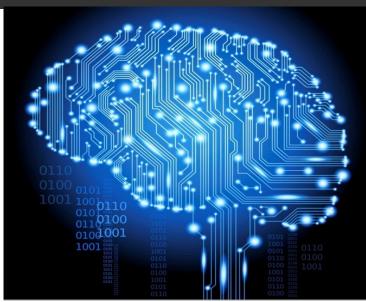


Information Technology

Recursive Sorts

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Overview

- To review what a "divide and conquer" algorithm is
- To review in more depth two different "divide and conquer" sorting algorithms:
 - Merge Sort
 - Quick Sort
- To be able to implement them and compare their efficiency for different classes of inputs





Divide and Conquer

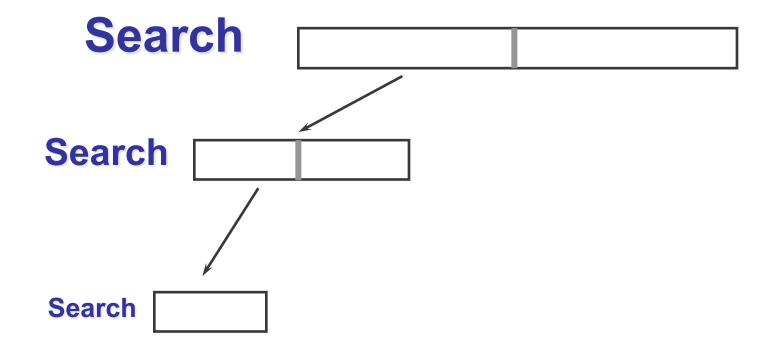
Divide-and-conquer algorithms

- One common recursive strategy: divide-and-conquer
 - Split the original problem into subproblems
 - Solve each subproblem independently
 - Combine their solutions to yield the final solution
- Such algorithms can be very efficient, especially when the subproblems have roughly the same size



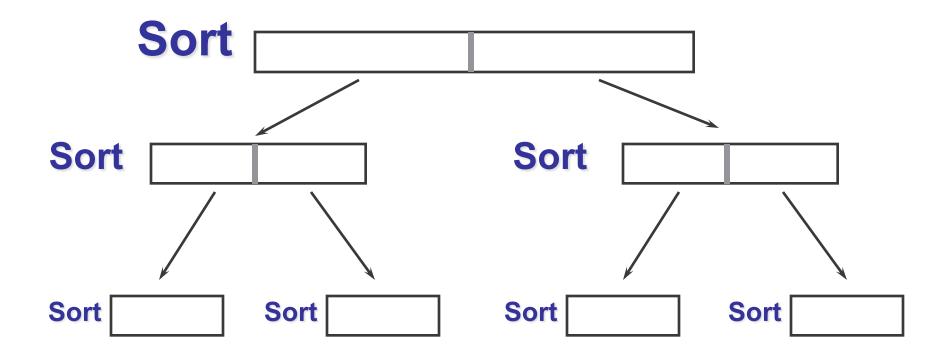
Searching by Divide and Conquer

Recall: Binary Search



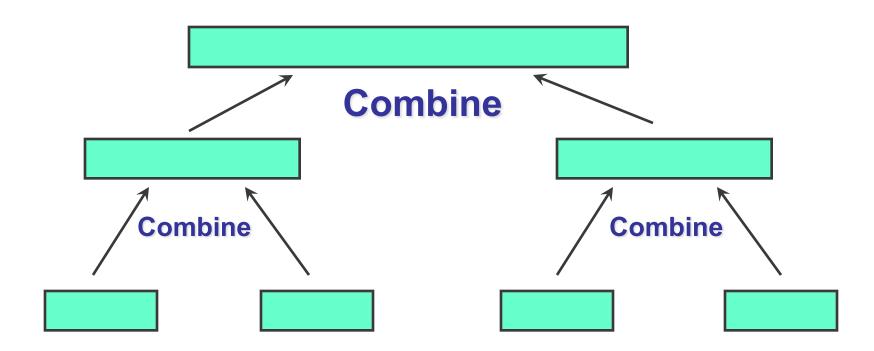


Sorting arrays by Divide and Conquer





Sorting arrays by Divide and Conquer





Sorting arrays by Divide and Conquer

The general strategy is (usually – not always -- for an array or a list implemented with arrays):

- Merge Sort has a simple split and a complex combine
- Quick Sort has a complex split and a simple combine

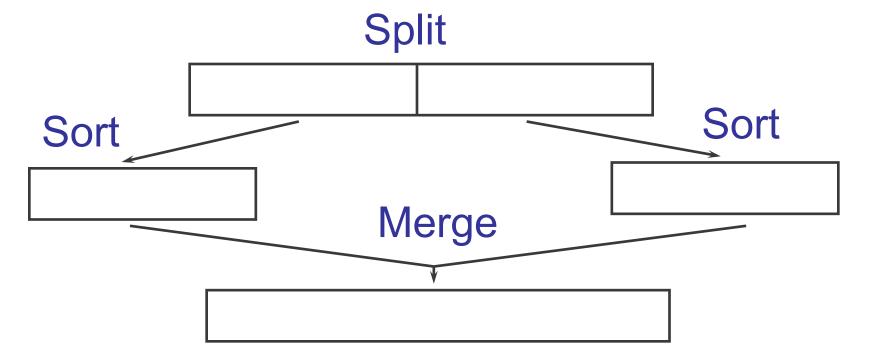




The split is trivial: splits the array into two halves

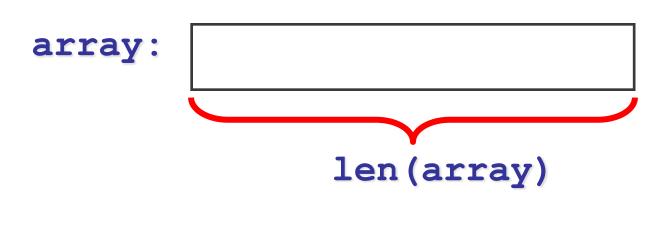
The combination is non-trivial: merges two sorted halves into a sorted

array

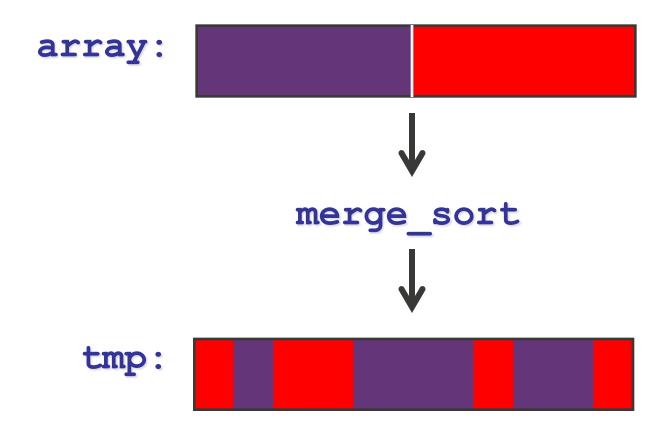


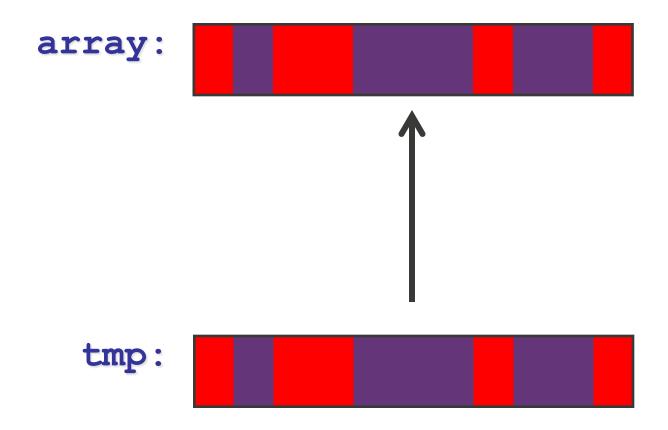


For this it requires a temporary array of the same size as the original array



tmp:









Merge Sort method

- So we need:
 - A temporary array to store the merged (partial) solution
 - Markers for the part of the array I am looking at (no need for new array)



Auxiliary method with all arguments I need

merge_sort_aux method

```
def merge_sort_aux(array: ArrayR, start: int, end: int, tmp: T) -> None:
    if not start == end: # 2 or more still to sort
        mid = (start + end) // 2
        # split into two halves
        merge_sort_aux(array, start, mid, tmp)
        merge_sort_aux(array, mid+1, end, tmp)
        # merge
        merge_arrays(array, start, mid, end, tmp)
        # copy tmp back into the original
        for i in range(start, end+1):
              array[i] = tmp[i]
```



a: 3 5 15 28 30 32 **b:** 10 14 22 43 50

start: 0 mid: 5 end: 10

tmp:

a: 3 5 15 28 30 32 b: 10 14 22 43 50 ia=0 ib=6





a: 3 5 15 28 30 32 b: 10 14 22 43 50 ia=2 ib=6



























Done!

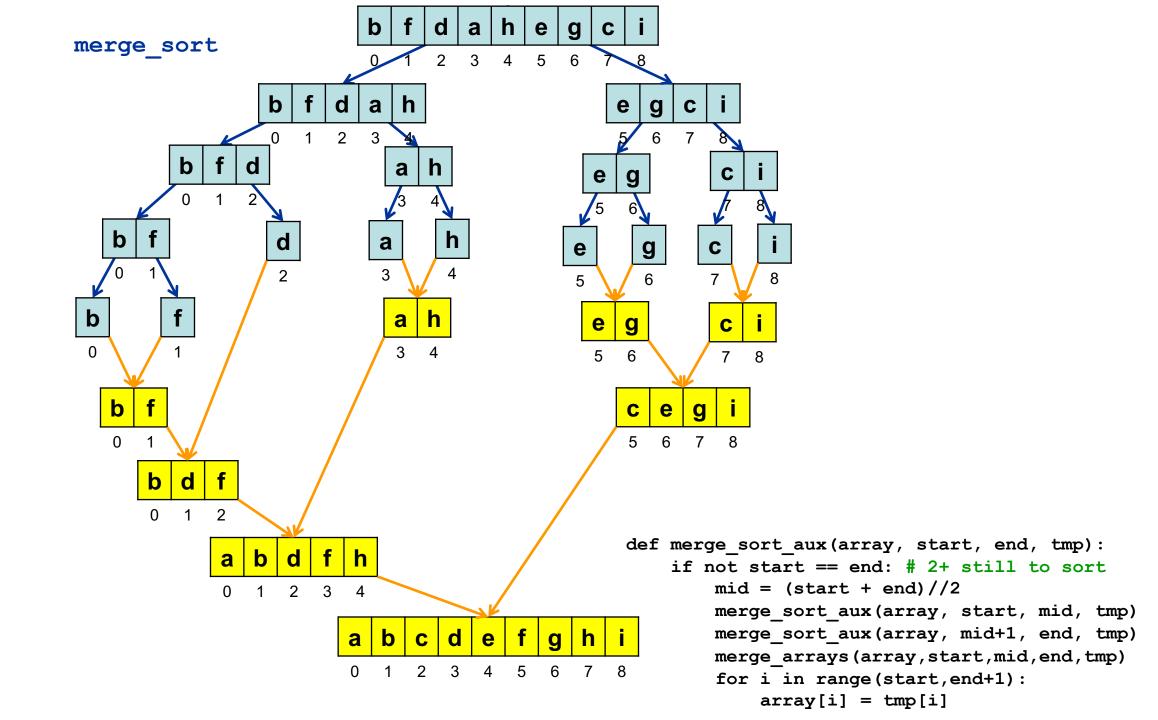


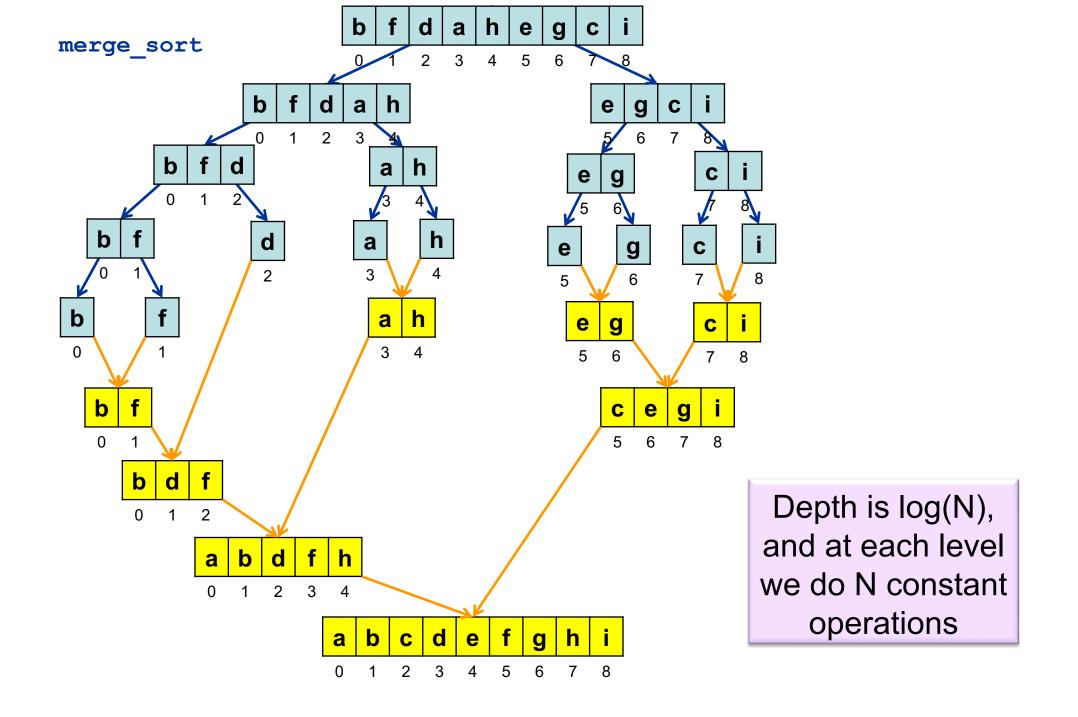


merge_arrays method

```
def merge arrays(a: ArrayR, start: int, mid: int, end: int, tmp: T) -> None:
    ia = start
    ib = mid+1
    for k in range(start, end+1):
        if ia > mid: # half1 finished, copy half2
           tmp[k] = a[ib]
           ib += 1
        elif ib > end: # half2 finished, copy half1
           tmp[k] = a[ia]
           ia += 1
        elif a[ia] <= a[ib]: # a[ia] is the item to copy
           tmp[k] = a[ia]
           ia += 1
        else:
           tmp[k] = a[ib] # a[ib] is the item to copy
           ib += 1
```







Merge Sort Analysis

- This is typically the method that you would naturally use when sorting a pile of books, CDs cards, etc.
- Most of the work is in the merging
- Uses more space than other sorts
- Takes O(n log(n))*CompEq, where n is the number of elements in the array and CompEq is the complexity of the comparison
- Close to optimal in number of comparisons
 - Good for languages where comparison is expensive

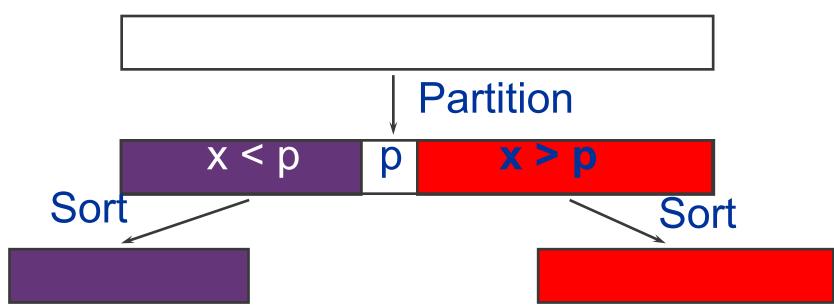




Quick Sort

Quick Sort

- The split is non-trivial: partitions the array into two:
 - One with all elements smaller than the pivot p
 - Another with all elements greater than the pivot
- Combination is trivial: append with pivot in the middle





Quick Sort: choosing the pivot

- It must be an element in the array
- Ideally, its value is the median of all values
 - This is does not mean is in the middle position! (unless the list is already sorted)
- Why: best when the subarrays are about the same size
- Hitch: how to pick the median element? (or similar)
- One strategy: pick a small sample of elements (e.g. first, last, mid) and choose the median of these
- Or you could just pick a random element
- First element is usually a bad idea: presorted input



Example: Quick Sort

array: 5 89 35 14 24 15 37 13 20 7 70

start:0 end:10

Example: Quick Sort

array: 5 89 35 14 24 15 37 13 20 7 70

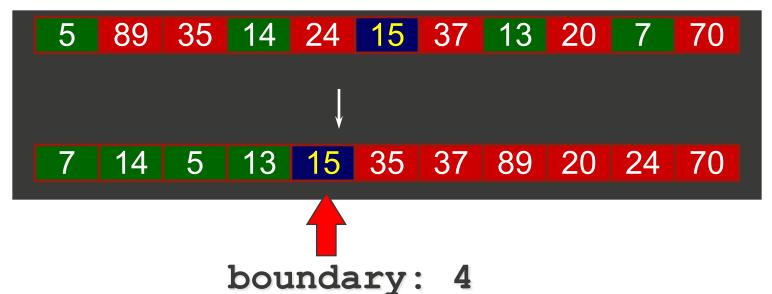
Say we randomly choose 15 as the pivot.

It happens to be in the middle of the list, but that's just a coincidence!

Example: Quick Sort

array: 5 89 35 14 24 15 37 13 20 7 70

partition:





Quick Sort method

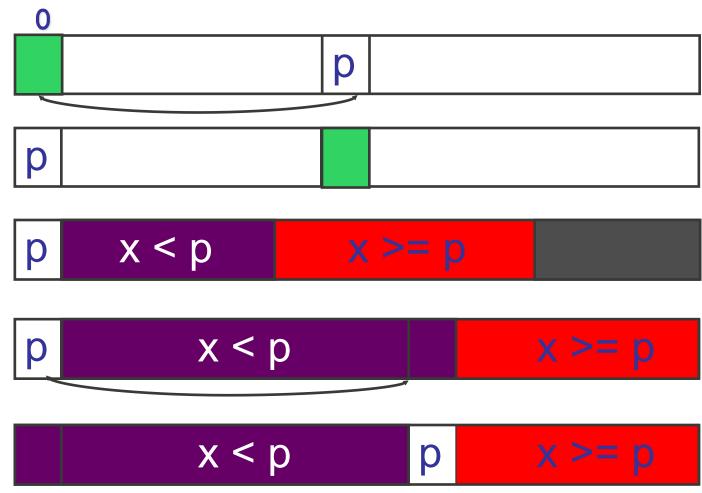
```
But no need for a temporary array this time

start = 0
end = len(array) -1
quick_sort_aux(array; start, end)

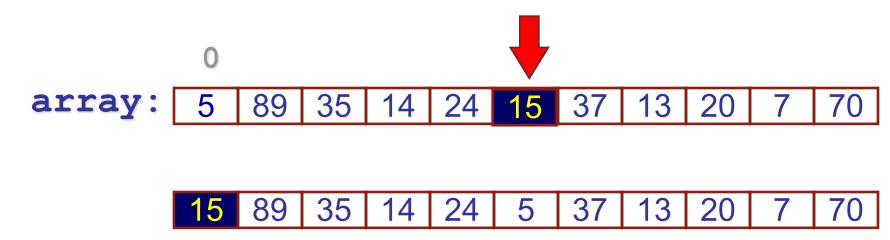
def quick_sort_aux(array; ArrayR, start: int, end: int) -> None: if start < end:
    boundary = partition(array, start, end)
    quick_sort_aux(array, start, boundary-1)
    quick_sort_aux(array, boundary+1, end)
```

def quick sort(array: ArrayR) -> None:

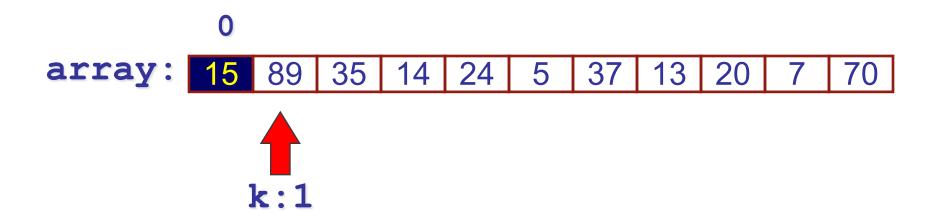
Partition: Checklist



randomly pick element in position 5









boundary:0

array: 15 89 35 14 24 5 37 13 20 7 70

k:1

boundary:0

array: 15 89 35 14 24 5 37 13 20 7 70

k:2



boundary:0

array: 15 89 35 14 24 5 37 13 20 7 70

k:3



boundary:1

array: 15 14 35 89 24 5 37 13 20 7 70

k:3

boundary:1

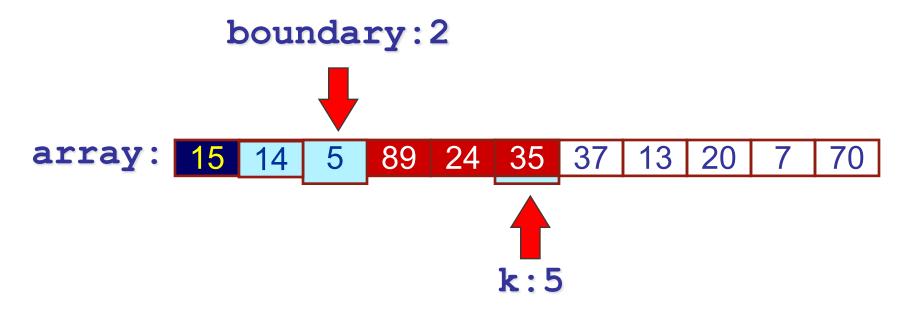
array: 15 14 35 89 24 5 37 13 20 7 70

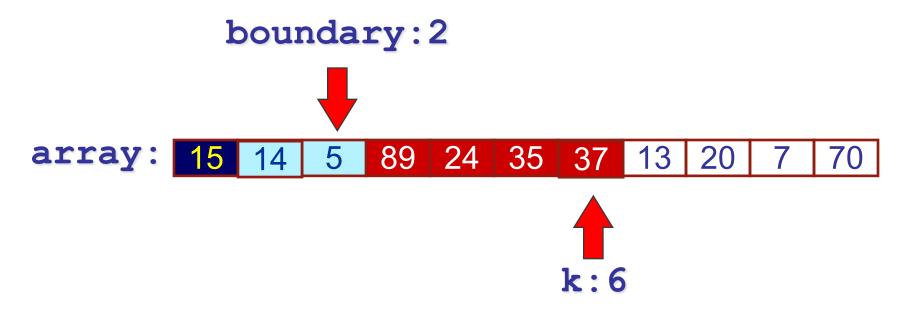
k:4

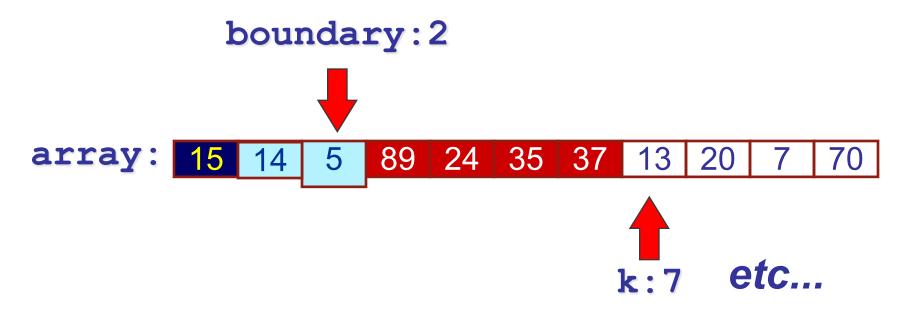
boundary:1

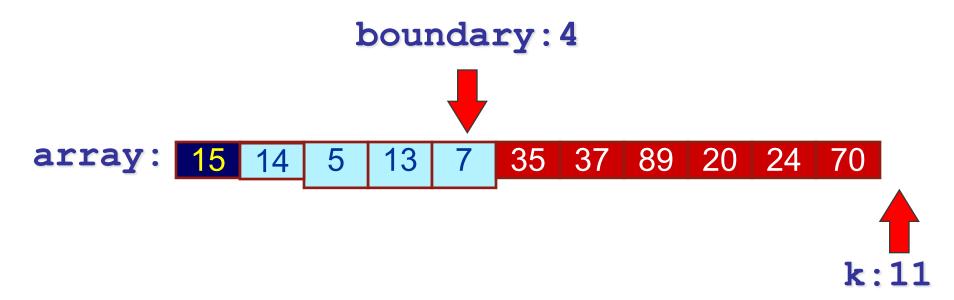
array: 15 14 35 89 24 5 37 13 20 7 70

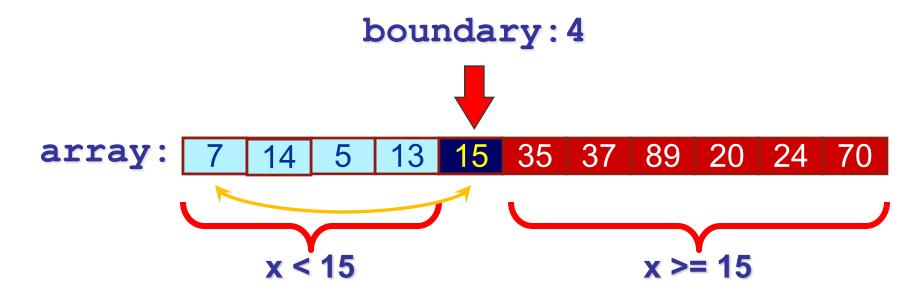
k:5



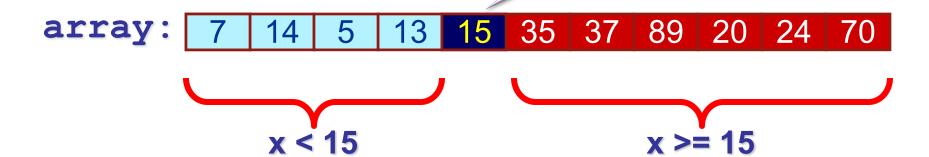




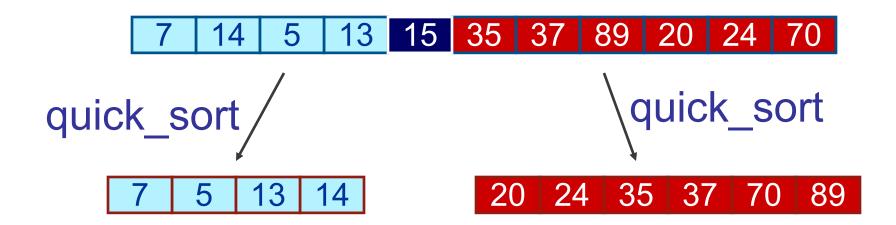




pivot now in correct position







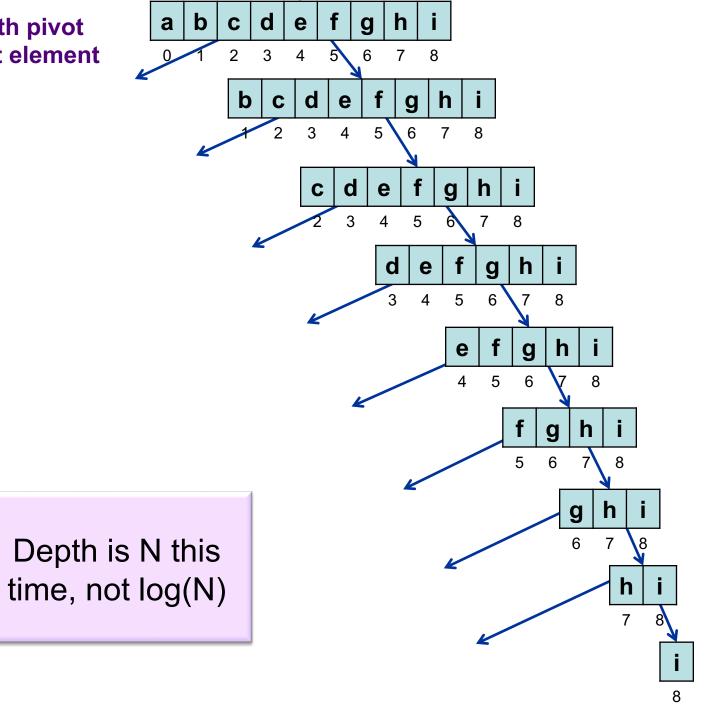


Partition Method

```
def partition(array: ArrayR, start: int, end: int) -> int:
    mid = (start+end)//2
                                        For simplicity, let
    pivot = array[mid] ----
                                         me choose mid
    swap(array, start, mid)
                                        element here...
    boundary = start
    for k in range(start+1, end+1):
        if array[k] < pivot:</pre>
            boundary += 1
            swap(array, k, boundary)
    swap(array, start, boundary)
    return boundary
```



quick_sort with pivot
being the first element



Quick Sort Analysis

- Most of the work done in partitioning
- Remember: careful about choice of pivot
 - Example: 2, 4, 6, 7, 3, 1, 5 what happens if we use the mid element as pivot?
 - Often used: median among 3 random
- Complexity?
 - Best case takes O(n log(n))*CompEq time
 - Always pick median as pivot
 - Worst case takes O(n²)*CompEq time
 - Always pick min/max as pivot
- Main advantage over merge ksort:
 - no need to copy back (so quicksort has a smaller constant if the pivot is good)
- Still, for small arrays (≤20) it has significant overhead (due to the recursion); for those better use insertion sort



Summary

- Divide and Conquer algorithms
 - Relationship with recursive algorithms
- Merge Sort
 - Easy split
 - Complex part: merge method
- Quick Sort
 - Complex split: partition method
 - Easy combination