

DM HW 4

211-6260

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Q1

(A)

$$\begin{matrix} & \begin{matrix} \bullet & \bullet & \bullet & \bullet & & & \bullet & \bullet & \bullet & \bullet \end{matrix} \\ & \begin{matrix} 1 & 2 & 3 & 4 & & & 7 & 8 & 9 & 10 \end{matrix} \\ \begin{pmatrix} 0 & 1 & 0 & 0 & \dots & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & \dots & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & \dots & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \dots & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & \dots & 1 & 0 & 1 \\ 0 & \dots & \dots & \dots & \dots & 0 & 1 & 0 \end{pmatrix} \end{matrix}$$

(B)



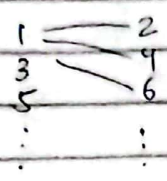
$$\begin{pmatrix} 0 & 1 & 1 & \dots & 1 & 1 \\ 1 & 0 & 0 & \dots & 0 & 0 \\ 1 & 0 & 0 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & 0 & 0 & \dots & 0 & 0 \end{pmatrix}$$

(c)



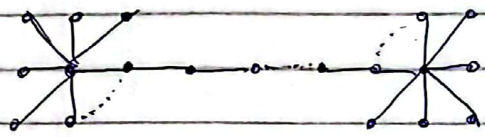
$$\begin{pmatrix} 0 & 1 & 0 & \dots & 0 & 1 \\ 1 & 0 & 1 & \dots & 0 & 0 \\ 0 & 1 & 0 & 1 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 1 & 0 & 1 \\ 1 & 0 & 0 & \dots & 0 & 1 & 0 \end{pmatrix}$$

(D)



$$\begin{bmatrix} 0 & 1 & 0 & 1 & 0 & \dots & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & \dots & 0 & 1 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 1 & 0 & 1 & \dots & \dots & 0 & 1 \\ 1 & 0 & 0 & 0 & \dots & \dots & 1 & 0 \end{bmatrix}$$

(F)



1 to n n+1 to m m+1 to p

$$\begin{bmatrix} \overbrace{\begin{matrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots \\ 1 & 0 & 0 & 0 \end{matrix}}^n & \overbrace{\begin{matrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{matrix}}^m & \begin{matrix} 000 \dots 0 \\ 000 \dots 0 \\ 000 \dots 0 \\ 000 \dots 0 \\ 000 \dots 0 \\ 000 \dots 0 \\ 000 \dots 0 \end{matrix} \end{bmatrix}$$

Q#2

$G(V, E)$ is unweighted and undirected graph.

$$N_v = |V|$$

$$N_e = |E|$$

A = adjacency Matrix.

We are given that.

$[A^k]_{ij}$ is $\Rightarrow k^{\text{th}}$ power of adjacency matrix.
 ij is number walks of length k .

A) Base Case $k=1$:

A being adjacency matrix of graph G .

$[A^k]_{ij}$ is walk length k from i to j vertex.

so $[A^k]_{ij}$, $k=1$ means $\{i, j\} \in E$ and

E are edges of G . and $[A^1]_{ij}$ represents walklength of 1 from i to j with direct edges between them.

Graph G has as i vertex to $v_2 j$ vertex and edge $\{i, j\} \in E$ then $k=1$ representing walk length of 1. Hence $k=1$ is true.

B) Inductive Step

For $n=k-1$;

$[A^k]_{ij}$ holds for $k=1$

we know

$$A^k = A^{k-1} \times A$$

$$A^k_{ij} = \sum_{p=1}^{k-1} A^p_{ip} \times A_{pj}$$

where p is some vertex cases.

$$i - \{p, i\} \in G \rightarrow A_{pi} = 1$$

$$\text{if } - \{p, i\} \notin G \rightarrow A_{pi} = 0.$$

so in particular all walks from i - j of length k are from some vertex p . therefore

$$A_{ij}^k = \sum_{p=1}^{k-1} A_{ip}^{k-1} \times A_{pj}$$

represents total i - j walks holds true for. $k-1$.

$$A^2 = \begin{bmatrix} 2 & 1 & 0 & 1 \\ 1 & 2 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \quad \text{2 means 2 step path.}$$

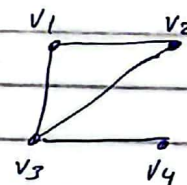
C) Generalization for Digraphs: Yes, the matrix multiplication result also applies for Digraphs where $[A^k]_{ij}$ gives the number of i - j walks of length k in graph.

$$D) |A^2|_{ij} = d_i, \quad |A^3|_{ij}/2 = \# \Delta.$$

$$\text{As } \frac{\text{tr}(A^2)}{2} = N_e \Rightarrow \text{no of edges}$$

$$\frac{\text{tr}(A^3)}{6} = \# \Delta$$

Suppose we take graph $G \Rightarrow$



$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

Now, computing A^2 and A^3 ~~computation~~

$$A^2 = A \cdot A$$

$$A^2 = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$A^2 = \begin{pmatrix} 2 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 \\ 1 & 1 & 2 & 0 \\ 1 & 1 & 0 & 2 \end{pmatrix}$$

$$\text{As } Ne = \frac{tn(A^2)}{2}$$

$$tn(A^2) = 2+2+2+2 \Rightarrow 8$$

$$Ne = 8/2 \Rightarrow 4.$$

We can see G has 4 edges.

$$\text{Now Similarly } A^3 = \begin{pmatrix} 2 & 3 & 4 & 1 \\ 3 & 2 & 4 & 1 \\ 3 & 3 & 2 & 2 \\ 1 & 1 & 3 & 0 \end{pmatrix}$$

$$\text{We know } \# \Delta = \frac{tn(A^3)}{6} \Rightarrow \frac{2+2+2+0}{6} \Rightarrow 1$$

We see that 1 triangle exists in G

Hence proved.

E) Checking connectivity, $S = A + A^2 + A^3$ would provide full matrix of non-zero values, showing multiple paths of varying lengths connecting every pair. If any vertex is not connected, it will have 0 values in S matrix as no walk $i-j$ exists.