Fibonacci numbers

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***Abstract*—This document depicts a way to setup a matrix equation to find the fibonacci sequence.**

Download all python codes from

https://github.com/Zeeshan−IITH/IITH−EE5609/new/master/codes

and latex-tikz codes from

https://github.com/Zeeshan−IITH/IITH−EE5609

1 Problem

Given a kxk matrix $\mathbf{A}$,find the powers of $\mathbf{A}^n$ within $O(logn)$ time.

2 Construction

For the sake of simplicity we will be calculating the powers given $n = 2^m$, where n is much larger than k.The required result will be of the form $\mathbf{A}^1, \mathbf{A}^2, \mathbf{A}^4, \mathbf{A}^8, \mathbf{A}^{16}..$

The first matrix multiplication will be

$$\mathbf{A}^2 = \mathbf{A}\mathbf{A} \tag{2.0.1}$$

Instead of using repeated multiplication by $\mathbf{A}$,we can use the previous result and square it to be closer to the result using less complutations

$$\mathbf{A}^4 = \mathbf{A}^2\mathbf{A}^2 \tag{2.0.2}$$

$$\mathbf{A}^8 = \mathbf{A}^4\mathbf{A}^4 \tag{2.0.3}$$

$$\mathbf{A}^{16} = \mathbf{A}^8\mathbf{A}^8 \tag{2.0.4}$$

So $\mathbf{A}^{2^m}$ need only m products of the resultant matrix from the previous computation.Since $m = log_2(n)$, the result can be computed in $O(log_2 n)$ time.

As a more general case any number n can be represented as a sum of powers of 2, just like binary numbers are represented with the radix 2 i.e when we represent n in binary form we get

$$n = b_k b_{k-1} b_{k-2} ... b_0 \tag{2.0.5}$$

where

$$k = \log_2 n$$

because we require that many number of bits to represent n.

Now to calculate

$$\boldsymbol{A}^n$$

we can use

$$\boldsymbol{A}^n = b_k \boldsymbol{A}^{2^k} + b_{k-1}\boldsymbol{A}^{2^{k-1}} + b_{k-2}\boldsymbol{A}^{2^{k-2}} ... + b_2\boldsymbol{A}^4 + b_1\boldsymbol{A}^2 + b_0\boldsymbol{A} \tag{2.0.6}$$

Now each of the $\boldsymbol{A}^{2^k}$ can be calculated by squaring the previous $\boldsymbol{A}^{2^{k-1}}$.Neglecting matrix addition as it only takes m^2 time to compute when compared to m^3 time of each multiplication,where m is the order of the square matrix, we can say that since $\boldsymbol{A}^{2^k}$ calculation takes $O(\log_2 n)$ time,the whole computation $\boldsymbol{A}^n$ takes $O(\log_2 n)$.

3 Fibonacci

Consider the special matrix which begins with fibonacci numbers in it

$$\mathbf{F} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \tag{3.0.1}$$

This matrix has the fibonaci numbers as its elements

$$\mathbf{F^2} = \begin{pmatrix} 2 & 1 \\ 1 & 0 \end{pmatrix} \tag{3.0.2}$$

$$\mathbf{F^3} = \begin{pmatrix} 3 & 2 \\ 2 & 1 \end{pmatrix} \tag{3.0.3}$$

Therefore the n^{th} fibonacci number F_n is the element in the first row,first column of $\mathbf{F^{n-1}}$,where $n \geq 2$.The general form will be

$$\mathbf{F^{n-1}} = \begin{pmatrix} F_n & F_{n-1} \\ F_{n-2} & F_{n-3} \end{pmatrix} \tag{3.0.4}$$