

Convex problems

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Overview

- 1 Convex optimization
- 2 Basic properties
- 3 Linear optimization
- 4 Quadratic optimization

Optimization in standard form

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p\end{array}$$

- $x \in \mathbb{R}^n$: optimization variables
- $f_0 : \mathbb{R}^n \rightarrow \mathbb{R}$: objective function
- $f_i : \mathbb{R}^n \rightarrow \mathbb{R}$: inequality constraint functions
- $h_j : \mathbb{R}^n \rightarrow \mathbb{R}$: equality constraint functions

Optimization in standard form

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p \end{array}$$

optimal value

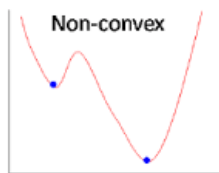
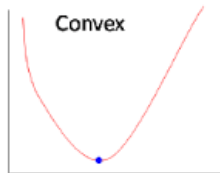
$$p^* = \inf \{ f_0(x) \mid f_i(x) \leq 0, h_j(x) = 0 \}$$

- $p^* = \infty$ problem is infeasible
- $p^* = -\infty$: problem is unbounded below

Optimal and locally optimal points

- x is feasible if $x \in \text{dom}(f_0)$ and it satisfies the constraints
- a feasible x is optimal if $f_0(x) = p^*$; X_{opt} is the set of optimal points
- x is locally optimal if there is an $\epsilon > 0$ such that x is optimal for

$$\begin{aligned}
 & \text{minimize} && f_0(z) \\
 & \text{subject to} && f_i(z) \leq 0, \quad i = 1, \dots, m \\
 & && h_j(z) = 0, \quad j = 1, \dots, p \\
 & && \|z - x\|_2 \leq \epsilon
 \end{aligned}$$



Examples

- $f_0(x) = 1/x$, $\text{dom}(f_0) = \mathbb{R}_{++}$

$$p^* = 0, \quad \text{no optimal } x$$

- $f_0(x) = x \log x$, $\text{dom}(f_0) = \mathbb{R}_{++}$

$$p^* = -1/e, \quad \text{optimal } x = 1/e$$

- $f_0(x) = x^3 - 3x$

$$p^* = -\infty, \quad \text{local optimal } x = 1$$

Implicit constraints

Above standard form optimization problem has an implicit constraints

$$x \in \mathcal{D} \doteq \cap \text{dom}(f_i) \cap \text{dom}(h_j)$$

- \mathcal{D} : domain of the problem
- f_i and h_j are the explicit constraints
- a problem is unconstrained if it has no explicit constraints

$$\text{minimize} \quad f_0(x) = - \sum \log(b_i - a_i^T x)$$

is an unconstrained problem with implicit constraint $a_i^T x \leq b_i$

Feasibility problem

$$\begin{array}{ll} \text{find} & x \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p \end{array}$$

can be considered a special case with $f_0(x) = 0$

$$\begin{array}{ll} \text{minimize} & 0 \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_j(x) = 0, \quad j = 1, \dots, p \end{array}$$

- $p^* = 0$ if constraints are feasible, any feasible x is optimal
- $p^* = \infty$ if constraints are infeasible

Convex optimization problem

Standard form convex optimization problem

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && f_i(x) \leq 0, \quad i = 1, \dots, m \\ & && a_j^T x = b_j, \quad j = 1, \dots, p \end{aligned}$$

- f_0 and f_i are convex,; equality constraints are convex
- feasible set of a convex optimization problem is convex
- equivalent form

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && f_i(x) \leq 0, \quad i = 1, \dots, m \\ & && Ax = b \end{aligned}$$

Example

$$\begin{aligned}
 &\text{minimize} && f_0(x) = x_1^2 + x_2^2 \\
 &\text{subject to} && f_1(x) = x_1/(1 + x_2^2) \leq 0 \\
 &&& h_1(x) = (x_1 + x_2)^2 = 0
 \end{aligned}$$

- f_0 is convex; f_1 is not convex, h_1 is not affine
- feasible set $\{(x_1, x_2) \mid x_1 = x_2 \leq 0\}$ is convex
- equivalent convex form

$$\begin{aligned}
 &\text{minimize} && x_1^2 + x_2^2 \\
 &\text{subject to} && x_1 \leq 0 \\
 &&& x_1 + x_2 = 0
 \end{aligned}$$

Local is global

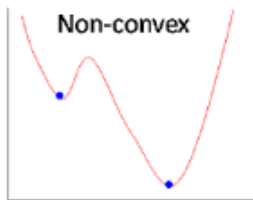
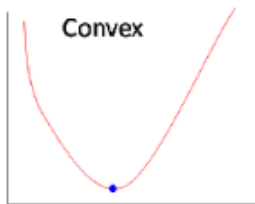
Any local optimal point of a convex problem is globally optimal.

- local minimizer

$$f(x^*) \leq f(y), \quad \forall y \in (x^* - \epsilon, x^* + \epsilon)$$

- global minimizer

$$f(x^*) \leq f(y), \quad \forall y \neq x$$



Local is global — proof

- let x be locally optimal, then there exists ϵ such that $f_0(z) \geq f_0(x)$ holds for any $\|x - z\|_2 \leq \epsilon$
- suppose there exists a feasible y with $f_0(y) < f_0(x)$
- consider $z = \alpha y + (1 - \alpha)x$ with $\alpha = \frac{\epsilon}{2\|y - x\|_2}$
- as $\|y - x\|_2 > \epsilon$, then $0 < \alpha < 1/2 < 1$
- z is convex combination of two feasible points, hence feasible
- then $\|z - x\|_2 = \alpha \|y - x\|_2 = \epsilon/2$ and

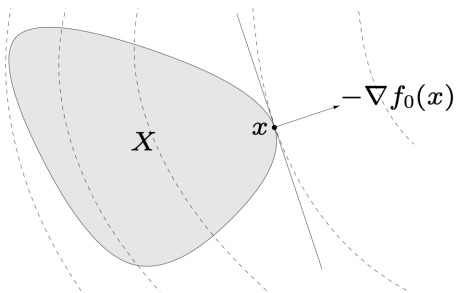
$$f_0(z) \leq \alpha f_0(y) + (1 - \alpha)f_0(x) < f_0(x)$$

- contradicts our assumption that x is locally optimal

Optimality condition for differentiable f_0

x is optimal if and only if it is feasible and

$$\nabla f_0(x)^T (y - x) \geq 0, \quad \forall \text{ feasible } y$$

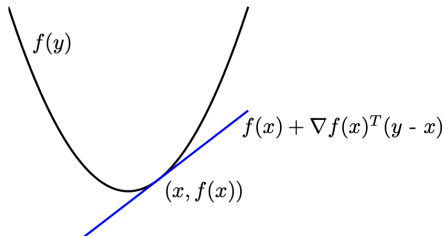


First order convexity conditions

Suppose $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is differentiable. Then f is convex if and only if

$$f(y) \geq f(x) + \nabla f(x)^T (y - x)$$

for all $x, y \in \text{dom}(f)$.



Optimality condition for differentiable f_0 — Proof

We need to prove

$$f_0(x) \leq f_0(y) \quad \forall y \in X \iff \nabla f_0(x)^T(y - x) \geq 0 \quad \forall y \in X$$

- \implies : Suppose x is optimal, but there exists y with $\nabla f_0(x)^T(y - x) < 0$. Consider point $z = \alpha y + (1 - \alpha)x$ with $\alpha \in \{0, 1\}$ is also feasible. As

$$\lim_{\alpha \rightarrow 0} \frac{f_0(z) - f_0(x)}{t} = \nabla f_0(x)^T(y - x) < 0$$

then $f_0(z) < f_0(x)$ for sufficiently small α , hence a contradiction.

- \impliedby : Suppose $\nabla f_0(x)^T(y - x) \geq 0$ for all y , since f_0 is convex

$$f_0(y) \geq f_0(x) + \nabla f_0(x)^T(y - x) \geq f_0(x)$$

Examples

- Unconstrained problem: x is optimal if and only if

$$x \in \text{dom}(f_0), \quad \nabla f_0(x) = 0$$

- equality constrained problem

$$\text{minimize} \quad f_0(x) \quad \text{subject to} \quad \nabla Ax = b$$

x is optimal if and only if there exists a v such that

$$x \in \text{dom}(f_0), \quad Ax = b, \quad \nabla f_0(x) + Av = 0$$

Linear programming (LP)

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Gx \leq h \\ & Ax = b\end{array}$$

- convex problem with affine objective and constraint functions
- feasible set is a polyhedron
- number of solutions: zero, one, infinity

Linear programming (LP) — examples

- ℓ_∞ norm

$$\begin{aligned}
 & \min_x \|Ax + b\|_\infty \\
 \Rightarrow & \min_x \max_i |a_i^T x + b_i| \\
 \Rightarrow & \min_{x,t} t \quad \text{s.t.} \quad -t \leq a_i^T x + b_i \leq t
 \end{aligned}$$

- ℓ_1 norm

$$\begin{aligned}
 & \min_x \|Ax + b\|_1 \\
 \Rightarrow & \min_x \sum_i |a_i^T x + b_i| \\
 \Rightarrow & \min_{x,t_i} \sum_i t_i \quad \text{s.t.} \quad -t_i \leq a_i^T x + b_i \leq t_i
 \end{aligned}$$

Quadratic programming (QP)

$$\begin{array}{ll}\text{minimize} & \frac{1}{2}x^T Px + q^T x + r \\ \text{subject to} & Gx \leq h \\ & Ax = b\end{array}$$

- $P \in \mathbf{S}_{+}^n$, objective is convex quadratic
- minimize a convex quadratic function over a polyhedron

Quadratic programming (QP) — examples

- Least squares

$$\begin{aligned}
 & \min_x \|Ax - b\|_2 \\
 \implies & \min_x \|Ax - b\|_2^2 \\
 \implies & \min_x x^T A^T A x - 2b^T A x + b^T b
 \end{aligned}$$

- Distance between two polyhedra

$$\begin{aligned}
 & \text{dist}(P_1, P_2) = \inf \{ \|x_1 - x_2\|_2 \mid x_1 \in P_1, x_2 \in P_2 \} \\
 \implies & \min_x \|x_1 - x_2\|_2^2 \\
 & \text{s.t. } A_1 x_1 \leq b_1, A_2 x_2 \leq b_2
 \end{aligned}$$

Quadratically constrained quadratic program (QCQP)

$$\begin{aligned} & \text{minimize} && \frac{1}{2}x^T P_0 x + q_0^T x + r_0 \\ & \text{subject to} && \frac{1}{2}x^T P_i x + q_i^T x + r_i \leq 0 \\ & && Ax = b \end{aligned}$$

- $P_i \in \mathbf{S}_{+}^n$, objective and constraints are convex quadratic

Second-order cone programming (SOCP)

$$\begin{aligned}
 & \text{minimize} && h^T x \\
 & \text{subject to} && \|A_i x + b_i\|_2 \leq c_i^T x + d_i \\
 & && Fx = g
 \end{aligned}$$

with $A \in \mathbb{R}^{n_i \times n}$, $F \in \mathbb{R}^{p \times n}$

- inequalities are second-order cone (SOC) constraints:

$$(A_i x + b_i, c_i^T x + d_i) \in \text{second-order cone in } \mathbb{R}^{n_i+1}$$

- if $n_i = 0$, reduces to an LP;
- if $c_i = 0$, reduces to a QCQP

Next Lecture

Convex problems