Follow the leader with Dropout perturbations - Additive versus multiplicative noise

Manfred K. Warmuth

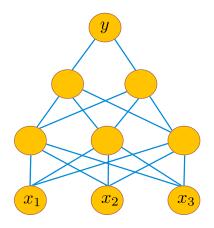
Simon's Institute Nov. 17, 2016

Joint work with Tim Van Erven and Wojciech Kotłowski

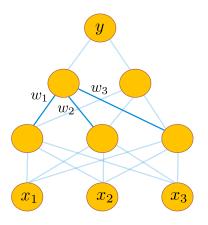
Major insights from [Devroye, Lugosi, Neu 2013]

Outline

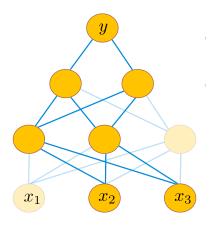
Feed forward neural net



Weights parameters - sigmoids at internal nodes



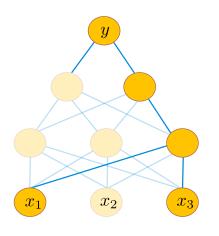
Dropout training



- Stochastic gradient descent
- Randomly remove every hidden/input node with prob. ½ before each gradient descent update

[Hinton et al. 2012]

Dropout training



- Very successful in image recognition & speech recognition
- Why does it work?

[Wagner, Wang, Liang 2013] [Helmbold, Long 2014]

What are we doing?

Prove bounds for dropout

- single neuron
- linear loss

Outline

	E_1	E_2	E_3	 E_n	prediction	label	loss
day 1	0	1	0	 0	0	1	1

	$ E_1 $	E_2	E_3	 E_n	prediction	label	loss
day 1	0	1	0	 0	0	1	1
day 2	1	1	0	 0	1	1	0

	E_1	E_2	E_3		E_n	prediction	label	loss
day 1	0	1	0		0	0	1	1
day 1 day 2	1	1	0	• • •	0	1	1	0
notation	x_1	x_1	x_2		x_n	\widehat{y}	y	$ \widehat{y} - y $

	$ E_1 $	E_2	E_3	 E_n	prediction	label	loss
day 1	0	1	0	 0	0	1	1
day 1 day 2	1	1	0	 0	1	1	0
notation	x_1	x_1	x_2	 x_n	\widehat{y}	y	$ \widehat{y} - y $
scope	$\in [0,1]$				$\in [0, 1]$	$\in \{0,1\}$	$\in [0, 1]$

	E_1	E_2	E_3	 E_n	prediction	label	loss
day 1	0	1	0	 0	0	1	1
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notation scope	$x_1 \in [0,1]$	x_1	x_2	 x_n	$\widehat{y} \in [0,1]$	$y \in \{0,1\}$	

- lacktriangle Algorithm maintains probability vector $m{w}$:
 - prediction $\widehat{y} = \boldsymbol{w} \cdot \boldsymbol{x}$

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day 1	0	1	0	 0	0	1	1
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notation scope	$\begin{vmatrix} x_1 \\ \in [0,1] \end{vmatrix}$	x_1	x_2	 x_n	$\widehat{y} \in [0,1]$	$y \in \{0,1\}$	$ \widehat{y} - y \in [0, 1]$

- lacksquare Algorithm maintains probability vector $oldsymbol{w}$:
 - prediction $\widehat{y} = {m w} \cdot {m x}$
- Loss linear because label $y \in \{0, 1\}$

Outline

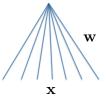
Predicting with expert advice

$$\hat{y} = \mathbf{w} \cdot \mathbf{x}$$
 loss $|\hat{y} - y|$



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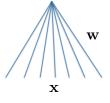


trial t

- get advice vector $oldsymbol{x}_t$
- predict $\widehat{y}_t = oldsymbol{w}_t \cdot oldsymbol{x}_t$
- get label y_t
- exp. losses $|x_{t,i} y_t|$
- alg. loss $|\widehat{y}_t y_t|$
- update $oldsymbol{w}_t
 ightarrow oldsymbol{w}_{t+1}$

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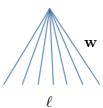


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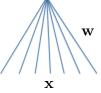
Hedge setting

loss $\mathbf{w} \cdot \ell$



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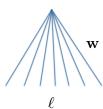


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- update w_t \to w_{t+1}
```

weights are implicit

Only works for linear loss

How do we measure performance

Worst-case regret

$$\sum_{t=1}^T oldsymbol{w}_t \cdot oldsymbol{\ell}_t \qquad - \inf_i oldsymbol{\ell}_{\leq T,i}$$
total expected loss of alg

Should be logarithmic in # of experts n

Outline

	E_1	E_2	E_3	E_4	E_5
	0	1	0	0	1
	1	1	0	1	1
$day\ t-1$	0	0	1	1	1
$\ell <_{t-1,i}$	1	2	1	2	3

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	0	1	0	0	1
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FL
$$\hat{i}_t = \operatorname{argmin}_i \ell_{\leq t-1,i}$$

ties broken uniformly

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$$\operatorname{Hedge}(\eta) \quad w_i = \frac{e^{-\eta \ell_{\leq t-1,i}}}{Z} \qquad \qquad \operatorname{Weighted Majority all}$$

Weighted Majority algorithm for pred. with Expert Advice Soft min

Dropout

	E_1	E_2	E_3	E_4	E_5
	0	1	0	0	1
	1	1	0	1	1
$day\ t-1$	0	0	1	1	1

$$\frac{\widehat{\ell}_{\leq t-1,i}}{\widehat{\ell}_{\leq t-1,i}}$$

Dropout

	E_1	E_2	E_3	E_4	E_5
	0	1	0	0	1
	1	1	0	1	1
$day\ t-1$	0	0	1	1	1
$\widehat{\ell}_{\leq t-1,i}$	1	1	0	1	2

$$\widehat{\ell}_{t,i} = \beta_{t,i} \ell_{t,i},$$

 $\widehat{\ell}_{t,i} = eta_{t,i} \ell_{t,i}, \qquad$ where $eta_{t,i}$ iid Bernoilli



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$$\begin{array}{ll} \mathsf{FL} \ \mathsf{on} \\ \mathsf{dropout} \end{array} \quad \widehat{i}_t = \underset{i}{\operatorname{argmin}} \ \widehat{\boldsymbol{\ell}}_{\leq t-1,i} \end{array}$$

How good?

Optimal worst case regret: $\sqrt{L^* \ln n} + \ln n$

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- In the meantime
 - new fancy algorithms by Haipeng Luo, Rob Schapire $\&\ {\sf Tim}\ {\sf van}\ {\sf Erven},$ Wouter Koolen
 - also no tuning, many other advantages

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- lacksquare Prob. vectors $oldsymbol{w}_t$ \longrightarrow density matrices $oldsymbol{\mathrm{W}}_t$
- lacksquare Hedge $w_{t,i} = rac{e^{-\eta\ell} \leq t-1,i}{Z} \longrightarrow lacksquare$ Matrix Hedge

$$\mathbf{W}_t = \frac{\exp\left(-\eta \mathbf{L}_{\leq t-1}\right)}{Z'}$$

■ Matrix Hedge $O(n^3)$ per update

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- Proof techniques break down
 - settled for vector case and independent multiplicative noise
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- Follow the skipping leader has linear regret [Lugosi, Neu2014]

What regularization?

 $\mathsf{Hedge}(\eta) \qquad \qquad \mathsf{relative} \ \mathsf{entropy}$

What regularization?

 $\begin{array}{ll} \operatorname{Hedge}(\eta) & \operatorname{relative\ entropy} \\ \operatorname{FPL}(\eta) & \operatorname{additive\ } \frac{1}{\eta} \ \log \ \operatorname{exponential\ noise} = \operatorname{Hedge}(\eta) \end{array}$

What regularization?

```
\mathsf{Hedge}(\eta) relative entropy
```

 $\mathsf{FPL}(\eta)$ additive $\frac{1}{\eta}$ log exponential noise = $\mathsf{Hedge}(\eta)$

FL on dropout tricky

Feed forward NN Logistic regression Linear loss case [Wagner, Wang, Liang 2013] [Helmbold, Long 2014] [ALST 2014]

Outline

Any deterministic alg. (such as FL) has huge regret

- \blacksquare For T trials: give algorithm's expert a unit of loss
- Loss of alg.: T loss of best: $\leq \frac{T}{n}$

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Recall optimum regret: $\sqrt{L^* \ln n} + \ln n$

FL with random ties

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FL with random ties

- Give every expert one unit of loss
 - iterate $L^* + 1$ times
- Loss per sweep $\frac{1}{n} + \frac{1}{n-1} + \ldots + \frac{1}{2} + 1 \approx \ln n$
- Loss of alg.: $(L^* + 1) \ln n$ loss of best: L^* regret: $L^* \ln n$

Our analysis of dropout

Unit rule

Adversary forces more regret by splitting loss vectors into units

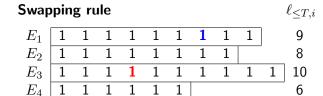
$$\begin{pmatrix} \mathbf{1} \\ 0 \\ \mathbf{1} \\ 1 \end{pmatrix} \longrightarrow \begin{pmatrix} \mathbf{1} \\ 0 \\ 0 \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ \mathbf{1} \\ 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

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Swapping rule

$$\ell_{\leq T,i}$$

E_1	1	1	1	1	1	1	1	1	1		9
E_2	1	1	1	1	1	1	1	1			8
E_3	1	1	1	1	1	1	1	1	1	1	10
E_4	1	1	1	1	1	1					6

- 1's occur in some order
- Worst case: 1 before 1
- Otherwise adversary benefits from swapping

Worst-case pattern

Assume we have s leaders

Assume we have s leaders

$$s$$
 leader get unit ignore non-leaders
$$\left\{ \begin{array}{c} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{array} \right.$$

Assume we have s leaders

$$s$$
 leader get unit ignore non-leaders
$$\begin{cases} & 1 \\ & 1 \\ & 1 \\ & 1 \end{cases}$$

FL

$$\frac{1}{s} + \frac{1}{s-1} + \frac{1}{s-2} + \frac{1}{s-3} + \ldots + \underbrace{\frac{1}{s-s-2}}_{2} + \underbrace{\frac{1}{s-s-1}}_{1}$$

 $\approx \ln s$

Assume we have s leaders

$$s$$
 leader get unit ignore non-leaders
$$\begin{cases} 1\\1\\1\\1\\1 \end{cases}$$

FL

$$\frac{1}{s} + \frac{1}{s-1} + \frac{1}{s-2} + \frac{1}{s-3} + \dots + \underbrace{\frac{1}{s-s-2}}_{2} + \underbrace{\frac{1}{s-s-1}}_{1}$$

$$\approx \ln s$$

Dropout

$$\frac{1}{s} + \frac{1}{s - 1/2} + \frac{1}{s - 2/2} + \frac{1}{s - 3/2} + \ldots + \frac{1}{s - (s - 2)/2} + \frac{1}{s - (s - 1)/2}$$

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$$\approx \ln s$$

Dropout

$$\frac{2}{2s} + \frac{2}{2s-1} + \frac{2}{2s-2} + \frac{2}{s-3} + \dots + \frac{2}{2s-(s-2)} + \frac{2}{2s-(s-1)}$$

$$\approx 2\left(\ln 2s - \ln s\right) = 2\ln 2$$

Assume we have s leaders

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$$\begin{cases} & 1 \\ & 1 \\ & 1 \\ & 1 \end{cases}$$

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$$\approx 2 \left(\ln 2s - \ln s \right) = 2 \ln 2$$

$L^* = 0$ - one expert incurs no loss

FL

One sweep

$$\frac{1}{n} + \frac{1}{n-1} + \ldots + \frac{1}{2} + 1 \approx (\ln n) - 1$$

Optimal

$L^* = 0$ - one expert incurs no loss

FL

One sweep

$$\frac{1}{n} + \frac{1}{n-1} + \ldots + \frac{1}{2} + 1 \approx (\ln n) - 1$$

Optimal

Dropout

- # of leaders reduced by half in each sweep

 $2 \ln n$

Overview of proof for noisy case

- \blacksquare Focus on first L sweeps
- lacksquare Only occurs constant regret if number of leaders >1

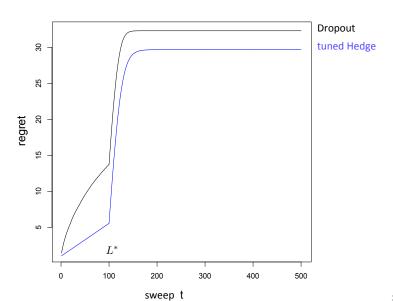
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- lacksquare Prob. that number of leaders >1 is at most $\sqrt{\frac{\ln n}{q+1}}$ for sweep q

Overview of proof for noisy case

- lacktriangle Focus on first L sweeps
- $lue{}$ Only occurs constant regret if number of leaders >1
- \blacksquare Prob. that number of leaders >1 is at most $\sqrt{\frac{\ln n}{q+1}}$ for sweep q
- For $\operatorname{Hedge}(\eta)$ and $\operatorname{FPL}(\eta)$ cost per sweep constant and dependent on η

Dropout versus Hedge



Outlook

- Combinatorial experts
- Matrix case
- Where else can dropout perturbations be used?
- Dropout for convex losses
- Dropout for neural nets

Outlook

- Combinatorial experts
- Matrix case
- Where else can dropout perturbations be used?
- Dropout for convex losses
- Dropout for neural nets
- Privacy

[Lugosi, Neu 2014] dense counter example

$$\begin{array}{c|cc}
0 & 1^* \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
1 & 0 \\
\hline
1 & 0 \\
\frac{n-1}{n} & \frac{1}{2}
\end{array}$$

Iterate this pattern n times:

$$\sum_{i=1}^{n} \left(\frac{n-i}{n-i+1} + \frac{1}{2} \right)$$

$$\approx n - \ln n + \frac{n}{2}$$

 $L^* = n$: Follow the Scipping Leader has linear regret

How does dropout ovoid this example?

0	1
1	0
1	0
1	0
1	0
1	0
n-1	1
\overline{n}	$\frac{n-1}{2}$

How does dropout ovoid this example?

0	1
1	0
1	0
1	0
1	0
1	0
$\frac{n-1}{n}$	$\frac{1}{\frac{n-1}{2}}$

It leaves the adversary clueless as to who the leader is i.e. privacy against adversary

sparse counter example

$$\begin{array}{|c|c|c|} \hline 0 & 1^* \\ \frac{1}{n-1} & 0 \\ \hline \frac{n-1}{n^2} & \frac{1}{2} \\ \hline \end{array}$$

Iterate this pattern n times:

$$\sum_{i=1}^{n} \left(\frac{n-i}{(n-i+1)^2} + \frac{1}{2} \right)$$

$$= \sum_{i=1}^{n} \left(\frac{1}{n-i+1} - \frac{1}{(n-i+1)^2} + \frac{1}{2} \right)$$

$$\approx \ln n - O(1) + \frac{n}{2}$$

 $L^* = \ln n$: Follow the Scipping Leader has linear regret