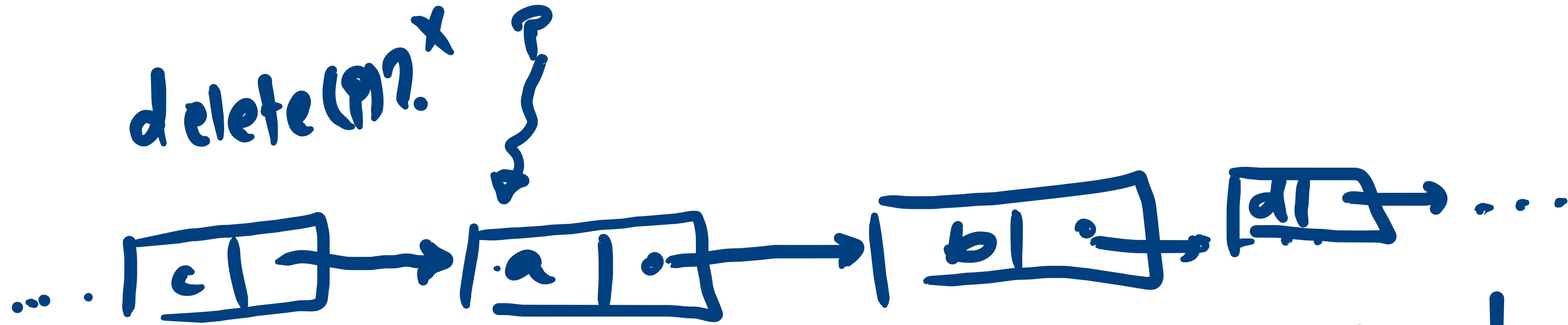
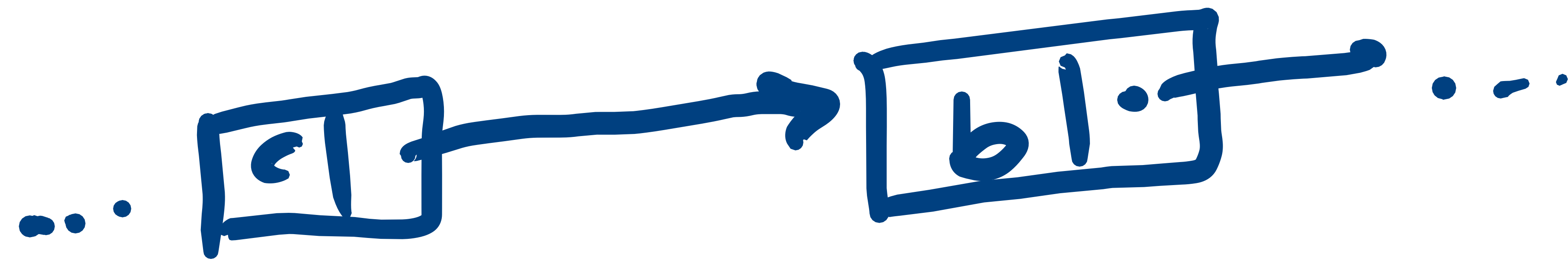


delete(P)?<sup>x</sup>



Not a good option.  
Need access to c.



delete\_next(P) ← delete the node following  
P from the list



delete\_next(P)

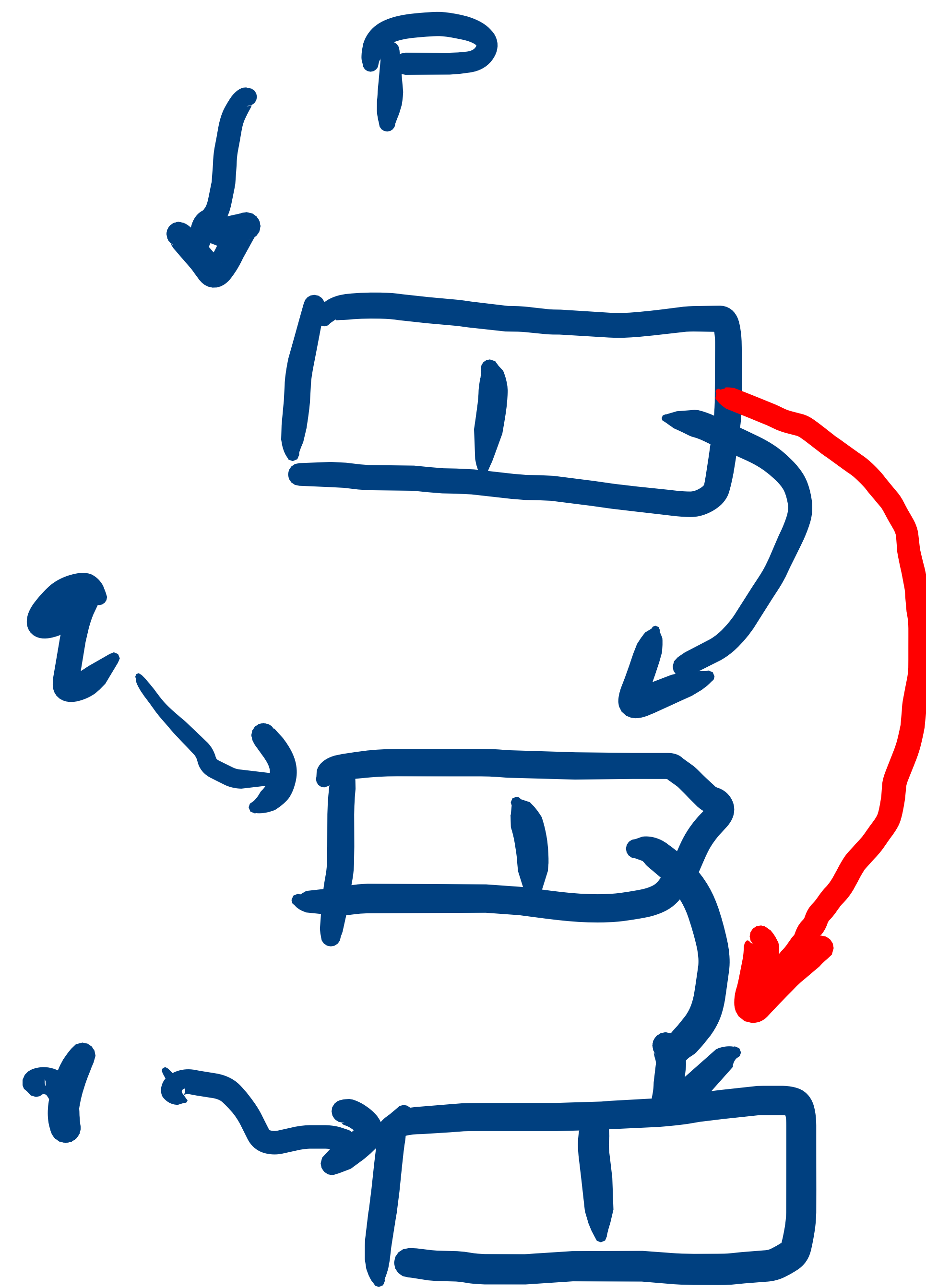
q = P.next

r = q.next

P.next = r

delete q

deallocate  
memory for q.



$O(1)$ -time

What happens if P is last elt?

## Two solutions

1. Require P is not last elt.
2. Check and do nothing if P is last.

1. is faster

2. is safer

# Stacks & Queues using Linked List

Stack

Push

Pop

is\_empty

empty\_stack()

} O(1)-time

How is empty stack represented? NULL

S = Empty()

Push(S, 1)

Push(S, 2)

Pop()

Pop()

S = NULL



S = NULL



Push(s, v)

t = new\_node(v)

t.next = s

return t { s = t  
return s

Pop(s)  
// may be deallocate  
return s.next

Empty()

return NULL

is\_empty(s)

return s == NULL

$S_1 = \text{copy-stack}(S_2)$

$S = \text{empty}()$

$\text{push}(S, 1)$

$t = \text{copy}(S)$

$\text{push}(t, 2)$

$\text{push}(S, 3)$

$\text{print}(S) \quad // \quad S = (3 \quad 1)$

$\text{print}(t) \quad // \quad t = (2 \quad 1)$

$\text{pop}(t)$

$\text{print}(S) \quad // (3 \quad 1)$

$\text{pop}(t)$

$\text{print}(S) \quad // (3 \quad 1)$

a = 2

b = a

a += 1

b += 2

Print(a) // 3

Print(b) // 4

Be like water  
C++ can talk  
"Be like int"



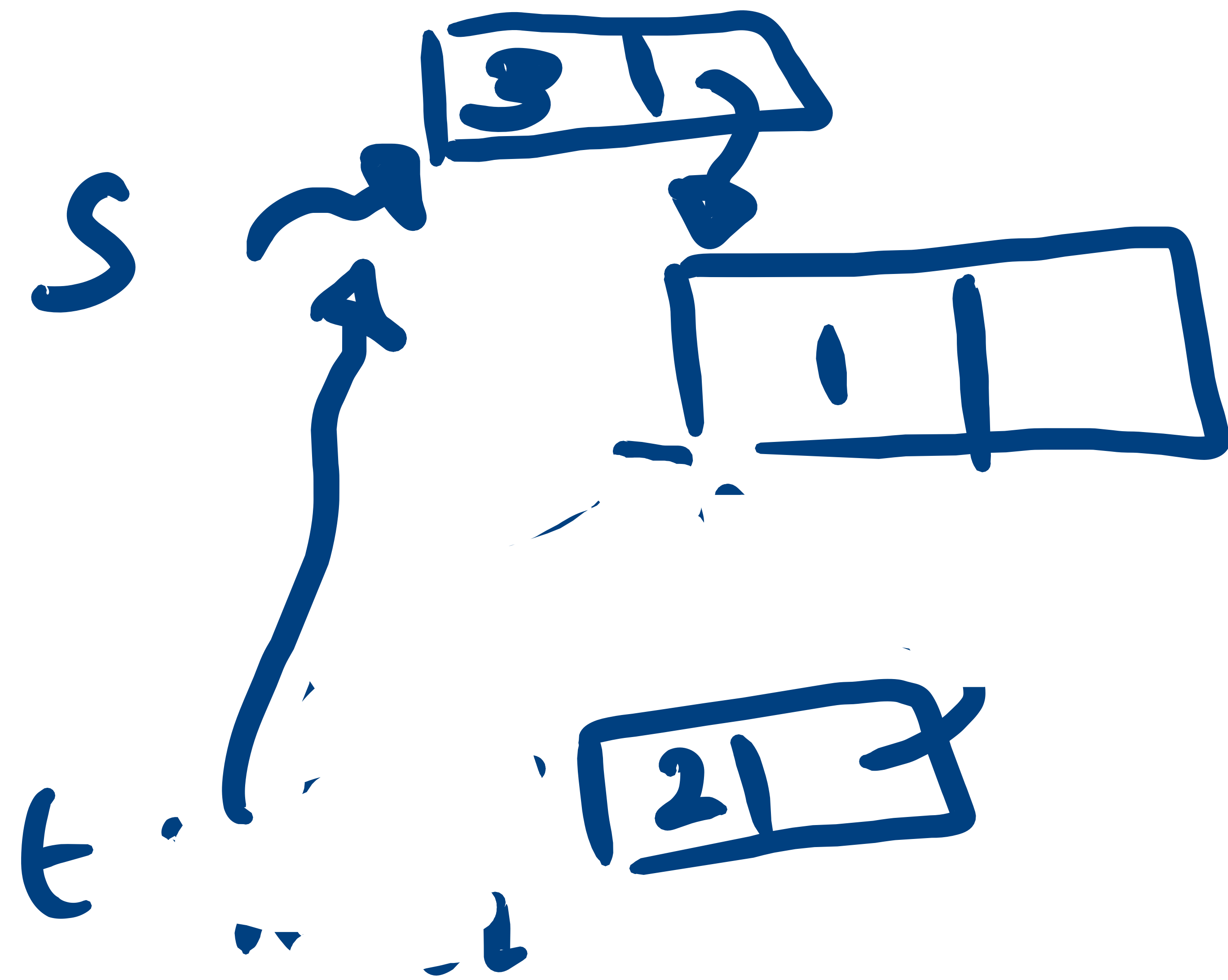
Copy-Stack ( )

array :  $\Omega(n)$

allocate a new array

copy all elts from old array  
to new array

Linked list ?



"tail-showing"

enables copy stack  
in  $O(1)$ -time.

Push( $t, 2$ )

Push( $S, 3$ )

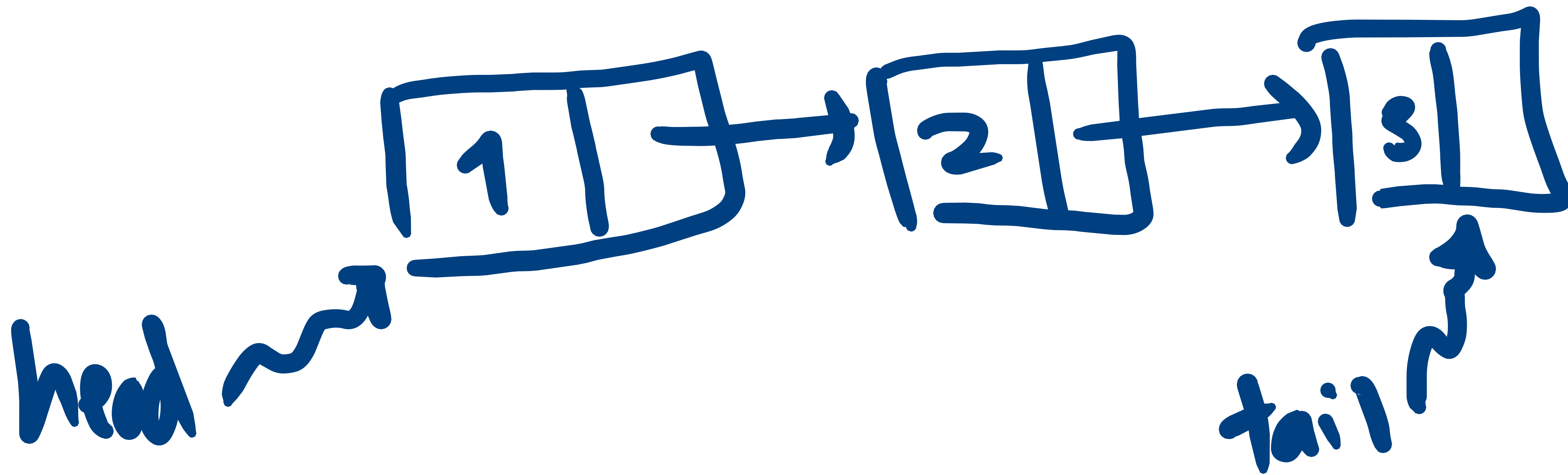
Pop( $t$ )

Pop( $t$ )

$t = \text{copy}(S)$

# Queue

enqueue() }  $O(1)$  - time  
dequeue()



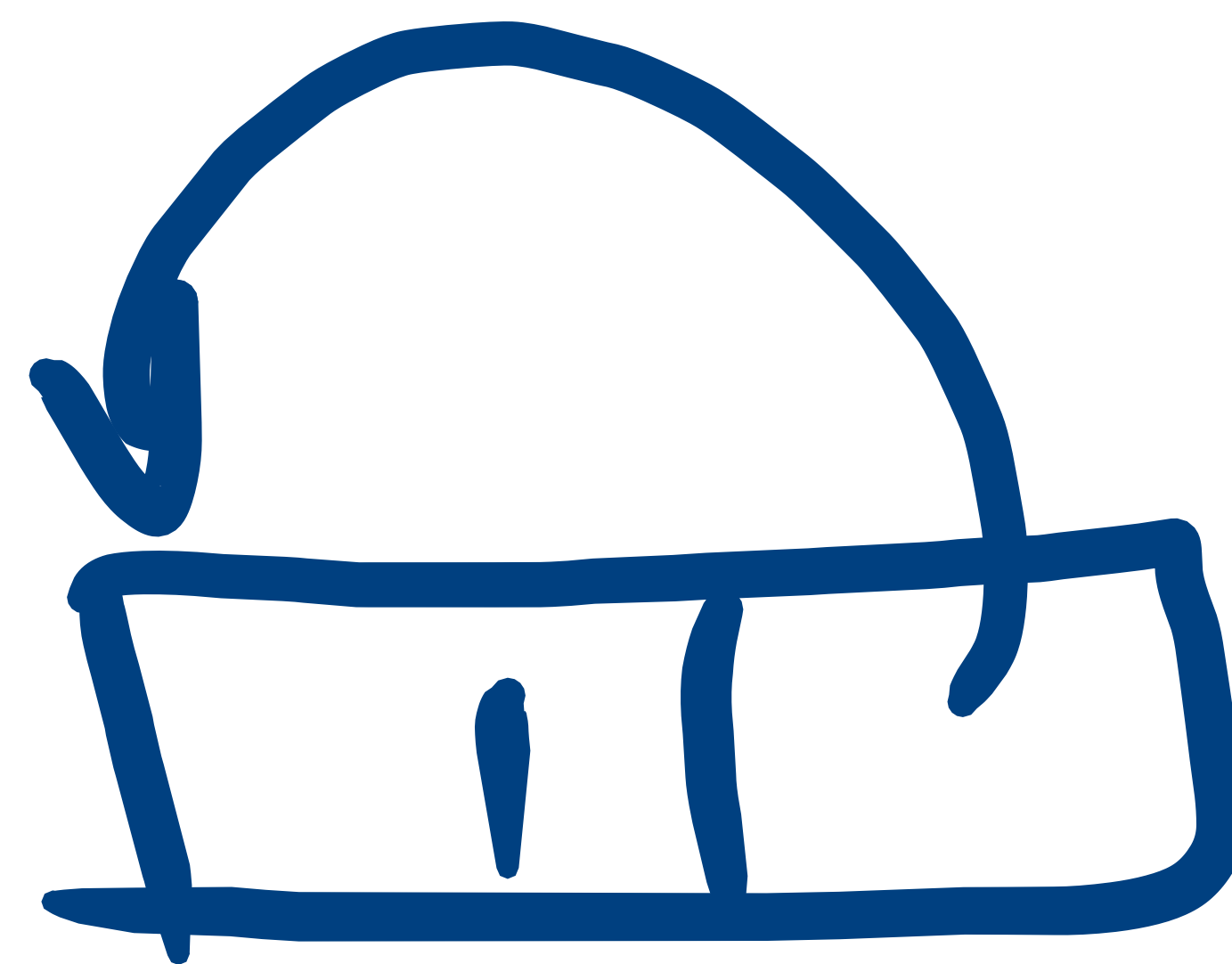
# Circular Singly linked list



tail

copy\_queue() ?

Corner cases



tail