

#### **Control Hijacking**

# Basic Control Hijacking Attacks

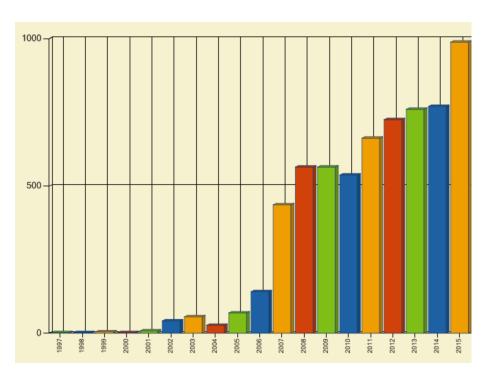
# Control hijacking attacks

- Attacker's goal:
  - Take over target machine (e.g. web server)
    - Execute arbitrary code on target by hijacking application control flow

- Examples:
  - Buffer overflow and integer overflow attacks
  - Format string vulnerabilities
  - Use after free

#### First example: buffer overflows

- Extremely common bug in C/C++ programs.
  - First major exploit: 1988 Internet Worm. fingerd.



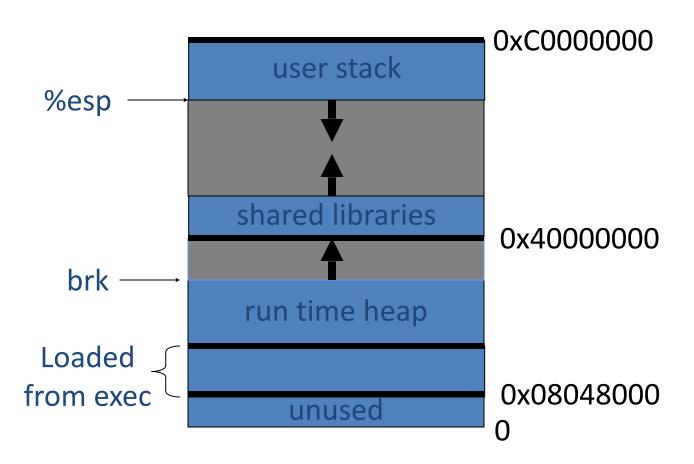
Source: web.nvd.nist.gov

#### What is needed

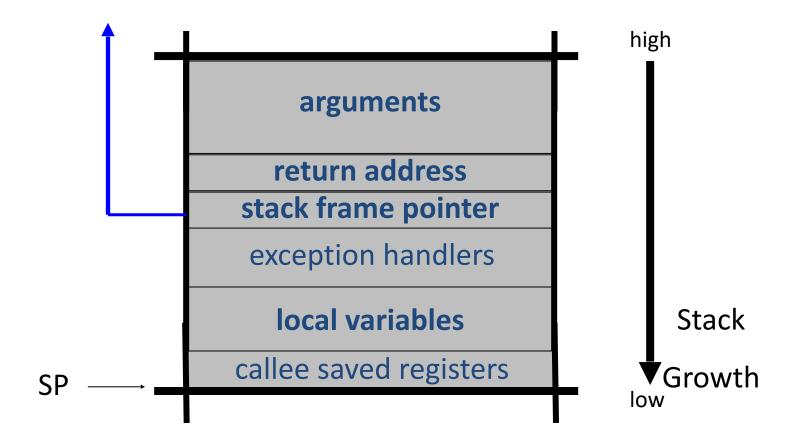
- Understanding C functions, the stack, and the heap.
- Know how system calls are made
- The exec() system call

- Attacker needs to know which CPU and OS used on the target machine:
  - Our examples are for x86 running Linux or Windows
  - Details vary slightly between CPUs and OSs:
    - Little endian vs. big endian (x86 vs. Motorola)
    - Stack Frame structure (Unix vs. Windows)

# Linux process memory layout



#### Stack Frame



#### What are buffer overflows?

Suppose a web server contains a function:

When func() is called stack looks like:

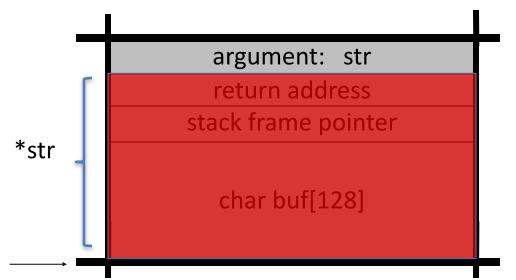
```
argument: str
return address
stack frame pointer

char buf[128]
```

```
void func(char *str) {
   char buf[128];
   strcpy(buf, str);
   do-something(buf);
}
```

#### What are buffer overflows?

```
What if *str is 136 bytes long?
After strcpy:
```



```
void func(char *str) {
   char buf[128];
   strcpy(buf, str);
   do-something(buf);
}
```

```
Problem: no length checking in strcpy()
```

# Basic stack exploit

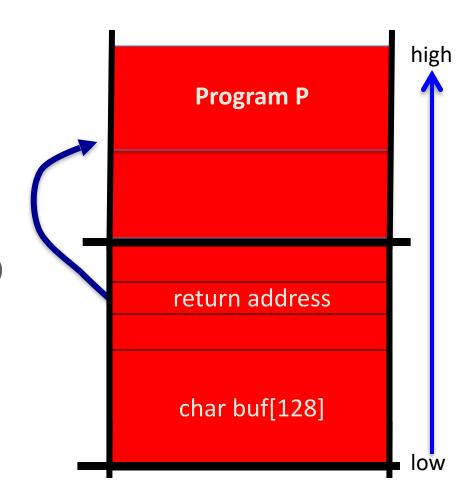
Suppose \*str is such that after strcpy stack looks like:

Program P: exec("/bin/sh")

(exact shell code by Aleph One)

When func() exits, the user gets shell!

Note: attack code P runs in stack.



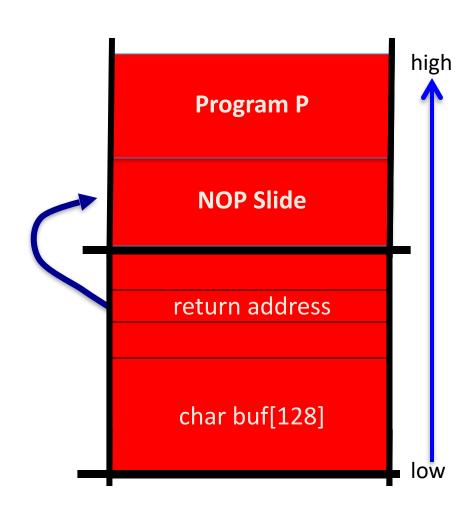
#### The NOP slide

Problem: how does attacker

determine ret-address?

Solution: NOP slide

- Guess approximate stack state when func() is called
- Insert many NOPs before program P:
   nop , xor eax,eax , inc ax



# Details and examples

- Some complications:
  - Program P should not contain the '\0' character.
  - Overflow should not crash program before func() exits.

- (in)Famous <u>remote</u> stack smashing overflows:
  - Overflow in Windows animated cursors (ANI). LoadAnilcon()
  - Buffer overflow in Symantec virus detection (May 2016)
     overflow when parsing PE headers ... kernel vuln.

# Many unsafe libc functions

```
strcpy (char *dest, const char *src)
strcat (char *dest, const char *src)
gets (char *s)
scanf ( const char *format, ... ) and many more.
```

- "Safe" libc versions strncpy(), strncat() are misleading
  - e.g. strncpy() may leave string unterminated.
- Windows C run time (CRT):
  - strcpy\_s (\*dest, DestSize, \*src): ensures proper termination

#### Buffer overflow opportunities

- Exception handlers: (Windows SEH attacks ... more on this later)
  - Overwrite the address of an exception handler in stack frame.

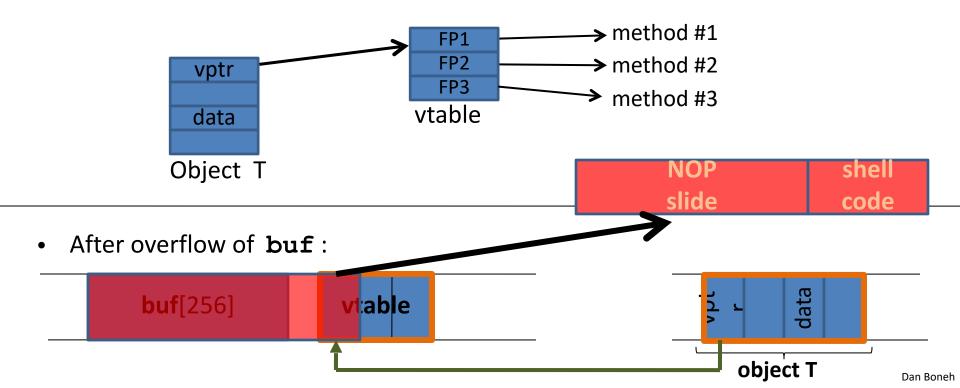
• Function pointers: (e.g. PHP 4.0.2, MS MediaPlayer Bitmaps)



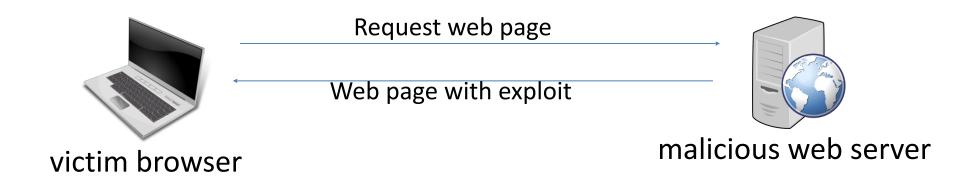
- Overflowing buf will override function pointer.
- Longimp buffers: longimp(pos) (e.g. Perl 5.003)
  - Overflowing buf next to pos overrides value of pos.

#### Heap exploits: corrupting virtual tables

Compiler generated function pointers (e.g. C++ code)



#### An example: exploiting the browser heap



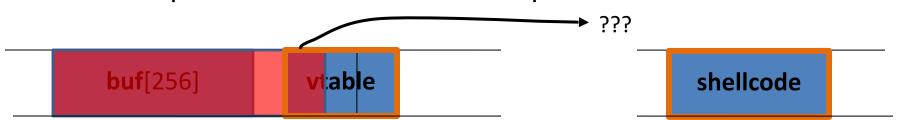
Attacker's goal is to infect browsers visiting the web site

How: send javascript to browser that exploits a heap overflow

# A reliable exploit?

```
<SCRIPT language="text/javascript">
shellcode = unescape("%u4343%u4343%...");  // allocate in heap
overflow-string = unescape("%u2332%u4276%...");
cause-overflow(overflow-string);  // overflow buf[]
</SCRIPT>
```

Problem: attacker does not know where browser places **shellcode** on the heap

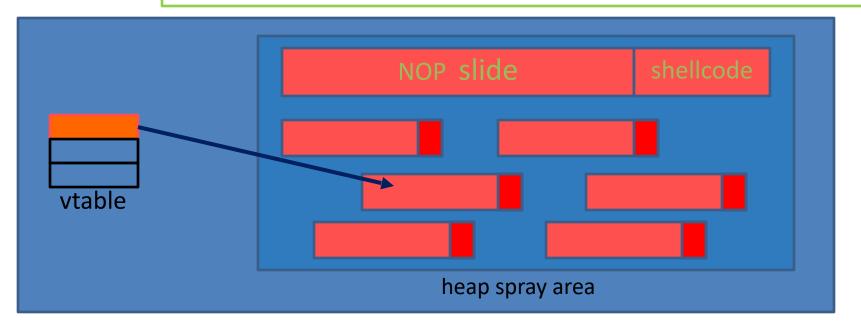


# Heap Spraying

[SkyLined 2004]

Idea:

- 1. use Javascript to spray heap with shellcode (and NOP slides)
- 2. then point vtable ptr anywhere in spray area



heap

# Javascript heap spraying

```
var nop = unescape("%u9090%u9090")
while (nop.length < 0x100000) nop += nop;

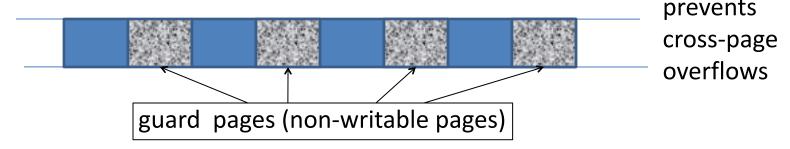
var shellcode = unescape("%u4343%u4343%...");

var x = new Array ()
for (i=0; i<1000; i++) {
    x[i] = nop + shellcode;
}</pre>
```

Pointing function-ptr almost anywhere in heap will cause shellcode to execute.

#### Ad-hoc heap overflow mitigations

- Better browser architecture:
  - Store JavaScript strings in a separate heap from browser heap
- OpenBSD and Windows 8 heap overflow protection:



Nozzle [RLZ'08]: detect sprays by prevalence of code on heap

# Finding overflows by fuzzing

- To find overflow:
  - Run web server on local machine
  - Issue malformed requests (ending with "\$\$\$\$")
    - Many automated tools exist (called fuzzers next week)
  - If web server crashes, search core dump for "\$\$\$\$" to find overflow location

Construct exploit (not easy given latest defenses)



#### **Control Hijacking**

# More Control Hijacking Attacks

# More Hijacking Opportunities

- Integer overflows: (e.g. MS DirectX MIDI Lib)
- **Double free**: double free space on heap
  - Can cause memory mgr to write data to specific location
  - Examples: CVS server
- Use after free: using memory after it is freed
- Format string vulnerabilities

# Integer Overflows

(see Phrack 60)

Problem: what happens when int exceeds max value?

$$c = 0x80 + 0x80 = 128 + 128 \qquad \Rightarrow c = 0$$

$$s = 0xff80 + 0x80 \qquad \Rightarrow s = 0$$

$$m = 0xffffff80 + 0x80 \qquad \Rightarrow m = 0$$

### An example

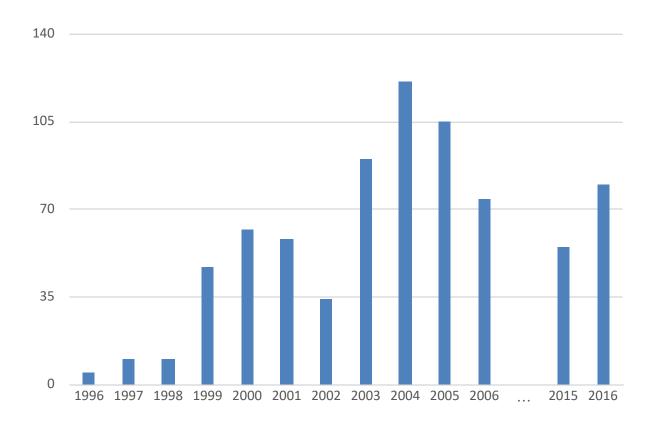
```
void func( char *buf1, *buf2, unsigned int len1, len2) {
    char temp[256];
    if (len1 + len2 > 256) {return -1}
                                              // length check
    memcpy(temp, buf1, len1);
                                               // cat buffers
    memcpy(temp+len1, buf2, len2);
    do-something(temp);
                                               // do stuff
```

```
What if len1 = 0x80, len2 = 0xffffff80 ?

\Rightarrow len1+len2 = 0
```

Second memcpy() will overflow heap!!

# Integer overflow exploit stats



Format string bugs

# Format string problem

```
int func(char *user) {
  fprintf( stderr, user);
}
```

- <u>Problem:</u> what if \*user = "%s%s%s%s%s%s%s" ??
  - Most likely program will crash: DoS.
  - If not, program will print memory contents. Privacy?
  - Full exploit using user = "%n"

#### Vulnerable functions

Any function using a format string.

```
Printing:

printf, fprintf, sprintf, ...

vprintf, vfprintf, vsprintf, ...
```

# Logging: syslog, err, warn

# **Exploit**

- Dumping arbitrary memory:
  - Walk up stack until desired pointer is found.
  - printf( "%08x.%08x.%08x.%08x|%s|")

- Writing to arbitrary memory:
  - printf( "hello %n", &temp) -- writes '6' into temp.
  - printf( "%08x.%08x.%08x.%08x.%n")

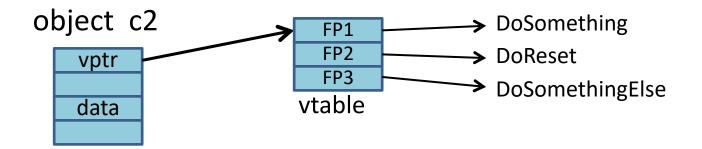
Use after free exploits

### IE11 Example: CVE-2014-0282 (simplified)

```
<form id="form">
 <textarea id="c1" name="a1" ></textarea>
 <input id="c2" type="text" name="a2" value="val">
</form>
                                                    Loop on form elements:
                                                       c1.DoReset()
<script>
                                                       c2.DoReset()
 function changer() {
    document.getElementById("form").innerHTML = "";
    CollectGarbage();
  document.getElementById("c1").onpropertychange = changer;
  document.getElementById("form").reset();
</script>
```

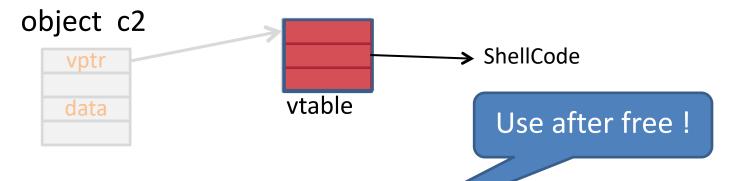
# What just happened?

c1.doReset() causes changer() to be called and free object c2



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c1.doReset() causes changer() to be called and free object c2



Suppose attacker allocates a string of time size as vtable

When c2.DoReset() is called, attacker gets shell

# The exploit

```
<script>
  function changer() {
    document.getElementById("form").innerHTML = "";
    CollectGarbage();
    --- allocate string object to occupy vtable location ---
  document.getElementById("c1").onpropertychange = changer;
  document.getElementById("form").reset();
</script>
```

Lesson: use after free can be a serious security vulnerability!!

## Next lecture ...

**DEFENSES** 

## THE END