

Syntax-Generic Operations, Reflectively Reified

Extended Abstract

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Hello everyone, my name is Tzu-Chi, I am a research assistant at Academia Sinica in Taiwan. Today I will be talking about the joint work of my advisor's and mine, the title is "Syntax-generic Operations, Reflectively Reified". So firstly we will see what are syntax-generic operations, then I will elaborate on what Reflectively reified means.

TYPICAL LANGUAGES IN LANGUAGES

Intrinsic typing is common for λ -calculus with De Bruijn indices.

data $\vdash_ : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**

$\backslash_ : \Gamma \ni A \rightarrow \Gamma \vdash A$

$\lambda_ : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$

$\cdot_ : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

data $\text{Ty} : \text{Set}$ **where**

$\alpha : \text{Ty}$

$\Rightarrow_ : \text{Ty} \rightarrow \text{Ty} \rightarrow \text{Ty}$

data $\text{Context} (A : \text{Set}) : \text{Set}$ **where**

$\emptyset : \text{Context } A$

$\rightarrow, - : (\Gamma : \text{Context } A) \rightarrow A \rightarrow \text{Context } A$

data $_ \ni _ : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**

$z : \Gamma, A \ni A$

$s : \Gamma \ni A \rightarrow \Gamma, B \ni A$

Syntax-Generic Operations, Reflectively Reified

└ Typical Languages in Languages

For starters, Let's take a look at a common thing dependently typed programmers would do, that is defining a programming language in a dependently typed programming language, such as Agda. Here is an example of the syntax of a language defined in Agda. It's a simply-typed lambda calculus. We define a datatype called `Ty` to describe the possible types for this language, a term has either a single type or a function type. Here *Context* is defined as the List of `Ty`. And we have a simple “has” relation that represents variables. For if a variable is typed it must be in the context. And we know that a simply typed lambda calculus has three constructions, each corresponding a a typing rule. Here are three constructors, each of them corresponds to the variable rule, the abstraction rule, and the application rule. We can see for a term to be constructed, it must be well-typed by definition, that's what we call intrinsic typing.

Intrinsic typing is common for λ -calculus with De Bruijn indices.

```
data Ty : Context → Ty → Set where
  v_  : F → A → F × A
  λ_  : F , A → B → F × A → B
  _·_ : F × A → B → F × A → F × B
```

```
data Ty : Set where
  a_  : Ty
  _·_ : Ty → Ty → Ty
data Context (A : Set) : Set where
  ()  : Context A
  _·_ : (F : Context A) → A → Context A
data _·_ : Context → Ty → Set where
  z : F , A → A
  s : F → A → F , B → A
```

Scope-safe syntax operations:

$$\begin{aligned} \text{rename} : \forall \{ \Gamma \ \Delta \} \rightarrow (\forall \{ A \} \rightarrow \Gamma \ni A \rightarrow \Delta \ni A) \\ \rightarrow (\forall \{ A \} \rightarrow \Gamma \vdash A \rightarrow \Delta \vdash A) \end{aligned}$$

$$\text{rename } \rho \ (\backslash x) \quad = \quad \backslash (\rho x)$$

$$\text{rename } \rho \ (\lambda N) \quad = \quad \lambda (\text{rename } (\text{ext } \rho) N)$$

$$\text{rename } \rho \ (L \cdot M) \quad = \quad (\text{rename } \rho L) \cdot (\text{rename } \rho M)$$

Syntax-Generic Operations, Reflectively Reified

Scope-safe syntax operations:

$$\begin{aligned} \text{rename} &: \forall \{F, A\} \rightarrow (\forall \{A\} \rightarrow F \supset A \rightarrow A \supset A) \\ &\quad \rightarrow (\forall \{A\} \rightarrow F \supset A \rightarrow A \supset A) \\ \text{rename } \rho \text{ (} x \text{)} &= `(\rho x) \\ \text{rename } \rho \text{ (} \lambda N \text{)} &= \lambda (\text{rename } (\text{ext } \rho) N) \\ \text{rename } \rho \text{ (} L \cdot M \text{)} &= (\text{rename } \rho L) \cdot (\text{rename } \rho M) \end{aligned}$$

Then let's define some operations on this syntax. Again, because of intrinsic typing, such operations are scope-safe. For example, a *rename* function says that if there's a mapping from every variable in a context Γ to the other context Δ , we can always find a mapping from every term that is typed in Γ to a term typed in Δ . Because we can remap every variable in that term. If we encounter a subterm that has a extended context, for example, the case of lambda abstraction, the variable mapping should be extended accordingly. We can see that rename is still called recursively, but the variable mapping function is modified with ext.

Intrinsic typing is common for λ -calculus with De Bruijn indices.

data $Ty : Set$ **where**

...

$\backslash \mathbb{N} : Ty$

data $\vdash_~ : Context \rightarrow Ty \rightarrow Set$ **where**

...

$\backslash zero : \Gamma \vdash \backslash \mathbb{N}$

$\backslash suc_~ : \Gamma \vdash \backslash \mathbb{N} \rightarrow \Gamma \vdash \backslash \mathbb{N}$

$case : \Gamma \vdash \backslash \mathbb{N} \rightarrow \Gamma \vdash A \rightarrow \Gamma, \backslash \mathbb{N} \vdash A \rightarrow \Gamma \vdash A$

$\mu_~ : \Gamma, A \vdash A \rightarrow \Gamma \vdash A$

Syntax-Generic Operations, Reflectively Reified

└ Motivation

Now what if we extend the language? we might want natural number primitives in our language, so we extend the Ty datatype with a new type, we add four typing rules, two of them, zero and suc are for constructing natural numbers, and another two rules are for branching and recursion.

Intrinsic typing is common for λ -calculus with De Bruijn indices.

data Ty : Set **where**

—

"H : Ty

data ϕ_n : Context \rightarrow Ty \rightarrow Set **where**

—

"zero : $F \vdash "H$

"suc_n : $F \vdash "H \rightarrow F \vdash "H$

case : $F \vdash "H \rightarrow F \vdash A \rightarrow F, "H \vdash A \rightarrow F \vdash A$

μ_n : $F, A \vdash A \rightarrow F \vdash A$

MOTIVATION

$$\begin{aligned} \text{rename} : \forall \{ \Gamma \Delta \} \rightarrow (\forall \{ A \} \rightarrow \Gamma \ni A \rightarrow \Delta \ni A) \\ \rightarrow (\forall \{ A \} \rightarrow \Gamma \vdash A \rightarrow \Delta \vdash A) \end{aligned}$$

...

$$\begin{aligned} \text{rename } \rho \text{ (`zero)} &= \text{'zero} \\ \text{rename } \rho \text{ (`suc } M) &= \text{'suc (rename } \rho \text{ } M) \\ \text{rename } \rho \text{ (case } L \text{ } M \text{ } N) &= \text{case (rename } \rho \text{ } L) \\ &\quad (\text{rename } \rho \text{ } M) \\ &\quad (\text{rename (ext } \rho) \text{ } N) \\ \text{rename } \rho \text{ (}\mu \text{ } N) &= \mu \text{ (rename (ext } \rho) \text{ } N) \end{aligned}$$

Syntax-Generic Operations, Reflectively Reified

└ Motivation

MOTIVATION

```

rename :  $\forall (F, A) \rightarrow (\forall \{A\} \rightarrow F \supset A \rightarrow A \supset A)$ 
         $\rightarrow (\forall \{A\} \rightarrow F \supset A \rightarrow A \supset A)$ 
...
rename  $\rho$  ("zero")      = "zero"
rename  $\rho$  ("var M")    = "var (rename  $\rho$  M)"
rename  $\rho$  ("case L M N") = "case (rename  $\rho$  L)
                        (rename  $\rho$  M)
                        (rename (ext  $\rho$ ) N)"
rename  $\rho$  ( $\mu$  N)      =  $\mu$  (rename (ext  $\rho$ ) N)

```

If we want a rename function for this extended language, we can copy and paste the rename we just defined and add a new clause for each constructor. The patterns still follow, a renaming always take a term from a constructor to one with the same constructor, if encounter subterm, call rename recursively, if there's context extension, extend the variable mapping as well. So no matter how many constructors there are we follow the same logic.

Change/extend the object language even further:

data $\vdash_{-} : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**

...

$\text{con} : \mathbb{N} \rightarrow \Gamma \vdash \text{Nat}$

$_ \ast _ : \Gamma \vdash \text{Nat} \rightarrow \Gamma \vdash \text{Nat} \rightarrow \Gamma \vdash \text{Nat}$

$\backslash \text{let} : \Gamma \vdash A \rightarrow \Gamma, A \vdash B \rightarrow \Gamma \vdash B$

$\backslash \langle _, _ \rangle : \Gamma \vdash A \rightarrow \Gamma \vdash B \rightarrow \Gamma \vdash A \backslash \times B$

$\backslash \text{proj}_1 : \Gamma \vdash A \backslash \times B \rightarrow \Gamma \vdash A$

$\backslash \text{proj}_2 : \Gamma \vdash A \backslash \times B \rightarrow \Gamma \vdash B$

$\text{case}\times : \Gamma \vdash A \backslash \times B \rightarrow \Gamma, A, B \vdash C \rightarrow \Gamma \vdash C$

Syntax-Generic Operations, Reflectively Reified

└ Motivation

What if we extend the language even further? Here we add even more rules. syntax sugaring, pairs and projections, etc.

Change/extend the object language even further:

```
data  $\phi_n$  : Context  $\rightarrow$  Ty  $\rightarrow$  Set where
  --
  cons      : M  $\rightarrow$  F  $\vdash$  Nat
   $\lambda_{\text{Nat}}$      : F  $\vdash$  Nat  $\rightarrow$  F  $\vdash$  Nat  $\rightarrow$  F  $\vdash$  Nat
  'let      : F  $\vdash$  A  $\rightarrow$  F, A  $\vdash$  B  $\rightarrow$  F  $\vdash$  B
  ' $\text{let}_A$ ,  $\text{let}_B$  : F  $\vdash$  A  $\rightarrow$  F  $\vdash$  B  $\rightarrow$  F  $\vdash$  A  $\times$  B
  'proj1    : F  $\vdash$  A  $\times$  B  $\rightarrow$  F  $\vdash$  A
  'proj2    : F  $\vdash$  A  $\times$  B  $\rightarrow$  F  $\vdash$  B
  cases     : F  $\vdash$  A  $\times$  B  $\rightarrow$  F, A, B  $\vdash$  C  $\rightarrow$  F  $\vdash$  C
```

Redefine/extend syntax operations:

...

$$\text{rename } \rho \text{ (con } n) = \text{con } n$$
$$\text{rename } \rho \text{ (M `* N) = rename } \rho \text{ M `* rename } \rho \text{ N}$$
$$\text{rename } \rho \text{ (`let M N) = `let (rename } \rho \text{ M) (rename (ext } \rho) \text{ N)}$$
$$\text{rename } \rho \text{ (`< M , N > = `< rename } \rho \text{ M , rename } \rho \text{ N >}$$
$$\text{rename } \rho \text{ (`proj}_1 \text{ L) = `proj}_1 \text{ (rename } \rho \text{ L)}$$
$$\text{rename } \rho \text{ (`proj}_2 \text{ L) = `proj}_2 \text{ (rename } \rho \text{ L)}$$
$$\begin{aligned} \text{rename } \rho \text{ (case}\times \text{ L M)} &= \text{case}\times \text{ (rename } \rho \text{ L)} \\ &\quad (\text{rename (ext (ext } \rho)) \text{ M)} \end{aligned}$$

Syntax-Generic Operations, Reflectively Reified

└ Motivation

The *rename* function must be extended as well, so everytime we make some changes to the object language, it becomes a very repeating work.

Redefine/extend syntax operations:

```

...
rename  $\rho$  (cons n)      = cons n
rename  $\rho$  (M ^* N)      = rename  $\rho$  M ^* rename  $\rho$  N
rename  $\rho$  ('let M N)     = 'let (rename  $\rho$  M) (rename (ext  $\rho$ ) N)
rename  $\rho$  ('< M , N >)   = '< rename  $\rho$  M , rename  $\rho$  N >
rename  $\rho$  ('proji L)    = 'proji (rename  $\rho$  L)
rename  $\rho$  ('projj L)    = 'projj (rename  $\rho$  L)
rename  $\rho$  (caseex L M)  = caseex (rename  $\rho$  L)
                        (rename (ext (ext  $\rho$ )) M)

```

Other repeating operations:

$$\begin{aligned} \textit{subst} : \forall \{ \Gamma \Delta \} \rightarrow (\forall \{ A \} \rightarrow \Gamma \ni A \rightarrow \Delta \vdash A) \\ \rightarrow (\forall \{ A \} \rightarrow \Gamma \vdash A \rightarrow \Delta \vdash A) \end{aligned}$$

$$\textit{print} : \Gamma \vdash A \rightarrow \textit{String}$$

...

└ Motivation

Other repeating operations:

$$\begin{aligned} \text{subst} &: \forall (\Gamma \ d) \rightarrow (\forall (A) \rightarrow \Gamma \ni A \rightarrow A \vdash A) \\ &\rightarrow (\forall (A) \rightarrow \Gamma \vdash A \rightarrow A \vdash A) \\ \text{print} &: \Gamma \vdash A \rightarrow \text{String} \\ &\dots \end{aligned}$$

There are other repeating operations that should be redefined for every change in the object language, you probably would want some kind of generic programs that work for every syntax, or a mechanism that generates these operations from any given datatypes of this kind.

WHERE WE ARE GOING...

1. Existing Work for Syntax-generic Operations
2. Elaborator Reflection to the Rescue
3. Discussion

└ Where we are going...

1. Existing Work for Syntax-generic Operations
2. Elaborator Reflection to the Rescue
3. Discussion

Now we know the problem we are dealing with. In the next section, I will introduce an existing work that eliminates the repetitions we just mentioned. After that I will illustrate how we develop an alternative approach based on this previous work, by incorporating elaborator reflection. We will make some comparisons along the way, and finally we raise some questions regarding our own work, and hopefully get some feedbacks from you.

Existing Work for Syntax-generic Operations

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Syntax-Generic Operations, Reflectively Reified

- └ Existing Work for Syntax-generic Operations

Existing Work for
Syntax-generic Operations

Let's start with the existing work.

There are generic libraries for a family/families of syntaxes with binders.

We improve upon Allais et al.'s approach presented at ICFP '18 (later published in JFP '21).

Syntax-Generic Operations, Reflectively Reified

└ Existing Work for Syntax-generic Operations

└ Existing work

There are generic libraries for a family/families of syntaxes with binders.

We improve upon Allais et al.'s approach presented at ICFP '18 (later published in JFP '21).

There have been some libraries that provide generic operations for the problem we introduced. Those programs are syntax-generic, that's the first part of our title. We focus on one of such works, that is the generic library by Allais et al. They have developed a framework with a variety of syntax-generic operations, such as renaming, substitution, printing and scope checking.

Allais et al.'s *Desc*:

```
data Desc (I : Set) : Set1 where  
  `σ : (A : Set) → (A → Desc I) → Desc I  
  `X : List I → I → Desc I → Desc I  
  `■ : I → Desc I
```

```
data `STLC : Set where  
  App Lam : Ty → Ty → `STLC  
STLCD : Desc Ty  
STLCD = `σ `STLC λ where  
  (App i j) →  
    `X [] (i ⇒ j) (`X [] i (■ j))  
  (Lam i j) →  
    `X (i :: []) j (■ (i ⇒ j))
```

Syntax-Generic Operations, Reflectively Reified

└ Existing Work for Syntax-generic Operations

└ Existing work by Allais et al.

How do they achieve it? They provide a universe of descriptions called Desc that describes a family of syntaxes. We won't go into its details here, but we can say that every inhabitant in this universe represents a syntax, and simply-typed lambda calculus is one of them. The parameter I in the Desc universe says that a syntax is intrinsically typed by I and has a context of List of I . So simply-typed lambda calculus can be encoded in the universe Desc

Allais et al.'s Desc:

```
data Desc (I : Set) : Set where
  'σ : (A : Set) → (A → Desc I) → Desc I
  'X : List I → I → Desc I → Desc I
  '■ : I → Desc I
```

```
data 'STLC : Set where
  App Lam : Ty → Ty → 'STLC
STLCD : Desc Ty
STLCD = 'σ 'STLC λ where
  (App i j) →
    'X [] (i ⇒ j) ('X [] i (■ j))
  (Lam i j) →
    'X (i :: []) j (■ (i ⇒ j))
```


Simply typed λ -calculus

data $\vdash_{-} : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**

$\backslash_{-} : \Gamma \ni A \rightarrow \Gamma \vdash A$

$\lambda_{-} : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$

$\cdot_{-} : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

Simply typed λ -calculus

data $\vdash_- : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**
 $_ : \Gamma \ni A \rightarrow \Gamma \vdash A$
 $\lambda_ : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$
 $_ : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

encoded in *Desc*:

$STLCD : \text{Desc Ty}$
 $STLCD = \backslash \sigma \backslash STLC \lambda$ **where**
 $(App\ i\ j) \rightarrow$
 $\backslash X\ []\ (i \Rightarrow j)\ (\backslash X\ []\ i\ (\blacksquare\ j))$
 $(Lam\ i\ j) \rightarrow$
 $\backslash X\ (i :: [])\ j\ (\blacksquare\ (i \Rightarrow j))$

Simply typed λ -calculus

data $\vdash_- : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**
 $_ : \Gamma \ni A \rightarrow \Gamma \vdash A$
 $\lambda_ : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$
 $_ : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

encoded in *Desc*:

$STLCD : \text{Desc Ty}$
 $STLCD = \backslash \sigma \backslash STLC \lambda$ **where**
 $(App\ i\ j) \rightarrow$
 $\backslash X\ []\ (i \Rightarrow j)\ (\backslash X\ []\ i\ (\blacksquare\ j))$
 $(Lam\ i\ j) \rightarrow$
 $\backslash X\ (i :: [])\ j\ (\blacksquare\ (i \Rightarrow j))$

$STLC' = Tm\ STLCD$

Simply typed λ -calculus

data $\vdash_- : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**
 $_ : \Gamma \ni A \rightarrow \Gamma \vdash A$
 $\lambda_- : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$
 $_ : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

encoded in *Desc*:

$STLCD : \text{Desc Ty}$
 $STLCD = \backslash \sigma \backslash STLC \lambda$ **where**
 $(App\ i\ j) \rightarrow$
 $\backslash X\ []\ (i \Rightarrow j)\ (\backslash X\ []\ i\ (\blacksquare\ j))$
 $(Lam\ i\ j) \rightarrow$
 $\backslash X\ (i :: [])\ j\ (\blacksquare\ (i \Rightarrow j))$

$STLC' = Tm\ STLCD$

$\vdash_- \cong STLC'$

Syntax-Generic Operations, Reflectively Reified

- Existing Work for Syntax-generic Operations

- Existing work by Allais et al.

EXISTING WORK BY ALLAIS ET AL.

Simply typed λ -calculus
$$\text{data } \mathcal{J}_\lambda : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set} \text{ where}$$

$$\begin{aligned} \mathcal{J}_\lambda &: F \triangleright A \rightarrow F \triangleright A \\ \mathcal{J}_\lambda &: F, A \triangleright B \rightarrow F \triangleright A \multimap B \\ \mathcal{J}_\lambda &: F \triangleright A \multimap B \rightarrow F \triangleright A \rightarrow F \triangleright B \end{aligned}$$

encoded in Dae:

 $STLCD : \text{DescTy}$ $STLCD = \text{'or' } STLC \text{ A where}$ $(App\ i\ j) \rightarrow$ $\text{'X [] (i \multimap j) (X [] i (■ j))}$ $(Lam\ i\ j) \rightarrow$ $\text{'X (i :: []) j (■ (i \multimap j))}$ $STLC' = Tm\ STLCD$ $\mathcal{J}_\lambda \approx STLC'$

A simply-typed lambda calculus here on the left hand side has just been rewritten as a description on the right hand side. In fact, to use such generic libraries, programmers must encode their syntax in the given description, because syntax-generic operations are defined on these descriptions. To acquire something structurally similar to the native syntax datatype, we can use the Tm type constructor, which takes the fixpoint of the functor of a description. We may prove that a syntax defined as a isolated datatype is isomorphic to the fixpoint of such functor.

GENERIC FUNCTIONS FOR THE WHOLE UNIVERSE

$$\text{Var } \sigma \Gamma = \Gamma \ni \sigma$$

$$\text{Renaming} : \forall \{d : \text{Desc } I\} \rightarrow \text{Semantics } d \text{ Var } (\text{Tm } d)$$

$$\text{rename} : \forall \{d : \text{Desc } I\} \rightarrow (\forall \{\sigma\} \rightarrow \text{Var } \sigma \Gamma \rightarrow \text{Var } \sigma \Delta)$$

$$\forall \{\tau\} \quad \rightarrow \text{Tm } d \tau \Gamma \rightarrow \text{Tm } d \tau \Delta$$

$$\text{rename} = \text{semantics Renaming}$$

GENERIC FUNCTIONS FOR THE WHOLE UNIVERSE

$$\text{Var } \sigma \Gamma = \Gamma \ni \sigma$$

$$\text{Renaming} : \forall \{d : \text{Desc } I\} \rightarrow \text{Semantics } d \text{ Var } (\text{Tm } d)$$

$$\text{rename} : \forall \{d : \text{Desc } I\} \rightarrow (\forall \{\sigma\} \rightarrow \text{Var } \sigma \Gamma \rightarrow \text{Var } \sigma \Delta)$$

$$\forall \{\tau\} \quad \rightarrow \text{Tm } d \tau \Gamma \rightarrow \text{Tm } d \tau \Delta$$

$$\text{rename} = \text{semantics Renaming}$$

Generic programs are *Semantics* records.

GENERIC FUNCTIONS FOR THE WHOLE UNIVERSE

$$\text{Var } \sigma \Gamma = \Gamma \ni \sigma$$

$$\text{Renaming} : \forall \{d : \text{Desc } I\} \rightarrow \text{Semantics } d \text{ Var } (Tm \ d)$$

$$\text{rename} : \forall \{d : \text{Desc } I\} \rightarrow (\forall \{\sigma\} \rightarrow \text{Var } \sigma \Gamma \rightarrow \text{Var } \sigma \Delta)$$

$$\forall \{\tau\} \qquad \qquad \qquad \rightarrow Tm \ d \ \tau \ \Gamma \rightarrow Tm \ d \ \tau \ \Delta$$

$$\text{rename} = \text{semantics Renaming}$$

Generic programs are *Semantics* records.

Functions are realized on fixpoints *Tm* via *semantics*.

GENERIC FUNCTIONS FOR THE WHOLE UNIVERSE

$$\text{Var } \sigma \Gamma = \Gamma \ni \sigma$$

$$\text{Renaming} : \forall \{d : \text{Desc } I\} \rightarrow \text{Semantics } d \text{ Var } (Tm \ d)$$

$$\text{rename} : \forall \{d : \text{Desc } I\} \rightarrow (\forall \{\sigma\} \rightarrow \text{Var } \sigma \Gamma \rightarrow \text{Var } \sigma \Delta)$$

$$\forall \{\tau\} \quad \quad \quad \rightarrow Tm \ d \ \tau \ \Gamma \rightarrow Tm \ d \ \tau \ \Delta$$

$$\text{rename} = \text{semantics Renaming}$$

Generic programs are *Semantics* records.

Functions are realized on fixpoints *Tm* via *semantics*.

rename can be applied to fixpoints of any description (e.g. *Tm STLCD*).

Syntax-Generic Operations, Reflectively Reified

└ Existing Work for Syntax-generic Operations

└ Generic Functions for the Whole Universe

$$\begin{aligned}
 \text{Var } \sigma \ F &= F \circ \sigma \\
 \text{Renaming} &: \forall (d : \text{Desc } 1) \rightarrow \text{Semantics } d \ \text{Var} \ (\text{Tm } d) \\
 \text{rename} &: \forall (d : \text{Desc } 1) \rightarrow (\forall (\sigma) \rightarrow \text{Var } \sigma \ F \rightarrow \text{Var } \sigma \ A) \\
 &\quad \forall (\tau) \rightarrow \text{Tm } d \ \tau \ F \rightarrow \text{Tm } d \ \tau \ A \\
 \text{rename} &= \text{semantics } \text{Renaming}
 \end{aligned}$$

Generic programs are *Semantics* records.

Functions are realized on fixpoints *Tm* via *semantics*.

rename can be applied to fixpoints of any description (e.g. *Tm STLCD*).

With such a universe, Allais they can and have defined some generic functions that work for a whole family of syntaxes. For example, they have defined a generic renaming function. A semantics datatype with the upper case *S* is a description of a generic function given a syntax description *d*. You can see *Renaming* here is generic as it's quantified over *d*. The semantics function with lower case *s* is used to obtain the actual generic rename function, we can see for every description *d*, this function works on *Tm d*. It essentially says that there's a rename for every given *d*, and of course there's a rename for the description of simply-typed lambda calculus.

PROBLEMS WITH SYNTAX UNIVERSES: READABILITY

data $\vdash_ : \text{Context} \rightarrow \text{Ty} \rightarrow \text{Set}$ **where**

$_ : \Gamma \ni A \rightarrow \Gamma \vdash A$

$\lambda_ : \Gamma, A \vdash B \rightarrow \Gamma \vdash A \Rightarrow B$

$_ : \Gamma \vdash A \Rightarrow B \rightarrow \Gamma \vdash A \rightarrow \Gamma \vdash B$

$STLCD : \text{Desc Ty}$

$STLCD = _ \sigma _ STLC \lambda$ **where**

$(App\ i\ j) \rightarrow$

$_X\ []\ (i \Rightarrow j)\ (_X\ []\ i\ (\blacksquare\ j))$

$(Lam\ i\ j) \rightarrow$

$_X\ (i :: [])\ j\ (\blacksquare\ (i \Rightarrow j))$

$STLC' = Tm\ STLCD$

Syntax-Generic Operations, Reflectively Reified

Existing Work for Syntax-generic Operations

Problems with Syntax Universes: Readability

PROBLEMS WITH SYNTAX UNIVERSES: READABILITY

```

data  $\mathcal{U}_\alpha$  : Context  $\rightarrow$  Ty  $\rightarrow$  Set where
   $\ulcorner$  :  $F \triangleright A \rightarrow F \vdash A$ 
   $\downarrow_\alpha$  :  $F, A \triangleright B \rightarrow F \vdash A \Rightarrow B$ 
   $\ulcorner_\alpha$  :  $F \vdash A \Rightarrow B \rightarrow F \triangleright A \rightarrow F \vdash B$ 

STLCD : Dec Ty
STLCD =  $\ulcorner$  or "STLC A where
  (App i j)  $\rightarrow$ 
    "X [] (i  $\Rightarrow$  j) ("X [] i (j))
  (Lam i j)  $\rightarrow$ 
    "X (i :: []) j ("X (i :: j))
STLC" =  $\ulcorner$  or STLCD

```

We quickly summerize what we think are the reasons these libraries are not widely adopted. First of all is readability. One good thing about of intrinsic typing is that types of constructors closely resemble typing rules, and typing rules are less obvious for syntaxes defined in generic universes.

```
data Desc (I : Set) : Set1 where  
  `σ : (A : Set) → (A → Desc I) → Desc I  
  `X : List I → I → Desc I → Desc I  
  `■ : I → Desc I
```

Syntax-Generic Operations, Reflectively Reified

└ Existing Work for Syntax-generic Operations

└ Problems with Syntax Universes: Burden on Programmers

```
data Desc (I : Set) : Set where
  'α : (A : Set) → (A → Desc I) → Desc I
  'X : List I → I → Desc I → Desc I
  '■ : I → Desc I
```

Secondly, to utilize such generic libraries, programmers are required to understand the generic universe instead of just defining syntaxes the way they want.

PROBLEMS WITH SYNTAX UNIVERSES: INTEROPERABILITY

$STLCD : Desc\ Ty$

$STLCD = \dots$

$STLCD' : Desc' ???$

$STLCD' = ???$

Syntax-Generic Operations, Reflectively Reified

└ Existing Work for Syntax-generic Operations

└ Problems with Syntax Universes: Interoperability

STLCD : Desc Ty
STLCD = ...
STLCD' : Desc' ???
STLCD' = ???

And they need to learn a new generic representation everytime they want some features that only exist in another generic library, the syntax they are working on must be redefined. This leads us to the third problem, interoperability. It would be hard to use two or more generic libraries at once.

Elaborator Reflection to the Rescue

Therefore, we want the best from both worlds, we want programmers and researchers to use native definitions whenever possible, while generic programs can still be invoked. We achieve this by elaborator reflection. Elaborator reflection is the metaprogramming mechanism provided by Agda, it allows us to read and define datatype and function definitions.

“Datatype-Generic Programming Meets Elaborator Reflection” by Josh Ko, Liang-Ting Chen, and Tzu-Chi Lin at 15:50, Tuesday.

“Datatype-Generic Programming Meets Elaborator Reflection” by Josh Ko, Liang-Ting Chen, and Tzu-Chi Lin at 15:50, Tuesday.

Syntax-generic operations *are* Datatype-generic programs with constraints.

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─Elaborator Reflection to the Rescue

I would like to shamelessly promote the other work we are presenting at ICFP, Datatype-Generic Programming Meets Elaborator Reflection. My advisor Josh will present it on Tuesday. We have demonstrated in that work that we can define programs on a family of datatypes with our program description, while using elaborator reflection to reify such programs as natural function definitions defined on native datatypes. What does this have to do with syntax-generic libraries that we spent so much time introducing? It turns out, syntax-generic programs sometimes datatype-generic programs. We can constrain a subset of all datatypes such that datatypes in this subset are also describable by a generic universe, in this case Allais their library.

1. The programmer defines a native datatype T .

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2. A metaprogram generates the description D of T .

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4. A metaprogram takes D and P , generates a native function accordingly.

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─The process

THE PROCESS

1. The programmer defines a native datatype T .
2. A metaprogram generates the description D of T .
3. The programmer chooses a description P from a set of pre-defined generic programs.
4. A metaprogram takes D and P , generates a native function accordingly.

Now let us skip the introduction of datatype-generic programming and metaprograms. To get the whole picture, let's do a rundown of our alternative process for a programmer to invoke generic programs. Firstly they define a native datatype T that we know to be a syntax, instead of relying on any generic description. Then by metaprograms in our datatype-generic library, they get the datatype description D of T . Then the programmer can choose a generic program P to reify. This generic program is pre-defined by the generic library. Lastly, another metaprogram takes D and P , and gives the programmer a native, reified function definition that works on S , as if everything is defined by hands.

1. The programmer defines a native datatype T .
2. A metaprogram generates the description D of T .
3. The programmer provides a proof S of D that says T is indeed a syntax.
4. The programmer chooses a description P from a set of pre-defined generic programs.
5. A metaprogram takes D , S , and P , generates a native function accordingly.

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─The process

THE PROCESS

1. The programmer defines a native datatype T .
2. A metaprogram generates the description D of T .
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4. The programmer chooses a description P from a set of pre-defined generic programs.
5. A metaprogram takes D , S , and P , generates a native function accordingly.

But this process is actually not sufficient, how does the generic library know S is a syntax? To pre-define a datatype-generic program that work specifically on a family of syntax, we must also pre-define a predicate that says a datatype is indeed a syntax. So the generic library should also provide predicates that constrain general datatypes, saying they are syntaxes. And the programmer must provide a proof of this predicate when choosing a generic program to reify. So the process we just showed actually requires an extra step, that is the programmer providing a proof of the datatype T being a syntax.

FLOW CHART

User $\text{Lam} : \text{Type} \rightarrow \text{List Type} \rightarrow \text{Set}$

Allais et al.'s
library (ported)

Metaprograms

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─Flow Chart

User
Lam : Type \rightarrow List Type \rightarrow SetAllais et al.'s
library (ported)

Metaprograms

Let's do a rundown again, but with our actual definitions. We have three worlds here, the user's own world, the syntax-datatype-generic library, and metaprograms. Say the user defines a language called Lam, indexed by the type and context of a term.

FLOW CHART

User

$\text{Lam} : \text{Type} \rightarrow \text{List Type} \rightarrow \text{Set}$

Allais et al.'s
library (ported)

$\text{LamD} : \text{DataD}$

genDataD

Metaprograms

2022-09-10

Syntax-Generic Operations, Reflectively Reified

└─ Elaborator Reflection to the Rescue

└─ Flow Chart

They can invoke a metaprogram `genDataD`, which generates a datatype description of type `DataD`, we call it `LamD`.

FLOW CHART



FLOW CHART

User

$\text{Lam} : \text{Type} \rightarrow \text{List Type} \rightarrow \text{Set}$

$\text{LamSyn} : \text{Syntax LamD}$

Allais et al.'s
library (ported)

$\text{LamD} : \text{DataD}$

genDataD

Metaprograms

2022-09-10

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

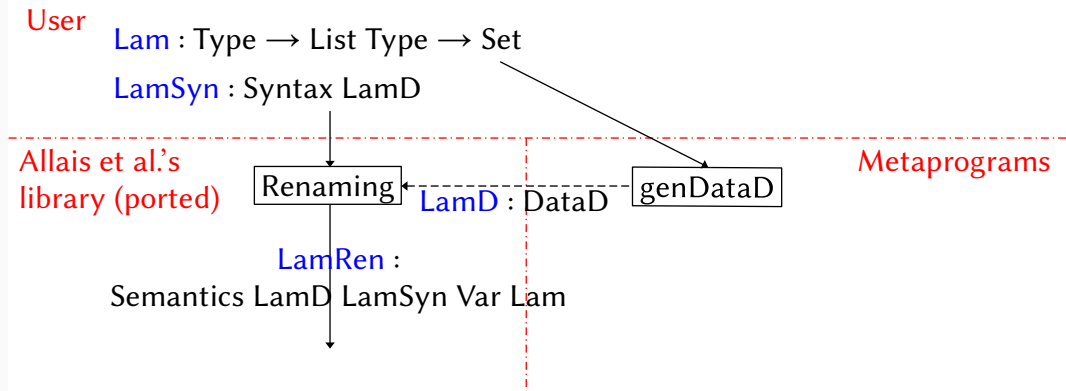
└─Flow Chart

Then to use the generic library, the user must provide a proof that the description we just generated is syntax. Here Syntax is the predicate, and we call the proof LamSyn.

FLOW CHART



FLOW CHART



2022-09-10

Syntax-Generic Operations, Reflectively Reified

└─ Elaborator Reflection to the Rescue

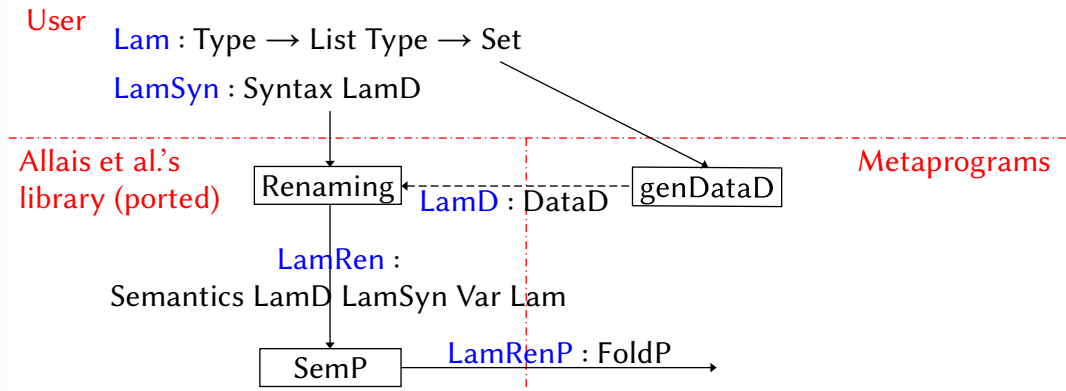
└─ Flow Chart

Then it's the job for the generic library. Renaming here is a function that generates syntax-generic programs. Syntax-generic programs are represented as Semantics.

FLOW CHART



FLOW CHART



2022-09-10

Syntax-Generic Operations, Reflectively Reified

└─ Elaborator Reflection to the Rescue

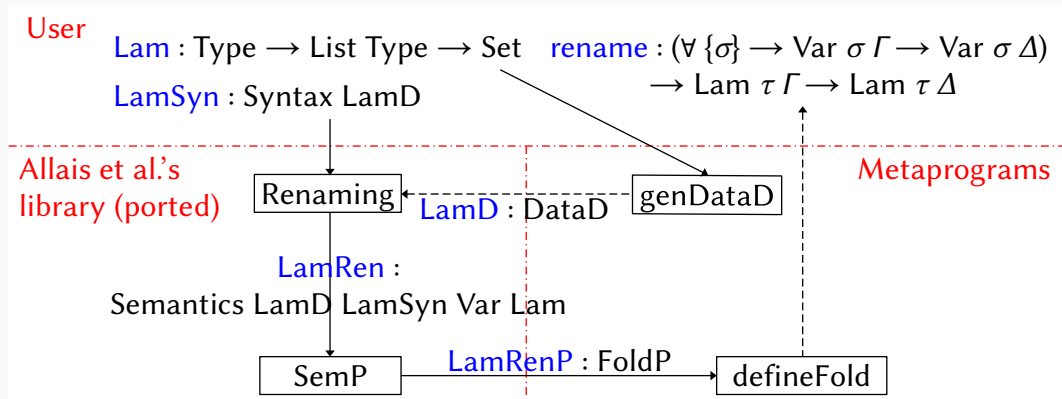
└─ Flow Chart

FLOW CHART



But our metaprograms that generates native functions are for datatype-generic programs. so we need another translation in the generic library that is from syntax-generic programs to datatype-generic programs. SemP is this translation and FoldP is the type of datatype-generic fold programs.

FLOW CHART

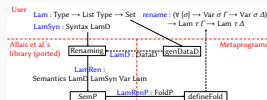


Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─Flow Chart

FLOW CHART



Finally, defineFold is a metaprogram that generates actual function definitions. So the user doesn't need to understand the detailed definitions in the generic library, they can write native datatype and get native function definitions if they know what metaprograms and functions to call.

In “Datatype-generic Programming Meets Elaborator Reflection”:

- *DataD*
- *FoldP* for folds (and *IndP* for inductions)
- metaprogram *genDataD*
- metaprogram *defineFold*

In “Datatype-generic Programming Meets Elaborator Reflection”:

- *DataD*
- *FoldP* for folds (and *IndP* for inductions)
- metaprogram *genDataD*
- metaprogram *defineFold*

In this work:

- predicate *Syntax* on *DataD* that captures *Desc*.
- function *SemP* that generates *FoldP* from *Syntax* proofs.

Syntax-Generic Operations, Reflectively Reified

└─ Elaborator Reflection to the Rescue

└─ Some of Our Contributions

The datatype and program descriptions are provided by our other work we mentioned, so are the metaprograms. In this work we provide the Syntax predicate and the translations built around it.

In "Datatype-generic Programming Meets Elaborator Reflection":

- *DataD*
- *FoldP* for folds (and *IndP* for inductions)
- metaprogram *genDataD*
- metaprogram *defineFold*

In this work:

- predicate *Syntax* on *DataD* that captures *Desc*.
- function *SemiP* that generates *FoldP* from *Syntax* proofs.

$Syntax : Set \ell \rightarrow DataD \rightarrow Set \omega$

Desc are captured by *Syntax* as each:

$Syntax : Set \ell \rightarrow DataD \rightarrow Set\omega$

Desc are captured by *Syntax* as each:

- has a variable rule,

$Syntax : Set \ell \rightarrow DataD \rightarrow Set\omega$

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Desc are captured by *Syntax* as each:

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- has two indices, I and *List* I , and

$Syntax : Set \ell \rightarrow DataD \rightarrow Set\omega$

Desc are captured by *Syntax* as each:

- has a variable rule,
- is not universe polymorphic,
- has two indices, I and $List\ I$, and
- supports context extensions.

Syntax-Generic Operations, Reflectively Reified

└─ Elaborator Reflection to the Rescue

└─ The *Syntax* Predicate

Syntax : $\text{Set } \ell \rightarrow \text{DataD} \rightarrow \text{Set} \omega$

Data are captured by *Syntax* as each:

- has a variable rule,
- is not universe polymorphic,
- has two indices, *I* and *List I*, and
- supports context extensions.

Unfortunately our time is limited, we can't go into the details of the Syntax predicate.

Does *PCF* satisfies *Syntax*?

```
data PCF : Ty → Context → Set where  
  `var  : Var σ Γ → PCF σ Γ  
  `app  : PCF (σ ⇒ τ) Γ → PCF σ Γ → PCF τ Γ  
  `lam  : PCF τ (σ :: Γ) → PCF (σ ⇒ τ) Γ  
  `zero : PCF `ℕ Γ  
  `suc_ : PCF `ℕ Γ → PCF `ℕ Γ
```

THE *Syntax* PREDICATE

Does *PCF* satisfies *Syntax*?

Proof of *PCF* being *Syntax*:

data *PCF* : *Ty* \rightarrow *Context* \rightarrow *Set* **where**

`var : *Var* σ *Γ* \rightarrow *PCF* σ *Γ*

`app : *PCF* ($\sigma \Rightarrow \tau$) *Γ* \rightarrow *PCF* σ *Γ* \rightarrow *PCF* τ *Γ*

`lam : *PCF* τ ($\sigma :: \Gamma$) \rightarrow *PCF* ($\sigma \Rightarrow \tau$) *Γ*

`zero : *PCF* \mathbb{N} *Γ*

`suc_ : *PCF* \mathbb{N} *Γ* \rightarrow *PCF* \mathbb{N} *Γ*

SyntaxPCF : *Syntax* *Ty* (*genDataD* *PCF*)

SyntaxPCF = _

, *refl*

, (*refl* , *refl*)

, _

, *refl*

, *refl*

, (_ , _ , _ , *refl* , (λ _ \rightarrow *refl*))

, (_ , _ , _ , *refl* , (λ _ \rightarrow *refl*))

, (_ , _ , _ , *refl* , (λ _ \rightarrow *refl*))

, (_ , _ , _ , *refl* , (λ _ \rightarrow *refl*))

, *tt*

Syntax-Generic Operations, Reflectively Reified

└─Elaborator Reflection to the Rescue

└─The *Syntax* Predicate

We were worried that the Syntax proof might be too complicated, because if it's too complicated it would not save any efforts and nobody would want to use it. Fortunately in this case it's pretty straightforward. Suppose we define a language PCF, which has five constructors. For this PCF datatype to be a syntax, a lot of things has to be considered, like the datatype's universe level, numbers of parameters and indices, and constraints on each field of each constructor. In this case the parameters must be empty and there must be exactly two indices, a type and a list of type. It turns out most of these are equality proofs, and if the datatype truly is a syntax, these can be proved by simply using the refl constructor. Proofs of any datatypes being syntaxes look pretty much the same, so it is possible to write yet another metaprogram for generating such proofs. Which is a future work we haven't done yet.

```

data PCF : Ty → Context → Set where
  'var  : Var σ F → PCF σ F
  'app  : PCF (σ → τ) F → PCF σ F → PCF τ F
  'lam  : PCF τ (σ :: F) → PCF (σ → τ) F
  'zero : PCF 'H F
  'sucn : PCF 'H F → PCF 'H F

SyntaxPCF : Syntax Ty (genDataD PCF)
SyntaxPCF =
  'var  , refl
  'app  , (refl , refl)
  'lam  , (refl , refl)
  'zero , refl
  'sucn , refl
  , (λ w → ..., refl , (λ w → refl))
  , (λ w → ..., refl , (λ w → refl))
  , (λ w → ..., refl , (λ w → refl))
  , (λ w → ..., refl , (λ w → refl))
  , H

```

Discussion

TOWARDS DATATYPE-GENERIC LIBRARIES FOR SYNTAXES?

Do we really need syntax-generic libraries?

Syntax-Generic Operations, Reflectively Reified

└ Discussion

└ Towards Datatype-generic Libraries for Syntaxes?

As you can see, our work has a lot to be done. Since we are running out of time, I would like to address one issue that's probably the elephant in the room. What if people actually don't want syntax-generic operations at all? Maybe they only define a language once in a while, and it's not worth the time looking up what libraries they can use. Or, since researchers define languages with new features all the time, maybe it's common for them to come up something no generic universes can cover. Our framework could still be help in that case, maybe our metaprograms can analyse the constructors in a datatype and determine which of them fit in a universe, and generate functions that are partially defined, then leave the uncertain parts to the programmer. So what do you think about it? Please share with us your concerns or what you think this framework can be going. Thank you all for listening.