 0; \\ 0, & -2 \end{bmatrix}$$

$$E = \begin{bmatrix} 1, & 2; \\ 1, & 3; \\ 3, & 4; \\ 3, & 5; \\ 5, & 6 \end{bmatrix}$$

directionality inherent
in column order

Needed operations:

1. Add a new vertex
2. Add a new edge
3. Backtrack the path

Trees are special because there is a ROOT node

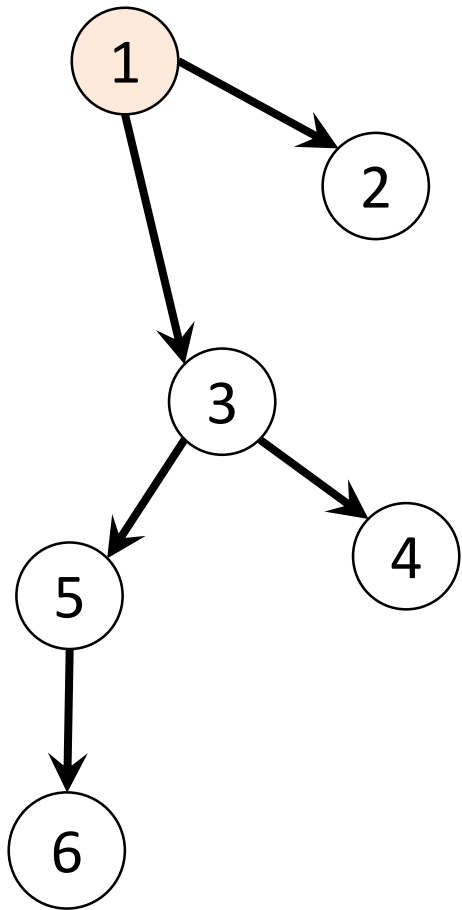
We are trying to find a path back to the root

Add some directionality to the representation!

$$E = \begin{bmatrix} 0, & 1, & 1, & 0, & 0, & 0; \\ -1, & 0, & 0, & 0, & 0, & 0; \\ -1, & 0, & 0, & 1, & 1, & 0; \\ 0, & 0, & -1, & 0, & 0, & 0; \\ 0, & 0, & -1, & 0, & 0, & 1; \\ 0, & 0, & 0, & 0, & -1, & 0 \end{bmatrix}; \quad E = \{ [2, 3]; \text{??} \\ [1]; \\ [1, 4, 5]; \\ [3]; \\ [3, 6]; \\ [5] \}$$

Use -1s for “from” node and 1s for “to” nodes

Tree Representations



```
Node {  
    coord: [x,y]  
    parent: i  
    children: [j,k,...]  
}
```

Implement this as a
struct in MATLAB

```
v(1) = {  
    coord: [3,7]  
    parent: NaN  
    children: [2,3]  
}
```

```
v(3) = {  
    coord: [4,2]  
    parent: 1  
    children: [4,5]  
}
```

To find a path back to the root, follow the parent pointer

Tree struct vs Adjacency Info

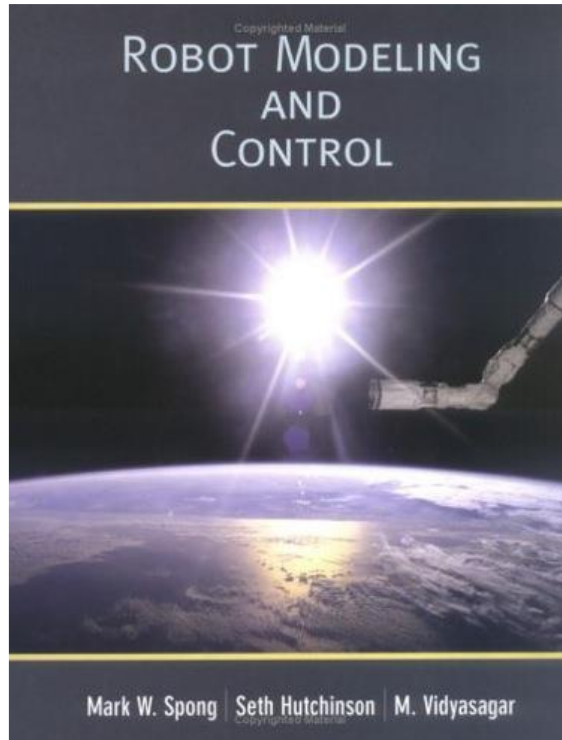
```
Node {  
    coord: [x,y]  
    parent: i  
    children: [j,k,...]  
    cost-to-goal: c  
}
```

```
V = [3, 7;  
      8, 5;  
      4, 2;  
      9, 1;  
      0, 0;  
      0, -2]  
E = [1, 2;  
      1, 3;  
      3, 4;  
      3, 5;  
      5, 6]
```

- Can store extra information in the struct
- Can store structs in data structure useful for search (lookup k-d trees)
- Matrix manipulation usually faster than struct manipulation

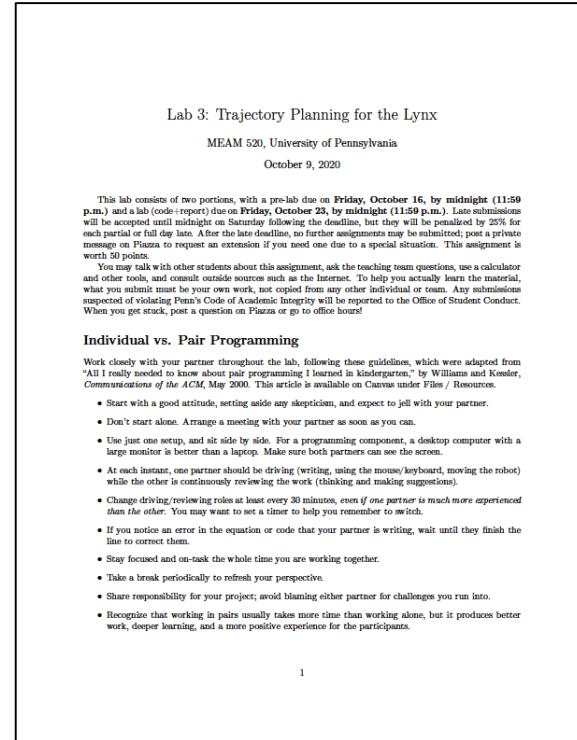
That brings us to the end of position/orientation analysis

Next time: Velocity



Chapter 4: Velocity Kinematics

- Read 4.intro – 4.4



Lab 3: Planning due 10/23