Arithmetic Expressions

The following Exp datatype encodes the abstract syntax for arithmetic expressions. Note that it is recursively defined: Exp occurs in the right-hand side of its own definition.

1. Interpreter

• Write an interpreter eval :: Exp -> Int that evaluates arithmetic expressions.

Examples

```
Main> eval (Add (Mul (Const 2) (Const 4)) (Const 3))

11

Main> eval (Sub (Const 42) (Mul (Const 6) (Const 7)))

0
```

2. Compiler

Instead of evaluating an arithmetic expression directly, we can also compile it to a program of a simple stack machine and subsequently execute the program. We represent a program as a list of instructions. The instructions IAdd, ISub, IMul take the two topmost elements from the stack, perform the corresponding operation, and push the result onto the stack. The IPush instruction pushes the given value onto the stack. A stack is modelled by a list of integers.

```
data Inst = IPush Int | IAdd | ISub | IMul
  deriving (Show, Eq)

type Prog = [Inst]
type Stack = [Int]
```

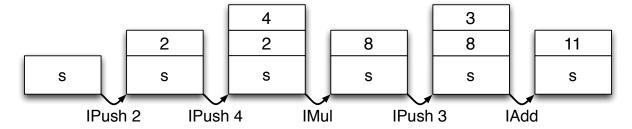
For example, the following arithmetic expression

```
Add (Mul (Const 2) (Const 4)) (Const 3)
```

is equivalent to the stack program

```
[IPush 2, IPush 4, IMul, IPush 3, IAdd]
```

The stack program leaves the result on top of the stack. The stack machine performs the following steps when executing the program on an initial stack s



• Write an execution function execute:: Inst -> Stack -> Stack that executes a single instruction. Since there are cases where the function can crash (e.g. a stack overflow in cases where we want to execute an IAdd instruction but the stack contains fewer than two elements), you can use the following exception-raising function where needed:

```
runtimeError :: Stack
runtimeError = error "Runtime error."
```

- Write a function run:: Prog -> Stack -> Stack that runs a whole program on a given initial stack.
- Write a compiler compile:: Exp -> Prog that compiles arithmetic expressions to stack machine programs. The compiled program should leave the result of the computation as the top element on the stack. Make sure that your compiler uses a left-to-right evaluation order and that it produces results equivalent to the interpreter, i.e. the following identity holds

```
forall (s :: Stack). run (compile e) s == (eval e) : s
```

Examples

```
Main> execute IAdd [4,5,6]
[9,6]

Main> execute ISub [4,5,6]
[1,6]

Main> execute (IPush 2) [4,5,6]
[2,4,5,6]

Main> run [IAdd, ISub] [4,5,6]
[-3]

Main> run [IAdd, ISub, IPush 7, IMul] [4,5,6,8]
[-21,8]
```

```
Main> run [IPush 1,IPush 2,IPush 3,IMul,ISub] []
[-5]
```

Main> compile (Sub (Const 1) (Mul (Const 2) (Const 3)))
[IPush 1,IPush 2,IPush 3,IMul,ISub]