

Homework #2: Operational Semantics for the WHILE Language

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1 Introduction to the WHILE Language

The “WHILE” language is a basic language that was defined in class. Figure 1 defines the expressions, values, and operators in this language. This notation for expressions (e), values (v), variables/addresses (x), and store (σ) applies to all sections of this document.

$e ::=$	x v $x := e$ $e; e$ $e \text{ op } e$ $\text{if } e \text{ then } e \text{ else } e$ $\text{while } (e) \ e$ $\text{not } e$ $\text{and } (e) \ (e)$ $\text{or } (e) \ (e)$	<i>Expressions</i> variables/addresses values assignment sequential expressions binary operations conditional expressions while expressions not expressions and expressions or expressions
$v ::=$	i b	<i>Values</i> integer values boolean values
$op ::=$	$+ \mid - \mid * \mid / \mid > \mid >= \mid < \mid <=$	<i>Binary operators</i>
σ		<i>Store</i>

Figure 1: The WHILE language

2 Base WHILE Language Small-Step Semantics Rules

The following figures enumerate the execution order, small-step semantics rules for the WHILE language expressions that were defined in class.

Variable Evaluation Rule:

$$[\text{SS-VAR}] \quad \frac{x \in \text{domain}(\sigma) \quad \sigma(x) = v}{x, \sigma \rightarrow v, \sigma}$$

Figure 2: Variable Small-Step Semantics Evaluation Order Rule**Set/Assignment Evaluation Rules:**

$$[\text{SS-ASSIGNCONTEXT}] \quad \frac{e, \sigma \rightarrow e', \sigma'}{x := e, \sigma \rightarrow x := e', \sigma'}$$

$$[\text{SS-ASSIGNREDUCTION}] \quad \frac{}{x := v, \sigma \rightarrow v, \sigma[x := v]}$$

Figure 3: Set/Assignment Small-Step Semantics Evaluation Order Rules**Binary Operator (op) Evaluation Rules:**

$$[\text{SS-OPCONTEXT1}] \quad \frac{e_1, \sigma \rightarrow e'_1, \sigma'}{e_1 \text{ op } e_2, \sigma \rightarrow e'_1 \text{ op } e_2, \sigma'}$$

$$[\text{SS-OPCONTEXT2}] \quad \frac{e, \sigma \rightarrow e', \sigma'}{v \text{ op } e, \sigma \rightarrow v \text{ op } e', \sigma'}$$

$$[\text{SS-OPREDUCTION}] \quad \frac{v_3 = v_1 \text{ op } v_2}{v_1 \text{ op } v_2, \sigma \rightarrow v_3, \sigma}$$

Figure 4: Binary Operator (op) Evaluation Order Rules**Sequence (;) Evaluation Rules:**

$$[\text{SS-SEQCONTEXT}] \quad \frac{e_1, \sigma \rightarrow e'_1, \sigma'}{e_1; e_2, \sigma \rightarrow e'_1; e_2, \sigma'}$$

$$[\text{SS-SEQREDUCTION}] \quad \frac{}{v; e, \sigma \rightarrow e, \sigma}$$

Figure 5: Sequence (;) Evaluation Order Rules

Conditional Statement (if) Evaluation Rules:

[SS-IFCONTEXT]	$\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3, \sigma \rightarrow \text{if } e'_1 \text{ then } e_2 \text{ else } e_3, \sigma'}$
[SS-IFTRUEREDUCTION]	$\frac{}{\text{if true then } e_1 \text{ else } e_2, \sigma \rightarrow e_1, \sigma}$
[SS-IFFALSEREDUCTION]	$\frac{}{\text{if false then } e_1 \text{ else } e_2, \sigma \rightarrow e_2, \sigma}$

Figure 6: Conditional (if) Small-Step Semantics Evaluation Order Rules

while Evaluation Rule:

[SS-WHILEREDUCTION]	$\frac{}{\text{while } (e_1) e_2, \sigma \rightarrow \text{if } e_1 \text{ then } (e_2; \text{while } (e_1) e_2) \text{ else false}, \sigma}$
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Figure 7: while Small-Step Semantics Evaluation Order Rule

3 Boolean Expressions Small-Step Semantics Rules

In following subsections, I describe three additional expression types in the updated the WHILE language namely: **not**, **and**, and **or**.

3.1 not Expression

not in my modified version of the WHILE language behaves as a standard Boolean **not**. It takes a single Boolean value and returns its complement. If an expression is passed, the language simplifies that expression until it is in normal form at which point it applies the Boolean **not**.

not Evaluation Rule:

[SS-NOTREDUCTION]	$\frac{}{\text{not } e, \sigma \rightarrow \text{if } e \text{ then false else true}, \sigma}$
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Figure 8: not Small-Step Semantics Evaluation Order Rule

3.2 and Expression

and is designed to mimic the Boolean **and** with the exception that it supports short circuit compare. Hence, if the first expression in the "**and**" evaluates to **false**, the second parameter is not evaluated at all.

and Evaluation Rules:	
[SS-ANDCONTEXT]	$\frac{e_1, \sigma \rightarrow e'_1, \sigma'}{\text{and } (e_1) (e_2), \sigma \rightarrow \text{and } (e'_1) (e_2), \sigma'}$
[SS-ANDREDUCTION]	$\frac{e' = \text{if } e \text{ then true else false}}{\text{and } (v) (e), \sigma \rightarrow \text{if } v \text{ then } e' \text{ else false}, \sigma}$

Figure 9: and Small-Step Semantics Evaluation Order Rules

3.3 or Expression

or is a composite of the expressions "**not**" and "**and**" described in sections 3.1 and 3.2 respectively.

or Evaluation Rule:	
[SS-ORREDUCTION]	$\frac{e'_1 = \text{not } e_1 \quad e'_2 = \text{not } e_2 \quad e_3 = \text{and } (e'_1) (e'_2)}{\text{or } (e_1) (e_2), \sigma \rightarrow \text{not } e_3, \sigma}$

Figure 10: or Small-Step Semantics Evaluation Order Rule