### CS252 - Final Exam Study Guide

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### Lecture #01 – General Introduction

### **Reasons for Different Programming Languages**

- 1. Different domains (e.g. web, security, bioinformatics)
- 2. Legacy code and libraries
- 3. Personal preference

### **Programming Language Design Choices**

- 1. Flexibility
- 2. Type safety
- 3. Performance
- 4. Build Time
- 5. Concurrency

### **Features of Good Programming Languages**

- 4. Safety (e.g. security and can errors be
  - caught at compile time)
  - 5. Machine independence
  - 6. Efficiency

### Goals almost always conflict

### **Conflict: Type Systems**

- Advantage: Prevents bad programs.
- Disadvantage: Reduces programmer flexibility.

Blub Paradox: Why do I need advanced programming language techniques (e.g. monads, closures, type inference, etc.)? My language does not have it, and it works just fine.

### **Current Programming Language Issues**

1. Simplicity

2. Readability

3. Learnability

- Multi-core "explosion"
- Big Data
- Mobile Devices

### **Advantages of Web and Scripting Languages**

- Examples: Perl, Python, Ruby, PHP, JavaScript
- · Highly flexible
- Dynamic typing
- · Easy to get started
- Minimal typing (i.e. type systems)

### **Major Programming Language Research Contributions**

- Garbage collection
- · Sound type systems
- Concurrency tools
- Closures

### **Programs that Manipulate Other Programs**

- Compilers & interpreters
- JavaScript rewriting
- Instrumentation
- Program Analyzers
- IDFs

### **Formal Semantics**

- Used to share information unambiguously
- Can formally prove a language supports a given property
- Crisply define how a language works

### **Types of Formal Semantics**

- Operational
  - Big Step "natural"
  - o Small Step "structural"
- Axiomatic
- Denotational

### Haskell

- Purely functional Define "what stuff is"
- No side effects
- Referential transparency A function with the same input parameters will always have the same result.
  - o A function call can be replaced with its value and nothing will change.
- Supports type inference.

**Duck Typing** – Suitability of an object for some function is determined not by its type but by presence of certain methods and properties.

- o More flexible but less safe.
- Supported by Haskell
- o Common in scripting languages (e.g. Python, Ruby)

### Side Effects in Haskell

- Generally not supported.
- Example of Support Side Effects: File IO
- Functions that do have side effects must be separated from other functions.

### **Lazy Evaluation**

- · Results are not calculated until they are needed
- Allows for the representation of infinite data structures

### Lecture #02 - Introduction to Haskell

### **Key Traits of Haskell**

- 1. Purely functional
- 2. Lazy evaluation
- 3. Statically typed
- 4. Type Inference
- 5. Fully curried functions

### ghci - Interactive Haskell.

let - Keyword required in ghci to set a variable value. Example:

> let f x = x + 1

> f 3 4

**Run Haskell from Command Line** Use runhaskell keyword.

Lists

> runhaskell <FileName>.hs

### **Hello World in Haskell**

main :: IO () main = do

> [1..4]

[1, 2, 3, 4]

> [1,2..6]

> [1,3..10]

[1, 3, 5, 7, 9]

putStrLn "Hello World"

Ranges • Can be infinite or bounded

• Use the "..." notation. Examples:

### **Primitive Classes in Haskell**

**Hello World in Haskell** 

putStrLn "Hello World"

- 1. Int Bounded Integers
- 2. Integer Unbounded
- 3.Float
- 4.Double
- 5.Bool 6.Char

main :: IO ()

main = do

- Base 0
- Comma separated in square brackets
- Operators
  - o: Prepend
  - ++ Concatenate
  - o!! Get element a specific index
  - o head First element in list
  - o tail All elements after head
- o last Last element in the list o init All elements in the list except

the last one

- o take n Take first n elements from a
- o replicate 1 m Create a list of length 1 containing only m
- o repeat m Create an infinite list

### containing only m

### > [5, 4..1] [5, 4, 3, 2, 1]

[2, 4, 6, 8, 10]

**Infinite List Example** > let even = [2,4..]> take 5 even

[1, 2, 3, 4, 5, 6]

### **List Examples**

> putStrLn \$ "Hello " ++ "World" "Hello World"

"abracadabra"

> let s = bra in s !! 2 : s ++ 'c' : last s : 'd' : s

1

```
List Comprehension
                                                                        A Simple Function
• Based off set notation.
                                                              > let inc x = x + 1
                                                              > inc 3
• Supports filtering as shown in second example
                                                                                                                  Pattern Matching
• If multiple variables (e.g. a, b, c) are specified, iterates through
                                                                                                    • Used to handle different input data
 them like nested for loops.
                                                              > inc 4.5
                                                                                                    • Guard uses the pipe ( ) operator
• Uses the pipe (|) operator. Examples:
                                                              5.5
                                                                                                    • Example:
> [ 2*x | x <- [1..5]]
                                                              > inc (-5) -- Negative
                                                                                                    inc :: Int -> Int
[2, 4, 6, 8, 10]
                                                                         Type Signature
                                                                                                      | x < 0 = error "invalid x"
> [(a, b, c) | a <- [1..10], b <-[1..10],
                                                              • Uses symbols ":: " and "->"
                                                                                                    inc x = x + 1
                  c \leftarrow [1..10], a^2 + b^2 = c^2]
                                                              • Example:
                                                              inc :: Int -> Int
 [(3, 4, 5), (4, 3, 5), (6, 8, 10), (8, 6, 10)]
                                                              inc x = x + 1
```

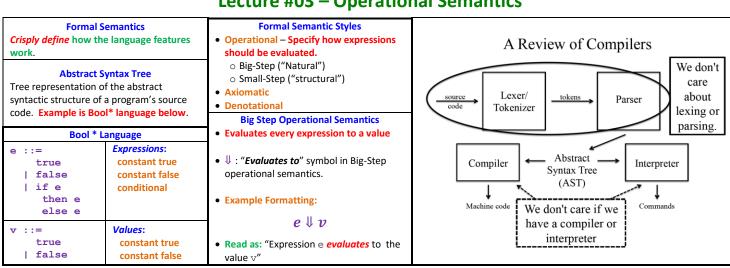
```
Recursion
• Base Case - Says when recursion should
                                                      Lab #01 – Max Number
                                                                                                       Reasons for a Large Number of
• Recursive Step - Calls the function with a
                                      > maxNum :: [Int] -> Int
                                                                                                          Programming Languages
 smaller version of the problem
                                      > maxNum [] = error "Invalid Input"

    Different domains

                                      > maxNum [x] = x
                                                                                                   • Different design choices
Example:
                                      > maxNum (x:xs) = if x > maxXs then x else maxXs
addNum :: [Int] -> Int
                                           where maxXs = maxNum xs
addNum [1] = 0
addNum (x:xs) = x + addNum xs
```

```
Recursion
                                                                  Haskell's Base Typeclasses
• :t or :type - Gets the type of a variable or function.
                                                          • Ord - Can be ordered
                                                          • Eq - Can perform equality check
Example:
                                                          • Show - Can convert to String
> :type 'A'
                                                          • Read - Can convert from String
'A' :: Char
                                                          • Enum - Sequentially Ordered
> :t "Hello"
                                                          • Bounded – Has upper and lower bound.
"Hello" :: [Char]
```

### **Lecture #03 – Operational Semantics**



Small-Step Operational Semantics	Bool* Small-Step Operational Semantics Rules		
Evaluate an expression until it is in normal form	E-IfTrue:	Example: Reduce the expression	
	<u></u>	if (if true then false else true) then true else false	
Normal Form – Any form that cannot be	<b>if</b> true <b>then</b> $e_2$ <b>else</b> $e_3 \rightarrow e_2$		
evaluated further.	E-IfFalse:	Step #1: Use rule "E-IfTrue" with "E-If"	
<ul> <li>→ : "Evaluates to" symbol in small step operational semantics. Example:</li> </ul>	if false then $e_2$ else $e_3 \rightarrow e_3$	<b>if</b> false <b>then</b> true <b>else</b> false	
e  ightarrow e''  ightarrow v	E-If:	Step #2: Use rule "E-IfFalse" (Now in normal form)	
$ullet$ $ o^*$ : Many evaluation steps required. Example: $oldsymbol{e}  o^* \ oldsymbol{v}$	$\frac{e_1 \to e_1'}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3 \to \text{if } e_1' \text{ then } e_2 \text{ else } e_3}$	false	

# • 0 : The Number "0"

**Bool\* Extension: Numbers** true false | if e then e else e • succ 0: Represents "1" 0 • succ succ 0: Represents "2" | succ e pred e • pred n: Gets the predecessor v ::= true | false

### **Literate Haskell**

• File Extension: ".lhs"

• Code lines begin with ">"

· All other lines are comments.

• "Essentially swaps code with comments."

```
Case Statement in Haskell
```

• Keywords: case, of, otherwise

Operator: ->

Example:

case x of val1 -> "Value 1" val2 -> "Value 2" otherwise -> "Everything else."

### Lab #02 Review

```
Bool Expression Type
                                       > data BoolVal = BVTrue
 data BoolExp = BTrue
        BFalse
>
        | Bif BoolExp BoolExp
                                            deriving Show
        | B0
        | Bsucc BoolExp
                                       > data BVInt = BV0
>
        | Bpred BoolExp
    deriving Show
                                       >
                                            deriving Show
```

```
BoolVal Type
  | BVFalse
  | BVNum BVInt
 | BVSucc BVInt
```

Type Constructors: BoolExp, BoolVal, BVInt

Non-nullary Value Constructors: Blf, Bsucc, Bpred, BVSucc, BVNum

Note: Even constants like BO, BTrue, BFalse, BVTrue, and BVFalse are nullary value constructors (since they take no arguments)

### **Lecture #04 – Higher Order Functions**

### Lambda

- Analogous to anonymous classes in Java.
- Based off Lambda calculus
- Example:

```
> (\x -> x + 1) 1
>(\x y -> x + y) 2 3
```

### **Function Composition**

Extended Bool \* Language

| IntV

IntV ::= 0 | succ IntV

- Uses the period (.)
- f(g(x)) can be rewritten (f . g) x

### **Point-Free Style**

- Pass no arguments to a function
- Example:

```
> let inc = (+1) - No args
> inc 3
```

### **Example: Lambda with Function** Composition $> let f = (\x -> x - 5)$ . $(\y -> y * 2)$ > f 7 9 > let f = ( $x y \rightarrow x - y$ ) $(\z -> z * (-1))$

#### Iterative vs. Recursive

- Iterative tends to be more efficient than recursive.
- Compiler can optimize tail recursive function.

Tail Recursive Function - The recursive call is the last step performed before returning a value.

### **Not Tail Recursive**

```
public int factorial(int n) {
  if (n==1) return 1;
  else {
    return n * factorial(n-1);
```

Last step is the multiplication so not tail recursive.

### **Tail Recursive Factorial**

```
public int factorialAcc(int n, int acc)
  if (n==1) return acc;
 else {
    return factorialAcc(n-1, n*acc);
}
```

Tail recursive code often uses the accumulator pattern like above.

### **Tail Recursion in Haskell** fact' :: Int -> Int -> Int fact' 0 acc = acc fact' n acc = fact' (n - 1) (n \* acc)

### **Higher Order Functions**

### **Functions in Functional Programming**

- Functional languages treat programs as mathematical functions.
- Mathematical Definition of a Function: A function f is a rule that associates to each x from some set X of values a unique y from a set of Y values.

### $(x \in X \land y \in Y) \rightarrow y = f(x)$

- f Name of the function
- X Independent variable
- y Dependent variable
- X Domain
- **Y** Range

### **Qualities of Functional Programming**

- Functions clearly distinguish:
  - Incoming values (parameters)
  - Outgoing Values (results)
- No (re)assignment
- No loops
- · Return values depend only on input parameters
- Functions are first class values; this means they can:
  - Passed as arguments to a function
  - o Be returned from a function
  - o Construct new functions dynamically

### **Higher Order Function**

Any function that takes a function as a parameter or returns a function as a result.

### **Function Currying**

Transform a function with multiple arguments into multiple functions that each take exactly one argument.

Named after Haskell Brooks Curry.

### **Currying Example**

addNums :: Num a => a -> a -> a

addNums is a function that takes in a number and returns a function that takes in another number.

### map

- Built in Haskell higher order function
- . Applies a function to all elements of a list.

#### filter

- Built in Haskell higher order function
- · Removes all elements from a list that do not satisfy (i.e. make true) some predicate.

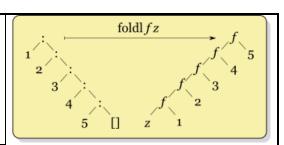
- Built in higher order function
- Does not support infinite lists.
- · Should only be used for special cases.

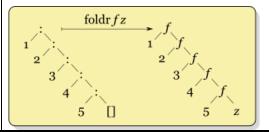
### Example:

> foldl (
$$x y -> x - y$$
) 0 [1, 2, 3, 4] -10 -- (((0-1) - 2) - 3) - 4

- Built in higher order function
- Supports infinite lists.
- "Usually the right fold to use"

#### Example:





Thunk – A delayed computation

Due to lazy evaluation, foldl and foldr build thunks rather than calculate the results as they go.

foldl'

> foldr (x y -> x + y) 0 [1, 2, 3, 4]

- Data.List.foldl' evaluates its results eagerly (i.e. does not use thunks)
- Good for large, but finite lists.

foldl in terms of foldr

myFoldl' f acc x = foldr (flip f) acc (reverse x)

### **Lecture #05 – Small-Step Operational Semantics**

### **WHILE Language**

• Unlike the Bool\* language, WHILE supports mutable references.

e ::= a	Variable/addresses
l v	Values
a:=e	Assignment
e;e	Sequence
e op e	Binary Operations
if e then e	Conditional
else e	
while (e) e	While Loops
v ::= i	Integers
b	Boolean
op ::= +   -   *	/

### **Small Step Semantics with State**

• Since the WHILE language supports mutable references, the grammar must be updated to support it.

#### While Relation:

$$e, \sigma \rightarrow e', \sigma'$$

• σ – Store. Maps references to values.

### **Example Operations:**

- $\sigma(a)$  Retrieves the value at address "a"
- $\sigma[a \coloneqq v]$  Identical to the original store with the exception that it now stores the value  $\boldsymbol{v}$  at address " $\boldsymbol{a}$ "

### **Evaluation Order Rules**

- Tend to be repetitive and clutter the semantics.
- Context based rules tend to represent the same information as evaluation order rules but more concisely.

### **Reduction Rule**

Rewrites the expression. Example:

### F-IfFalse:

if false then e2 else e3  $\rightarrow$  e3

### **Context Rule**

Specify the order for evaluating expressions. Example:

if  $e_1$  then  $e_2$  else  $e_3 \rightarrow$  if  $e'_1$  then  $e_2$  else  $e_3$ 

Reducible Expression (Redex) - Any expression that can be transformed (reduced) in one step.

### **Example: Redex**

if true then (if true then false else false) else true

This reduces to "if true then false else false"

### **Example: Not a Redex**

if (if true then false else false) then true else true

Not a redex as expression "if true then false else false" must be evaluated first.

### **Evaluation Contexts**

- Alternative to evaluation order rules.
- Marker (•) / hole indicate the next place for evaluation (i.e. where we will do the work).

### Example:

C[r]

= if (if true then false else false) then true else true

r = if true then false else false

**C** = **if** • **then** true else true

C[r] is the original expression.

### **Rewriting Evaluation Order Rules** Context based rules only apply to reducible

expressions (redexs). Example:

 $C[if false then e_2 else e_3] \rightarrow C[e_3]$ 

### **Context Syntax**

| if C then e else e | C op e | v op C

### Data.Map

- Library: import Data.Map as Map
- Immutable
- Example Methods:
  - o Map.empty Creates and returns an empty map
  - o Map.insert k v m Inserts a value "v" at key "k" into map "m". Returns a new, updated map.
  - $\circ$  Map.lookup k m Returns the value at key "k" in map "m". Wrapped in a Maybe.
  - Map.member k m Returns true if k is in map "m" and false otherwise.

Precondition - Text above the line in a rule.

**Context Rule for Binary Op:** 

 $v_3 = v_1 \text{ op } v_2$  $C[v_1 \text{ op } v_2] \rightarrow C[v_3]$ 

How to Read a Small Step Semantic Rule: "Given < Precondition >, then <LeftSideArrow> evaluates to <RightSideArrow>."

### Lecture #06 - LaTeX

#### TeX

- Created by Donald Knuth
- Domain specific language for typesetting documents.
- · Precisely controls the interface of content.
- Type of Literate Programming - Logic is in natural language and code is interspersed.
- "Mark code instead of marking comments."

#### LaTeX

- Developed by Leslie Lamport. Derives from TeX.
- Type of Domain Specific Language (DSL) A computer language that is specialized for a particular application domain.
- Enforces separation of concerns Design principle for separating a computer program into different sections, such that each section addresses a separate
  - o Example: LaTeX separates formatting from content.
- Literate Programming

### **Specify Document Type** \documentclass{article}

**Specify Title Block Content** \title{Hello World!}

> **Start Document** \begin{document}

**Generate Title from Title Information** \title{Hello World!}

> **Close the Document** \end{document}

```
Cross-Reference
\ref{<referenceName>}
```

Reference a Bibliography Citation \cite{<citationName>}

**Create a Reference** \label{<referenceName>}

Create a Bibliography \bibliography{<bibFileName>}

> Create a List \begin{itemize} \item Text for #1 \item Text for #2 \end{itemize}

### **Create Section with Label** \section{Section #1}

\label{sec:one}

### **Create Subsection with Label**

\subsection{<SubsectionName>} \label{sec:<refName>}

### Use of Tilde (~)

Creates an undividable space so the text "Section~\ref{sec:one}" will appear on one line

#### **BibTeX**

- · References are tedious to reformat and renumber.
- Reference details shorted in a "\*.bib" file.

Create a Bibliography \bibliography{biblio}

BibTeX filename for the example would be "biblio.bib"

**Define Bibliography Style** \bibliographystyle{plainurl}

# **BibTeX Article Reference Example**

```
@article{citationName,
   author = {Donald Knuth},
   title = {Literate Programming},
   journal = {},
   year = \{1984\},
   volume = {27},
   number = \{2\},
   pages = {97-111},
```

### **Lecture #07 – Types and Typeclasses**

#### Maybe Type

### · Example of an algebraic data type

- Enables behavior similar to null in Java
- Can be used to provide context.
- Used when:
  - o A function may not return a value
- o A caller may not pass an argument
- Definition:

```
data Maybe a = Nothing
             I Just a
```

Algebraic Data Type

A composite data

type (i.e. a type

Created via the

Keyword: data

made from other

### Maybe "Divide" Example

```
divide :: Int -> Int -> Maybe Int
divide _ 0 = Nothing
divide x y = Just $ x 'div' y
> divide 5 2
2
> divide 4 0
Nothing
```

DO NOT FORGET THE Just IN CORRECT SOLUTION

### Maybe Map Example

```
import Data.Map
m = Map.empty
m' = Map.insert "a" 42 m
case (Map.lookup "a") of
  Nothing -> error "Element not in map"
  Just x -> putStrLn $ show x
```

Since element may not be in the map, you need to use a maybe

### **Example Algebraic Data Type**

data Tree k = EmptyTree | Node (Tree k) (Tree k) val deriving (Show)

k - Type parameter. Specifies a type not a value.

Node: Value Constructor that creates values of type "Tree k"

# • Tree and Tree Int have no types since they themselves form a concrete

• Node does have a type:

> :t 3

```
> :t Node
Node :: (Tree k) \rightarrow (Tree k) \rightarrow k \rightarrow (Tree k)
```

Explanation: To make a complete Node object, you pass it two objects of type "Tree k" and another object of type "k" and that returns a "Tree k" object.

### **Partially Applying a Value Constructor**

- Value constructors can be partially applied similar to functions.
- > let leaf = Node EmptyTree EmptyTree
- > Node (leaf 3) (leaf 7) 5

This creates a three node tree with value 5 at the root and values 3 and 7 at the leaves.

Type of the "+" Operator

```
> :t (+)
(+) :: (Num \ a) => a -> a -> a
```

Explanation: The plus sign takes two numbers of type "a" and returns an object of type "a".

```
Type of a Number
3 :: (Num \ a) => a
```

**Explanation:** Since "3" has no explicit type, it can for now be any type that satisfies the "Num" type class.

### • Examples:

types).

- o Either
- o Maybe
- o Tree

```
Typeclasses
                            Kinds
                                                                                                            Example: Make Maybe an Instance of Eq

    Similar to interfaces in Java.

                                                                                                            instance (Eq a) => Eq (Maybe a) of

    Like a contract.

                                                                                                                   (==) Nothing Nothing = true
                                                                   o Implementation details can be included
                                        String Kind
                                                                                                                   (==) (Just x) (Just y) = x == y
                                                                     in typeclass definition.
                             > :kind String
                                                                                                                                                = false
• "The type of types".
                             String:: *
                                                                • No relation to classes in object-oriented
                                                                                                            Need to ensure type "a" supports "Eq" so add that as
• Concrete types have a kind
                                                                  programming.
                                        Map Kind
                                                                                                            a class constraint.
                             > :k Map
                                                                   o Example: Do not have any data
                             Map :: * -> * -> *
                                                                     associated with them.
• Keyword :k, :kind
                                                                                                            Class Constraint
                                       Maybe Kind
• Example:

    Simplify polymorphism.

                                                                                                            Operator: =>
                             > :k Maybe
                                                                                                            • Ensures that a type parameter satisfies some
                             Map:: * -> *
> :k Tree
                                                                Example: Eq Typeclass
                                                                                                              typeclass requirement.
Tree :: * -> *
                                     Map String Kind
                                                                class Eq a where
                                                                                                                           Kind of Typeclasses
                             > :kind (Map String)
Explanation: A Tree requires
                                                                     (==) :: a -> a -> Bool
                             (Map String) :: * -> *
one type parameter (e.g.
                                                                     (/=) :: a -> a -> Bool
                                                                                                            > :k Eq
Int) to be made a concrete
                                                                     x == y = not (x /= y)
                                                                                                            Eq :: * -> Constraint
                             Explanation: Map String is has one
                                                                     x \neq y = not (x == y)
type.
                             of the two type parameters filled so
                                                                                                            > :k Num
                             it has one less asterisk.
                                                                The last two lines in the type class definition
                                                                                                            Num :: * -> Constraint
                                                                allow the developer to program either (==) or
                                                                (/=) but not necessarily both.
                                                                                                            Note: Typeclasses are a class constaint (not a type)
                                                                                                            so their kind is different.
```

### **Lecture #08 – Functors**

```
Functor – Something that can be mapped over.
                                                                                          Examples: map and fmap on Lists
        Functor Type Class Definition
                                             • Handles things "inside a box"
                                                                                                                           Examples: fmap on Maybes
                                                                                          > map (+1) [1, 2, 3]
class Functor f where
                                              Example: List ([]) as an instance of Functor
                                                                                          [2, 3, 4]
  fmap :: (a \rightarrow b) \rightarrow f a \rightarrow f b
                                                                                                                           > fmap (+1) (Just 3)
                                                                                                                           Just 4
                                                                                          > fmap (+1) [1, 2, 3]
                                             instance Functor [] where
This is very similar to the definition of the
                                                                                          [2, 3, 4]
                                                fmap = map
                                                                                                                           > fmap (+1) Nothing
higher order function "map"
                                                                                                                           Nothing
                                                                                          > fmap (+1) []
                                             Explanation: map is a specialized version of
map :: (a -> b) -> [a] -> [b]
                                                                                          []
                                             fmap for lists.
```

```
Example: Either as an Instance of Functor
                                                     Either Algebraic Data Type
Example: Maybe as an Instance of Functor
                                                                                       instance Functor (Either a) where
                                           data Either a b = Left a
                                                                                           fmap _ (Left x) = Left x
fmap f (Right y) = Right (f y)
                                                               | Right b
instance Functor Maybe where
                                                    deriving (Eq,Ord,Read,Show)
   fmap _ Nothing = Nothing
   fmap f (Just x) = Just (f x)
                                                                                       > fmap (+1) Leftt 20
                                           • Left - Error type that is not mappable.
                                                                                       20 -- No Change
DO NOT FORGET THE Just IN VALID SOLUTION
                                                                                       > fmap (+1) Right 20

    Right - Expected type

                                                                                       21 -- Changed
```

### IO in Haskell

<ul> <li>Haskell avoids side effects but they are</li> </ul>	Type Signature of the main Function in	<ul> <li>do – Allows for the chaining of multiple</li> </ul>	
inevitable in real programs.	Haskell	IO/Monad commands together. Syntactic	
	main :: IO ()	sugar for bind ">>="	
• Monads			
<ul> <li>Related to Functors</li> </ul>	Hello World in Haskell	<ul> <li>&lt;- Extracts data out of an IO/Monad</li> </ul>	
<ul> <li>Compartmentalize side effects.</li> </ul>	main = putStrLn "Hello World"	"Box"	
• ()	Type Signature of getLine	<ul> <li>return – Places data into an IO/Monad</li> </ul>	
<ul> <li>Unit type in Haskell</li> </ul>	getLine :: IO String	"Box"	

```
do Example
main = do
    line <- getLine</pre>
    if null line -- Checks for empty str
       then return ()
       else putStrLn $ reverseWords line
reverseWords :: String -> String
reverseWords = unwords .
               map reverse . words
```

```
return in Haskell
• Unrelated to "return" in other
  languages
```

• Better described as "wrap" or "box"

### Summary:

return - Boxes an IO (since IO is a monad)

Unboxes an IO

```
Type of the Unit Type ()

    Base type

> :t ()
() :: ()
                     Type of return
> :t (return ())
(return ()) :: Monad m => m ()
```

### Using IO as a Functor

```
main = do
       line <- fmap (++"!!!") getLine
       putStrLn line
```

**Explanation:** This function takes a string input from standard in and appends "!!!" at which point it prints it to the console.

#### Definition of IO as a Functor

```
instance Functor IO where
 fmap f action = do
                result <- action
                 return (f result)
```

**Explanation:** The action object is taken out of the IO box, the function "f" applied to it, and then returned to the IO box.

#### id Function

• Takes one input parameter and returns that input parameter unmodified. Examples:

```
> id 3
```

Monad is a typeclass.

> id "Hello World" "Hello World"

### **Functor Laws**

Functor Law #1: If we map the id function over a Functor, the Functor that we get back should be the same as the original Functor.

```
Examples:
> fmap id (Just 3)
Just 3
> fmap id Nothing
```

Nothing > fmap id [1, 2, 3] [1, 2, 3]

Functor Law #2: Composing two functions and then mapping the resulting (composed) function over a Functor should be the same as first mapping one function over the Functor and then mapping the other one.

```
Law #2 Written Formally
fmap (f . g) = fmap f . fmap g
```

The Functor laws are NOT enforced. They are good practice that makes the code easier to reason about.

### **Lecture #09 – Applicative Functors**

Functor - Something that can be mapped over. Allow you to map functions over different data types. Examples:

- Maybe
- Either
- IO
- Lists
- <\*>

Functors return boxed up values.

### **Functor Example**

```
> fmap (+1) [1, 2, 3]
[2, 3, 4]
> let x = fmap (+) [1, 2, 3]
```

Explanation: In this case x is: [(1+), (2+), (3+)]

### **Applicative Functor**

• Requires the importing of a special library as shown below:

import Control.Applicative

Functions in Applicative Typeclass:

- pure Wraps/boxes a value
- <\*> Infix version of fmap. Is itself a Functor.

Example Uses of pure > pure 7

> pure 7 :: Maybe Int Just 7

> pure 7 :: [Int]

### Type Class Definition of Applicative

```
class (Functor f) => Applicative f where
     pure :: a -> f a
     <*> :: f (a -> b) -> f a -> f b
```

Only difference between <\*> and fmap is that the function in <\*> is boxed while it is not in fmap (see the green f).

### Make Maybe an Instance of Applicative

```
instance Applicative Maybe where
     pure = Just
     Nothing <*> = Nothing
(Just f) <*> x = fmap f x
```

Explanation: pure simply wraps the value in  ${\tt Just}$ . No need to explicitly check if " ${\tt x}$ " is maybe as **fmap** will do that for you.

### Examples of Applicative Maybe

```
> Just (+3) <*> Just 4
Just 7
> pure (+3) <*> Just 4
Just 7
> pure (+) <*> Just 3 <*> Just 4
Just 7
> (+) <$> Just 3 <*> Just 4
Just 7
```

Explanation: x <\$> is fmap as an infix operator. It is NOT necessarily the same as pure x <\*>. It should be based off Applicative Functor Law #1.

### Making [] an Instance of Applicative

```
instance Applicative [] where
  pure x = [x]
  fs <*> xs = [f x | f <- fs, x <- xs]
```

**Explanation:** The function is actually a list of functions so list comprehension is needed.

```
Example Use of Applicative on Lists
> (*) <$> [1, 2, 3] <*> [1,0,0,1]
[1,0,0,1,2,0,0,2,3,0,0,3]
```

```
> pure 7
7 -- No change
> pure 7 :: [Int]
[71
```

### Definition of IO as an Instance of Applicative

```
instance Applicative IO where
   pure = return
    a <*> b = do
             f <- a
              x <- b
              return (f x)
```

Example of Applicative IO	11000
import Control.Applicative	A function that simplifies the application of a normal function to two Functors.
<pre>main = do      a &lt;- (++) &lt;\$&gt; getLine &lt;*&gt; getLine      putStrLn a</pre>	<pre>liftA2 :: (Applicative f) =&gt; (a -&gt; b -&gt; c) -&gt; f a -&gt; f b -&gt; fc liftA2 f x y = f &lt;\$&gt; a &lt;*&gt; b</pre>

Example of liftA2	Applicative Functor Definition
> (:) <\$> Just 3 <*> Just [4]	
Just [3, 4]	A functor you can apply to
> liftA2 (:) (Just 3) (Just [4])	The same apply to
Just [3, 4]	other Functors.

### **Applicative Functor Laws**

<pre>Law 1:     pure f &lt;*&gt; x = fmap f x</pre>	Law 2: pure id <*> v = v	Law 3: pure (.) <*> u <*> v <*> w = u <*> (v <*> w)
<pre>Law 4:   pure f &lt;*&gt; pure x = pure (f x)</pre>	Law 5: u <*> pure y = pure (\$y) <*> u	Similar to Functor Laws, these are not strictly enforced but are good practice to make it easier to reason about the code.

### Monoids

Monoid: An associative binary function and a value that acts as an identity with respect to that function.		Definition of Monoid Typeclass	
	•	class Monoid m where	
• x * 1 • lst ++ [] • x + 0	Examples Identity of Multiplication Identity of Concatenation Identity of Addition	<pre>mempty :: m mappend :: m -&gt; m -&gt; m mconcat :: [m] -&gt; m mconcat = foldr mappend mempty</pre>	

### **Monoid Rules**

Rule #1:	Rule #2:	Rule #3:
mempty `mappend` x = x	x `mappend` mempty = x	<pre>(x `mappend` y) `mappend` z = x `mappend` (y `mappend` z)</pre>

### Lecture #10 - Monads

Problem with Functors: Do not support chaining of	Applicative Functor: A Functor that can be applied to other
multiple commands. Example:	Functors.
> fmap (+) (Just 3) (Just 4)	<pre>class (Functor f) =&gt; Applicative f where</pre>
	(<*>) :: f (a -> b) -> f a -> f b
Returns an error since it cannot resolve (Just 3+)	
and (Just 4)	Requires library Control. Applicative
	multiple commands. Example:  > fmap (+) (Just 3) (Just 4)  Returns an error since it cannot resolve (Just 3+)

```
Comparing <*> and >>=
                                                                       Example of <$>, <*> and >>=
Functor:
                                                              > (\x -> x + 1) < > Just 3
                                                                                                            Example: Implement applyMaybe that applies a
(<*>) :: Applicative f => f (a -> b) -> f a -> f b
                                                              Just 4
                                                                                                            function to a Maybe
Monad:
                                                              > Just (x -> x + 1) <*> Just(3)
(>>=) :: Monad m => m a -> (a \rightarrow m b) -> m b
                                                                                                            applyMaybe :: Maybe a -> (a -> b) -
                                                               Just 4
                                                                                                            > (Maybe b)
                                                                                                            applyMaybe Nothing _ = Nothing
applyMaybe (Just x) f = Just (f x)
Differences:
1. Order of the arguments changed.
                                                               > (Just 3) >>= (\x -> Just(x+1))
2. The function is boxed in Functor but not Monad
3. Monad function returns a boxed result.
                                                               Just 4
```

```
> (Just 3) `applyMaybe` (\_ -> Nothing)
                                                                                             • "Applicative Functors you can chain."
applyMaybe Nothing _ = Nothing
                                                         `applyMaybe` (\y -> Just (y-1))
applyMaybe (Just x)^{-}f = Just (f x)
                                             Nothing
          Monad Typeclass Definition
                                                                   Example a Robot Moving Towards a Goal (Not Failure)
                                                                                -- Define Operator and start location
                                                                               x -: f = f x
class Monad m where
                                             --Location
      return :: a -> m a
                                             type Robot = (Int, Int)
                                                                               start = (0, 0)
      (>>=) :: m a -> (a -> m b) -> m b
                                             -- Functions
                                                                               > start -: up -: right
      (>>) :: m a -> m b -> m b
                                             up (x,y) = (x, y+1)
                                                                               (1, 1)
      x \gg y = x \gg (\ -> y) --Lamda
                                             down (x,y) = (x, y-1)
                                             left (x,y) = (x-1, y)
                                                                               > start -: up -: left -: left -: right -: down
      fail :: String -> m a
                                             right (x,y) = (x+1, y)
      fail msg = error msg
```

Chaining applyMaybe

**Additional Names for Monoids** 

• "Programmable Semicolons"

Just 5

Example: Implement applyMaybe that applies a

applyMaybe :: Maybe a -> (a -> Maybe b)

-> (Maybe b)

function to a Maybe

```
Example a Robot Moving Towards a Goal (with Failure)
                                      -- Once the goal is reached,
                                      -- the robot stops
                                      goal := Map.empty
                                                                                 start = (0, 0)
                                              -: (Map.insert (0, 2) True)
Maybe as an Instance of the Monad Typeclass
                                              -: (Map.insert (-1, 3) True)
                                              -: (Map.insert (-3, -8) True)
                                                                                 > return start >>= up >>= left >>= left
instance Monad Maybe where
                                                                                                >>= right >>= down
                                     moveTo :: Pos -> Maybe Pos
                                                                                 Just (-1, 0)
     return = Just
                                     moveTo p = if Map.member p goal
                                                                                 > return start >>= left >>= left >>= up
                                                       then Nothing
     (>>=) Nothing
                      = Nothing
                                                                                                else Just p
     (>>=) (Just x) \overline{f} = f x
                                                                                                >>= right >>= right >>= down
                                                                                 Nothing
                                      -- Since these are in bind, no need
     fail _
                      = Nothing
                                      -- to handle Nothing. Bind handles it.
                                     up(x,y) = moveTo(x, y+1)
                                                                                 Explanation: Reached one of the goals (-1, 3) at the red up
                                      down (x,y) = moveTo (x, y-1)
                                      left (x,y) = moveTo (x-1, y)
                                      right(x,y) = moveTo(x+1, y)
```

### **Integer Division Using Monads**

```
Integer Division with Bind with "do"
                                                                                                            Integer Division with Bind with "do" and return
       Integer Division with Bind and No "do"
                                                    mydiv :: Maybe Int -> Maybe Int -> Maybe Int
                                                                                                        mydiv :: Maybe Int -> Maybe Int -> Maybe Int
mydiv :: Maybe Int -> Maybe Int -> Maybe Int
                                                    mydiv x y = do
                                                                                                         mydiv x y = do
mydiv x y = x >>= ( numer ->
                                                                 numer <- x
                                                                                                                      numer <- x
             y >>= (\denom ->
                                                                 denom <- y
                                                                                                                      denom <- y
             if denom > 0
                                                                 if denom > 0
                                                                                                                      if denom > 0
                 then Just (div numer denom)
else fail "Div by zero"))
                                                                       then Just (div numer denom)
                                                                                                                         then return $ div numer denom
                                                                       else fail "Div by 0"
                                                                                                                         else fail "Div by 0"
```

### **List Monad**

```
Making List an Instance of Monad
                                                                Example Use of List as a Monad
instance Monad [] where
                                                        listOfTuples :: [(Int, Char)]
        return x = [x]
                                                                                                                Combining a Maybe and a List Monad
                                                        listOfTuples = do
         (>>=) xs f = concat(map f xs)
                                                                        n <- [1, 2]
        fail _
                     = []
                                                                                                         > Just [2,3] >>= (\x -> Just(fmap (+1) x))
                                                                        ch <- ['a', 'b']
                                                                                                         [3, 4]
                                                                        return (n, ch)
Explnation: concat is needed here as f returns elements
                                                        > listOfTuples
already in a list. As such, concat merges the individual lists
                                                        [(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]
(from each call to f) into a single list.
```

### **Lecture #11 – Parsing Combinators**

Semantics: Enumerate what a program means. Defined by the interpreter or compiler.  Compilation Flow Step #1: Tokenizer/lexer generates a set of tokens.		Converts the characters of the program into words of the language.	
	Step #2: Parser turns the tokens into an abstract syntax tree.	Examples:	
Syntax: Enumerate how a program Is structured. Defined by the lexer and parser.	Step #3: Compilers and interpreters convert the AST into machine code or commands respectively.	Lex/Flex (C/C++) ANTLR & JavaCC (Java) Parsec (Haskell)	

Lover

#### **Categories of Tokens**

- Reserved Words/Keywords.
  - o Examples: while, if, then, else
- Literals/Constants.
- o Examples: 123, "Hello World!"
- · Special symbols.
  - o Examples: ";", "=>", "&&"
- Identifiers.
  - o Examples: "balance", "myFunction"

#### **Parsing**

- · Parser converts tokens to abstract syntax trees.
- Defined by context free grammars (CFG)
- Types of Parsers:
  - o Bottom-up/Shift-Reduce Parsers
  - o Top-down parsers

### **Context Free Grammars**

- · Grammars specify the language.
- Specified in Backus-Naur form format. Example:

```
Expr -> Number
    Number + Expr
```

- Terminal Cannot be broken down further.
- Non-terminals Can be broken down further.

Example: "0", "1", "2", ..., "9" are terminals but digit, number, and expression are not.

#### **Example Grammar**

```
expr -> expr + expr
        expr - expr
        (expr)
        number
number -> number digit
        | digit
digit -> 0 | 1 | 2 | ... | 9
```

### Bottom-Up / Shift-Reduce Parser

- Shift tokens onto a stack
- Reduce the stack to a non-terminal.
- LR Left to right, Rightmost derivation
- LALR Look-Ahead LR parsers are the most popular type of LR parsers.
- o Examples: YACC/Bison
- · Fading from popularity

#### **Top-Down Parser**

- Non-terminals are expanded to match tokens.
- LL <u>Left</u> to right, <u>Leftmost derivation</u>
- LL(k) Parser Looks ahead up to k elements. **Examples:** Java CC, ANTLR
  - o The higher the k, the more difficult language is to parse. k can be arbitrary.
  - o LL(1) Easy to parse using either LL or recursive descent parsers. Many computer languages are designed to be LL(1).

### **Parser Combinator**

Combine simpler parsers to make a more complex parser.

Example: Parsec

### **Useful Parsec Functions**

- many Parses zero or more occurrences of the given parser.
- many1 Parses 1 or more occurrences of the given parser.
- noneOf Anything but the specified value
- spaces Whitespace characters
- **char** The specific specified character
- **string** The specific specified string.
- sepBy Separate tokens by some token.

```
import Text.ParserCombinators.Parsec
num :: GenParser st String
num = many1 digit
main = do
       print $ parse num "Hello" "42"
```

```
Example Parsec Code
import Text.ParserCombinators.Parsec
```

num :: GenParser st Integer str <- many1 digit return \$ read str

main = do

print \$ parse num "World" "42"

- st "State." Always required for our purposes.
- String/Integer Parser return type
- digit 0, 1, 2, 3, ..., 9 (terminal)
- num Parser entry function
- "Hello"/"World" Debug string.

```
Example with try, <|>, and <?>
```

```
eol = try (string "\n")
   <|> string "\n\r"
   <?> "end of line"
```

- try If an incomplete match is found, rewind.
- <|> "Or" Operator for matching tokens.
- <?> Otherwise with an accompanying error message.

• many1 - Select one of more digits.

• "42" - String to parse.

### **Practice Midterm and Review Notes**

Question #1	Question #2	Question #2 Question #3		Question #5
a. True	a. True	a. False – Big step	a. False – Imperative	a. True
b. False – Lazy evaluation	b. False – Applicative functor	b. True	b. True	b. False – Typeclass
c. False – Lazy evaluation	c. True	c. False – Use store	c. False	c. True
d. False – Statically type	d. True	d. True	d. True	d. False
e. True	e. True	e. False	e. True	e. False – Algebraic data type

d. False – Statically type	l. True	d. Irue	d. True	d. False
e. True	. True	e. False	e. True	e. False – Algebraic data type
Haskell  Purely Functional  Lazy evaluation  Fully Curried Language  Statically Typed  Type Inference – Via context, Haskell can deduce the type.	Purely Function  Referential Transparency call can be replaced with i value without affecting th  No (re)assignment  No loop  No side effects	<ul><li>A function</li><li>ts equivalent</li></ul>	Functional Languages  • Functions are first class objects meaning they can be passed to a function, returned from it, or created on the fly.  • Higher order function support	Operational Semantics  Small Step – Structural Semantics  Big Step – Natural Semantics  "Get stuck" – When a function is encountered that does not have an associated rule.

### **CSV Parser Example**

```
Verbose Approach
import Text.ParserCombinator.Parsec
import System.Environment
csvFile :: GenParser st [[String]]
csvFile = do
          arr <- many line
          char eof
          return arr
line :: GenParser st [String]
line = do
       result <- many1 cell
       char '\n'
       return result
cells :: GenParser st [String]
cells = do
        firstCell <- cellContents
nextCells <- remainingCells</pre>
        return (firstCell:nextCells)
cellContent :: GenParser st String
cellContent = many $ noneOf ", \n" -- Two characters
remainingCells :: GenParser st [String]
remainingCells = do
                  (char "," >> cells)
                  <|> return []
main = do
       args <- getArgs
       p <- parseFromFile csvFile "example 1" (head args)
       case p of
           Left msg -> error msg
           Right csv -> print csv
```

### Miscellaneous

### Kind of Show and show > :k Show Show :: \* -> Constraint Type and Kind of show > :k show Error (A function not a type) > :t show

```
Lambda and ADT Combined
> (\x -> Just (x+1)) 1
Just 2
        Creating Type Alias
```

type String = [Char]

Allows for more readable code as developer can use a type name that makes more sense for a given application.

```
Example: applyMaybe that takes a (Maybe a) and applies to it a function
            that takes a normal a and returns a (Maybe b)
```

```
applyMaybe :: (Maybe a) -> (a -> Maybe b) -> (Maybe b)
applyMaybe Nothing _ = Nothing
applyMaybe (Just x) f = f x
```

Explanation: Since the function "f" already returns a Maybe, you do not need to re-box it. However, since it does not take a Maybe, you need to unbox the first input parameter.

```
Applying return to Items
> return 7 :: Maybe Int
Just 7
> return 7 :: [Int]
```

show :: (Show a) => a -> String

List comprehension is syntactic sugar for using lists as monads.

Conclusion: Behavior for return is the same as pure. Both put the object in the minimum default context that still yields that value.

[7] -- Need Int or get an error

### Monads and Lambda

When trying to chain multiple functions together in a Monad, remember the Monad must return a boxed value. Hence, Lambda often work well as they simplifying boxing.

Applicative Typeclass - Allows you to use normal functions on values that have a context (i.e. are inside a Functor).

Monad: Given a value of type, a, in a context, m, apply a function that takes a normal value of type a and returns a value in the context m.

(>>=) :: (Monad m) => m a -> (a -> m b) -> m b

Monads are just applicative functors that support bind (>>=).

Key Difference: Applicative functors support normal functions that take and return unboxed values while Monads return boxed values.

return - Monad equivalent of "pure" for Applicative Functors.

Cannot use fmap in the definition of a Monad since fmap returns a boxed value while the function of the Monad returns a boxed value. Hence, if you used fmap with a Monad, you would return a double boxed value.

### **Functor Definitions**

```
TO
                                                          Maybe
                 Lists
                                                                                  instance Functor IO where
                                           instance Functor Maybe where
instance Functor [] where
                                                                                    fmap f a = do
                                                   Nothing = Nothing
                                             fmap
 fmap = map
                                                                                               x <- a
                                             fmap f (Just x) = Just (f x)
                                                                                               return (f x)
```

### **Applicative Functor Definitions**

Lists
<pre>instance Applicative [] where pure x = [x]</pre>
(<*>) fs xs = [ f x   f <- fs, x <- xs]

```
Maybe
instance Applicative Maybe where
 pure x = Just x
  (<*>) Nothing _ = Nothing
  (<*>) (Just f) x = fmap f x
```

```
instance Applicative IO where
  a <*> b = do
            f <- a
            x <- b
            return (f x)
```

### **Monad Definitions**

```
instance Monad Maybe where
                                                                                    instance Monad IO where
instance Monad [] where
                                              return x = Just x
                                                                                        (>>=) a f = do
return x = [x]
                                              (>>=) Nothing
                                                              = Nothing
                                                                                                   x <- a
 (>>=) xs f = concat $ map f x
                                              (>>=) (Just x) \overline{f} = f x
                                                                                                   fх
 fail _ = []
                                              fail
                                                        = Nothing
                                                                                       fail s = ioerror (userError s)
```

### Lecture #12 - Introduction to JavaScript

#### **JavaScript**

- Developed at Netscape by Brendan Eichs in 10 days
- · Originally named "Mocha"
- Syntax similar to Java

• Standard Approach:

### Multi-paradigm JavaScript Supported programming paradigms:

- Imperative
- Functional
- Object-Oriented (through prototypes)

#### Where JavaScript is Run

- Client Side Versions
- o Runs on user machine
- Server-side Versions o IVM: Rhino & Nashorn
- Node is

```
var = i, sum = 0;
for( i = 0; i < list.length ; i++){</pre>
  sum += list[i];
return sum;
```

var x = 42; // Create with var

y = 7; // No error without var

function noReturnAdd(a, b) {

var c = noReturnAdd(x, y)

function add(a, b){

return a + b;

//Lambda Function

a + b;

### **Example: Functional JavaScript**

```
function addList(list) {
 if(list.length == 0) {
   return 0;
 return list[0]
         + addList(list.slice(1));
}
```

slice(begin[, end]) - Takes a subset of an array from the "begin" index to the "end" (exclusive). If no "end" is specified, it takes all elements to the end of the list.

**Printing to the Console in JavaScript** 

console.log("...")

o Not supported by all implementations.

print

var print = console.log

• Solution to Support a Single Interface:

• JVM based JavaScript Approach:

### **Example: Object-Oriented JavaScript**

```
function Adder (amount) {
  this.amount = amount;
Adder.prototype.add = function(x){
             return this.amount + x;
var myAdder = new Adder(1)
var y = myAdder.add(7)
```

Adder - Name of a new constructor. Convention is to start constructors with a capital letter.

this - Not optional in JavaScript.

```
function addList(list) {
```

Example: Quirks of JavaScript

**Example: Imperative JavaScript** 

• Functions whose inner variables refer to independent (free) variables.

### Closure Example

```
function getNextInt(){
  var nextInt = 0;
  return function(){
             return nextInt++;
        }() // Double paren
             // run the function
console.log(getNextInt()); // print "0"
console.log(getNextInt()); // print "1"
console.log(getNextInt()); // print "2"
```

### var myLambda = function(x) {return x \* x;} Node.is

- JavaScript runtime environment and library designed to run outside the browser.
- Based off Google's V8 engine.
- npm Package manager to get new packages.

// c is "undefined" since no return

### **Callback Function**

- Functions in JavaScript are first class objects of type "Object".
- Not executed immediately.

JavaScript supports both "null" and "undefined"

### Reading from a File with Callbacks in Node.js

```
var fs = require('fs')
fs.readFile('myFile.txt',
    function(err, data) {
      if(err)
        throw err:
        console.log("" + data);
console.log("All done")
"All done" prints before the file contents due
to callbacks.
require - Includes the JavaScript package "fs"
```

```
Synchronous File IO in Node
var fs = require('fs')
var data =
fs.readFileSync('myFile.txt');
console.log("All done")
```

To eliminate callbacks, most function names can be appended with "Sync"

### **Undeclared Object Fields**

Any undeclared object fields or uninstantiated variables are undefined.

```
var y; // Uninstantiated
// Both print 'undefined'
console.log(y)
console.log(myDog.name)
```

### **Creating a JavaScript Object**

```
var myDog = {age : 3,
               weight: 100}
Every object is a map.
```

### Adding a Field to a JavaScript Object

```
myDog['height'] = 45 // Add a new height field
                     // Note the single quotes
```

Adding a Function to a JavaScript Object's Prototype

```
myDog.speak = function() { console.log("Grr"); }
```

Delete a Function from a JavaScript Object's Prototype

delete myDog.speak

### **Prototypes**

### **Object Prototypes**

```
JavaScript prototypes are just like any other object.
```

```
var dogPrototype = {
    speak: function(){
      console.log("bark!");
```

### **Defining an Object's Prototype**

```
var rex = { name: "Rex",
            proto : dogPrototype}
```

Prototypical Inheritance: If an object does not have a method of field, JavaScript looks to the object's \_\_proto\_\_ object.

```
Add a Special "speak" Method to Rex
                                                                                  Effect of the "new" Keyword
                                                          function Cat(name, breed) {
rex.speak = function() {
                                                            var this = {}; // Add when new is used
               console.log("Grr");
                                                            this.prototype = Cat.prototype; // Also comes from new
             };
rex.speak(); // Prints "Grr"
                                                            this.name = name;
                                                            this.breed = breed;
                                                            this.speak = function(){console.log("meow");};
delete rex.speak;
// Prints "Bark!" from __proto__
                                                            return this; // Also comes from new
rex.speak();
                                                                                  No "return" in a Function
delete rex.speak; // Does nothing
rex.speak(); // Prints "Bark!"
                                                          function noReturnAdd(x, y){
                                                            x + y; // without "return"
Unspecified Function Arguments: In JavaScript, any unspecified function
                                                          // c is "undefined" since no return
argument defaults to "undefined".
                                                          var c = noReturnAdd(x, y)
                                                          console.log(c); // Prints "undefined"
```

```
Top Prototypes
                                                                                 require
                                   Iterating Using "forEach"
                                                                                                              Running from the Command Line
                            var arr = [1, 2, 3];
                                                                    • Used to import an external module in
Object.prototype - Top
                                                                      Node.is
                             // Print each element in array
                                                                                                          • Use the keyword "node" for Node.js.
of all object prototypes
                            arr.forEach(function(val){
                                                                                                            Example:
                                            console.log(val);
                                                                     • Can be stored in a variable. Example:
Function.prototype -
                                                                                                                $ node my_program.js
                            }
Top of all function prototypes.
                                                                       var net = require('net');
```

```
Example: Currying in JavaScript
                                                     Function.prototype.curry = function(){
        Example: Create an Object with a Constructor
                                                       // Take slice from the Array class' prototype
                                                       var slice = Array.prototype.slice;
var Droid = {
                                                       // Convert arguments to an array
 speak: function() {
                                                       var args = slice.apply(arguments);
            console.log("I am "
                                                       var that = this;
                         + this.name);
                                                       return function(){
                                                         return that.apply(null,
 create: function(name) {
               var clone = Object.create(this);
clone.name = name;
                                                                             arg.concat(slice.apply(arguments));
                                                     };
               return clone;
                                                     function add(x, y){
        },
                                                       return x + y;
};
                                                     var addOne = add.curry(1);
                                                     console.log(addOne(3)); // Prints "4"
```

### Lecture #13 – Lambda Calculus

### **Expressions**

(Variables, immutable) e ::= x (Lambda abstraction)

(Function application)

Note: Lambda ( $\lambda$ ) is simply a function.

 $v ::= (\lambda x.e)$  (Lambda abstraction)

### **Function Application**

Given a function where **E** is a complex expression:

 $\lambda(x.E)v$ 

Then:

 $\lambda(x.E)v \to E[x \mapsto v]$ 

Hence, "x" replaces "v" in "E".

Lambda Calculus is a simple, Turing complete language. Hence it is equal in power to a Turing Machine.

Lambda calculus stops evaluating when the result is in **normal form**.

### Small-Step Evaluation Order Rules for Lambda Calculus

Rule: SS-E1

$$\frac{e_1 \rightarrow e_1'}{e_1 e_2 \rightarrow e_1' e_2}$$

 $oldsymbol{e_2}
ightarrowoldsymbol{e_2}'$  $(\lambda x.e) e_2 \rightarrow (\lambda x.e) e_2'$  **Rule: SS-Lambda Context** 

 $(\lambda x. e) v \rightarrow e[x \vdash > v]$ 

Optional Rule: Lazy SS-Lambda Context

 $(\lambda x. e) e_2 \rightarrow e[x \vdash > e_2]$ 

### **Evaluation Strategies**

Strict Evaluat	ion Strategies	Lazy Evaluation Strategies	
Call by Value: Pass a copy of a parameter	Call by Reference: Implicit reference (e.g., pointer) to the parameter is passed	Call By Name: Re-evaluate the argument each time it is used.	Call by Need: Memoizes parameter value after first use.

Language Equivalents of $(\lambda x. e)$		True and False in Lambda Calculus		
		True in Lambda Calculus:	True in Lambda Calculus:	
JavaScript:	Haskell:	$getFirstParam = tru = (\lambda x. \lambda y. x)$	$getSecondParam = fls = (\lambda x. \lambda y. y)$	
<pre>function(x) {return e;}</pre>	(\x -> e)	Note: This returns the <i>first</i> parameter in the pair of values.	Note: This returns the <i>second</i> parameter in the pair of values.	

### **Conditional in Lambda Calculus**

### $test = \lambda cond \cdot \lambda then \cdot \lambda els. (cond then els)$

Example #1:

test(tru tru fls)

 $\lambda$  cond.  $\lambda$  then.  $\lambda$  els. (cond then els)(tru tru fls)

*λthen* . *λels*.(tru then els)(tru fls)

 $\lambda$ els. (tru tru els)(fls)

(tru tru fls)

 $(\lambda x. \lambda y. x)(tru fls)$ 

 $(\lambda y. tru)(fls)$ 

tru

Example #2:

test(fls tru fls)

 $\lambda$  cond.  $\lambda$  then.  $\lambda$  els. (cond then els)(fls tru fls)

 $\lambda$ then.  $\lambda$ els. (fls then els)(tru fls)

 $\lambda$ els. (fls tru els)(fls)

(fls tru fls)

 $\lambda x. \lambda y. y (tru fls)$ 

 $\lambda y. y (fls)$ 

**Boolean And** 

and  $d = \lambda b \cdot \lambda c \cdot (b c f l s)$ 

Pair

 $pair = \lambda f. \lambda s. \lambda b. (b f s)$ 

Pair – A tuple-like data structure in Lambda Calculus.

### Working with a Pair in Lambda Calculus

First Element in a Pair

Second Element in a Pair

 $first = \lambda p. (p tru)$ 

 $second = \lambda p.(p fls)$ 

Note #1: In the case of both first and second, the term p must be a pair.

Note #2: Both of these rely on the tru or fls being substituted for the "b" in the pair data structure in term selecting either the first or second element.

### **Church Encoding Numerals**

 $zero = \lambda s. \lambda z. z$   $one = \lambda s. \lambda z. s z$  $two = \lambda s. \lambda z. s s z$ 

### **Successor Function**

 $scc = \lambda n. \lambda s. \lambda z. s(n z)$ 

### Example:

one' = scc zerotwo' = scc one' = scc(scc zero)

### Plus in Lambda Calculus

### Lecture #14 - JavaScript Scoping

### **Example: First Class Function**

```
function makeAdder(x) {
  return function(y) {
      return x + y;
      };
}
var addOne = makeAdder(1);
// Prints "11"
console.log(addOne(10));
```

#### **Example: Function Application**

JavaScript lacks block scope for the closure to be right, must create the function inside another function.

Block Scope – The scope (i.e. visibility) of a variable is limited to a specific block (e.g., for loop, if statement, etc.).

- Unlike most languages, JavaScript does not have block scope.
- To create a new scope, use an anonymous function.

Variable Hoisting – All variable declarations (i.e., use of "var") are treated as if they are at the beginning of the function.

#### "this" in JavaScript

this – Refers to the scope where the function is called.

- In Normal Function Calls this refers to the global "this"
- Object Methods The object itself.
- Constructor (using "new") The newly created object.
- Exceptions: apply, call, and bind. Inline event handles on DOM elements

Any time a new function is created, the other "this" is no longer in scope

### **Execution Context**

### Consists of three part:

- A Variable Object Container for variables and functions.
- Scope Chain Variable object plus parent scopes
- Context Object this

### **Global Context**

- Top Level Context
- Variable object is known as the "global object"
- this Refers to the global object.

Any variable declared without var is added to the global context.

### **Function Contexts**

- Activation or Variable Objects which include:
  - o Arguments passed to the function
  - o A special arguments object
  - o Local variables

```
apply, bind, call Example
```

```
x = 3;
function foo(y) {
    console.log(this.x + y);
}
foo(100); // Prints "103"
// Array passed for args
foo.apply(null, [100]);
// Update the context
foo.apply({x:4}, [100]);
// No array needed
foo.call({x:4}, 100);
// Create a new function
var bf = foo.bind({x:5});
bf(100);
```

- apply Calls a function with the arguments passed as an array.
- call Calls a function with the arguments passed in comma separated.
- **bind** Used to create a new function with a custom context.

### Lecture #14.5 - JSLint and TypeScript

### Issues in JavaScript

- No block scope
- Forgetting var can lead to unexpected behavior since variables become global.
- Operator "==" is not transitive.
- Switch/case statements require "break"

```
JavaScript Automatically Inserts Semicolons
```

```
function makeObject () {
  return // Semicolon inserted here
  {
    madeBy: 'Austin Tech. Sys.'
  }
}
var o = makeObject();
```

console.log(o.madeBy); // error

# Behavior of "typeOf"

typeOf 5 // "number"

typeOf — Returns a string. May yield unexpected results.

```
typeOf "hi" // "string"
typeOf NaN // "number"
```

typeOf null // "object"

### Behavior of "typeOfChar"

typeChar — Returns a string. Classifies letters as "digits".

```
typeOfChar "5" // "digit"
typeOfChar "q" // "digit"
// "Other character"
typeOfChar " "
```

### JSLint

- A tool to write cleaner and safer JavaScript.
- Requires that "use strict" (with quotes) be added at the beginning of all functions.
- · Performs static code analysis.
- Helps catch common programming errors by requiring:
- o Variables declared before they are used.
- o Semicolons are always used.
- o Double equals never used.
- Inspired by the "lint" tool

### **Benefits of Type Systems**

- Tips for compilers
- Hints for IDEs
- Enforced documentation
- Prevent code with errors from running.

### **TypeScript**

- Developed by Microsoft
- Static type checker for JavaScript.
   A new "superset" language
- of JavaScript with:

  O Type annotations
- O Type annotationsO Classes
- Compiles to JavaScript

### Function Type Annotations in TypeScript

```
function greet(person: string) {
  console.log("Hello " +
  person);
}
var user : string = "Vlad";
// Prints "Hello Vlad"
```

greet (user);

### Types in TypeScript

```
number (var pi : number = 3.14)
boolean (var b : boolean = true)
string (var greet : string = "hi")
array (var lst : number[] = [1, 2])
enum
any (var a : any = 3; var b : any = "hi")
```

```
TypeScript Class
```

```
class Employee{
  name : string;
  salary : number;

constructor(name : string, salary : number){
    this.name = name;
    this.salary = salary;
  }
  display(){ console.log(this.name); }
}

var emp = new Employee('Jon', 50000);
emp.display();
```

### **Lecture #15 – Event-Based Programming and Cryptocurrencies**

### JavaScript Embedded in HTML

Create a button on a website that prints

#### Improved JavaScript in HTML

### **Adding an Event Listener**

 If clicking a button should perform multiple functions, then an event listener should be used.

#### **Removing an Event Listener**

Event listeners can be removed by function name.
 Example:

### **Events in JavaScript**

· JavaScript is single threaded.

**Types of Keys** 

• Private Key: Known

only by the owner

• Public Key: Known

by everyone

An event must be run to completion before the next event handler can run.

#### **Event Emitter**

• Import the "events" module using the syntax

```
var ee =
require('events').EventEmitter;
```

Used to create event via the keyword "on".
 Example:

```
ee.on('die', function() {
          console.log("Died");
     });
```

 Invoking (emitting) an event using the keyword "emit" Example:

### **Create a TCP Server in Node.js Using Event Listeners**

```
var net = require('net');
var eol = require('os').EOL;
var srvr = net.createServer();

// Add an event listener
srvr.on('connection', function(client) {
    client.write('Hello there!' + eol);
    client.end();
});
srvr.listen(9000);
```

telnet – Used to connect to a TCP server on the command line.

127.0.0.1 - IP address of localhost

### Cryptocurrencies

### **Digital Signature**

### Non-Repudiation – Involves associating actions or changes to a unique individual.

- Solution in Cryptocurrency: Digital signature.
- Procedure:
  - Step #1: Owner encrypts the message with his private key
  - Step #2: Use the public key to decrypt the message.
- Analogy: Enclosed Bulletin Board

### **Private Key Encryption**

Used to transmit sensitive data to a specific recipient.

- Procedure:
- Step #1: A user encrypts his data using the recipient's public key.
- Step #2: The intended recipient decrypts the data using his private key.
- Analogy: A public mailbox. Anyone can put letters in, but only the mailman has the key to open the box.

- update Used to update the signature with the specified message contents. Each signature object can only be updated once.
- hex Specifies that the output should be in hexadecimal format.
- Sync Ensures that the file read is done immediately without relying on a callback.
- SHA "Secure Hash Algorithm"
- RSA Signature algorithm

### **Example: JavaScript Signer Example**

```
var crypto = require('crypto');
var fs = require('fs');

// Constructor for a "Signer" object
function Signer(privKeyFile) {
    this.privKey = fs.readFileSync(privKeyFile).toString('ascii');
}

// Add a "signMessage" function to the Signer prototype
Signer.prototype.signMessage = function(msgFileName) {
    var msg = fs.readFileSync(msgFileName).toString('ascii');
    var sign = crypto.createSign('RSA-SHA256');

    return sign.update(msg).sign(this.privKey, 'hex');
}
```

Double Spending – Spend the same funds in multiple places.

### **Solutions to Prevent Double Spending:**

- Centralized Authority Disadvantages include that the central authority would charge a fee and not everyone trusts central authorities.
- Decentralized Authority Broadcast transactions to everyone.

**Ledger** – Used to keep a **history of all transactions** and the funds held by all users.

```
Example: JavaScript Verifier Example
var crypto = require('crypto');
var fs = require('fs');
// Constructor for a "Verifier" object
function Verifier(publicKeyFile) {
  this.publicKey = fs.readFileSync(privKeyFile).toString('ascii');
}
// Add a "verifySignature" function to the Verifier prototype
Verifier.prototype.verifySignature = function (msgFileName,
                                                signature) {
            var msg = fs.readFileSync(msgFileName).toString('ascii');
            // Create a verifier
            var ver = crypto.createVerifier('RSA-SHA256').update(msg);
           // Verify signature matches the hash
            var legit = ver.verify(this.publicKey, signature, 'hex');
            return legit;
```

### **Bitcoin Mining**

- Block Chain Defines the transaction history.
   Used to prevent double spending.
- Proof of Work Verification of the block chain.
- Miners hash transaction details plus a "proof" (i.e. nonce)
   Reward: New bitcoins are mined for the first to find a proof.
- Cost to *Derive* a Proof: 2<sup>N</sup> where N is the number of the initial bits that must be "0" for the proof to be valid.
- Cost to Verify a Proof: A single hash
- Bitcoin protocol is designed to make mining more profitable than cheating.

### **Attributes of a Good Hash Function**

Role of a Hash Function:	One Way: Given an output "y", it is	Collision Resistant: It is infeasible to find		
Compress arbitrary length inputs	infeasible to find an "x" such that:	any "x" and "y" such that:	Compression	Efficient
to small, fixed length outputs.	h(x) = y	h(x) == h(y)		

### Lecture #17 - Macros and Sweet.js

**Example: C Preprocessor Example** 

• Short for "macroinstruction"

**Basic Compiler Structure with C-Style Macros** 

#### #define PT 3.14159 Expanded Source **Tokens** #define SWAP(a,b) {int tmp=a;a=b;b=tmp;} Code Code Pre-Lexer/ • Rule specifies how an input Parser Processo Tokenizer sequence maps to a int main(void) { replacement sequence. int x = 4, y=5, diam = 7; double circum = diam \* PI; SWAP (x,y) Macros in C • Performed by a preprocessor C Preprocessor Output **Abstract** Compiler Interpretter Syntax • Rely on text substitution. int main(void) { Tree int x = 4, y = 5, diam = 7;• Embedded languages like PHP, double circum = diam \* 3.14159; Ruby, etc. use a similar approach. {int tmp=x;x=y;y=tmp;} **Machine Code** Interpretter Problem with C Macros (Input) **Macros** in JavaScript Hygienic Macro - Any macro whose expansion is guaranteed not · No standard macro system for JavaScript // Macro should be on one line to cause the accidental capture of identifiers. • Sweet.js has been gaining interest. #define SWAP(a,b) {int tmp=a; • Recently redesigned. a=b: **Syntactic Macros** b=tmp;} Sweet, is • Derive from Lisp since Lisp programs are essentially one big AST. Borrows concepts from Racket. int main(void){ • Source-to-source compiler (i.e., transpiler) for int x = 4, tmp = 5; Work at the level of abstract syntax trees. JavaScript. SWAP (x, tmp) • Examples of other JavaScript transpilers: Powerful by expensive. o TypeScript **Problem with C Macros (Output)** o CoffeeScript Hygiene easier to address at the AST level. o Dart (includes its own VM) int main(void){ · Project backed by Mozilla int x = 4, tmp = 5; • Essentially a source-to-source compiler. $\{ int tmp = x; \}$ a = tmp;**Basic Compiler Structure with Syntactic Macros Invoking Sweet.js** tmp = tmp; From command line: 1 \$ sjs myfile.js -d out/ } Abstract **Abstract** Macro Hence, a variable name collision between Syntax **Syntax Expander** • Compiled files run normally (as shown below for Node): the two variables named "tmp". This is Tree Tree \$ node out/myfile.js known as "inadvertent variable capture"

```
Keywords in Sweet.js
                                                                      Concatenating Multiple Result Strings
            Writing a Swap Function in Sweet.js
                                                         This function squares a set of input variables.
                                                                                                                    • let - Create a Sweet.js
syntax swap = function(ctx){
                                                                                                                      variable.
  let innerCtx = ctx.next().value().inner();
                                                         syntax square = function(){
  let first = innerCtx.next().value();
                                                           var innerCtx = ctx.next().value().inner();
                                                                                                                    • ctx.next().value() -
  // Eat the comma
                                                           // Start with empty results
                                                                                                                      Get the next value from the
  innerCtx.next(); // No need for "value()"
                                                           result = #``;
                                                                                                                      context
                                                           while(let stx of innerCtx) {
  // Get the second parameter
                                                             result =
  let second = innerCtx.next().value();
                                                                                                                    • #`...` - Used to define a
                                                                result.concat(#`${stx}=${stx}*${stx};`);
  return #`var tmp = ${first};
                                                                                                                      result string.
                                                              // Eat comma
             ${first} = ${second};
                                                              // Ignored if no comma present
             ${second} = tmp; ;

    concat – Used to combine

                                                              innerCtx.next();
                                                                                                                      two result strings.
                                                           }
swap(a, b); // Invokes the macro
                                                         square(a, b, c); // Invokes the macro
                                                                                                                    • let xxx of yyy - Iterate
                                                                                                                      over a list of tokens.
Note #1: The returned string is preceded by a pound (#) sign and
                                                         Note #1: Use ".concat" to concatenate multiple result strings.
is enclosed in backticks (`).
                                                                                                                    • isIndentifier - Used to
                                                         Note #1: If a token is not present, ".next()" does not cause an
                                                                                                                      check if a Sweet.js variable
Note #2: Sweet.js variables are declared with "let".
                                                         error.
                                                                                                                      matches some string.
```

### A class in Sweet.js

```
syntax class = function(ctx){
 let className = ctx.next().value();
 let bodyCtx = ctx.next().value().inner();
  // By default assume empty constructor
 let construct = #`function() { }`;
let result = #``;
 while( let item of bodyCtx ){
    // Check if constructor
   if(item.isIndenifier('constructor')){
      // Get arguments then function code
      construct = #`function ${className}
                    ${bodyCtx.next().value()}
                    ${bodyCtx.next().value()}`;
   else {
     // Add the function to the class prototype
     result = result.concat(
         #`${className}.prototype.${item} =
             function ${bodyCtx.next.value()}
             ${bodyCtx.next.value()}`);
    // Return the constructor and methods
```

return construct.concat(result);

### Input JavaScript class Code to be Parsed by Sweet.js

# Lecture #18 – Simply Typed Lambda Calculus

### **Lecture #19 – Metaprogramming and JS Proxies**

**Metaprogramming**: Writing programs that manipulate other programs.

• Proposed in ECMAScript 6 for JavaScript.

### **Terminology in Reflection**

- Introspection: Ability to examine (but not modify) the structure of a program.
- Self-modification: Ability to modify the structure of a program.

Introspection
Ability to examine (but not modify) the structure of a program.

### **JavaScript Examples**

```
Property Lookup
"x" in o; //o is an object

Iterate Over All Properties of an Object
for ( prop in o ) {
```

// Do something

}

```
Self-modification
Ability to modify the structure of a program.
```

# JavaScript Examples

```
o["x"]; // Computed property
o.y = 42; // Add new property
delete o.y; // Delete property
// Reflected method call
O["m"].apply(null, [38]);
```

### Proxies in JavaScript

- Metacircular Interpretation The language is able to understand its own language.
- Until recently, JavaScript did not support
- o Javascript proxies are intended to fix that.
- Node.js' implementation of proxies lags behind the standard.

intercession.

Proxies only exist for objects and functions.
 Proxies do not exist for primitives.

### **Proxies and Common Lisp**

- Developed before object oriented languages were popular.
- Many libraries were create with non-standard OO systems.
- Common Lisp Object System (CLOS) – Standard object oriented system for Lisp.

### Achieving Lisp Object Backwards Compatibility

Option #1: Rewrite all libraries using CLOS. Disadvantage

- Huge number of libraries.
- Not feasible to rewrite them all.

### Option #2: Make a complex API

- API difficult to understand.
- Systems had conflicting features.

Option #3: Keep API simple and modify object behavior to fit different systems.

• This approach relies on **metaobject protocols**.

### **Proxies and Handlers**

- The behavior of a proxy is determined by traps specified in its handler (i.e., the metaobject).
- Trap Methods that intercept an operation.
- Handler The metaobject that specifies the details of the trap. The handler itself is usually a normal object.
- Using proxies in node requires a special flag "--harmony-proxies". Example:
- \$ node --harmony-proxies prog.js

### **Kinds of JavaScript Proxies**

• Object Proxies - Defined with:

Proxy.create(handler, proto)

• Functions (with extra traps) - Defined with:

• Proxies do not exist for primitives.

### **Read Only Handler**

Information Control – Share a reference to an object, but do allow it to be modified.
 Example: Reference to the DOM.

```
function ReadOnlyHandler(obj){
  delete : function(name){
    return obj[name];
}

// rcvr can be ignored
  set : function(rcvr, name, val){
    return true;
}
```

```
A Noop Proxy – All Operations Passed through Unchanged
```

```
function handlerMaker(obj){
  // Delete a property from an object
 delete : function(name) { return obj[name]; },
  // Check if object has the specified property
 has : function(name) { return name in obj;},
  // Check if object (not prototype chain) has property
 hasOwn : function(name) {return Object.property
                                   .hasOwn(obj, name);},
  // Get a property value
 get : function(name) { return obj[name]; },
 // Set a property value
 set : function(rcvr, name, val) { obj[name] = val; },
  // Get all properties of an object
 enumerate : function(){
   var props = [];
    var prop:
   for(prop in obj) { props.push(prop); }
   return props;
 },
  // Get all of the keys of an object
 keys: function() { return Object.keys(obj); }
}
```

### **Aspect Oriented Programming**

- Some code not well organized into objects. Example:
  - Cross-cutting concern where code is spread throughout a program.
- Canonical Example: Logging Statements
  - $\,\circ\,$  Littered throughout the code
- Swapping out a logger requires massive code changes.
- Solution: Use a proxy

### **Lecture #20 – Introduction to Ruby**

### **Influences of Ruby**

- SmallTalk
- Everything is an object
- Blocks
- Metaprogramming
- Perl
  - Regular Expressions
  - o Function names

### **Ruby on Rails**

- "Killer" app for Ruby
   Lightweight web
   framework
  - "Convention over configuration" – If use standard configuration, very little configuration required.
- Initial framework was PHP, but that was abandoned.

### **Basic Ruby Syntax**

```
puts "Hello World"

a = [1, 2, 3]

m = { 'a' => "Apple",
    'b' => "Bear",
    'c' => "Cat" }

# Prints "1"
```

# Prints "Apple"

puts a[0]

puts a['m']

### Keywords

@ - Represents an object property

### **Returning From a Function**

- Every function in Ruby returns a value, even if return is not used.
- If no return is specified, a function returns the last used value.

### **Basic Ruby Class**

```
class Person
  # Constructor
 def initialize name # Parameter
    # Attribute
   @name = name
 end
 # Getter
 def name
   return @name
 end
  # Setter
 def name = newName
   @name = name
 end
 # Method
 def say_hi
   puts "Hi my name is #{@name}"
 end
```

```
Using Metaprogramming for Getters and Setters

class Person

# Replaces getters and setters

# Uses metaprogramming
attr_accessor :name

# Constructor
def initialize name

# Attribute
@name = name
End

# Method
def say_hi
   puts "Hi my name is #{@name}"
end
end
```

### **Using a Class in Ruby**

```
p = Person.new "Joe"
puts "Name is #{p.name}"
p.say_hi
#{...} - Embeds a variable in a Ruby String
```

### **Getters and Setters the Ruby Way**

### Relies on metaprogramming

- attr\_reader Getter only
- attr writer-Setter only
- attr\_accessor Getter and setter

### **Reopening a Class in Ruby**

- Class definitions can be changed during runtime in Ruby.
- This is known as "reopening the class"

#### **Parent Class**

end

```
class Dog
  # Parentheses optional
  def initialize(name)
     @name = name
  end
  def speak
     puts "#{@name} says bark"
  end
end
```

## Inheritance in Ruby

```
class GuardDog < Dog
  attr_accessor :breed
  def initialize(name, breed)
    # Use parent constructor
    super(name)
    @breed = breed
end
  def attack
    puts "Grrr"
end
end</pre>
```

**Child Class** 

Note: Inheritance is doing using the less than (<) operator.

#### Mixin

- Add features to a class
- Similar to interfaces in Java with the exception that they <u>can</u> include functionality.
- module Keyword to define a Mixin.
- include Keyword to include a Mixin into a class.

### **Blocks in Ruby**

### Superficially similar to blocks in other languages.

- Create custom control structures.
- Can be represented with curly brackets ({...}) or do/end.

### File IO without Blocks

```
file = File.open('test.txt','r')
file.each_line do |line|
  puts line
end
file.close
```

Note #1: Similar "boilerplate" code of open and closing the file.

Note #2: It is possible one may forget to the close the file.

### File IO without Blocks

```
File.open('test.txt','r') do |file|
  file.each_line { |line|
    puts line
  }
end
```

Note #1: Eliminates the "boilerplate" code.

Note #3: When using a block (both do/end, and curly brackets), surround the variable names in pipes (1).

### Example: Mixin

```
# Define the mixin
module RevString
  def to_rev_s
    # Object is implicit
    to_s.reverse
    end
end

# Reopen the Person Class
class Person include RevSting
  def to_s
    # Returns the value
    @name
    end
end
```

### Dynamic Code Evaluation (eval)

- Executes source code dynamically
  - Code passed as either a string (or a block of code)
- Popular feature in JavaScript
  - Early usage was to convert JSON strings to variables since not supported by JavaScript.
- Source of security concerns.

### Additional Ruby eval Methods

- instance\_eval Evaluates code within an object's body.
  - o Access the internals of an object.
- class\_eval Evaluates code within a class' body.
- o Modifies the class' definition.
- Takes either a string or block of code. Block of code is more secure.

# Example: Use instance\_eval to Change an Object's Value

```
# Create with the name Bob
bob = Person.new "Bob"
```

- # Change his name
  bob.instance\_eval do
   @name = "Steve"
  end
- # Prints "Steve"
  puts bob.name

### **Regular Expressions in Ruby**

- sub Replaces the first instance of a string match.
  - To perform the modification in place, must include an exclamation point (!) after sub.
- gsub Replaces all instance of a string match.
   To perform the modification in place, must include an exclamation point (!) after sub.

```
Example: class eval in Ruby
                                                                           Example: Using Regular Expressions in Ruby
# Applies to all classes
class Class
                                                         s = "Hi, I'm Larry; this is my" +
  # Simulate the "attr accessor" function
                                                              " brother Darryl, and this" +
  def my_attr_accessor(args)
                                                              " is my other brother Darryl."
    args.each do |prop|
                                                         s.sub(/Larry/,'Laurent')
      # Create getter
      self.class_eval("def #{prop};
                                                          # Prints s unchanged
                         return @#{prop};
                        end")
      # Create setter
                                                         # Changes first "Larry" to "Laurent"
      self.class eval("def #{prop} = v;
                          @\#\{prop\} = v;
                                                         s.sub! (/Larry/, 'Laurent')
                        end")
   end
                                                          # Prints first "brother" replaced with
  end
                                                          # "frere". s is unchanged, bt it did
end
                                                         # return the modified string.
# Use the new attribute
                                                         puts s.sub(/brother/, 'frère')
class Musician
 my attr accessor :name, :genre
                                                          # Same as previous except all where
                                                         # changed when printing.
end
                                                         puts s.gsub(/brother/, 'frère')
m = Musician.new
m.name = "Bob Marley"
puts m.name # Prints "Bob Marley"
```

### **Regular Expression Symbols in Ruby**

/ . / - Any character except a newline	/\w/ - Any word character: [a-zA-Z0-9_]	/\d/ – Any digit character: [0–9]	/\w/ – Any whitespace character: [ \t\r\n\f]	
	/\w/ – Any non-word character:	/\D/ – Any non-digit character:	/\w/- Any non-whitespace character:	
	[^a-zA-Z0-9_]	[^0-9]	[^ \t\r\n\f]	
- Zero or more times + - One or more times		? – Zero or one times (optional)		

### **Important Syntax in Ruby**

For Each Loop  object.each do  val  end	Create a Mixin module Name end	Return from a Block def block_name yield x end	<pre>Single Line If Statement  x = 5 # Does nothing x = 3 if (x &gt; 10) puts x # Prints "5"</pre>	# Create array from 1 to 5  x = (15)  Note: Uses parentheses.
irb – Command line for Ruby similar to GHCi.				

### Lecture #21 - Blocks and Messages

### Influence of Smalltalk on Ruby

- Everything is an object
- Blocks
- Message passing

### Benefits of Blocks in Ruby

- Create custom control structures
- Eliminate boilerplate code.
- Ruby blocks are closures, but they are different than JavaScript blocks.

### Example: do\_noisy Block

```
def do_noisy
  puts "About to call block"
  yield # Calls block code
  puts "Just called block
end
```

Note: Called with a do/end or with curly brackets.

# **Example:** Extend Array Class to Return Lowercase Version of Every Element

```
# Reopen the Array class
class Array
def each_downcase
    self.each do |val|
        yield val.downcase
    end
end
end
```

#### Example: Probabilistic Run Block

```
# Probabilistic Run Block
def with_prof(prob)
  yield if (Random.rand < prob)
end
with_prob 0.42 do
  puts "Prints 42% of time."</pre>
```

### **Example: Passing Code to a Block**

```
def with_prob2(prob, &blk)
  blk.call if (Random.rand < prob)
end</pre>
```

 $\mbox{\bf blk}$  – Block of code passed to the function.

Note #1: Argument name has an ampersand (&) before it.

Note #2: No ampersand is used when calling the block.

### Example: Sharing Code Between Blocks

```
def half_the_time(prob, &blk)
  with_prob2(0.5, &blk)
end
```

Note: Need to pass argument to the function with the ampersand (&).

### Example: with prob in JavaScript

```
function with_prob(prob, f) {
  if(Math.random() < prob) {
    return f();
  }
}</pre>
```

**Note:** The JavaScript implementation relies on **callbacks**.

### Example: Difference Between Ruby and JavaScript Blocks

```
Ruby

def coin_flip
  with_prob 0.5 do
    return "Heads"
  end
  return "Tails"
end
```

Note: This returns "Heads" half the time and "Tails" half the time.

This is because a return in a Ruby block returns for the

 This is because a return in a Ruby block returns for the entire function.

# JavaScript function coin\_flip() { with\_prob(0.5, function() { return "Heads";} return "Tails"; }

Note: This always returns "Tails"

• This is because even if "with\_prob" runs, the return only occurs within the anonymous function.

### **Example: Probabilistic Run Block**

```
# Probabilistic Run Block
def with_prof(prob)
  yield if (Random.rand < prob)
end
with_prob 0.42 do
  puts "Prints 42% of time."
end</pre>
```

### **Example: Passing Code to a Block**

```
def with_prob2(prob, &blk)
  blk.call if (Random.rand < prob)
end</pre>
```

**blk** – Block of code passed to the function.

Note #1: Argument name has an ampersand (&) before it.

Note #2: No ampersand is used when calling the block.

### **Example: Sharing Code Between Blocks**

```
def half_the_time(prob, &blk)
with_prob2(0.5, &blk)
```

Note: Need to pass argument to the function with the ampersand (&).

### Singleton Classes

- In Ruby, every object has its own singleton class.
- This class holds methods and fields unique to that object.
- This is different from Singleton Objects in design patterns.

### Example: Adding a Property to a Variable in JavaScript

```
function Employee(name, salary) {
  this.name = name;
  this.salary = salary;
}
var a = new Employee("Alice", 500);
var b = new Employee("Bob", 1000);

// Add a signing bonus to "Alice"
a.signingBonus = 2000;
```

### **Accessing Singleton Classes in Ruby**

- To open an object's singleton class, use double less than symbols ("<<").</li>
- Code only added to the specific object being reference.

### Example: Adding a Property to an Object in Ruby

```
class Employee
  attr accessor :name,:salary
  def initialize(name, salary)
    @name = name
    @salary = salary
  end
  def to s
    @name # No return required
  end
end
# Create the Objects
a = Employee.new("Alice", 500)
b = Employee.new("Bob", 1000)
# Access the singleton class of "a"
class << a
  def signing bonus
   2000
  end
end
```

### **Example:** Using a Singleton Class to Create Static Methods

```
# Add Static Methods to Employee Class
class Employee
  class << self</pre>
    def get_employee_by_name(name)
   @employee[name] # No return needed
    end
    # Called in constructor
    def add(emp)
      puts "Adding #{emp}"
      # Create map if not exist
      @employee = Hash.new unless @employee
      @employee[emp.name] = emp
    end
  end
end
```

### **Message Passing**

- Represents inter-object interaction.
- o Method name
- o Data: Method parameters (if any)
- Receiver:
  - o Processes the message
- Receiver may not understand the message.

### method missing

- overridden.
- o Smalltalk Name: doesNotUnderstand
- O Ruby Name: method\_missing
- Invoked whenever an unknown method is called.

- Sender Sends:

- o (Optionally) returns data

- Method that is part of every class. Can by

```
Example: missing method in Ruby
```

```
class Person
 attr accessor :name
 def initialize(name)
   @name = name
  # Called when method unknown
 def method missing(m)
   puts "Didn't understand #{m}"
  end
end
```

### **Active Record and Message Passing**

- Relational database tool in Ruby.
- Specify fields in the database to be extracted based off method names. Example:

Person.find by first name "John"

### Lecture #22 - Virtual Machines and Just-In Time Compilation

### Virtual Machine Overview

- Code is compiled to bytecode o Byte code is low level o Platform independent
- The VM interprets the bytecode

### Scheme

- Similar to an AST. Uses parentheses.
- Relies on a stack.

### **Supported VM Operations**

- PUSH Adds an argument to the stack
- PRINT Pops an argument off the stack and prints it.
- ADD Pops two elements off the stack, adds them, and places result on the
- SUB Similar to add but for subtraction. If "A" is on the top of the stack and "B" is below it, the result is B - A
- MUL Similar to add but with multiplication.

### Compilers vs. Interpreters vs. JIT

- Compiler
- o Efficient execution

#### Interpreter

- o Runtime flexibility
- o Efficient execution with runtime flexibility.

### **Just-In-Time Compliers**

- Interpret code
- "Hot" (i.e., heavily-used) sections are compiled at runtime.

### Advantages

- o Speed of compiled code
- o Flexibility of interpreter

### Disadvantages

- o Overhead of compiler and interpreter
- o Complex implementation

### **Dynamic Recompilation**

- JIT pursues aggressive optimizations o Makes assumptions about the code o Guard conditions verify assumptions
- Unexpected cases are interpreted (i.e., not compiled)
- Can in some corner cases outperform static compilation.

### **Types of JITs**

- Method Based Compile Methods
- Trace Based Compile loops

### **How to Support JITs for a Language**

- Option #1: Build your own JIT.
- o Study the latest techniques
- o Build large code bases to test.
- o Profile the code execution
- Option #2: Use someone else's Just-In-Time VM.

### **Final Exam Review Notes**

#### Parsec

- Parser combinator in Haskell.
- Example Question: Write a grammar a the level of the CSV parser.

### **JavaScript**

- Prototype based language.
- o Inherit from an object not from a class.
- o Add properties and methods on the fly.
- Closure By wrapping an inner function with an outer function, the inner function can encapsulate variables.
  - o Have a scope chain.

```
Example: Write a Function that Toggles a Variable Each Time the Function is Called
```

#### Example #1: Use a Global

```
var b = true;
function flip(){
 if(b) b = false;
 else b = true;
 return b;
```

```
var flip = function(){
 var b = true;
 return function(){
   if(b) b = false;
   else b = true;
   return b:
    // Needed to call out function
}();
```

Example #2: Use a Closure

### JavaScript - Multiparadigm Language

- Imperative
- Object Oriented
- Functional
  - o Supports higher order functions but not purely functional since not immutable.

### **JavaScript Scoping**

- No block scope
- Variable declaration hoisting to the top of the function.

### this in JavaScript

- In a method, this refers to the associated object.
- In a non-method, this refers to the global
- Constructor (with new) Refers to the object being created.
- DOM for Event Listeners Refers to the DOM
- apply, call, bind User can define what "this" refers to.

### Scope Precedence in JavaScript

#### From highest to lowest precedence

- 1. Variable object (i.e., local variables)
- 2. Scope Chain
- 3. Global object (i.e., this)

### Example: Scope Chain

```
function my hello(){
 var x = 5;
  // Scope Chain
 function print hello(){
   console.log("Hello " + x);
 print_hello();
```

### Quirks of JavaScript

- Semicolon insertion
- typeOf
  - o typeOf NaN Number
- o typeOf null Object
- == Not Transitive
  - $\circ$  If a == b and b == c, it is not guaranteed that a == c

- Designed to catch common JavaScript
- Based off lint for the C language.
- Performs static code analysis

### **TypeScript**

- Developed by Microsoft
- Source-to-source compiler (i.e., transpiler)
- Compiles to JavaScript
- Provides type annotation and classes.

### **Event Based Programming**

• Relies on listeners

my\_hello();

- Events are placed into an event queue.
- No concurrency.
- emit Invokes an event
- on Registers a listener
- Client-Based Programming Often used in GUIs.

### Metaprogramming

- Reflection Two primary categories. Both occur at runtime.
  - o Introspection Examine program's execution at runtime.
- o Self-modification Modify a program execution at runtime.
- Intercession Trigger or control interaction at runtime.
- Reflection is more common, but intercession is more powerful.

### Aspect-Oriented **Programming**

Designed to address crosscutting code (i.e., code that is interspersed everywhere in a program.

### **Metaobject Protocols**

- Metaobject Any object that can reason about the behavior of other objects.
  - o Example: Proxies
- Handler
- o Type of Meta Object.
- o Defines traps.
- Trap Methods that intercept an operation.

var b = true; function flip(){ if(b) b = false; else b= true; return b: var flip = function(){ var b = true; return function(){ if(b) b = false; else b = true; return b; }; };