

Blond's Rule

Choosing always the first non-basic variable with negative reduced cost will prevent any loop.

Th.

If we follow blond's rule we will not end up in loops

Proof:

Consider an instance of the tableau (T) in which the worst possible choice following Blond's rule

	...	$x_{p[i]}$	x_n	x_n
=		0	-	0
$x_{p[i]}$	0	0	-	0
.		1	-	0
.		0	...	0
.		...	+	1
$x_{p[t]}$	0	0	-	0

row t

↳ $\bar{b} = 0$ for the objective function to be constant

$$\bar{c}_n < 0, \bar{c}_{p[i]} = 0 \quad \forall i, \quad \bar{a}_{in} \leq 0 \quad \forall i \neq t, \quad \bar{a}_{tn} > 0$$

Let's assume that this table will loop: at some point x_n will enter the basis (\tilde{T}):

	\dots	$x_{b[i]}$	\dots	x_n	\dots	x_n
=		+	0	+	\dots	-

$\rightarrow \tilde{c}_i \geq 0 \ \forall i \neq n, \tilde{c}_n < 0$
 (this must happen to choose x_n as the new basis variable)

$$[\text{row } 0 \text{ of } \tilde{T}] = [\text{row } 0 \text{ of } T] + \sum_{i=1}^m \mu_i [\text{row } i \text{ of } T]$$

• $\tilde{c}_{b[i]} = \bar{c}_{b[i]} + \mu_i$ \rightarrow elements of the basis belong to the identity matrix

\downarrow
 $\mu_i \geq 0 \ \forall i \neq t$

• $\tilde{c}_{b[t]} = \tilde{c}_n = \bar{c}_{b[t]} + \mu_t$

\downarrow
 $\mu_t < 0$

• $\tilde{c}_n = \bar{c}_n + \sum_{i=1}^m \mu_i \bar{c}_{in} = \bar{c}_n + \sum_{i \neq t} \mu_i \bar{c}_{in} + \mu_t \bar{c}_{tn}$

$\Rightarrow \begin{cases} \tilde{c}_n \geq 0 \\ \tilde{c}_n < 0 \end{cases}$

\rightarrow impossible to reach \tilde{T} from T , hence following Bland's rule prevents any loop

Dealing with bounded variables

Always worked with: $0 \leq x_i \quad \forall_i$

We might have an upper bound: $0 \leq x_i \leq q_i \quad \forall_i$

or also a different lower bound: $l_{bi} \leq x_i \leq u_{bi} \quad \forall_i$

$$(l_{bi} \leq x_i \leq u_{bi}) - l_{bi}$$

$$0 \leq x'_i \leq q_i, \quad x'_i = x_i - l_{bi}$$

$$q_i = u_{bi} - l_{bi}$$

We could add this as new constraints to the matrix

↓

HUGE MATRIX

These constraints can be implicitly deduced while producing the theta using complemented variables, that make an upper bound a non-negativity constraint

$$x_j^c = q_j - x_j \geq 0 \quad (\text{yet no one implemented this})$$

$$x = \begin{bmatrix} x_B \\ x_F \end{bmatrix} = \begin{bmatrix} x_B \\ x_L \\ x_U \end{bmatrix} \rightarrow \begin{array}{l} \text{at lower bound} \\ \text{at upper bound} \end{array}$$

↓

$$Ax = b \leadsto \begin{bmatrix} B & L & U \end{bmatrix} \begin{bmatrix} x_B \\ x_L \\ x_U \end{bmatrix} = Bx_B + \underbrace{Lx_L}_F + Ux_U$$

$$\text{Basic solutions: } x_L = 0, x_U = q_U \Rightarrow x_B = B^{-1}b - B^{-1}Uq_U$$

Optimality test

$$\bar{c}_h \geq 0 \quad \forall x_h \text{ non-basic at the lower bound}$$

$$\bar{c}_h \leq 0 \quad \forall x_h \text{ non-basic at the upper bound}$$

Optimality test fails in two cases:

1) $\bar{c}_h < 0$ and x_h is non-basic at the lower bound

$$x_h: 0 \rightarrow \min \{\vartheta, \vartheta'\}$$

$$\vartheta = \min \left\{ \frac{\bar{b}_i}{\bar{a}_{ih}} : \bar{a}_{ih} > 0 \right\}$$

$$\vartheta' = \dots \quad (\text{to avoid going over the upper bound})$$

2) $\bar{c}_h > 0$ and x_h is non-basic at the upper bound

$$x_h: q_h \rightarrow w_h$$

↳ similar procedure than ϑ