STA261 MANAJEMEN DATA RELASIONAL

Model Data Relasional

DEPARTEMEN STATISTIKA
FAKULTAS MATEMATIKA DAN ILMU PENGETAHUAN ALAM
INSTITUT PERTANIAN BOGOR

SEMESTER GANJIL 202!/2022

Relational Data model

- Introduced by Ted Codd of IBM Research in 1970
- The model uses the concept of a mathematical relation
- First commercial implementations available in early 1980s by IBM and Oracle
- Has been implemented in a large number of commercial system
- Popular Relational DBMS: Oracle, DB2, MySQL, PostgreSQL,
 SQLite
- Preceded by hierarchical and network models

Relational Model Concept

- Represents database as a collection of relations
- Each relation resembles a table of values
 - o Row
 - ✓ Represents a collection of related data values
 - ✓ Represents a fact that typically corresponds to a real world entity or relationship
 - Table name and column names
 - ✓ Interpret the meaning of the values in each row
 - Formal Terminology
 - \checkmark Row \rightarrow Tuple
 - ✓ Column header → Attribute
 - ✓ Table → Relation

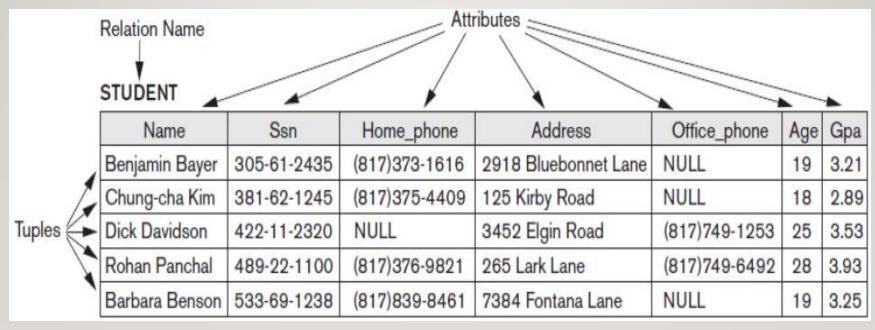


Figure 5.1 The attributes and tuples of a relation STUDENT

Equivalent Terminology

Formal (Relational Model)	Non Formal
Relation	Table (File)
Tuple	Row (Record)
Attribute	Column Header (Field)
Domain	All possible column values
Schema of a relation	Table definition
State of the relation	Populated table

Characteristics of Relations

- Ordering of **tuples in a relation**
 - Relation defined as a set of tuples
 - o Elements have no order among them

STUDENT

Name	Ssn	Home_phone	Address	Office_phone	Age	Gpa
Dick Davidson	422-11-2320	NULL	3452 Elgin Road	(817)749-1253	25	3.53
Barbara Benson	533-69-1238	(817)839-8461	7384 Fontana Lane	NULL	19	3.25
Rohan Panchal	489-22-1100	(817)376-9821	265 Lark Lane	(817)749-6492	28	3.93
Chung-cha Kim	381-62-1245	(817)375-4409	125 Kirby Road	NULL	18	2.89
Benjamin Bayer	305-61-2435	(817)373-1616	2918 Bluebonnet Lane	NULL	19	3.21

Characteristics of Relations

- Ordering of values within a tuple
 - Order of attributes and values is not important
 - As long as correspondence between attributes and values maintained
- Alternative definition of a relation
 - Tuple considered as a set of (<attribute>, <value>) pairs
 - Each pair gives the value of the mapping from an attribute A_i to a value v_i from dom (A_i)

Relational Model Constraints

- Classification of *Relational Integrity Constraints*
 - Key constraint
 - Entity integrity constraint
 - o Referential Integrity constraint
 - Domain constraint

Key Constraints and Constraints on NULL Values

Key Constraints : **No two tuples** can have the same combination of values for all their attributes.

Superkey (SK)
 an attribute or set of attributes that guarantee that no two distinct tuples R
 can have the same value for SK

Key Constraints and Constraints on NULL Values

Key Constraints: No two tuples can have the same combination of values for all their attributes.

Key

- Superkey of R
- Key satisfies two properties:
 - 1. Two distinct tuples in any state of relation **cannot** have identical values for (all) attributes in key
 - 2. Minimal superkey: cannot be removed and still have **uniqueness** constraint

Key Constraints and Constraints on NULL Values

Key Constraints: No two tuples can have the same combination of values for all their attributes.

- Candidate key
 - Relation schema may have more than one key
- **Primary key** of the relation
 - Designated among candidate keys
 - Underline attribute

SSN	FName	LName	BirthDate	Sex	Address
0606007800	Ahmad	Zakky	10-4-87	L	Jakarta
0607001123	Gede	Saraswati	19-9-87	P	Denpasar
0607120012	Bayu	Wirawan	12-12-86	L	Jimbaran
0607121023	Satya	Wirawan	12-12-86	L	Jimbaran
0607131240	Fira	Bahira	1-3-87	P	Jakarta
0607132222	Nayla	Putri	1-9-86	P	Depok

Super key: SSN, {SSN, Lname}, {FName, BirthDate}, {FName, Sex}, ...

Candidate key: SSN, Fname

Primary Key: SSN

Alternate Key: FName

CAR

License_number	Engine_serial_number	Make	Model	Year
Texas ABC-739	A69352	Ford	Mustang	02
Florida TVP-347	B43696	Oldsmobile	Cutlass	05
New York MPO-22	X83554	Oldsmobile	Delta	01
California 432-TFY	C43742	Mercedes	190-D	99
California RSK-629	Y82935	Toyota	Camry	04
Texas RSK-629	U028365	Jaguar	XJS	04

Figure 5.4 The CAR relation, with two candidate keys: License_number and Engine_serial_number

Relational database schema

EMPLOYEE									
Fname	Minit	Lname	San	Bdate	Address	Sex	Salary	Super_ssn	Dno
DEPARTM	DEPARTMENT								
Dname	Dname Dnumber Mgr_ssn Mgr_start_date								
DEPT_LO	CATION	s							
Dnumbe	r Dloo	cation							
PROJECT					_				
Pname Pnumber Plocation Dnum									
WORKS_ON									
Essn Pno Hours									
DEPENDENT						S	igure 5.5 chema diagram for the OMPANY relational		
<u>Essn</u>	Depend	lent_name	Sex	Bdate	Relations	ship			atabase schema.

Integrity, Referential Integrity, and Foreign Keys

Entity integrity constraint

No primary key value can be NULL

Referential integrity constraint

Specified between two relations

Maintains consistency among tuples in two relations

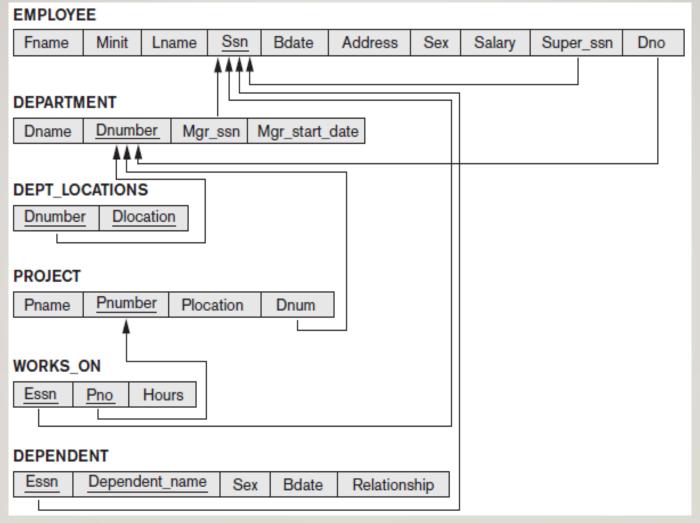
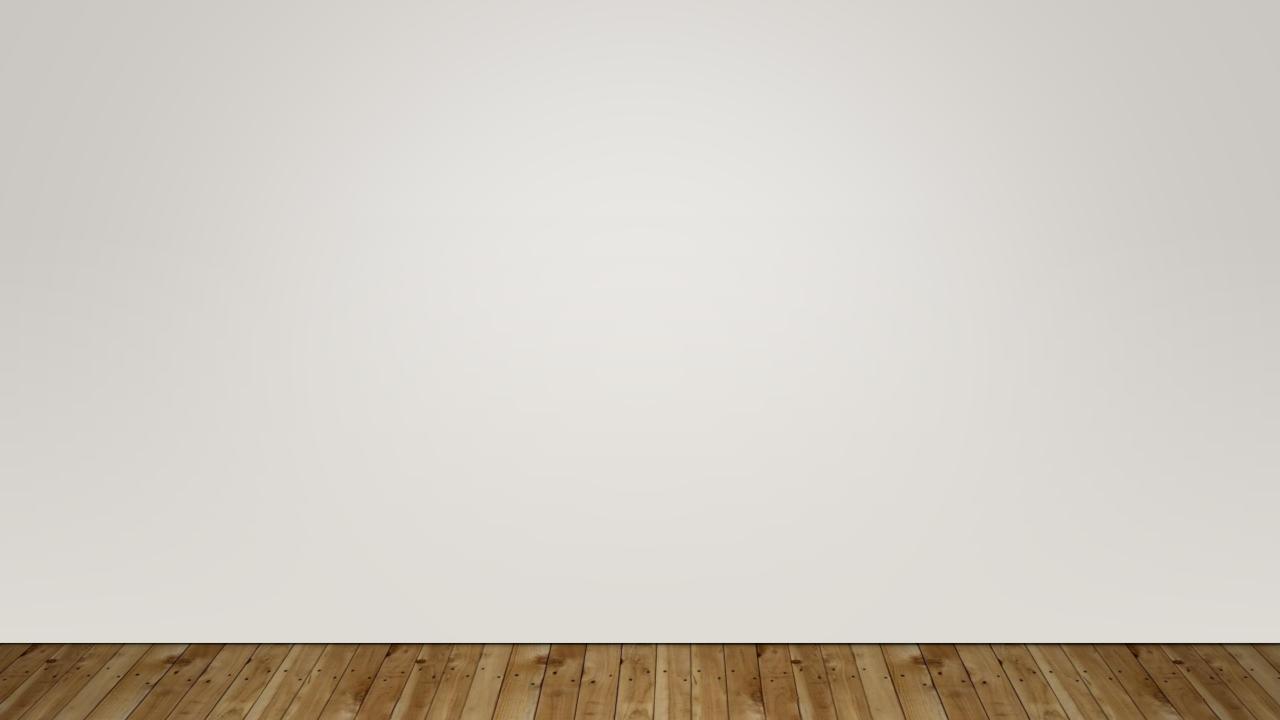


Figure 5.7 Referential integrity constraints displayed on the COMPANY relational database schema



Diketahui:

Basis data yang memproses order (pemesanan) pada sebuah perusahaan memiliki 6 relasi berikut:

```
CUSTOMER (Cust#, Cname, City)
ORDER (Order#, Odate, Cust#, Ord_Amt)
ORDER_ITEM (Order#, Item#, Qty)
ITEM (Item#, Unit_price)
SHIPMENT (Order#, Warehouse#, Ship_date)
WAREHOUSE (Warehouse#, City)
```

- I. Ord_Amt mengacu pada jumlah harga pada satu kali order.
- 2. Odate menyatakan tanggal pemesanan dilakukan
- 3. Ship_date menyatakan tanggal pengiriman barang yang dipesan customer dari gudang.
- a) Buatlah entity relationship untuk basis data tersebut.
- b) Buatlah **referential integrity constraints** untuk basis data tersebut.

 Asumsikan bahwa suatu order dapat mengambil barang dari beberapa gudang (warehouse).

 Nyatakan **foreign key** yang mungkin untuk skema basis data ini.

Model Data Relasional