

## Assignment 4

due Monday, May 7, 2018

1. Illustrate the Floyd-Warshall algorithm on the graph described by the following weight matrix:

$$W = \begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ 1 & 0 & \infty & 2 & \infty & \infty \\ \infty & 2 & 0 & \infty & \infty & -8 \\ -4 & \infty & \infty & 0 & 3 & \infty \\ \infty & 7 & \infty & \infty & 0 & \infty \\ \infty & 5 & 10 & \infty & \infty & 0 \end{pmatrix}$$

[6 points]

2. We are given a weighted graph  $G = (V, E)$  with weights given by  $W$ . The nodes represent cities and the edges are (positive) distances between the cities. We want to get a car that can travel between any two cities: it can leave with a full tank of gas, but cannot purchase gas until it reached its destination. The *zero-stop-capacity* of the car we purchase is defined to be the longest distance between any two cities that we may travel (but of course, between any two cities we would take the shortest path). A simple way to compute this is to run Floyd-Warshall ( $O(n^3)$ ), and then look at all pairs of vertices/cities and return the greatest distance ( $O(n^2)$ ).

A bit more formally, Floyd-Warshall returns  $W^*$ , where  $W^*[u, v]$  is the length of the shortest route from city  $u$  to  $v$ . The *zero-stop-capacity* is thus

$$\max\{ W^*[u, v] \mid u, v \in V \}.$$

Now suppose that you want to calculate the necessary *one-stop-capacity* of a potential car: we are allowed to stop at just **one** intermediate city to purchase gas for the car. Give an algorithm to determine what capacity would be needed for our car to travel between any two cities with at most one refueling stop. [6 points]

The first two steps of the development of a dynamic programming algorithm for a problem are

**step 1** describe the structure of the subproblem

**step 2** find a recurrence for the optimal value of the subproblem in terms of smaller subproblems

Perform just these two steps for the two problems listed below. Do not write (pseudo) code - just the subproblem and recurrence structure.

3. (exercise 6.2 from the DPV text) You are going on a long trip. You start on the road at mile post 0. Along the way there are  $n$  hotels, at mile posts  $a_1 < a_2 < \dots < a_n$ , where each  $a_i$  is measured from the starting point. The only places you are allowed to stop are at these hotels, but you can choose which of the hotels you stop at. You must stop at the final hotel (at distance  $a_n$ ), which is your destination.

You'd ideally like to travel 200 miles a day, but this may not be possible (depending on the spacing of the hotels). If you travel  $x$  miles during a day, the penalty for that day is  $(200 - x)^2$ . You want to plan your trip so as to minimize the total penalty - that is, the sum, over all travel days, of the daily penalties. Perform the two steps above to start the process of determining the minimum possible penalty. **[8 points]**

4. There is a small company which in some months maintains a tourist store in Waldport (code W) and in others in Yachats (code Y), and moves back and forth between these two cities (they can only afford to have one store operating at a time). This company wants to have the cheapest possible location plan - the two cities have different operating costs and these costs can change from month to month.

We are given  $M$ , a fixed cost of moving between the two cities, and lists  $W = (w_1, \dots, w_n)$  and  $Y = (y_1, \dots, y_n)$ . Here  $w_i$  is the cost of operating out of Waldport in month  $i$ , and  $y_i$  is the cost of being in Yachats that month. Suppose that  $M = 10$ ,  $W = (1, 3, 20, 30)$ , and  $Y = (50, 20, 2, 4)$ . If the location plan is (W, W, Y, Y), its cost will be  $1 + 3 + 10 + 2 + 4 = 20$ . On the other hand, the cost of the plan (Y, Y, W, Y) is  $50 + 20 + 10 + 20 + 10 + 4 = 114$ . The goal here is to (start to) devise a dynamic programming algorithm which, given  $M$ ,  $W$ , and  $Y$ , determines the *cost* of the optimal plan. The plan can start in either city, and end in either city. Note that you will likely need two subproblems, which will be mutually recursive. **[8 points]**

**Total: 27 points**