CS252 - Final Exam Study Guide

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Lecture #01 – General Introduction

Reasons for Different Programming Languages

- 1. Different domains (e.g. web, security, bioinformatics)
- 2. Legacy code and libraries
- 3. Personal preference

Programming Language Design Choices

- 1. Flexibility
- 2. Type safety
- 3. Performance
- 4. Build Time
- 5. Concurrency

Features of Good Programming Languages

- 4. Safety (e.g. security and can errors be
 - caught at compile time)
 - 5. Machine independence
 - 6. Efficiency

Goals almost always conflict

Conflict: Type Systems

- Advantage: Prevents bad programs.
- Disadvantage: Reduces programmer flexibility.

Blub Paradox: Why do I need advanced programming language techniques (e.g. monads, closures, type inference, etc.)? My language does not have it, and it works just fine.

Current Programming Language Issues

1. Simplicity

2. Readability

3. Learnability

- Multi-core "explosion"
- Big Data
- Mobile Devices

Advantages of Web and Scripting Languages

- Examples: Perl, Python, Ruby, PHP, JavaScript
- · Highly flexible
- Dynamic typing
- · Easy to get started
- Minimal typing (i.e. type systems)

Major Programming Language Research Contributions

- Garbage collection
- · Sound type systems
- Concurrency tools
- Closures

Programs that Manipulate Other Programs

- Compilers & interpreters
- JavaScript rewriting
- Instrumentation
- Program Analyzers
- IDFs

Formal Semantics

- Used to share information unambiguously
- Can formally prove a language supports a given property
- Crisply define how a language works

Types of Formal Semantics

- Operational
 - Big Step "natural"
 - o Small Step "structural"
- Axiomatic
- Denotational

Haskell

- Purely functional Define "what stuff is"
- No side effects
- Referential transparency A function with the same input parameters will always have the same result.
 - o A function call can be replaced with its value and nothing will change.
- Supports type inference.

Duck Typing – Suitability of an object for some function is determined not by its type but by presence of certain methods and properties.

- o More flexible but less safe.
- Supported by Haskell
- o Common in scripting languages (e.g. Python, Ruby)

Side Effects in Haskell

- Generally not supported.
- Example of Support Side Effects: File IO
- Functions that do have side effects must be separated from other functions.

Lazy Evaluation

- · Results are not calculated until they are needed
- Allows for the representation of infinite data structures

Lecture #02 - Introduction to Haskell

Key Traits of Haskell

- 1. Purely functional
- 2. Lazy evaluation
- 3. Statically typed
- 4. Type Inference
- 5. Fully curried functions

ghci - Interactive Haskell.

let - Keyword required in ghci to set a variable value. Example:

> let f x = x + 1

> f 3

4

Run Haskell from Command Line Use runhaskell keyword.

Lists

> runhaskell <FileName>.hs

Hello World in Haskell

main :: IO () main = do

putStrLn "Hello World"

Primitive Classes in Haskell

- 1. Int Bounded Integers
- 2. Integer Unbounded
- 3.Float
- 4.Double
- 5.Bool 6.Char

- Base 0
- Comma separated in square brackets
- Operators
 - o: Prepend
 - ++ Concatenate
 - o!! Get element a specific index
 - o head First element in list
 - o tail All elements after head
- the last one

o last Last element in the list

o take n Take first n elements from a

o init All elements in the list except

- o replicate 1 m Create a list of length 1 containing only m
- o repeat m Create an infinite list containing only m

• Can be infinite or bounded • Use the "..." notation. Examples:

> [1..4]

Ranges

[1, 2, 3, 4]

> [1,2..6] [1, 2, 3, 4, 5, 6]

> [1,3..10] [1, 3, 5, 7, 9]

> [5, 4..1] [5, 4, 3, 2, 1]

Hello World in Haskell

putStrLn "Hello World"

main :: IO () main = do

> putStrLn \$ "Hello " ++ "World" "Hello World"

> let s = bra in s !! 2 : s ++ 'c' : last s : 'd' : s "abracadabra"

List Examples

Infinite List Example

> let even = [2,4..]> take 5 even

[2, 4, 6, 8, 10]

```
List Comprehension
                                                                        A Simple Function
• Based off set notation.
                                                              > let inc x = x + 1
                                                              > inc 3
• Supports filtering as shown in second example
                                                                                                                  Pattern Matching
• If multiple variables (e.g. a, b, c) are specified, iterates through
                                                                                                    • Used to handle different input data
 them like nested for loops.
                                                              > inc 4.5
                                                                                                    • Guard uses the pipe ( ) operator
• Uses the pipe (|) operator. Examples:
                                                              5.5
                                                                                                    • Example:
> [ 2*x | x <- [1..5]]
                                                              > inc (-5) -- Negative
                                                                                                    inc :: Int -> Int
[2, 4, 6, 8, 10]
                                                                         Type Signature
                                                                                                      | x < 0 = error "invalid x"
> [(a, b, c) | a <- [1..10], b <-[1..10],
                                                              • Uses symbols ":: " and "->"
                                                                                                    inc x = x + 1
                  c \leftarrow [1..10], a^2 + b^2 = c^2]
                                                              • Example:
                                                              inc :: Int -> Int
 [(3, 4, 5), (4, 3, 5), (6, 8, 10), (8, 6, 10)]
                                                              inc x = x + 1
```

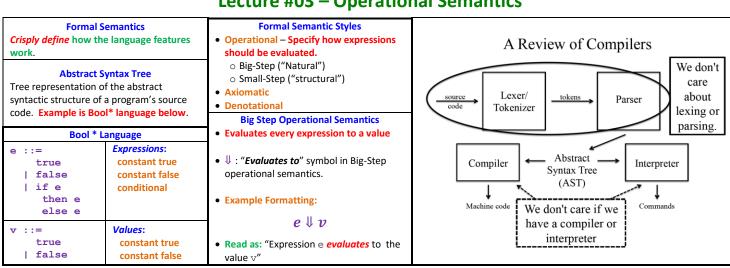
```
Recursion
• Base Case - Says when recursion should
                                                      Lab #01 – Max Number
                                                                                                       Reasons for a Large Number of
• Recursive Step - Calls the function with a
                                      > maxNum :: [Int] -> Int
                                                                                                          Programming Languages
 smaller version of the problem
                                      > maxNum [] = error "Invalid Input"

    Different domains

                                      > maxNum [x] = x
                                                                                                   • Different design choices
Example:
                                      > maxNum (x:xs) = if x > maxXs then x else maxXs
addNum :: [Int] -> Int
                                           where maxXs = maxNum xs
addNum [1] = 0
addNum (x:xs) = x + addNum xs
```

```
Recursion
                                                                  Haskell's Base Typeclasses
• :t or :type - Gets the type of a variable or function.
                                                          • Ord - Can be ordered
                                                          • Eq - Can perform equality check
Example:
                                                          • Show - Can convert to String
> :type 'A'
                                                          • Read - Can convert from String
'A' :: Char
                                                          • Enum - Sequentially Ordered
> :t "Hello"
                                                          • Bounded – Has upper and lower bound.
"Hello" :: [Char]
```

Lecture #03 – Operational Semantics



Small-Step Operational Semantics	Bool* Small-Step Operational Semantics Rules		
Evaluate an expression until it is in normal form	E-IfTrue:	Example: Reduce the expression	
	<u></u>	if (if true then false else true) then true else false	
Normal Form – Any form that cannot be	if true then e_2 else $e_3 \rightarrow e_2$		
evaluated further.	E-IfFalse:	Step #1: Use rule "E-IfTrue" with "E-If"	
 → : "Evaluates to" symbol in small step operational semantics. Example: 	if false then e_2 else $e_3 \rightarrow e_3$	if false then true else false	
e ightarrow e'' ightarrow v	E-If:	Step #2: Use rule "E-IfFalse" (Now in normal form)	
$ullet$ $ o^*$: Many evaluation steps required. Example: $oldsymbol{e} o^* \ oldsymbol{v}$	$\frac{e_1 \to e_1'}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3 \to \text{if } e_1' \text{ then } e_2 \text{ else } e_3}$	false	

• 0 : The Number "0"

Bool* Extension: Numbers true false | if e then e else e • succ 0: Represents "1" 0 • succ succ 0: Represents "2" | succ e pred e • pred n: Gets the predecessor v ::= true | false

Literate Haskell

• File Extension: ".lhs"

• Code lines begin with ">"

· All other lines are comments.

• "Essentially swaps code with comments."

```
Case Statement in Haskell
```

• Keywords: case, of, otherwise

Operator: ->

Example:

case x of val1 -> "Value 1" val2 -> "Value 2" otherwise -> "Everything else."

Lab #02 Review

```
Bool Expression Type
                                       > data BoolVal = BVTrue
 data BoolExp = BTrue
        BFalse
>
        | Bif BoolExp BoolExp
                                            deriving Show
        | B0
        | Bsucc BoolExp
                                       > data BVInt = BV0
>
        | Bpred BoolExp
    deriving Show
                                       >
                                            deriving Show
```

```
BoolVal Type
  | BVFalse
  | BVNum BVInt
 | BVSucc BVInt
```

Type Constructors: BoolExp, BoolVal, BVInt

Non-nullary Value Constructors: Blf, Bsucc, Bpred, BVSucc, BVNum

Note: Even constants like BO, BTrue, BFalse, BVTrue, and BVFalse are nullary value constructors (since they take no arguments)

Lecture #04 – Higher Order Functions

Lambda

- Analogous to anonymous classes in Java.
- Based off Lambda calculus
- Example:

```
> (\x -> x + 1) 1
>(\x y -> x + y) 2 3
```

Function Composition

Extended Bool * Language

| IntV

IntV ::= 0 | succ IntV

- Uses the period (.)
- f(g(x)) can be rewritten (f . g) x

Point-Free Style

- Pass no arguments to a function
- Example:

```
> let inc = (+1) - No args
> inc 3
```

Example: Lambda with Function Composition $> let f = (\x -> x - 5)$. $(\y -> y * 2)$ > f 7 9 > let f = ($x y \rightarrow x - y$) $(\z -> z * (-1))$

Iterative vs. Recursive

- Iterative tends to be more efficient than recursive.
- Compiler can optimize tail recursive function.

Tail Recursive Function - The recursive call is the last step performed before returning a value.

Not Tail Recursive

```
public int factorial(int n) {
  if (n==1) return 1;
  else {
    return n * factorial(n-1);
```

Last step is the multiplication so not tail recursive.

Tail Recursive Factorial

```
public int factorialAcc(int n, int acc)
  if (n==1) return acc;
 else {
    return factorialAcc(n-1, n*acc);
}
```

Tail recursive code often uses the accumulator pattern like above.

Tail Recursion in Haskell fact' :: Int -> Int -> Int fact' 0 acc = acc fact' n acc = fact' (n - 1) (n * acc)

Higher Order Functions

Functions in Functional Programming

- Functional languages treat programs as mathematical functions.
- Mathematical Definition of a Function: A function f is a rule that associates to each x from some set X of values a unique y from a set of Y values.

$(x \in X \land y \in Y) \rightarrow y = f(x)$

- f Name of the function
- X Independent variable
- y Dependent variable
- X Domain
- **Y** Range

Qualities of Functional Programming

- Functions clearly distinguish:
 - Incoming values (parameters)
 - Outgoing Values (results)
- No (re)assignment
- No loops
- · Return values depend only on input parameters
- Functions are first class values; this means they can:
 - Passed as arguments to a function
 - o Be returned from a function
 - o Construct new functions dynamically

Higher Order Function

Any function that takes a function as a parameter or returns a function as a result.

Function Currying

Transform a function with multiple arguments into multiple functions that each take exactly one argument.

Named after Haskell Brooks Curry.

Currying Example

addNums :: Num a => a -> a -> a

addNums is a function that takes in a number and returns a function that takes in another number.

map

- Built in Haskell higher order function
- . Applies a function to all elements of a list.

filter

- Built in Haskell higher order function
- · Removes all elements from a list that do not satisfy (i.e. make true) some predicate.

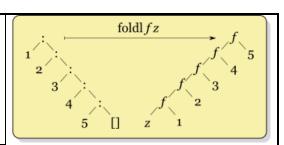
- Built in higher order function
- Does not support infinite lists.
- · Should only be used for special cases.

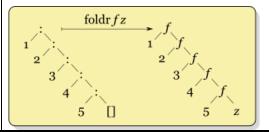
Example:

> foldl (
$$x y -> x - y$$
) 0 [1, 2, 3, 4] -10 -- (((0-1) - 2) - 3) - 4

- Built in higher order function
- Supports infinite lists.
- "Usually the right fold to use"

Example:





Thunk – A delayed computation

Due to lazy evaluation, foldl and foldr build thunks rather than calculate the results as they go.

foldl'

> foldr (x y -> x + y) 0 [1, 2, 3, 4]

- Data.List.foldl' evaluates its results eagerly (i.e. does not use thunks)
- Good for large, but finite lists.

foldl in terms of foldr

myFoldl' f acc x = foldr (flip f) acc (reverse x)

Lecture #05 – Small-Step Operational Semantics

WHILE Language

• Unlike the Bool* language, WHILE supports mutable references.

e ::= a	Variable/addresses
l v	Values
a:=e	Assignment
e;e	Sequence
e op e	Binary Operations
if e then e	Conditional
else e	
while (e) e	While Loops
v ::= i	Integers
b	Boolean
op ::= + - *	/

Small Step Semantics with State

• Since the WHILE language supports mutable references, the grammar must be updated to support it.

While Relation:

$$e, \sigma \rightarrow e', \sigma'$$

• σ – Store. Maps references to values.

Example Operations:

- $\sigma(a)$ Retrieves the value at address "a"
- $\sigma[a \coloneqq v]$ Identical to the original store with the exception that it now stores the value \boldsymbol{v} at address " \boldsymbol{a} "

Evaluation Order Rules

- Tend to be repetitive and clutter the semantics.
- Context based rules tend to represent the same information as evaluation order rules but more concisely.

Reduction Rule

Rewrites the expression. Example:

F-IfFalse:

if false then e2 else e3 \rightarrow e3

Context Rule

Specify the order for evaluating expressions. Example:

if e_1 then e_2 else $e_3 \rightarrow$ if e'_1 then e_2 else e_3

Reducible Expression (Redex) - Any expression that can be transformed (reduced) in one step.

Example: Redex

if true then (if true then false else false) else true

This reduces to "if true then false else false"

Example: Not a Redex

if (if true then false else false) then true else true

Not a redex as expression "if true then false else false" must be evaluated first.

Evaluation Contexts

- Alternative to evaluation order rules.
- Marker (•) / hole indicate the next place for evaluation (i.e. where we will do the work).

Example:

C[r]

= if (if true then false else false) then true else true

r = if true then false else false

C = **if** • **then** true else true

C[r] is the original expression.

Rewriting Evaluation Order Rules Context based rules only apply to reducible

expressions (redexs). Example:

 $C[if false then e_2 else e_3] \rightarrow C[e_3]$

Context Syntax

| if C then e else e | C op e | v op C

Data.Map

- Library: import Data.Map as Map
- Immutable
- Example Methods:
 - o Map. empty Creates and returns an empty map
 - o Map.insert k v m Inserts a value "v" at key "k" into map "m". Returns a new, updated map.
 - \circ Map.lookup k m Returns the value at key "k" in map "m". Wrapped in a Maybe.
 - Map.member k m Returns true if k is in map "m" and false otherwise.

Precondition - Text above the line in a rule.

Context Rule for Binary Op:

 $v_3 = v_1 \text{ op } v_2$ $C[v_1 \text{ op } v_2] \rightarrow C[v_3]$

How to Read a Small Step Semantic Rule: "Given < Precondition >, then <LeftSideArrow> evaluates to <RightSideArrow>."

Lecture #06 - LaTeX

TeX

- Created by Donald Knuth
- Domain specific language for typesetting documents.
- · Precisely controls the interface of content.
- Type of Literate Programming - Logic is in natural language and code is interspersed.
- "Mark code instead of marking comments."

LaTeX

- Developed by Leslie Lamport. Derives from TeX.
- Type of Domain Specific Language (DSL) A computer language that is specialized for a particular application domain.
- Enforces separation of concerns Design principle for separating a computer program into different sections, such that each section addresses a separate
 - o Example: LaTeX separates formatting from content.
- Literate Programming

Specify Document Type \documentclass{article}

Specify Title Block Content \title{Hello World!}

> **Start Document** \begin{document}

Generate Title from Title Information \title{Hello World!}

> **Close the Document** \end{document}

```
Cross-Reference
\ref{<referenceName>}
```

Reference a Bibliography Citation \cite{<citationName>}

Create a Reference \label{<referenceName>}

Create a Bibliography \bibliography{<bibFileName>}

> Create a List \begin{itemize} \item Text for #1 \item Text for #2 \end{itemize}

Create Section with Label \section{Section #1}

\label{sec:one}

Create Subsection with Label

\subsection{<SubsectionName>} \label{sec:<refName>}

Use of Tilde (~)

Creates an undividable space so the text "Section~\ref{sec:one}" will appear on one line

BibTeX

- · References are tedious to reformat and renumber.
- Reference details shorted in a "*.bib" file.

Create a Bibliography \bibliography{biblio}

BibTeX filename for the example would be "biblio.bib"

Define Bibliography Style \bibliographystyle{plainurl}

BibTeX Article Reference Example

```
@article{citationName,
   author = {Donald Knuth},
   title = {Literate Programming},
   journal = {},
   year = \{1984\},
   volume = {27},
   number = \{2\},
   pages = {97-111},
```

Lecture #07 – Types and Typeclasses

Maybe Type

· Example of an algebraic data type

- Enables behavior similar to null in Java
- Can be used to provide context.
- Used when:
 - o A function may not return a value
- o A caller may not pass an argument
- Definition:

```
data Maybe a = Nothing
             I Just a
```

Algebraic Data Type

A composite data

type (i.e. a type

Created via the

Keyword: data

made from other

Maybe "Divide" Example

```
divide :: Int -> Int -> Maybe Int
divide _ 0 = Nothing
divide x y = Just $ x 'div' y
> divide 5 2
2
> divide 4 0
Nothing
```

DO NOT FORGET THE Just IN CORRECT SOLUTION

Maybe Map Example

```
import Data.Map
m = Map.empty
m' = Map.insert "a" 42 m
case (Map.lookup "a") of
  Nothing -> error "Element not in map"
  Just x -> putStrLn $ show x
```

Since element may not be in the map, you need to use a maybe

Example Algebraic Data Type

data Tree k = EmptyTree | Node (Tree k) (Tree k) val deriving (Show)

k - Type parameter. Specifies a type not a value.

Node: Value Constructor that creates values of type "Tree k"

• Tree and Tree Int have no types since they themselves form a concrete

• Node does have a type:

> :t 3

```
> :t Node
Node :: (Tree k) \rightarrow (Tree k) \rightarrow k \rightarrow (Tree k)
```

Explanation: To make a complete Node object, you pass it two objects of type "Tree k" and another object of type "k" and that returns a "Tree k" object.

Partially Applying a Value Constructor

- Value constructors can be partially applied similar to functions.
- > let leaf = Node EmptyTree EmptyTree
- > Node (leaf 3) (leaf 7) 5

This creates a three node tree with value 5 at the root and values 3 and 7 at the leaves.

Type of the "+" Operator

```
> :t (+)
(+) :: (Num a) => a -> a -> a
```

Explanation: The plus sign takes two numbers of type "a" and returns an object of type "a".

```
Type of a Number
3 :: (Num \ a) => a
```

Explanation: Since "3" has no explicit type, it can for now be any type that satisfies the "Num" type class.

• Examples:

types).

- o Either
- o Maybe
- o Tree

```
Typeclasses
                            Kinds
                                                                                                            Example: Make Maybe an Instance of Eq

    Similar to interfaces in Java.

                                                                                                            instance (Eq a) => Eq (Maybe a) of

    Like a contract.

                                                                                                                   (==) Nothing Nothing = true
                                                                   o Implementation details can be included
                                        String Kind
                                                                                                                   (==) (Just x) (Just y) = x == y
                                                                     in typeclass definition.
                             > :kind String
                                                                                                                                                = false
• "The type of types".
                             String:: *
                                                                • No relation to classes in object-oriented
                                                                                                            Need to ensure type "a" supports "Eq" so add that as
• Concrete types have a kind
                                                                  programming.
                                        Map Kind
                                                                                                            a class constraint.
                            > :k Map
                                                                   o Example: Do not have any data
                            Map:: * -> * -> *
                                                                     associated with them.
• Keyword :k, :kind
                                                                                                            Class Constraint
                                       Maybe Kind
• Example:

    Simplify polymorphism.

                                                                                                            Operator: =>
                             > :k Maybe
                                                                                                            • Ensures that a type parameter satisfies some
                             Map:: * -> *
> :k Tree
                                                                Example: Eq Typeclass
                                                                                                              typeclass requirement.
Tree :: * -> *
                                     Map String Kind
                                                                class Eq a where
                                                                                                                           Kind of Typeclasses
                             > :kind (Map String)
Explanation: A Tree requires
                                                                     (==) :: a -> a -> Bool
                             (Map String) :: * -> *
one type parameter (e.g.
                                                                     (/=) :: a -> a -> Bool
                                                                                                            > :k Eq
Int) to be made a concrete
                                                                     x == y = not (x /= y)
                                                                                                            Eq :: * -> Constraint
                             Explanation: Map String is has one
                                                                     x \neq y = not (x == y)
type.
                             of the two type parameters filled so
                                                                                                            > :k Num
                             it has one less asterisk.
                                                                The last two lines in the type class definition
                                                                                                            Num :: * -> Constraint
                                                                allow the developer to program either (==) or
                                                                (/=) but not necessarily both.
                                                                                                            Note: Typeclasses are a class constaint (not a type)
                                                                                                            so their kind is different.
```

Lecture #08 – Functors

```
Functor – Something that can be mapped over.
                                                                                          Examples: map and fmap on Lists
        Functor Type Class Definition
                                             • Handles things "inside a box"
                                                                                                                           Examples: fmap on Maybes
                                                                                          > map (+1) [1, 2, 3]
class Functor f where
                                              Example: List ([]) as an instance of Functor
                                                                                          [2, 3, 4]
  fmap :: (a \rightarrow b) \rightarrow f a \rightarrow f b
                                                                                                                           > fmap (+1) (Just 3)
                                                                                                                           Just 4
                                                                                          > fmap (+1) [1, 2, 3]
                                             instance Functor [] where
This is very similar to the definition of the
                                                                                          [2, 3, 4]
                                                fmap = map
                                                                                                                           > fmap (+1) Nothing
higher order function "map"
                                                                                                                           Nothing
                                                                                          > fmap (+1) []
                                             Explanation: map is a specialized version of
map :: (a -> b) -> [a] -> [b]
                                                                                          []
                                             fmap for lists.
```

```
Example: Either as an Instance of Functor
                                                     Either Algebraic Data Type
Example: Maybe as an Instance of Functor
                                                                                       instance Functor (Either a) where
                                           data Either a b = Left a
                                                                                           fmap _ (Left x) = Left x
fmap f (Right y) = Right (f y)
                                                               | Right b
instance Functor Maybe where
                                                    deriving (Eq,Ord,Read,Show)
   fmap _ Nothing = Nothing
   fmap f (Just x) = Just (f x)
                                                                                       > fmap (+1) Leftt 20
                                           • Left - Error type that is not mappable.
                                                                                       20 -- No Change
DO NOT FORGET THE Just IN VALID SOLUTION
                                                                                       > fmap (+1) Right 20

    Right - Expected type

                                                                                       21 -- Changed
```

IO in Haskell

 Haskell avoids side effects but they are 	Type Signature of the main Function in	 do – Allows for the chaining of multiple 	
inevitable in real programs.	Haskell	IO/Monad commands together. Syntactic	
	main :: IO ()	sugar for bind ">>="	
• Monads			
 Related to Functors 	Hello World in Haskell	 <- Extracts data out of an IO/Monad 	
 Compartmentalize side effects. 	main = putStrLn "Hello World"	"Box"	
• ()	Type Signature of getLine	 return – Places data into an IO/Monad 	
 Unit type in Haskell 	getLine :: IO String	"Box"	

```
do Example
main = do
    line <- getLine</pre>
    if null line -- Checks for empty str
       then return ()
       else putStrLn $ reverseWords line
reverseWords :: String -> String
reverseWords = unwords .
               map reverse . words
```

```
return in Haskell
• Unrelated to "return" in other
  languages
```

• Better described as "wrap" or "box"

Summary:

return - Boxes an IO (since IO is a monad)

Unboxes an IO

```
Type of the Unit Type ()

    Base type

> :t ()
() :: ()
                     Type of return
> :t (return ())
(return ()) :: Monad m => m ()
```

Using IO as a Functor

```
main = do
       line <- fmap (++"!!!") getLine
       putStrLn line
```

Explanation: This function takes a string input from standard in and appends "!!!" at which point it prints it to the console.

Definition of IO as a Functor

```
instance Functor IO where
 fmap f action = do
                result <- action
                 return (f result)
```

Explanation: The action object is taken out of the IO box, the function "f" applied to it, and then returned to the IO box.

id Function

• Takes one input parameter and returns that input parameter unmodified. Examples:

```
> id 3
```

Monad is a typeclass.

> id "Hello World" "Hello World"

Functor Laws

Functor Law #1: If we map the id function over a Functor, the Functor that we get back should be the same as the original Functor.

```
Examples:
> fmap id (Just 3)
Just 3
> fmap id Nothing
```

Nothing > fmap id [1, 2, 3] [1, 2, 3]

Functor Law #2: Composing two functions and then mapping the resulting (composed) function over a Functor should be the same as first mapping one function over the Functor and then mapping the other one.

```
Law #2 Written Formally
fmap (f . g) = fmap f . fmap g
```

The Functor laws are NOT enforced. They are good practice that makes the code easier to reason about.

Lecture #09 – Applicative Functors

Functor - Something that can be mapped over. Allow you to map functions over different data types. Examples:

- Maybe
- Either
- IO
- Lists
- <*>

Functors return boxed up values.

Functor Example

```
> fmap (+1) [1, 2, 3]
[2, 3, 4]
> let x = fmap (+) [1, 2, 3]
```

Explanation: In this case x is: [(1+), (2+), (3+)]

Applicative Functor

• Requires the importing of a special library as shown below:

import Control.Applicative

Functions in Applicative Typeclass:

- pure Wraps/boxes a value
- <*> Infix version of fmap. Is itself a Functor.

Example Uses of pure > pure 7

> pure 7 :: Maybe Int Just 7

> pure 7 :: [Int]

Type Class Definition of Applicative

```
class (Functor f) => Applicative f where
     pure :: a -> f a
     <*> :: f (a -> b) -> f a -> f b
```

Only difference between <*> and fmap is that the function in <*> is boxed while it is not in fmap (see the green f).

Make Maybe an Instance of Applicative

```
instance Applicative Maybe where
     pure = Just
     Nothing <*> = Nothing
(Just f) <*> x = fmap f x
```

Explanation: pure simply wraps the value in ${\tt Just}$. No need to explicitly check if " ${\tt x}$ " is maybe as **fmap** will do that for you.

Examples of Applicative Maybe

```
> Just (+3) <*> Just 4
Just 7
> pure (+3) <*> Just 4
Just 7
> pure (+) <*> Just 3 <*> Just 4
Just 7
> (+) <$> Just 3 <*> Just 4
Just 7
```

Explanation: x <\$> is fmap as an infix operator. It is NOT necessarily the same as pure x <*>. It should be based off Applicative Functor Law #1.

Making [] an Instance of Applicative

```
instance Applicative [] where
  pure x = [x]
  fs <*> xs = [f x | f <- fs, x <- xs]
```

Explanation: The function is actually a list of functions so list comprehension is needed.

```
Example Use of Applicative on Lists
> (*) <$> [1, 2, 3] <*> [1,0,0,1]
[1,0,0,1,2,0,0,2,3,0,0,3]
```

```
> pure 7
7 -- No change
> pure 7 :: [Int]
[71
```

Definition of IO as an Instance of Applicative

```
instance Applicative IO where
   pure = return
    a <*> b = do
             f <- a
              x <- b
              return (f x)
```

Example of Applicative IO	11000
import Control.Applicative	A function that simplifies the application of a normal function to two Functors.
<pre>main = do a <- (++) <\$> getLine <*> getLine putStrLn a</pre>	<pre>liftA2 :: (Applicative f) => (a -> b -> c) -> f a -> f b -> fc liftA2 f x y = f <\$> a <*> b</pre>

Example of liftA2	Applicative Functor Definition
> (:) <\$> Just 3 <*> Just [4]	
Just [3, 4]	A functor you can apply to
> liftA2 (:) (Just 3) (Just [4])	The same apply to
Just [3, 4]	other Functors.

Applicative Functor Laws

<pre>Law 1: pure f <*> x = fmap f x</pre>	Law 2: pure id <*> v = v	Law 3: pure (.) <*> u <*> v <*> w = u <*> (v <*> w)
<pre>Law 4: pure f <*> pure x = pure (f x)</pre>	Law 5: u <*> pure y = pure (\$y) <*> u	Similar to Functor Laws, these are not strictly enforced but are good practice to make it easier to reason about the code.

Monoids

Monoid: An associative binary function and a value that acts as an identity with respect to that function.		Definition of Monoid Typeclass	
	•	class Monoid m where	
• x * 1 • lst ++ [] • x + 0	Examples Identity of Multiplication Identity of Concatenation Identity of Addition	<pre>mempty :: m mappend :: m -> m -> m mconcat :: [m] -> m mconcat = foldr mappend mempty</pre>	

Monoid Rules

Rule #1:	Rule #2:	Rule #3:
mempty `mappend` x = x	x `mappend` mempty = x	<pre>(x `mappend` y) `mappend` z = x `mappend` (y `mappend` z)</pre>

Lecture #10 - Monads

Problem with Functors: Do not support chaining of	Applicative Functor: A Functor that can be applied to other
multiple commands. Example:	Functors.
> fmap (+) (Just 3) (Just 4)	<pre>class (Functor f) => Applicative f where</pre>
	(<*>) :: f (a -> b) -> f a -> f b
Returns an error since it cannot resolve (Just 3+)	
and (Just 4)	Requires library Control. Applicative
	multiple commands. Example: > fmap (+) (Just 3) (Just 4) Returns an error since it cannot resolve (Just 3+)

```
Comparing <*> and >>=
                                                                       Example of <$>, <*> and >>=
Functor:
                                                              > (\x -> x + 1) < > Just 3
                                                                                                            Example: Implement applyMaybe that applies a
(<*>) :: Applicative f => f (a -> b) -> f a -> f b
                                                              Just 4
                                                                                                            function to a Maybe
Monad:
                                                              > Just (x -> x + 1) <*> Just(3)
(>>=) :: Monad m => m a -> (a \rightarrow m b) -> m b
                                                                                                            applyMaybe :: Maybe a -> (a -> b) -
                                                               Just 4
                                                                                                            > (Maybe b)
                                                                                                            applyMaybe Nothing _ = Nothing
applyMaybe (Just x) f = Just (f x)
Differences:
1. Order of the arguments changed.
                                                               > (Just 3) >>= (\x -> Just(x+1))
2. The function is boxed in Functor but not Monad
3. Monad function returns a boxed result.
                                                               Just 4
```

```
> (Just 3) `applyMaybe` (\_ -> Nothing)
                                                                                             • "Applicative Functors you can chain."
applyMaybe Nothing _ = Nothing
                                                         `applyMaybe` (\y -> Just (y-1))
applyMaybe (Just x)^{-}f = Just (f x)
                                             Nothing
          Monad Typeclass Definition
                                                                   Example a Robot Moving Towards a Goal (Not Failure)
                                                                                -- Define Operator and start location
                                                                               x -: f = f x
class Monad m where
                                             --Location
      return :: a -> m a
                                             type Robot = (Int, Int)
                                                                               start = (0, 0)
      (>>=) :: m a -> (a -> m b) -> m b
                                             -- Functions
                                                                               > start -: up -: right
      (>>) :: m a -> m b -> m b
                                             up (x,y) = (x, y+1)
                                                                               (1, 1)
      x \gg y = x \gg (\ -> y) --Lamda
                                             down (x,y) = (x, y-1)
                                             left (x,y) = (x-1, y)
                                                                               > start -: up -: left -: left -: right -: down
      fail :: String -> m a
                                             right (x,y) = (x+1, y)
      fail msg = error msg
```

Chaining applyMaybe

Additional Names for Monoids

• "Programmable Semicolons"

Just 5

Example: Implement applyMaybe that applies a

applyMaybe :: Maybe a -> (a -> Maybe b)

-> (Maybe b)

function to a Maybe

```
Example a Robot Moving Towards a Goal (with Failure)
                                      -- Once the goal is reached,
                                      -- the robot stops
                                      goal := Map.empty
                                                                                 start = (0, 0)
                                              -: (Map.insert (0, 2) True)
Maybe as an Instance of the Monad Typeclass
                                              -: (Map.insert (-1, 3) True)
                                              -: (Map.insert (-3, -8) True)
                                                                                 > return start >>= up >>= left >>= left
instance Monad Maybe where
                                                                                                >>= right >>= down
                                     moveTo :: Pos -> Maybe Pos
                                                                                 Just (-1, 0)
     return = Just
                                     moveTo p = if Map.member p goal
                                                                                 > return start >>= left >>= left >>= up
                                                       then Nothing
     (>>=) Nothing
                      = Nothing
                                                                                                else Just p
     (>>=) (Just x) \overline{f} = f x
                                                                                                >>= right >>= right >>= down
                                                                                 Nothing
                                      -- Since these are in bind, no need
     fail _
                      = Nothing
                                      -- to handle Nothing. Bind handles it.
                                     up(x,y) = moveTo(x, y+1)
                                                                                 Explanation: Reached one of the goals (-1, 3) at the red up
                                      down (x,y) = moveTo (x, y-1)
                                      left (x,y) = moveTo (x-1, y)
                                      right(x,y) = moveTo(x+1, y)
```

Integer Division Using Monads

```
Integer Division with Bind with "do"
                                                                                                            Integer Division with Bind with "do" and return
       Integer Division with Bind and No "do"
                                                    mydiv :: Maybe Int -> Maybe Int -> Maybe Int
                                                                                                        mydiv :: Maybe Int -> Maybe Int -> Maybe Int
mydiv :: Maybe Int -> Maybe Int -> Maybe Int
                                                    mydiv x y = do
                                                                                                         mydiv x y = do
mydiv x y = x >>= (\numer ->
                                                                 numer <- x
                                                                                                                      numer <- x
             y >>= (\denom ->
                                                                 denom <- y
                                                                                                                      denom <- y
             if denom > 0
                                                                 if denom > 0
                                                                                                                      if denom > 0
                 then Just (div numer denom)
else fail "Div by zero"))
                                                                       then Just (div numer denom)
                                                                                                                         then return $ div numer denom
                                                                       else fail "Div by 0"
                                                                                                                         else fail "Div by 0"
```

List Monad

```
Making List an Instance of Monad
                                                                Example Use of List as a Monad
instance Monad [] where
                                                        listOfTuples :: [(Int, Char)]
        return x = [x]
                                                                                                                Combining a Maybe and a List Monad
                                                        listOfTuples = do
         (>>=) xs f = concat(map f xs)
                                                                        n <- [1, 2]
        fail _
                     = []
                                                                                                         > Just [2,3] >>= (\x -> Just(fmap (+1) x))
                                                                        ch <- ['a', 'b']
                                                                                                         [3, 4]
                                                                        return (n, ch)
Explnation: concat is needed here as f returns elements
                                                        > listOfTuples
already in a list. As such, concat merges the individual lists
                                                        [(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]
(from each call to f) into a single list.
```

Lecture #11 – Parsing Combinators

Semantics: Enumerate what a program means. Defined by the interpreter or compiler. Compilation Flow Step #1: Tokenizer/lexer generates a set of tokens.		Converts the characters of the program into words of the language.	
	Step #2: Parser turns the tokens into an abstract syntax tree.	Examples:	
Syntax: Enumerate how a program Is structured. Defined by the lexer and parser.	Step #3: Compilers and interpreters convert the AST into machine code or commands respectively.	Lex/Flex (C/C++) ANTLR & JavaCC (Java) Parsec (Haskell)	

Lover

Categories of Tokens

- Reserved Words/Keywords.
 - o Examples: while, if, then, else
- Literals/Constants.
- o Examples: 123, "Hello World!"
- · Special symbols.
 - o Examples: ";", "=>", "&&"
- Identifiers.
 - o Examples: "balance", "myFunction"

Parsing

- · Parser converts tokens to abstract syntax trees.
- Defined by context free grammars (CFG)
- Types of Parsers:
 - o Bottom-up/Shift-Reduce Parsers
 - o Top-down parsers

Context Free Grammars

- · Grammars specify the language.
- Specified in Backus-Naur form format. Example:

```
Expr -> Number
    Number + Expr
```

- Terminal Cannot be broken down further.
- Non-terminals Can be broken down further.

Example: "0", "1", "2", ..., "9" are terminals but digit, number, and expression are not.

Example Grammar

```
expr -> expr + expr
        expr - expr
        (expr)
        number
number -> number digit
        | digit
digit -> 0 | 1 | 2 | ... | 9
```

Bottom-Up / Shift-Reduce Parser

- Shift tokens onto a stack
- Reduce the stack to a non-terminal.
- LR Left to right, Rightmost derivation
- LALR Look-Ahead LR parsers are the most popular type of LR parsers.
- o Examples: YACC/Bison
- · Fading from popularity

Top-Down Parser

- Non-terminals are expanded to match tokens.
- LL <u>Left</u> to right, <u>Leftmost derivation</u>
- LL(k) Parser Looks ahead up to k elements. **Examples:** Java CC, ANTLR
 - o The higher the k, the more difficult language is to parse. k can be arbitrary.
 - o LL(1) Easy to parse using either LL or recursive descent parsers. Many computer languages are designed to be LL(1).

Parser Combinator

Combine simpler parsers to make a more complex parser.

Example: Parsec

Useful Parsec Functions

- many Parses zero or more occurrences of the given parser.
- many1 Parses 1 or more occurrences of the given parser.
- noneOf Anything but the specified value
- spaces Whitespace characters
- **char** The specific specified character
- **string** The specific specified string.
- sepBy Separate tokens by some token.

```
import Text.ParserCombinators.Parsec
num :: GenParser st String
num = many1 digit
main = do
       print $ parse num "Hello" "42"
```

```
Example Parsec Code
import Text.ParserCombinators.Parsec
```

num :: GenParser st Integer str <- many1 digit return \$ read str

main = do

print \$ parse num "World" "42"

- st "State." Always required for our purposes.
- String/Integer Parser return type
- digit 0, 1, 2, 3, ..., 9 (terminal)
- num Parser entry function
- "Hello"/"World" Debug string.

```
Example with try, <|>, and <?>
```

```
eol = try (string "\n")
   <|> string "\n\r"
   <?> "end of line"
```

- try If an incomplete match is found, rewind.
- <|> "Or" Operator for matching tokens.
- <?> Otherwise with an accompanying error message.

• many1 - Select one of more digits.

• "42" - String to parse.

Practice Midterm and Review Notes

Question #1	Question #2	Question #2 Question #3		Question #5
a. True	a. True	a. False – Big step	a. False – Imperative	a. True
b. False – Lazy evaluation	b. False – Applicative functor	b. True	b. True	b. False – Typeclass
c. False – Lazy evaluation	c. True	c. False – Use store	c. False	c. True
d. False – Statically type	d. True	d. True	d. True	d. False
e. True	e. True	e. False	e. True	e. False – Algebraic data type

d. False – Statically type	l. True	d. Irue	d. True	d. False
e. True	. True	e. False	e. True	e. False – Algebraic data type
Haskell Purely Functional Lazy evaluation Fully Curried Language Statically Typed Type Inference – Via context, Haskell can deduce the type.	Purely Function Referential Transparency call can be replaced with i value without affecting th No (re)assignment No loop No side effects	A functionts equivalent	Functional Languages • Functions are first class objects meaning they can be passed to a function, returned from it, or created on the fly. • Higher order function support	Operational Semantics Small Step – Structural Semantics Big Step – Natural Semantics "Get stuck" – When a function is encountered that does not have an associated rule.

CSV Parser Example

```
Verbose Approach
import Text.ParserCombinator.Parsec
import System.Environment
csvFile :: GenParser st [[String]]
csvFile = do
          arr <- many line
          char eof
          return arr
line :: GenParser st [String]
line = do
       result <- many1 cell
       char '\n'
       return result
cells :: GenParser st [String]
cells = do
        firstCell <- cellContents
nextCells <- remainingCells</pre>
        return (firstCell:nextCells)
cellContent :: GenParser st String
cellContent = many $ noneOf ", \n" -- Two characters
remainingCells :: GenParser st [String]
remainingCells = do
                  (char "," >> cells)
                  <|> return []
main = do
       args <- getArgs
       p <- parseFromFile csvFile "example 1" (head args)
       case p of
           Left msg -> error msg
           Right csv -> print csv
```

Miscellaneous

Kind of Show and show > :k Show Show :: * -> Constraint Type and Kind of show > :k show Error (A function not a type) > :t show

```
Lambda and ADT Combined
> (\x -> Just (x+1)) 1
Just 2
        Creating Type Alias
```

type String = [Char]

Allows for more readable code as developer can use a type name that makes more sense for a given application.

```
Example: applyMaybe that takes a (Maybe a) and applies to it a function
            that takes a normal a and returns a (Maybe b)
```

```
applyMaybe :: (Maybe a) -> (a -> Maybe b) -> (Maybe b)
applyMaybe Nothing _ = Nothing
applyMaybe (Just x) f = f x
```

Explanation: Since the function "f" already returns a Maybe, you do not need to re-box it. However, since it does not take a Maybe, you need to unbox the first input parameter.

```
Applying return to Items
> return 7 :: Maybe Int
Just 7
> return 7 :: [Int]
```

show :: (Show a) => a -> String

List comprehension is syntactic sugar for using lists as monads.

Conclusion: Behavior for return is the same as pure. Both put the object in the minimum default context that still yields that value.

[7] -- Need Int or get an error

Monads and Lambda

When trying to chain multiple functions together in a Monad, remember the Monad must return a boxed value. Hence, Lambda often work well as they simplifying boxing.

Applicative Typeclass - Allows you to use normal functions on values that have a context (i.e. are inside a Functor).

Monad: Given a value of type, a, in a context, m, apply a function that takes a normal value of type a and returns a value in the context m.

(>>=) :: (Monad m) => m a -> (a -> m b) -> m b

Monads are just applicative functors that support bind (>>=).

Key Difference: Applicative functors support normal functions that take and return unboxed values while Monads return boxed values.

return - Monad equivalent of "pure" for Applicative Functors.

Cannot use fmap in the definition of a Monad since fmap returns a boxed value while the function of the Monad returns a boxed value. Hence, if you used fmap with a Monad, you would return a double boxed value.

Functor Definitions

```
TO
                                                          Maybe
                 Lists
                                                                                  instance Functor IO where
                                           instance Functor Maybe where
instance Functor [] where
                                                                                    fmap f a = do
                                                   Nothing = Nothing
                                             fmap
 fmap = map
                                                                                               x <- a
                                             fmap f (Just x) = Just (f x)
                                                                                               return (f x)
```

Applicative Functor Definitions

Lists				
	<pre>instance Applicative [] where pure x = [x]</pre>			
	(<*>) fs xs = [f x f <- fs, x <- xs]			

```
Maybe
instance Applicative Maybe where
 pure x = Just x
  (<*>) Nothing _ = Nothing
  (<*>) (Just f) x = fmap f x
```

```
instance Applicative IO where
  a <*> b = do
            f <- a
            x <- b
            return (f x)
```

Monad Definitions

```
instance Monad Maybe where
                                                                                    instance Monad IO where
instance Monad [] where
                                              return x = Just x
                                                                                        (>>=) a f = do
return x = [x]
                                              (>>=) Nothing
                                                              = Nothing
                                                                                                   x <- a
 (>>=) xs f = concat $ map f x
                                              (>>=) (Just x) \overline{f} = f x
                                                                                                   fх
 fail _ = []
                                              fail
                                                        = Nothing
                                                                                       fail s = ioerror (userError s)
```

Lecture #12 - Introduction to JavaScript

JavaScript

- Developed at Netscape by Brendan Eichs in 10 days
- · Originally named "Mocha"
- Syntax similar to Java

Multiparadigm JavaScript Supported programming paradigms:

- Imperative
- Functional
- Object-Oriented (through prototypes)

Where JavaScript is Run

- Client Side Versions
- o Runs on user machine
- Server-side Versions
- o IVM: Rhino & Nashorn ○ Node is

```
Example: Imperative JavaScript
sum += list[i];
```

Example: Functional JavaScript

```
function addList(list) {
  if(list.length == 0){
    return 0;
  return list[0]
addList(list.slice(1));
slice(begin[, end]) - Takes a subset of an
array from the "begin" index to the "end"
```

Example: Object-Oriented JavaScript

```
function Adder (amount) {
  this.amount = amount;
Adder.prototype.add = function(x){
            return this.amount + x;
var myAdder = new Adder(1)
var y = myAdder.add(7)
```

Adder - Name of a new constructor. Convention is to start constructors with a capital letter.

this - Not optional in JavaScript.

```
function addList(list) {
 var = i, sum = 0;
  for( i = 0; i < list.length ; i++){</pre>
 return sum;
```

Example: Functional JavaScript

var x = 42; // Create with var

y = 7; // No error without var

function noReturnAdd(a, b) {

var c = noReturnAdd(x, y)

//Lambda Function

function add(a, b){ return a + b;

a + b:

• Standard Approach:

elements to the end of the list.

```
console.log("...")
```

Printing to the Console in JavaScript

(exclusive). If no "end" is specified, it takes all

Not supported by all implementations.

- JVM based JavaScript Approach: print
- Solution to Support a Single Interface: var print = console.log

Closures

• Functions whose inner variables refer to independent (free) variables.

Closure Example

```
function getNextInt(){
 var nextInt = 0;
 return function(){
           return nextInt++;
       }() // Double paren
           // run the function
console.log(getNextInt()); // print "0"
console.log(getNextInt()); // print "1"
console.log(getNextInt()); // print "2"
```

var myLambda = function(x) {return x * x;} Node.js

- JavaScript runtime environment and library designed to run outside the browser.
- Based off Google's V8 engine.
- npm Package manager to get new packages.

// c is "undefined" since no return

Callback Function

- Functions in JavaScript are first class objects of type "Object".
- Not executed immediately.

JavaScript supports both "null" and "undefined"

Reading from a File with Callbacks in Node.js

```
var fs = require('fs')
fs.readFile('myFile.txt',
    function(err, data){
      if (err)
         throw err;
        console.log("" + data);
}
console.log("All done")
"All done" prints before the file contents due
to callbacks.
require - Includes the JavaScript package "fs"
```

Synchronous File IO in Node var fs = require('fs')

var data = fs.readFileSync('myFile.txt'); console.log("All done")

To eliminate callbacks, most function names can be appended with "Sync"

Undeclared Object Fields

Any undeclared object fields or uninstantiated variables are undefined.

```
var y; // Uninstantiated
// Both print 'undefined'
console.log(y)
console.log(myDog.name)
```

Creating a JavaScript Object

```
var myDog = {age : 3,
               weight: 100}
Every object is a map.
```

Adding a Field to a JavaScript Object

```
myDog['height'] = 45 // Add a new height field
                     // Note the single quotes
```

Adding a Function to a JavaScript Object's Prototype

```
myDog.speak = function() { console.log("Grr"); }
```

Delete a Function from a JavaScript Object's Prototype

delete myDog.speak

Prototypes

Object Prototypes

JavaScript prototypes are just like any other object.

```
var dogPrototype = {
    speak: function(){
      console.log("bark!");
```

Defining an Object's Prototype

```
var rex = { name: "Rex",
            __proto__ : dogPrototype}
```

Prototypical Inheritance: If an object does not have a method of field. JavaScript looks to the object's proto object.

```
Add a Special "speak" Method to Rex
                                                                                  Effect of the "new" Keyword
rex.speak = function(){
                                                          function Cat(name, breed) {
                                                            var this = {}; // Add when new is used
               console.log("Grr");
                                                             this.prototype = Cat.prototype; // Also comes from new
                                                            this.name = name;
rex.speak(); // Prints "Grr"
                                                            this.breed = breed;
                                                            this.speak = function() {console.log("meow");};
delete rex.speak;
                                                            return this; // Also comes from new
// Prints "Bark!" from __proto__
rex.speak();
                                                                                  No "return" in a Function
delete rex.speak; // Does nothing
rex.speak(); // Prints "Bark!"
                                                          function noReturnAdd(x, y) {
                                                            x + y; // without "return"
Unspecified Function Arguments: In JavaScript, any unspecified function
                                                          // c is "undefined" since no return
argument defaults to "undefined".
                                                          var c = noReturnAdd(x, y)
                                                          console.log(c); // Prints "undefined"
```

```
Top Prototypes
                                                                                   require
                                    Iterating Using "forEach"
                                                                                                               Running from the Command Line
                                                                     • Used to import an external module in
                             var arr = [1, 2, 3];
Object.prototype - Top
                             // Print each element in array
                                                                       Node.js
                                                                                                           • Use the keyword "node" for Node.js.
of all object prototypes
                             arr.forEach(function(val){
                                                                                                             Example:
                                             console.log(val);
                                                                     • Can be stored in a variable. Example:
                                           };
Function.prototype -
                                                                                                                 $ node my_program.js
Top of all function prototypes.
                                                                        var net = require('net');
      Create an Object with Methods without a Constructor
```

Lecture #13 - Lambda Calculus

Expressions

(Variables, immutable) e ::= x | λ x.e (Lambda abstraction)

(Function application)

Note: Lambda (λ) is simply a function. $v ::= (\lambda x.e)$ (Lambda abstraction)

Given a function where ${\bf E}$ is a

complex expression:

 $\lambda(x.E)v$

Function Application

Then:

 $\lambda(x.E)v \to E[x \vdash > v]$

Hence, "v" replaces "x" in "E".

Small-Step Evaluation Order Rules for Lambda Calculus

Rule: SS-E1	Rule: SS-E2	Rule: SS-Lambda Context	Optional Rule: Lazy SS-Lambda Context
$\boldsymbol{e_1} \rightarrow \boldsymbol{e_1'}$	$\boldsymbol{e}_2 \rightarrow \boldsymbol{e}_2'$		
$\overline{\boldsymbol{e_1}\boldsymbol{e_2} \rightarrow \boldsymbol{e_1'}\boldsymbol{e_2}}$	$(\lambda x. e) e_2 \rightarrow (\lambda x. e) e_2'$	$(\lambda x. e) \ v \to e[x \vdash > v]$	$(\lambda x. e) e_2 \to e[x \vdash > e_2]$

Lecture #14 - JavaScript Scoping

Example: First Class Function

```
function makeAdder(x) {
  return function(y) {
            return x + y;
            };
}
var addOne = makeAdder(1);
// Prints "11"
console.log(addOne(10));
```

Example: Function Application

```
function makeAdderList(arr) {
  var i;
  var output = [];
  for(i = 0; i < output.length; i++) {
     // Need to create a new scope
     function() {
          // Can add to arr without append
          output[i] = function(y) {
                return arr[i] + y;
                };
        }
    }
}</pre>
```

JavaScript lacks block scope for the closure to be right, must create the function inside another function.

Block Scope – The scope (i.e. visibility) of a variable is limited to a specific block (e.g., for loop, if statement, etc.).

- Unlike most languages, JavaScript does not have block scope.
- To create a new scope, use an anonymous function.

Variable Hoisting – All variable declarations (i.e., use of "var") are treated as if they are at the beginning of the function.

"this" in JavaScript

this – Refers to the scope where the function is called.

- In Normal Function Calls this refers to the global "this"
- Object Methods The object itself.
- Constructor (using "new") The newly created object.
- Exceptions: apply, call, and bind. Inline event handles on DOM elements

Any time a new function is created, the other "this" is no longer in scope

Execution Context

Consists of three part:

- A Variable Object Container for variables and functions.
- Scope Chain Variable object plus parent scopes
- Context Object this

Global Context

- Top Level Context
- Variable object is known as the "global object"
- this Refers to the global object.

Any variable declared without var is added to the global context.

Function Contexts

- Activation or Variable Objects which include:
 - o Arguments passed to the function
 - o A special arguments object
 - o Local variables

```
apply, bind, call Example
```

```
x = 3;
function foo(y) {
    console.log(this.x + y);
}
foo(100); // Prints "103"
// Array passed for args
foo.apply(null, [100]);
// Update the context
foo.apply({x:4}, [100]);
// No array needed
foo.call({x:4}, 100);
// Create a new function
var bf = foo.bind({x:5});
bf(100);
```

- apply Calls a function with the arguments passed as an array.
- call Calls a function with the arguments passed in comma separated.
- **bind** Used to create a new function with a custom context.

Lecture #14.5 - JSLint and TypeScript

Issues in JavaScript

- No block scope
- Forgetting var can lead to unexpected behavior since variables become global.
- Operator "==" is not transitive.
- Switch/case statements require "break"

```
JavaScript Automatically Inserts Semicolons
```

```
function makeObject () {
  return // Semicolon inserted here
    madeBy: 'Austin Tech. Sys.'
}
var o = makeObject();
```

console.log(o.madeBy); // error

```
Function "parseInt" can Yield Unexpected Results
// Drops the " tons"
console.log("what do you get? "
             + parseInt("16 tons"));
// Prints just "1" due to the "Oh"
parseInt("101");
```

Behavior of "typeOf" typeOf - Returns a string. May yield unexpected results. typeOf 5 // "number" typeOf "hi" // "string" typeOf NaN // "number"

typeOf null // "object"

```
Behavior of "typeOfChar"
```

typeChar - Returns a string. Classifies letters as "digits".

```
typeOfChar "5" // "digit"
typeOfChar "q" // "digit"
// "Other character"
typeOfChar " "
```

- A tool to write cleaner and safer JavaScript.
- Requires that "use strict" (with quotes) be added at the beginning of all
- · Performs static code analysis.
- Helps catch common programming errors by requiring:
- o Variables declared before they are used.
- o Semicolons are always used.
- o Double equals never used.
- Inspired by the "lint" tool

Benefits of Type Systems

- Tips for compilers
- Hints for IDEs
- Enforced documentation
- Prevent code with errors from running.

TypeScript

- Developed by Microsoft
- . Static type checker for JavaScript.
- A new "superset" language of JavaScript with:
 - Type annotations o Classes
- Compiles to JavaScript

Function Type Annotations in TypeScript

```
function greet(person: string) {
  console.log("Hello " +
person);
var user : string = "Vlad";
```

greet (user);

```
// Prints "Hello Vlad"
```

Types in TypeScript

```
• number
           (var pi : number = 3.14)
          (var b : boolean = true)
• boolean
• string
           (var greet : string = "hi")
• array
           (var lst : number[] = [1, 2])
• enum
• any (var a : any = 3; var b : any = "hi")
```

```
TypeScript Class
```

```
class Employee{
 name : string;
 salary : number;
 constructor(name : string, salary : number) {
    this.name = name;
   this.salary = salary;
 display() { console.log(this.name); }
var emp = new Employee('Jon', 50000);
emp.display();
```

Lecture #15 – Event-Based Programming and Cryptocurrencies

JavaScript Embedded in HTML

Create a button on a website that prints

Improved JavaScript in HTML

Adding an Event Listener

 If clicking a button should perform multiple functions, then an event listener should be used.

Removing an Event Listener

• Event listeners can be removed by function name. Example:

Events in JavaScript

· JavaScript is single threaded.

Types of Keys

• Private Key: Known

only by the owner

• Public Key: Known

by everyone

An event must be run to completion before the next event handler can run.

Event Emitter

• Import the "events" module using the syntax

```
var ee =
require('events').EventEmitter;
```

Used to create event via the keyword "on".
 Example:

```
ee.on('die', function() {
          console.log("Died");
     });
```

 Invoking (emitting) an event using the keyword "emit" Example:

Create a TCP Server in Node.js Using Event Listeners

```
var net = require('net');
var eol = require('os').EOL;
var srvr = net.createServer();

// Add an event listener
srvr.on('connection', function(client) {
    client.write('Hello there!' + eol);
    client.end();
});
srvr.listen(9000);
```

telnet – Used to connect to a TCP server on the command line.

127.0.0.1 - IP address of localhost

Cryptocurrencies

Digital Signature

Non-Repudiation – Involves associating actions or changes to a unique individual.

- Solution in Cryptocurrency: Digital signature.
- Procedure:
 - Step #1: Owner encrypts the message with his private key
 - Step #2: Use the public key to decrypt the message.
- Analogy: Enclosed Bulletin Board

Private Key Encryption

Used to transmit sensitive data to a specific recipient.

- Procedure:
- Step #1: A user encrypts his data using the recipient's public key.
- Step #2: The intended recipient decrypts the data using his private key.
- Analogy: A public mailbox. Anyone can put letters in, but only the mailman has the key to open the box.
- update Used to update the signature with the specified message contents. Each signature object can only be updated once.
- hex Specifies that the output should be in hexadecimal format.
- Sync Ensures that the file read is done immediately without relying on a callback.
- SHA "Secure Hash Algorithm"
- RSA Signature algorithm

Example: JavaScript Signer Example

```
var crypto = require('crypto');
var fs = require('fs');

// Constructor for a "Signer" object
function Signer(privKeyFile){
    this.privKey = fs.readFileSync(privKeyFile).toString('ascii');
}

// Add a "signMessage" function to the Signer prototype
Signer.prototype.signMessage = function(msgFileName){
    var msg = fs.readFileSync(msgFileName).toString('ascii');
    var sign = crypto.createSign('RSA-SHA256');

    return sign.update(msg).sign(this.privKey, 'hex');
}
```

Double Spending – Spend the same funds in multiple places.

Solutions to Prevent Double Spending:

- Centralized Authority Disadvantages include that the central authority would charge a fee and not everyone trusts central authorities.
- Decentralized Authority Broadcast transactions to everyone.

Ledger – Used to keep a **history of all transactions** and the funds held by all users.

```
Example: JavaScript Verifier Example
var crypto = require('crypto');
var fs = require('fs');
// Constructor for a "Verifier" object
function Verifier(publicKeyFile) {
  this.publicKey = fs.readFileSync(privKeyFile).toString('ascii');
}
// Add a "verifySignature" function to the Verifier prototype
Verifier.prototype.verifySignature = function (msgFileName,
                                                signature) {
            var msg = fs.readFileSync(msgFileName).toString('ascii');
            // Create a verifier
            var ver = crypto.createVerifier('RSA-SHA256').update(msg);
           // Verify signature matches the hash
            var legit = ver.verify(this.publicKey, signature, 'hex');
            return legit;
```

Bitcoin Mining

- Block Chain Defines the transaction history.
 Used to prevent double spending.
- Proof of Work Verification of the block chain.
- Miners hash transaction details plus a "proof" (i.e. nonce)
 Reward: New bitcoins are mined for the first to find a proof.
- Cost to *Derive* a Proof: 2^N where N is the number of the initial bits that must be "0" for the proof to be valid.
- Cost to Verify a Proof: A single hash
- Bitcoin protocol is designed to make mining more profitable than cheating.

Attributes of a Good Hash Function

Role of a Hash Function:	One Way: Given an output "y", it is	Collision Resistant: It is infeasible to find		
Compress arbitrary length inputs	infeasible to find an "x" such that:	any "x" and "y" such that:	Compression	Efficient
to small, fixed length outputs.	h(x) = y	h(x) == h(y)		

Lecture #17 - Macros and Sweet.js

Basic Compiler Structure with C-Style Macros Example: C Preprocessor Example Macros • Short for "macroinstruction" #define PT 3.14159 Source Expanded Tokens #define SWAP(a,b) {int tmp=a;a=b;b=tmp;} Code Code Lexer/ • Rule specifies how an input Processo Tokenizer sequence maps to a int main(void) { replacement sequence. int x = 4, y=5, diam = 7; double circum = diam * PI; SWAP (x,y) Macros in C • Performed by a preprocessor C Preprocessor Output **Abstract** Compiler **Syntax** • Rely on text substitution. int main(void) { Tree int x = 4, y = 5, diam = 7;• Embedded languages like PHP, double circum = diam * 3.14159; Ruby, etc. use a similar approach. {int tmp=x;x=y;y=tmp;} **Machine Code** Problem with C Macros (Input) **Macros** in JavaScript Hygienic Macro - Any macro whose expansion is guaranteed not · No standard macro system for JavaScript // Macro should be on one line to cause the accidental capture of identifiers. • Sweet.js has been gaining interest. #define SWAP(a,b) {int tmp=a; • Recently redesigned. a=b: **Syntactic Macros** b=tmp;} Sweet, is • Derive from Lisp since Lisp programs are essentially one big AST. Borrows concepts from Racket. int main(void){ • Source-to-source compiler (i.e., transpiler) for int x = 4, tmp = 5; Work at the level of abstract syntax trees. JavaScript. SWAP (x, tmp) • Examples of other JavaScript transpilers: Powerful by expensive. o TypeScript **Problem with C Macros (Output)** o CoffeeScript Hygiene easier to address at the AST level. o Dart (includes its own VM) int main(void){ int x = 4, tmp = 5; · Project backed by Mozilla • Essentially a source-to-source compiler. { int tmp = x; a = tmp;**Basic Compiler Structure with Syntactic Macros Invoking Sweet.js** tmp = tmp;

Abstract

Syntax

Tree

}

Hence, a variable name collision between

the two variables named "tmp". This is

known as "inadvertent variable capture"

```
Keywords in Sweet.js
                                                                     Concatenating Multiple Result Strings
            Writing a Swap Function in Sweet.js
                                                                                                                    • let - Create a Sweet.js
                                                         This function squares a set of input variables.
syntax swap = function(ctx){
                                                                                                                      variable
  let innerCtx = ctx.next().value().inner();
                                                         syntax square = function(){
  let first = innerCtx.next().value();
                                                           var innerCtx = ctx.next().value().inner();
                                                                                                                    • ctx.next().value() -
  // Eat the comma
                                                           // Start with empty results
  innerCtx.next(); // No need for "value()"
                                                                                                                      Get the next value from the
                                                           result = #``;
                                                                                                                      context
                                                           while(let stx of innerCtx) {
  // Get the second parameter
                                                             result =
  let second = innerCtx.next().value();
                                                                                                                    • #`...` - Used to define a
                                                                result.concat(#`${stx}=${stx}*${stx};`);
  return #`var tmp = ${first};
                                                                                                                      result string.
                                                              // Eat comma
             ${first} = ${second};
                                                              // Ignored if no comma present
             ${second} = tmp; ;
                                                                                                                    • concat - Used to combine
                                                              innerCtx.next();
                                                                                                                      two result strings.
swap(a, b); // Invokes the macro
                                                         }
                                                         square(a, b, c); // Invokes the macro
                                                                                                                    • let xxx of yyy-lterate
                                                                                                                      over a list of tokens.
Note #1: The returned string is preceded by a pound (#) sign and
                                                         Note #1: Use ".concat" to concatenate multiple result strings.
is enclosed in backticks (`).
                                                                                                                    • isIndentifier - Used to
                                                         Note #1: If a token is not present, ".next()" does not cause an
                                                                                                                      check if a Sweet.is variable
Note #2: Sweet.js variables are declared with "let".
                                                         error.
                                                                                                                      matches some string.
```

Macro

Expander

Abstract

Syntax

Tree

Parser

Interpretter

Interpretter

From command line:\$ sjs myfile.js -d out/

• Compiled files run normally (as shown below for Node):

\$ node out/myfile.js

let className

```
A Class in Sweet.js
```

```
Input JavaScript Code to be Parsed by Sweet.js
```

```
syntax class = function(ctx){
 let className = ctx.next().value();
 let bodyCtx = ctx.next().value().inner();
  // By default assume empty constructor
 let construct = #`function() { }`;
let result = #``;
 while( let item of bodyCtx ){
    // Check if constructor
   if(item.isIndenifier('constructor')){
      // Get arguments then function code
      construct = #`function ${className}
                    ${bodyCtx.next().value()}
                    ${bodyCtx.next().value()}`;
   else {
     // Add the function to the class prototype
     result = result.concat(
         #`${className}.prototype.${item} =
             function ${bodyCtx.next.value()}
             ${bodyCtx.next.value()}`);
    // Return the constructor and methods
   return construct.concat(result);
```

Lecture #18 – Simply Typed Lambda Calculus

Lecture #19 - Metaprogramming and JS Proxies

Metaprogramming: Writing programs that manipulate other programs.

• Proposed in ECMAScript 6 for JavaScript.

Terminology in Reflection

- Introspection: Ability to examine (but not modify) the structure of a program.
- Self-modification: Ability to modify the structure of a program.

```
Introspection
Ability to examine (but not modify) the structure of a program.
```

JavaScript Examples

```
Property Lookup
"x" in o; //o is an object

Iterate Over All Properties of an Object
```

```
Self-modification
Ability to modify the structure of a program.
```

JavaScript Examples

```
o["x"]; // Computed property
o.y = 42; // Add new property
delete o.y; // Delete property
// Reflected method call
O["m"].apply(null, [38]);
```

Proxies in JavaScript

- Metacircular Interpretation The language is able to understand its own language.
- Until recently, JavaScript did not support
- o Javascript proxies are intended to fix that.
- Node.js' implementation of proxies lags behind the standard.

intercession.

Proxies only exist for objects and functions.
 Proxies do not exist for primitives.

Proxies and Common Lisp

- Developed before object oriented languages were popular.
- Many libraries were create with non-standard OO systems.
- Common Lisp Object System (CLOS) – Standard object oriented system for Lisp.

Achieving Lisp Object Backwards Compatibility

Option #1: Rewrite all libraries using CLOS. Disadvantage

• Huge number of libraries.

for(prop in o){

}

// Do something

• Not feasible to rewrite them all.

Option #2: Make a complex API

- API difficult to understand.
- Systems had conflicting features.

Option #3: Keep API simple and modify object behavior to fit different systems.

• This approach relies on **metaobject protocols**.

Proxies and Handlers

- The behavior of a proxy is determined by traps specified in its handler (i.e., the metaobject).
- Trap Methods that intercept an operation.
- Handler The metaobject that specifies the details of the trap. The handler itself is usually a normal object.
- Using proxies in node requires a special flag "--harmony-proxies". Example:
- \$ node --harmony-proxies prog.js

Kinds of JavaScript Proxies

• Object Proxies - Defined with:

Proxy.create(handler, proto)

• Functions (with extra traps) - Defined with:

• Proxies do not exist for primitives.

Read Only Handler

Information Control – Share a reference to an object, but do allow it to be modified.
 Example: Reference to the DOM.

```
function ReadOnlyHandler(obj){
  delete : function(name) {
    return obj[name];
}
// rcvr can be ignored
  set : function(rcvr, name, val) {
    return true;
}
```

```
A Noop Proxy – All Operations Passed through Unchanged
```

```
function handlerMaker(obj){
  // Delete a property from an object
 delete : function(name) { return obj[name]; },
  // Check if object has the specified property
 has : function(name) { return name in obj;},
  // Check if object (not prototype chain) has property
 hasOwn : function(name) {return Object.property
                                   .hasOwn(obj, name);},
  // Get a property value
 get : function(name) { return obj[name]; },
 // Set a property value
 set : function(rcvr, name, val){ obj[name] = val; },
  // Get all properties of an object
 enumerate : function(){
   var props = [];
    var prop:
   for(prop in obj) { props.push(prop); }
   return props;
 },
  // Get all of the keys of an object
 keys: function() { return Object.keys(obj); }
}
```

Aspect Oriented Programming

- Some code not well organized into objects. Example:
 - Cross-cutting concern where code is spread throughout a program.
- Canonical Example: Logging Statements
 - o Littered throughout the code
 - Swapping out a logger requires massive code changes.
- Solution: Use a proxy

Lecture #20 – Introduction to Ruby

Influences of Ruby

- SmallTalk
- Everything is an object
- Blocks
- Metaprogramming
- Perl
 - Regular Expressions
 - o Function names

Ruby on Rails

- "Killer" app for Ruby
 Lightweight web
 framework
 - "Convention over configuration" – If use standard configuration, very little configuration required.
- Initial framework was PHP, but that was abandoned.

Basic Ruby Syntax

```
puts "Hello World"

a = [1, 2, 3]

m = { 'a' => "Apple",
    'b' => "Bear",
    'c' => "Cat" }

# Prints "1"
```

Prints "Apple"
puts a['m']

puts a[0]

- Keywords
- @ Represents an object property

Returning From a Function

- Every function in Ruby returns a value, even if return is not used.
- If no return is specified, a function returns the last used value

Basic Ruby Class

```
class Person
  # Constructor
 def initialize name # Parameter
    # Attribute
   @name = name
 end
 # Getter
 def name
   return @name
 end
  # Setter
 def name = newName
   @name = name
 end
 # Method
 def say_hi
   puts "Hi my name is #{@name}"
 end
```

```
Using Metaprogramming for Getters and Setters
class Person
  # Replaces getters and setters
  # Uses metaprogramming
  attr accessor :name
  # Constructor
  def initialize name
    # Attribute
    @name = name
  End
  # Method
  def say_hi
   puts "Hi my name is #{@name}"
  end
end
          Using a Class in Ruby
```

osing a class in Ruby

```
p = Person.new "Joe"
puts "Name is #{p.name}"
p.say_hi
#{...} - Embeds a variable in a Ruby String
```

Getters and Setters the Ruby Way

Relies on metaprogramming

- attr_reader Getter only
- attr writer-Setter only
- attr_accessor Getter and setter

Reopening a Class in Ruby

- Class definitions can be changed during runtime in Ruby.
- This is known as "reopening the class"

Parent Class

end

```
class Dog
# Parentheses optional
def initialize(name)
    @name = name
end
def speak
   puts "#{@name} says bark"
end
end
```

Inheritance in Ruby

```
class GuardDog < Dog
  attr_accessor :breed
  def initialize(name, breed)
    # Use parent constructor
    super(name)
    @breed = breed
end
  def attack
    puts "Grrr"
end
end</pre>
```

Child Class

Note: Inheritance is doing using the less than (<) operator.

Mixin

- Add features to a class
- Similar to interfaces in Java with the exception that they <u>can</u> include functionality.
- module Keyword to define a Mixin.
- include Keyword to include a Mixin into a class.

Blocks in Ruby

Superficially similar to blocks in other languages.

- Create custom control structures.
- Can be represented with curly brackets ({...}) or do/end.

File IO without Blocks

```
file = File.open('test.txt','r')
file.each_line do |line|
  puts line
end
file.close
```

Note #1: Similar "boilerplate" code of open and closing the file.

Note #2: It is possible one may forget to the close the file.

File IO without Blocks

```
File.open('test.txt','r') do |file|
  file.each_line { |line|
    puts line
  }
end
```

Note #1: Eliminates the "boilerplate" code.

Note #3: When using a block (both do/end, and curly brackets), surround the variable names in pipes (1).

Example: Mixin

```
# Define the mixin
module RevString
def to_rev_s
    # Object is implicit
    to_s.reverse
end
end

# Reopen the Person Class
class Person include RevSting
def to_s
    # Returns the value
    @name
end
end
```

Dynamic Code Evaluation (eval)

- Executes source code dynamically
 - Code passed as either a string (or a block of code)
- Popular feature in JavaScript
 - Early usage was to convert JSON strings to variables since not supported by JavaScript.
- Source of security concerns.

Additional Ruby eval Methods

- instance_eval Evaluates code within an object's body.
 - o Access the internals of an object.
- class_eval Evaluates code
 within a class' body.
- $\,\circ\,$ Modifies the class' definition.
- Takes either a string or block of code. Block of code is more secure.

Example: Use instance_eval to Change an Object's Value

```
# Create with the name Bob
bob = Person.new "Bob"
```

- # Change his name
 bob.instance_eval do
 @name = "Steve"
 end
- # Prints "Steve"
 puts bob.name

Regular Expressions in Ruby

- sub Replaces the first instance of a string match.
- To perform the modification in place, must include an exclamation point (!) after sub.
- gsub Replaces all instance of a string match.
 To perform the modification in place, must include an exclamation point (!) after sub.

```
Example: class eval in Ruby
                                                                           Example: Using Regular Expressions in Ruby
# Applies to all classes
class Class
                                                         s = "Hi, I'm Larry; this is my" +
  # Simulate the "attr accessor" function
                                                              " brother Darryl, and this" +
  def my attr accessor(args)
                                                              " is my other brother Darryl."
    args.each do |prop|
                                                         s.sub(/Larry/,'Laurent')
      # Create getter
      self.class_eval("def #{prop};
                                                         # Prints s unchanged
                         return @#{prop};
                        end")
      # Create setter
                                                         # Changes first "Larry" to "Laurent"
      self.class eval("def #{prop} = v;
                         @\#\{prop\} = v;
                                                         s.sub! (/Larry/, 'Laurent')
                       end")
    end
                                                         # Prints first "brother" replaced with
  end
                                                          # "frere". s is unchanged, bt it did
end
                                                         # return the modified string.
# Use the new attribute
                                                         puts s.sub(/brother/, 'frère')
class Musician
 my attr accessor :name, :genre
                                                          # Same as previous except all where
end
                                                         # changed when printing.
                                                         puts s.gsub(/brother/, 'frère')
m = Musician.new
m.name = "Bob Marley"
puts m.name # Prints "Bob Marley"
```

Regular Expression Symbols in Ruby

/ . / - Any character except a newline	/\w/ – Any word character: [a-zA-Z0-9_]	/\d/ – Any digit character: [0-9]	/\w/ - Any whitespace character: [\t\r\n\f]
	/\w/ – Any non-word character:	/\D/ – Any non-digit character:	/\w/- Any non-whitespace character:
	[^a-zA-Z0-9_]	[^0-9]	[^ \t\r\n\f]
* – Zero or more times	+ – One or more times	? – Zero or one times (optional)	

```
File.open('test.txt','r') do |file|
  file.each line { |line| puts line }
end
puts "Name is #{p.name}"
class GuardDog < Dog
  attr_accessor :breed
  def initialize (name, breed)
    super (name)
    @breed = breed
  end
  def attack
   puts "Grrr"
  end
end
p.say_hi
a = [1, 2, 3]
m = { 'a' => "Apple",
      'b' => "Bear",
      'c' => "Cat"}
// Prints "1"
puts a[0]
class Dog
  # Parentheses optional
  def initialize(name)
    @name = name
  end
```

```
def speak
   puts "#{@name} says bark"
  end
end
// Prints "Apple"
puts a['m']
o.y = 42; // Add new property
delete o.y; // Delete a property
// Reflected method call
O["m"].apply(null, [38]);
  // By default assume empty constructor
  let construct = #`function() { }`;
  let result = #``;
  while( let item of bodyCtx ){
    // Check if constructor
    if(item.isIndenifier('constructor')){
     // Get arguments then function code
     construct = #`function ${className}
                    ${bodyCtx.next().value()}
                    ${bodyCtx.next().value()};
    else {
     // Add the function to the class prototype
     result = result.concat(
         #`${className}.prototype.${item} =
             function ${bodyCtx.next.value()}
             ${bodyCtx.next.value()}`);
    // Return the constructor and methods
    return construct.concat(result);
});
srvr.listen(9000);
var Droid = {
  speak: function() {
            console.log("I am "
                        + this.name);
        },
  create: function(name) {
              var clone = Object.create(this);
              clone.name = name;
              return clone;
        } ,
};
function makeAdder(x){
 return function(y){
           return x + y;
var addOne = makeAdder(1);
// Prints "3"
console.log(addOne(2));
function Cat(name, breed) {
  var this = {}; // Done when new keyword is used
  this.prototype = Cat.prototype; // Also comes from new
  this.name = name;
  this.breed = breed;
  this.speak = function() {console.log("meow");};
  return this; // Also comes from new
}
//Lambda Function
var myLambda = function(x) { return x * x; }
```