System Call

System call is a way to shift from user mode to kernel mode, and access operating System.

At Kernel in the Sense Operating System. The main part or heart of 0.5.

A System call is a way your a work the o.s. all is a way your a

· File Related =>

Open(), Read(), Worite(), close()
(reate file et.c.

- · Device Related =>
- Read, write, Reposition, iocti, fonth
- · Information =>

9et Pid attributes 3 get System Time

jarnory '21
Su Mo Tu We Th Fr So

31 1 2

3 4 5 6 7 8 9

10 11 12 13 14 15 16

17 18 19 20 21 22 23

24 25 26 27 28 29 30

roces (ontral =)

(ommunication =)

0-> Child

-> Parent

7			-0	4				
2×!-	Main () { Pri	nt	2	+	im	C.	S	
	fork(); Hel	lo,	6	4 (Shi	11	4	
	Printf ("Hello"); P	ישיפר	# (C	<u>m</u>	وب	ve	atte
17	8					()		1
NOTES	P				feb	ירטכ	ary	21
		Su	Мо	Tu	We	Th	Fr	Sa
4	fork	F 7	1 8	2	3		5 12	6 13
	o / tve	P ₁₄						
	C P	21						
	učelo Hello	28			泛流	17 h	V	5-4
	Dreams are illustration from the book your soul is writing about	ıt you.	- Mai	rsha	Non	nan		

Load, Execute, abort, Fork, Signal, wait,

Pipe 1), (reate Idelete Connections, &

A fork() generally returns value

No child created

System Call

s used to create child Procey.

			S. Carlot
Ex3;-	Main () {		
	1)	ar a chainma a marian a mara a mara para na marian dia cata na paradian a da maria	and the same of th
10	fork();		
10	forkes;		
	Pointf ("Hel	10").	
11	3		
12		(, a, 1	
12			-
à	. 20/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/ 1/	Adam 1 7 may a market	
	P	The street of th	
	tork	2.00	
		Character	
		The same of the same	
	T PULL THE SE	P	
			1-2, -1
	tork		
10	Car	0	-7
		1	
4	Chald		
Hello	- CIVIC HO	Parent	Tolars 4
,	Heleo He	lo Kell	
<u> </u>			D
Tormula!	- 2"	The second secon	M. C. C. A.
	•	The state of the s	走到.
Child Po	DC 000 1 - 21	THE ELECTION OF THE	
TITLE IS	ocers :- 2h ==		
	the state of the s		
1 1 1 1 1		Bles water	11
A C			0.00

include < stdio-h>

main ()

int

include < Unistd-h>

							`
一个一个一个一个一个	7-1-	147		يفت ا	ei'		
Reads ist fork () of it stat	eme	fas		a		}	
Parent Program's Child is creat	ed.		1				
	# 1	المراكب الم			J.	avi	ж.
(1) P (3)	4	7	3	1			
10-11-500-	ÿ	1.00		_			
Ctown P						1	
		10	S	UNI	DAY	' 1	. /
44 is operator AND, Now #	6	ch:	d	Pr	500	61	λ
oretwins a so off o/1	Stot	9.07	1/1/2		<u>ල</u>		
ie it oreturns false. Therefor	10	H.	2	10	n c	9-1	
forke) Statement will not in	اردا		<u></u>	111	1	21	
		رف			a		
print the Hello Statement.	15 7	-	+ /h				
				feb)(LIC)(III	21
NOTES O THE	Su	Мо	- Tu		Th	•	Sc
<u> </u>		1.1	2	3	4	5	6
	7	8	9	10	11	12	13
Hello	21	15	16 23	17 24	18 25	19 26	20
		22	77.7				

Hello

GGK-04 (01 ₀
But the parent 'p' gretwing the means
1 ff1 = 1 -: it invokes pext
torke) do parent will (reate is
and the bacen.
12
tve
CI P
1
mello of ptue
But of all
But child process has veturn o.
Latent's type
Child's a greturing false
i.e. je (former)
5 (ORC) IT FORKI)
100
· Hence Skipss fork() inside
die all Jorke inside it
7 Print Hello.
But a
But Parent's oreturn the i.e the forther
The two honce
Jooks inlide la parents child execut
january 121
Su Mo Tu We Th Fr Sa
NOTES
10 11 12 12 1 1 1 1 1 1 1 1 1 1 1 1 1 1
17 18 19 20 21 22 22
24 25 26 27 28 29 30 Hello
Your many

You must pay the price if you wish to secure the blessing. - Andrew Jackson Hello Nello

UserMode	Modebit-1
User Proces Get System Executing Call	Return from System call
trap	
0	
Kernel Mode Call	A Carlotte
	Mode bit=0
Hand Disk.	
Hard Disk.	
Hand Disk.	
Hard Disk.	
Hand Disk. A Switching of bits called trap A Every application on process of modes	rang on User
Hand Disk. A Switching of bits called trap A Every application on process of model	rang on User righ Kernel february 12
Hard Disk. A Switching of bits called trap A Every application on process of Model wing System (all.	rang on User

* Proces and	Threads in Operating System
10	
Process	Threads (Userleve)
1) System call involved	1) There is no system [Call involved.
(2) OS treats different process differently.	2) All user level threads treated as slingle task for O.S.
(3) Different process Lave different Copies of Doda, files, Code.	3) Threads I have Same Copy of Code and data But
(4) Context Switching is slower.	Stack and origiston. 4) Context Switching IS faster.
(5) Blocking a process will not block the another	5) Blocking a thread will block entire
January 21 Su Mo Tu We Th Fr Sa 31 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23	dent (6) Interdependent
24 25 26 27 28 29 30 Writing books is the closest men ever of	come to childbearing Norman Mailer

January 21	71	
Week-04 (021-344)	THURSDAY 🚣 🗘	
Thoread Thoread2	Provent child	
TI T2 doubt not	Stack 12	
Stack Stack	Register R Stack	
10 Reg. Reg.	Regulor	
Code	e e	
11	Code S Code	
Data/	files files	
files	files files	
100		ME
	forhe)	
A CALL THE ALL AND A CONTROL OF THE ALL AND A	11 1.2	FEBRUARY
2 There are too level	s of threads:	ΑY
***	treanel level Thread	-
3 Uson level Thereach	hegines level invitate	
	(1) Kernel level threads	
1) User level Hureads	are managed by	
are managed by the	00	
5 User level- library.	The state of the s	MARCH
(2) (1907 level threads	(2) Kernel level threads	H
9) 55	are slower than User leve	P.
are typically fast.	supplied and the second	=
3) Context Switching	(3) (ontext Switching is	
is faster,	Slower -	
13 Just	and the second of	
(4) If one user level	(a) If one Kernel level	APRIL
NOTES flored perform	Haread blocked Mo To We The Fr Sa	F
blaking operation than	No offect 1 2 3 4 5 6	
entire Process get	On others. 7 8 9 10 11 12 13	1
blocked	21 22 23 24 25 26 27	
	28 Shiy Khera	-
Your positive action combined with	positive thinking results in success Shiv Khera	, i