# 表达式语法分析实验报告

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# 1 实验题目

表达式语法分析

## 2 实验目的

熟悉并设计一个表达式的语法分析器

# 3 实验内容

- 1. 设计表达式的语法分析器算法
- 2. 编写代码并上机调试运行通过

# 4 概要设计

本语法分析器实现了LR(1)和LL(1)语法分析方法,对给定文法自动生成LR(1)和LL(1)分析表,在几个简单文法的测试下正常运行。

下面介绍具体设计实现。

SyntaxAnalyzer 类定义 (代码清单 1) 如下:

```
class SyntaxAnalyzer
{
public:
    SyntaxAnalyzer(const string& filePath);

Grammar G;

void display(const LROItem& item);
void display(const LR1Item& item);
void display(const set<LR0Item>& I);
void display(const set<LR1Item>& I);
void display(const set<LR1Item>& I);
```

```
13
    vector < set < LR1Item >> lr1ItemSetFamily;
14
    vector < vector < LR1ActionTabItem >> lr1ActionTab;
    vector < vector < unsigned int >> lr1GoToTab;
16
    vector < unsigned int > lr1Stack;
17
    string lr1Input;
18
19
    set < LROItem > LROClosure(const set < LROItem > & I);
20
    set < LR1Item > LR1Closure(const set < LR1Item > & I);
    set < LR1Item > LR1GoTo(const set < LR1Item >& I, char X);
22
    void calLR1ItemSetFamily();
23
    void buildLR1ParseTab();
24
    bool lr1SyntaxAnalyze(const string& myLR1Input);
25
26
27
    vector < LL1Item >> ll1ParseTab;
28
    vector < char > ll1Stack;
29
    string ll1Input;
31
    void buildLL1ParseTab();
32
    bool ll1SyntaxAnalyze(const string& myLL1Input);
33
34
 private:
35
36
 };
37
```

Listing 1: SyntaxAnalyzer 类定义代码清单

构造函数和 display 多态函数不再赘述,下面分别介绍LR(1)和LL(1)分析器的构造。

## 4.1 LR(1)语法分析器的设计实现

#### 4.1.1 LR1项的数据结构设计

代码如下(代码清单2)如下:

```
class LR1Item
{
public:
    LR1Item() = default;
    LR1Item(const LR0Item& myLR0Item, char
    myLookaheadSymbol);
LR1Item(unsigned int myLProductionRuleIdx, unsigned
    int myRProductionRuleIdx, unsigned int myRProductionRuleIdx, unsigned int myDotPos,
    char myLookaheadSymbol);
```

```
unsigned int lProductionRuleIdx;
    unsigned int rProductionRuleIdx;
    unsigned int dotPos;
    char lookaheadSymbol;
    friend bool operator < (const LR1Item& i1, const
13
     LR1Item& i2);
    friend bool operator ==(const LR1Item& i1, const
     LR1Item& i2);
    friend bool operator > (const LR1Item& i1, const
     LR1Item& i2);
    friend bool operator !=(const LR1Item& i1, const
     LR1Item& i2);
 private:
18
19
20 };
22 bool operator <(const LR1Item& i1, const LR1Item& i2)
    return i1.1ProductionRuleIdx < i2.1ProductionRuleIdx
      true: (i1.lProductionRuleIdx > i2.
     lProductionRuleIdx?
        false: (i1.rProductionRuleIdx < i2.</pre>
     rProductionRuleIdx?
          true: (i1.rProductionRuleIdx > i2.
     rProductionRuleIdx?
            false: (i1.dotPos < i2.dotPos?</pre>
28
              true: (i1.lookaheadSymbol < i2.
29
     lookaheadSymbol))));
30 }
 bool operator ==(const LR1Item& i1, const LR1Item& i2)
32
33 {
    return i1.lProductionRuleIdx == i2.
     lProductionRuleIdx
      && i1.rProductionRuleIdx == i2.rProductionRuleIdx
35
      && i1.dotPos == i2.dotPos
      && i1.lookaheadSymbol == i2.lookaheadSymbol;
37
bool operator >(const LR1Item& i1, const LR1Item& i2)
41 {
   return !(i1 < i2) && !(i1 == i2);
```

```
43 }
44
bool operator !=(const LR1Item& i1, const LR1Item& i2)
    return !(i1 == i2);
47
 }
49
 LR1Item::LR1Item(const LR0Item& myLR0Item, char
     myLookaheadSymbol):
    lProductionRuleIdx(myLROItem.lProductionRuleIdx),
     rProductionRuleIdx(myLROItem.rProductionRuleIdx),
     dotPos(myLROItem.dotPos),
    lookaheadSymbol(myLookaheadSymbol)
52
53 {
54
 }
55
56
 LR1Item::LR1Item(unsigned int myLProductionRuleIdx,
     unsigned int myRProductionRuleIdx, unsigned int
     myDotPos, char myLookaheadSymbol):
    lProductionRuleIdx(myLProductionRuleIdx),
     rProductionRuleIdx(myRProductionRuleIdx), dotPos(
     myDotPos), lookaheadSymbol(myLookaheadSymbol)
 {
59
60
 }
61
```

Listing 2: LR1项的数据结构设计代码清单

## 4.1.2 LR1Closure 成员函数的设计实现

LR1Closure 类定义(代码清单3)如下:

```
13
      for(auto ptr = ret.begin(); ptr != ret.end(); ++
     ptr)
        if(ptr->dotPos < G.P[ptr->lProductionRuleIdx].
16
     rPartSet[ptr->rProductionRuleIdx].size()
          && G.isNonTerminal(G.P[ptr->lProductionRuleIdx
     ].rPartSet[ptr->rProductionRuleIdx][ptr->dotPos]))
18
          ProductionRule pr;
          int prIdx = G.getProductionRuleByLeft(pr, G.P[
20
     ptr->1ProductionRuleIdx].rPartSet[ptr->
     rProductionRuleIdx][ptr->dotPos]);
          for(size_t i = 0; i < pr.rPartSet.size(); ++i)</pre>
21
          {
            string str;
23
            if(ptr->dotPos + 1 < G.P[ptr->
24
     lProductionRuleIdx].rPartSet[ptr->
     rProductionRuleIdx].size())
25
               str = G.P[ptr->1ProductionRuleIdx].
     rPartSet[ptr->rProductionRuleIdx][ptr->dotPos + 1]
     + ch2str(ptr->lookaheadSymbol);
            }
27
            else
28
            {
29
               str = ch2str(ptr->lookaheadSymbol);
            set < char > s = G.getFirst(str);
32
            int oldSize = ret.size();
            for(auto& sPtr: s)
34
               if(sPtr != 'e')
36
                 ret.insert(LR1Item(prIdx, i, 0, sPtr));
38
            }
40
            int newSize = ret.size();
            if(oldSize != newSize)
42
               updated = true;
44
46
        }
47
      }
48
    }
```

```
50
51
    // cout << "LR1Closure:\n"; display(ret); cout << endl;
52
53
    return ret;
54
}
```

Listing 3: LR1Closure 成员函数代码清单

## 4.1.3 LR1GoTo 成员函数的设计实现

代码(代码清单4)如下:

```
set < LR1Item > SyntaxAnalyzer::LR1GoTo(const set < LR1Item</pre>
     >& I, char X)
2 {
    if(find(G.VN.begin(), G.VT.begin(), X) != G.VN.end()
      || find(G.VT.begin(), G.VT.end(), X) != G.VT.end()
     )
    {
      set < LR1Item > ret;
      for(auto ptr = I.begin(); ptr != I.end(); ++ptr)
        if(ptr->dotPos < G.P[ptr->lProductionRuleIdx].
     rPartSet[ptr->rProductionRuleIdx].size()
           && G.P[ptr->lProductionRuleIdx].rPartSet[ptr->
11
     rProductionRuleIdx][ptr->dotPos] == X)
           ret.insert(LR1Item(ptr->lProductionRuleIdx,
13
     ptr->rProductionRuleIdx, ptr->dotPos + 1, ptr->
     lookaheadSymbol));
        }
      }
      // cout << "LR1GoTo:\n"; display(LR1Closure(ret));</pre>
17
      cout << endl;</pre>
18
      return LR1Closure(ret);
    }
20
21
    else
      return set < LR1Item > ();
23
24
25 }
```

#### 4.1.4 calLR1ItemSetFamily 成员函数的设计实现

代码(代码清单5)如下:

```
void SyntaxAnalyzer::calLR1ItemSetFamily()
2 {
    G.P.push_back(ProductionRule('S', vector<string>({
     ch2str(G.S)})));
    sort(G.P.begin(), G.P.end(),
      [](const ProductionRule& pr1, const ProductionRule
     & pr2)
        return pr1.1Part < pr2.1Part?</pre>
          true: (pr1.1Part > pr2.1Part?
             false: pr1.rPartSet[0] < pr2.rPartSet[0]);</pre>
      });
    G.VN.push_back('S');
    sort(G.VN.begin(), G.VN.end());
13
14
    G.S = 'S';
    G.nullable.clear();
    G.first.clear();
18
    G.follow.clear();
19
    G.nullable.assign(G.VN.size(), false);
20
    G.first.assign(G.VN.size() + G.VT.size(), set<char</pre>
21
    G.follow.assign(G.VN.size(), set<char>());
22
23
    G.display();
24
25
    G.calNullableFirstFollow();
26
    lr1ItemSetFamily.push_back(
28
      LR1Closure(
        set < LR1Item > ({
30
          LR1Item(
31
             find_if(
               G.P.begin(), G.P.end(),
33
               [](const ProductionRule& pr)
34
```

```
return pr.lPart == 'S';
36
               }) - G.P.begin(),
37
             0, 0, '$'
38
           })
40
        )
41
      );
42
43
    bool updated = true;
44
    while (updated)
45
46
      updated = false;
47
48
      for(auto& itemSetPtr: lr1ItemSetFamily)
49
50
        for(auto& symbolPtr: G.VN)
51
           set < LR1Item > itemSet = LR1GoTo(itemSetPtr,
53
     symbolPtr);
           if(!itemSet.empty()
             && find_if(lr1ItemSetFamily.begin(),
55
     lr1ItemSetFamily.end(),
                [itemSet](const set<LR1Item>& s)
57
                  if(s.size() != itemSet.size())
58
59
                    return false;
60
61
                  else
62
63
                    auto ptr1 = itemSet.begin();
64
                    auto ptr2 = s.begin();
65
                    for(; ptr1 != itemSet.end(); ++ptr1,
66
     ++ptr2)
67
                      if(!(*ptr1 == *ptr2))
68
69
                        return false;
71
                    }
                    return true;
73
                 }
               }) == lr1ItemSetFamily.end())
           {
76
             lr1ItemSetFamily.push_back(itemSet);
77
             updated = true;
```

```
}
79
80
81
         for(auto& symbolPtr: G.VT)
83
            set < LR1Item > itemSet = LR1GoTo(itemSetPtr,
84
      symbolPtr);
            if(!itemSet.empty()
85
              && find_if(lr1ItemSetFamily.begin(),
86
      lr1ItemSetFamily.end(),
                [itemSet](const set<LR1Item>& s)
87
88
                  if(s.size() != itemSet.size())
89
                  {
90
                     return false;
91
                  }
92
                  else
93
94
                     auto ptr1 = itemSet.begin();
95
                     auto ptr2 = s.begin();
96
                     for(; ptr1 != itemSet.end(); ++ptr1,
97
      ++ptr2)
                       if(!(*ptr1 == *ptr2))
99
                       {
100
                         return false;
                       }
                     return true;
104
                }) == lr1ItemSetFamily.end())
106
              lr1ItemSetFamily.push_back(itemSet);
108
              updated = true;
         }
111
       }
     }
113
     // sort lr1ItemSetFamily
115
     \verb|sort(lr1ItemSetFamily.begin(), lr1ItemSetFamily.end|\\
      (),
       [](const set<LR1Item>& itemSet1, const set<LR1Item
117
      >& itemSet2)
       {
118
         auto ptr1 = itemSet1.begin();
119
```

```
auto ptr2 = itemSet2.begin();
120
         for(; ptr1 != itemSet1.end() && ptr2 != itemSet2
121
      .end(); ++ptr1, ++ptr2)
            if(*ptr1 > *ptr2)
            {
124
              return false;
            }
126
            else if(*ptr1 < *ptr2)</pre>
            {
              return true;
            }
130
131
         if(ptr1 != itemSet1.end() && ptr2 == itemSet2.
132
      end())
            return false;
134
         }
136
         else
            return true;
         }
139
       });
     cout << "item set family" << endl;</pre>
142
     for(size_t i = 0; i < lr1ItemSetFamily.size(); ++i)</pre>
143
       cout << "I" << i << endl;
145
       display(lr1ItemSetFamily[i]);
146
       cout << endl;</pre>
147
     }
148
149
     cout << "Goto graph" << endl;</pre>
     for(size_t i = 0; i < lr1ItemSetFamily.size(); ++i)</pre>
       for(size_t j = 0; j < G.VT.size(); ++j)</pre>
153
         set < LR1Item > Ij = LR1GoTo(lr1ItemSetFamily[i], G
155
      .VT[j]);
         if(!Ij.empty())
            cout << "I" << i << endl;
            display(lr1ItemSetFamily[i]);
            cout << "receives " << G.VT[j] << ", goes to"</pre>
160
      << endl;
            display(Ij);
161
```

```
cout << endl;</pre>
          }
       }
164
       for(size_t j = 0; j < G.VN.size(); ++j)</pre>
167
          set < LR1Item > Ij = LR1GoTo(lr1ItemSetFamily[i], G
168
       .VN[j]);
          if(!Ij.empty())
          {
            cout << "I" << i << endl;
            display(lr1ItemSetFamily[i]);
172
            cout << "receives " << G.VN[j] << ", goes to"</pre>
       << endl;
            display(Ij);
            cout << endl;</pre>
177
     }
178
179
```

Listing 5: calLR1ItemSetFamily代码清单

我们使用增广文法,即添加产生式 $S' \to S$ 。这时,文法的终结符集,开始符,nullable 集,first 集和 follow 集均需重新计算。这些计算仅与文法相关,文法及其上操作的设计实现见 7 附录。

#### 4.1.5 buildLR1ParseTab 成员函数的设计实现

代码(代码清单6)如下:

```
void SyntaxAnalyzer::buildLR1ParseTab()
// notice that the lookahead symbol may be e!!!

{
    lr1ActionTab.assign(lr1ItemSetFamily.size(), vector<
        LR1ActionTabItem>(G.VT.size() + 1, LR1ActionTabItem
        (ERR, 0, 0, 0)));
    lr1GoToTab.assign(lr1ItemSetFamily.size(), vector<
        unsigned int>(G.VN.size(), lr1ItemSetFamily.size())
    );
    // if an item in lr1GoToTab equals to the value of
    lr1ItemSetFamily.size(), it means this item is
    invalid

for(size_t i = 0; i < lr1ItemSetFamily.size(); ++i)
    {
}</pre>
```

```
for(auto itemSetPtr = lr1ItemSetFamily[i].begin();
      itemSetPtr != lr1ItemSetFamily[i].end(); ++
     itemSetPtr)
        auto VTPtr = find(G.VT.begin(), G.VT.end(),
     itemSetPtr->lookaheadSymbol);
        auto VNPtr = find(G.VN.begin(), G.VN.end(),
13
     itemSetPtr ->lookaheadSymbol);
        if(itemSetPtr->dotPos < G.P[itemSetPtr->
14
     lProductionRuleIdx].rPartSet[itemSetPtr->
     rProductionRuleIdx].size())
15
          if (G.P[itemSetPtr->lProductionRuleIdx].
     rPartSet[itemSetPtr->rProductionRuleIdx] == "e")
            lr1ActionTab[i][(VTPtr == G.VT.end())? G.VT.
18
     size(): (VTPtr - G.VT.begin())] =
              LR1ActionTabItem(REDUCE, itemSetPtr->
19
     lProductionRuleIdx, itemSetPtr->rProductionRuleIdx,
      0);
          }
          auto ptr = find(G.VT.begin(), G.VT.end(),
            G.P[itemSetPtr->lProductionRuleIdx].rPartSet
     [itemSetPtr->rProductionRuleIdx][itemSetPtr->dotPos
     ]);
          if(ptr != G.VT.end())
            set < LR1Item > Ij = LR1GoTo(lr1ItemSetFamily[i
26
     ],
              G.P[itemSetPtr->lProductionRuleIdx].
27
     rPartSet[itemSetPtr->rProductionRuleIdx][itemSetPtr
     ->dotPosl):
            int lr1ActionTabCol = ptr - G.VT.begin();
            lr1ActionTab[i][lr1ActionTabCol] =
29
     LR1ActionTabItem(
              SHIFT, 0, 0,
30
              find_if(lr1ItemSetFamily.begin(),
31
     lr1ItemSetFamily.end(),
                [Ij](const set < LR1Item > & s)
                   if(s.size() != Ij.size())
35
                     return false;
37
                  else
```

```
39
                      auto ptr1 = Ij.begin();
40
                      auto ptr2 = s.begin();
41
                      for(; ptr1 != Ij.end(); ++ptr1, ++
     ptr2)
                        if(!(*ptr1 == *ptr2))
44
45
                          return false;
46
                      }
48
                      return true;
49
                 }) - lr1ItemSetFamily.begin()
51
               ):
52
          }
        }
        else
55
        {
56
           if (G.P[itemSetPtr->lProductionRuleIdx].lPart
57
     != 'S')
58
             int lr1ActionTabCol = (VTPtr == G.VT.end())?
      G.VT.size(): (VTPtr - G.VT.begin());
             lr1ActionTab[i][lr1ActionTabCol] =
     LR1ActionTabItem(REDUCE, itemSetPtr->
     lProductionRuleIdx, itemSetPtr->rProductionRuleIdx,
      0);
          }
61
           else
62
63
             lr1ActionTab[i][G.VT.size()] =
64
     LR1ActionTabItem(ACCEPT, 0, 0, 0);
        }
      }
67
    }
68
    for(size_t i = 0; i < lr1ItemSetFamily.size(); ++i)</pre>
70
71
      for(size_t A = 0; A < G.VN.size(); ++A)</pre>
        set < LR1Item > Ij = LR1GoTo(lr1ItemSetFamily[i], G
74
     . VN[A]);
        lr1GoToTab[i][A] = find_if(lr1ItemSetFamily.
75
     begin(), lr1ItemSetFamily.end(),
```

```
[Ij](const set<LR1Item>& s)
76
              if(s.size() != Ij.size())
                return false;
80
              }
81
              else
83
                auto ptr1 = Ij.begin();
84
                auto ptr2 = s.begin();
                for(; ptr1 != Ij.end(); ++ptr1, ++ptr2)
86
87
                   if(!(*ptr1 == *ptr2))
88
89
                     return false;
91
                }
92
                return true;
93
              }
94
            }) - lr1ItemSetFamily.begin();
95
       }
96
     }
97
98
     cout << "lr1ActionTab" << endl;</pre>
99
     for(size_t i = 0; i < lr1ActionTab.size(); ++i)</pre>
100
       for(size_t j = 0; j < lr1ActionTab[i].size(); ++j)</pre>
102
         cout << "lr1ActionTab[" << i << "," << ((j == G.
104
      VT.size())? '$': G.VT[j]) << "] = ";</pre>
         switch(lr1ActionTab[i][j].action)
106
            case SHIFT: cout << "SHIFT " << lr1ActionTab[i</pre>
      ][j].itemSetIdx << endl; break;</pre>
            case REDUCE: cout << "REDUCE "; G.P[</pre>
108
      lr1ActionTab[i][j].lProductionRuleIdx].display();
      cout << endl; break;</pre>
            case ACCEPT: cout << "ACCEPT" << endl; break;</pre>
            default: cout << "ERR" << endl;</pre>
         }
111
       }
     }
     cout << endl;</pre>
114
     cout << "lr1GoToTab" << endl;</pre>
     for(size_t i = 0; i < lr1GoToTab.size(); ++i)</pre>
117
```

Listing 6: buildLR1ParseTab 成员函数代码清单

### 4.1.6 lr1SyntaxAnalyze 成员函数的设计实现

代码(代码清单7)如下:

```
bool SyntaxAnalyzer::lr1SyntaxAnalyze(const string&
     myLR1Input)
2 {
    if (myLR1Input.size() == 0)
      cout << "Input is empty!" << endl;</pre>
      return false;
    cout << "Parsing " << myLR1Input << "..." << endl;</pre>
    lr1Stack.push_back(
11
      find_if(
12
        lr1ItemSetFamily.begin(), lr1ItemSetFamily.end()
13
        [this](const set < LR1Item > & s)
14
15
          return find_if(s.begin(), s.end(),
16
             [this](const LR1Item& i)
17
18
               return G.P[i.1ProductionRuleIdx].1Part ==
     'S';
             }) != s.end();
          - lr1ItemSetFamily.begin()
22
      );
23
24
    string str = myLR1Input;
25
    reverse(str.begin(), str.end());
```

```
lr1Input = "$" + str;
27
28
    int idx = lr1Input.size() - 1;
29
    char a = lr1Input[idx];
    for(;;)
31
32
      unsigned int s = lr1Stack[lr1Stack.size() - 1];
33
      auto ptr = find(G.VT.begin(), G.VT.end(), a);
34
      int aIdx = (ptr == G.VT.end())? G.VT.size(): ptr -
      G.VT.begin();
      LR1ActionTabItem curAction = lr1ActionTab[s][aIdx
36
     ];
      cout << "Begin a new pass..." << endl;</pre>
37
      cout << "Current state(top of the stack): "</pre>
38
            << lr1Stack[lr1Stack.size() - 1] << endl;
      display(lr1ItemSetFamily[lr1Stack[lr1Stack.size()
40
     - 1]]);
      cout << "Current input symbol: " << a << endl <<</pre>
41
     endl;
      if(curAction.action == SHIFT)
42
43
        lr1Stack.push_back(curAction.itemSetIdx);
        cout << "Push back state " << curAction.</pre>
     itemSetIdx << endl;</pre>
        display(lr1ItemSetFamily[curAction.itemSetIdx]);
        --idx;
47
        a = lr1Input[idx];
        \verb"cout" << "Shift" and now the input symbol is " <<
49
     a << endl << endl;
      else if(curAction.action == REDUCE)
51
        cout << "Reduce using production rule "</pre>
53
              << G.P[curAction.lProductionRuleIdx].lPart
54
     << "->"
              << G.P[curAction.lProductionRuleIdx].
     rPartSet[curAction.rProductionRuleIdx] << endl;
        int num = (G.P[curAction.lProductionRuleIdx].
     rPartSet[curAction.rProductionRuleIdx] == "e")?
          0: G.P[curAction.lProductionRuleIdx].rPartSet[
     curAction.rProductionRuleIdx].size();
        for(int i = 0; i < num; ++i)</pre>
        {
59
          lr1Stack.erase(lr1Stack.end() - 1);
60
61
        cout << "Pop " << num << " state(s) out of the</pre>
```

```
stack" << endl;
        cout << "Now the stack top is state " <<</pre>
     lr1Stack[lr1Stack.size() - 1] << endl;</pre>
        display(lr1ItemSetFamily[lr1Stack[lr1Stack.size
     () - 1]]);
        int t = lr1Stack[lr1Stack.size() - 1];
        lr1Stack.push_back(lr1GoToTab[t][find(G.VN.begin
     (), G.VN.end(), G.P[curAction.lProductionRuleIdx].
     1Part) - G.VN.begin()]);
        cout << "lr1GoToTab[" << t << "," << G.P[</pre>
     curAction.lProductionRuleIdx].lPart << "] = "</pre>
              << ((lr1GoToTab[t][find(G.VN.begin(), G.VN.
68
     end(), G.P[curAction.lProductionRuleIdx].lPart) - G
     .VN.begin()] == lr1ItemSetFamily.size())?
                   " ": to_string(lr1GoToTab[t][find(G.VN
     .begin(), G.VN.end(), G.P[curAction.
     lProductionRuleIdx].lPart) - G.VN.begin()]));
        cout << ", so we push back state "
70
              << lr1GoToTab[t][find(G.VN.begin(), G.VN.
     end(), G.P[curAction.lProductionRuleIdx].lPart) - G
     .VN.begin()] << endl;
        display(lr1ItemSetFamily[lr1GoToTab[t][find(G.VN
     .begin(), G.VN.end(), G.P[curAction.
     lProductionRuleIdx].lPart) - G.VN.begin()]]);
        cout << endl;</pre>
73
      }
74
      else if(curAction.action == ACCEPT)
76
        cout << "Parsing succeed!" << endl << endl;</pre>
77
        return true;
      }
79
      else
80
81
        cout << "Parsing error!" << endl;</pre>
        cout << "In state " << lr1Stack[lr1Stack.size()</pre>
83
     - 1] << endl;
        display(lr1ItemSetFamily[lr1Stack[lr1Stack.size
84
     () - 1]]); cout << endl;</pre>
        return false;
85
      }
86
    }
87
 }
```

Listing 7: lr1SyntaxAnalyze 成员函数代码清单

我们使用\$作为分析栈初始栈顶符号。

# 5 LL(1)语法分析器的设计实现

#### 5.0.7 buildLL1ParseTab 成员函数的设计实现

代码(代码清单8)如下:

```
void SyntaxAnalyzer::buildLL1ParseTab()
2 {
    ll1ParseTab.assign(G.VN.size(), vector<LL1Item>(G.VT
     .size() + 1, LL1Item(G.P.size(), 0)));
   for(size_t i = 0; i < G.P.size(); ++i)</pre>
      unsigned int AIdx = find(G.VN.begin(), G.VN.end(),
      G.P[i].lPart) - G.VN.begin();
      for(size_t j = 0; j < G.P[i].rPartSet.size(); ++j)</pre>
        set < char > firstAlpha = G.getFirst(G.P[i].
11
     rPartSet[j]);
        for(auto& p: firstAlpha)
14
          unsigned int aIdx = find(G.VT.begin(), G.VT.
     end(), p) - G.VT.begin();
          if(aIdx < G.VT.size())</pre>
            ll1ParseTab[AIdx][aIdx] = LL1Item(i, j);
19
        }
20
21
        if(find(firstAlpha.begin(), firstAlpha.end(), 'e
22
     ') != firstAlpha.end())
23
          set < char > followA = G.getFollow(G.P[i].1Part);
24
25
          for(auto& p: followA)
26
27
            unsigned int bIdx = find(G.VT.begin(), G.VT.
     end(), p) - G.VT.begin();
            ll1ParseTab[AIdx][bIdx] = LL1Item(i, j);
          }
30
          if(find(followA.begin(), followA.end(), '$')
     != followA.end())
          {
```

```
ll1ParseTab[AIdx][G.VT.size()] = LL1Item(i,
34
      j);
         }
       }
37
    }
38
39
    for(size_t i = 0; i < ll1ParseTab.size(); ++i)</pre>
40
41
       for(size_t j = 0; j < ll1ParseTab[i].size(); ++j)</pre>
43
         cout << "ll1ParseTab[" << G.VN[i] << ", " << ((j</pre>
44
       == G.VT.size())? '$': G.VT[j]) << "] = ";
         if(ll1ParseTab[i][j].i < G.P.size())</pre>
45
46
           cout << G.P[ll1ParseTab[i][j].i].1Part << "->"
47
       << G.P[ll1ParseTab[i][j].i].rPartSet[ll1ParseTab[i</pre>
      ][j].j];
         }
         cout << endl;</pre>
49
       }
50
    }
51
  }
```

Listing 8: buildLL1ParseTab 成员函数代码清单

#### 5.0.8 lllSyntaxAnalyze 成员函数的设计实现

代码(代码清单9)如下:

```
bool SyntaxAnalyzer::ll1SyntaxAnalyze(const string&
     myLL1Input)
2 {
    if(myLL1Input.size() == 0)
      cout << "Input is empty!" << endl;</pre>
      return false;
    }
    cout << "Parsing " << myLL1Input << "..." << endl;</pre>
    string str = myLL1Input;
    reverse(str.begin(), str.end());
12
    111Input = '$' + str;
13
    ll1Stack.push_back('$');
14
    111Stack.push_back(G.S);
```

```
int ip = ll1Input.size() - 1;
16
    char x = ll1Stack[ll1Stack.size() - 1];
17
18
    while(x != '$')
20
      if(x == ll1Input[ip])
        111Stack.erase(111Stack.end() - 1);
23
24
      else if(find(G.VT.begin(), G.VT.end(), x) != G.VT.
26
27
        cout << "Parsing error!" << endl;</pre>
28
        return false;
29
30
      else
32
        unsigned int row = find(G.VN.begin(), G.VN.end()
     , x) - G.VN.begin();
        unsigned int col = (ll1Input[ip] == '$')? G.VT.
     size(): (find(G.VT.begin(), G.VT.end(), ll1Input[ip
     ]) - G.VT.begin());
35
        if(ll1ParseTab[row][col].i >= G.P.size())
36
          cout << "Parsing error!" << endl;</pre>
          return false;
39
        }
40
        else
41
42
          cout << G.P[ll1ParseTab[row][col].i].1Part <<</pre>
43
     "->" << G.P[ll1ParseTab[row][col].i].rPartSet[
     ll1ParseTab[row][col].j] << endl;</pre>
          111Stack.erase(111Stack.end() - 1);
44
          if(!(G.P[ll1ParseTab[row][col].i].rPartSet[
45
     ll1ParseTab[row][col].j].size() == 1 &&
            G.P[ll1ParseTab[row][col].i].rPartSet[
     ll1ParseTab[row][col].j][0] == 'e'))
          {
            for(int cnt = G.P[ll1ParseTab[row][col].i].
48
     rPartSet[ll1ParseTab[row][col].j].size() - 1; cnt
     >= 0; --cnt)
               111Stack.push_back(G.P[111ParseTab[row][
     col].i].rPartSet[ll1ParseTab[row][col].j][cnt]);
```

Listing 9: 111SyntaxAnalyze 成员函数代码清单

我们使用\$作为分析栈初始栈顶符号。

# 6 测试数据及运行结果

该语法分析器对下面的3个测试语法正确运行。 文法 1:

$$\begin{split} E &\to TA \\ A &\to +TA|-TA|\varepsilon \\ T &\to FB \\ B &\to *FB|/FB|\varepsilon \\ F &\to (E)|0|1|2|3|4|5|6|7|8|9 \end{split} \tag{1}$$

文法 2:

$$E \to E + T|E - T|T$$

$$T \to T * F|T/F|F$$

$$F \to (E)|0|1|2|3|4|5|6|7|8|9$$
(2)

文法 3:

$$Z \to AA$$

$$A \to aA|b \tag{3}$$

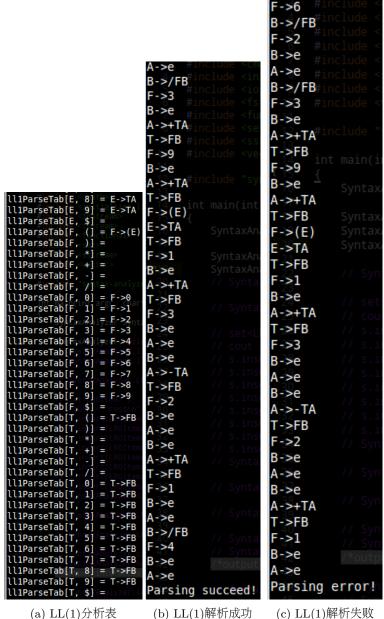
对于文法 3 和含左递归的文法 2, LR(1)语法分析器可正确生成分析表,对于不含左递归的文法 1, LL(1)语法分析器可正确生成分析表。利用自动生成的分析表,LR(1)和LL(1)语法分析器均可正确解析下面的书写正确的复杂的一位数含括号四则运算表达式:

均不能正确解析书写错误的表达式(最后一个4之前缺少一个运算符):

运行结果截图如下:

```
nullable
A: nullable periment-report-on-
B: nullable
E: not nullablerammarl
F: not nullablerammar2
T: not nullablerammar3
first
            fig: screenshot3fig
            fig: screenshot4fig
         BLOCKS
   {0}
   \{1\}
2: {2}
          ▶ LR(1)语法分析器的i
3: {3}
       ▼ [S] LL(1)语法分析器的i
4: {4}
5: {5}
6: {6}
7: {7}
         S 附录: 文法的设计与
8: {8}
            产生式数据结构的设
9: {9}
A: {+,-,e} 文法类的设计
B: {*,/,e}
E: {(,0,1,2,3,4,5,6,7,8,9}
F: {(,0,1,2,3,4,5,6,7,8,9}
T: {(,0,1,2,3,4,5,6,7,8,9}
follow
A: {$,)}
B: {$,),+,-}
E: {$,)}
F: {$,),*,+,-,/}
   {$,),+,-}
```

Figure 1: 文法



(a) LL(1)分析表

(b) LL(1)解析成功

Figure 2: LL(1)语法分析器运行结果截图

```
lr1GoToTab[0,A]
lr1GoToTab[0,B]
lr1GoToTab[0,E]
lr1GoToTab[0,F]
lr1GoToTab[0,S]
lr1GoToTab[0,T]
lr1GoToTab[1,A]
lr1GoToTab[1,B]
lr1GoToTab[1,E]
lr1GOToTab[1,E]
                                                                                                                                                                                                                              [T∹>.FB,)]
                                                                                                                                                                                                                               T->.FB,+
                                                                                                                                                                                                                            T->.FB,-
                                                                                                                                                                                                                           receives T,
                                                                                                                                                                                                                                                                                    goes to
                                                                                                                                                                                           ⊫111
                                                                                                                                                                                                                         [A->.+TA,)]
[A->.-TA,)]
[A->.e,)]
[E->T.A,)]
                                                                                                                    lr1GoToTab[1,F]
lr1GoToTab[1,S]
                                                                                                                  lr1GoToTab[1,5]
lr1GoToTab[2,A]
lr1GoToTab[2,B]
lr1GoToTab[2,E]
lr1GoToTab[2,F]
                                                                                                                                                                                          = 29
                                                                                                                                                                                                                          I33
                                                                                                                                                                                                                              F->(E.),
                                                                                                                  | CTGOTOTAD [2,F]
| CTGOTOTAD [2,F]
| CTGOTOTAD [2,T]
| CTGOTOTAD [3,A]
| CTGOTOTAD [3,B]
| CTGOTOTAD [3,E]
| CTGOTOTAD [3,F]
| CTGOTOTAD [3,F]
| CTGOTOTAD [3,F]
                                                                                                                                                                                                                            [F->(E.),
                                                                                                                                                                                                                             [F->(E.),/
                                                                                                                                                                                                                           receives ),
                                                                                                                                                                                                                                                                                    goes to
                                                                                                                                                                                                                            [F->(E).,$]
REDUCE F->(E) |011|2|3|4|5|6|7|8|9
                                                                                                                                                                                                                            [F->(E).,*
                                                                                                                                                                                           = 16
                                                                                                                                                                                                                          [F->(E).,+]
[F->(E).,-]
                                                                                                                   lr1GoToTab[3,T]
lr1GoToTab[4,A]
lr1GoToTab[4,B]
                                                                                                                                                                                                                          [F->(E).,/]
                                                                                                                   lr1GoToTab[4,E]
lr1GoToTab[4,F]
                                                                                                                                                                                                                           I35
                                                                                                                                                                                                                          [F->(E.),)]
                                                                                                                  lr1GoToTab[4,F]
lr1GoToTab[4,T]
lr1GoToTab[5,A]
lr1GoToTab[5,B]
lr1GoToTab[5,E]
                                                                                                                                                                                                                             [F->(E.),
                                                                                                                                                                                                                              F->(E.),+
                                                                                                                                                                                           = 9
                                                                                                                                                                                                                             [F->(E.),
                                                                                                                                                                                                                            [F->(E.),/]
                                                                                                                                                                                                                           receives ),
[F->(E).,)]
                                                                                                                                                                                                                                                                                    goes to
                                                                                                                  lr1GoToTab[5,E]
lr1GoToTab[5,S]
lr1GoToTab[5,T]
lr1GoToTab[6,A]
lr1GoToTab[6,B]
                                                                                                                                                                                                                              E->(E).,*]
                                                                                                                                                                                                                              F->(E).,+
                                                                                                                                                                                            = 13
                                                                                                                                                                                                                              F->(E).,
```

(a) LR(1)Action分析表

(b) LR(1)Goto分析表

(c) LR(1)项集族

Figure 3: LR(1)语法分析器运行结果截图1

(a) LR(1)解析成功

(b) LR(1)解析失败

Figure 4: LR(1)语法分析器运行结果截图2

# 7 附录: 文法的设计与实现

## 7.1 产生式数据结构的设计

代码(代码清单10)如下:

```
class ProductionRule
3 public:
    ProductionRule();
    ProductionRule(const char& 1, const vector<string>&
     r);
    void display();
    char lPart;
    vector < string > rPartSet;
11
12 private:
13
14 };
16 ProductionRule::ProductionRule()
17
18
19 }
21 ProductionRule::ProductionRule(const char& 1, const
     vector < string > & r): lPart(l), rPartSet(r)
22 {
23
24 }
void ProductionRule::display()
27 {
    cout << 1Part << "->";
    for(size_t i = 0; i < rPartSet.size(); ++i)</pre>
29
30
      cout << rPartSet[i];</pre>
31
      if(i != rPartSet.size() - 1)
32
33
        cout << "|";
34
35
    }
36
37 }
```

Listing 10: 产生式数据结构的设计代码清单

### 7.2 文法类的设计

代码(代码清单11)如下:

```
class Grammar
3 public:
    Grammar(const string& filePath);
    void calNullableFirstFollow();
    void display();
    bool isNonTerminal(char ch);
    int getProductionRuleByLeft(ProductionRule& pr, char
      1Part); // return index of production rule which
     has lPart as left part
    set < char > getFirst(char ch);
    set < char > getFirst(string str);
12
    set < char > getFollow(char ch);
13
14
    vector < char > VN;
    vector < char > VT;
16
    vector < ProductionRule > P;
17
    char S:
18
19
    vector < bool > nullable; // for symbols in VN and VT,
20
     elements in VT are followed by elements in VN
    vector<set<char>> first; // for symbols in VN and VT
     , elements in VT are followed by elements in VN
    vector<set<char>> follow; // only for symbols in VN
22
23
24 private:
25
 };
```

Listing 11: 文法类的设计代码清单

我们仅介绍构造函数 calNullableFirstFollow 成员函数的设计实现,其他函数仅用于输出信息和帮助对类的封装,不再赘述。

#### 7.2.1 构造函数的设计实现

代码(代码清单12)如下:

```
Grammar::Grammar(const string& filePath)
2
```

```
ifstream f(filePath);
    assert(f);
    // file format
    // VNSize VN
    // VTSize VT
    // PSize
    // P
    // S
12
    int VNSize;
13
    f >> VNSize;
14
    VN.resize(VNSize);
    for(int i = 0; i < VNSize; ++i)</pre>
16
      f >> VN[i];
18
19
20
    int VTSize;
21
    f >> VTSize;
22
    VT.resize(VTSize);
    for(int i = 0; i < VTSize; ++i)</pre>
24
25
      f >> VT[i];
26
27
28
    struct functor: public binary_function <</pre>
     ProductionRule, char, bool>{
    public:
30
      bool operator ()(const ProductionRule pr, const
31
     char ch) const {
        return (pr.1Part == ch);
      }
33
    };
34
35
    int PSize;
36
    f >> PSize;
37
    string str;
    for(int i = 0; i < PSize; ++i)</pre>
39
      f >> str;
41
      vector<string> v_pr = split(str, '>');
      v_pr[0].erase(v_pr[0].end() - 1);
43
      vector<string> v_pr_r = split(v_pr[1], '|');
45
      functor myFunctor;
```

```
auto itr = find_if(P.begin(), P.end(), bind2nd(
47
     myFunctor, *(v_pr[0].begin()));
      if(itr == P.end())
48
        ProductionRule pr;
        pr.lPart = *(v_pr[0].begin());
        for(size_t j = 0; j < v_pr_r.size(); ++j)</pre>
53
           pr.rPartSet.push_back(v_pr_r[j]);
54
        P.push_back(pr);
56
      }
57
      else
58
59
      {
        for(size_t j = 0; j < v_pr_r.size(); ++j)</pre>
60
61
           P[itr - P.begin()].rPartSet.push_back(v_pr_r[j
     ]);
      }
65
    }
66
67
    f >> S;
68
69
    f.close();
71
    sort(VN.begin(), VN.end());
72
    sort(VT.begin(), VT.end());
73
    for(size_t i = 0; i < P.size(); ++i)</pre>
74
75
      sort(P[i].rPartSet.begin(), P[i].rPartSet.end());
76
77
    sort(P.begin(), P.end(),
78
      [](ProductionRule pr1, ProductionRule pr2)
79
80
        return pr1.1Part < pr2.1Part?</pre>
81
           true: (pr1.1Part > pr2.1Part?
             false: pr1.rPartSet[0] < pr2.rPartSet[0]);</pre>
83
      });
84
85
    nullable.assign(VN.size(), false);
    first.assign(VN.size() + VT.size(), set<char>());
87
    follow.assign(VN.size(), set<char>());
88
89
    display();
```

```
calNullableFirstFollow();
3 }
```

Listing 12: 构造函数的设计实现代码清单

该文法类从格式化文本中读入信息并解析后完成初始化。

## 7.2.2 calNullableFirstFollow 成员函数的设计实现

代码(代码清单13)如下:

```
void Grammar::calNullableFirstFollow()
2 {
    // calculate nullable
    for(size_t i = 0; i < P.size(); ++i)</pre>
      for(size_t j = 0; j < P[i].rPartSet.size(); ++j)</pre>
        if(P[i].rPartSet[j].size() == 1 && *(P[i].
     rPartSet[j].begin()) == 'e')
           nullable[i] = true;
           break;
12
13
      }
14
    }
16
    bool updated = true;
17
    while(updated)
18
19
      updated = false;
20
21
      for(size_t i = 0; i < P.size(); ++i)</pre>
23
        for(size_t j = 0; j < P[i].rPartSet.size(); ++j)</pre>
           bool allNullable = true;
26
           for(size_t k = 0; k < P[i].rPartSet[j].size();</pre>
27
      ++k)
             auto ptr = find(VN.begin(), VN.end(), P[i].
29
     rPartSet[j][k]);
             if(ptr == VN.end() || !nullable[ptr - VN.
30
     begin()])
```

```
31
                allNullable = false;
32
                break;
             }
           }
35
           if(allNullable && !nullable[i])
37
38
             updated = true;
39
             nullable[i] = true;
             break;
41
           }
42
         }
43
      }
44
    }
45
46
    // calculate first
47
48
    for(size_t i = 0; i < VT.size(); ++i)</pre>
49
      first[i].insert(VT[i]);
51
    updated = true;
    while(updated)
55
56
      updated = false;
58
      for(size_t i = 0; i < P.size(); ++i)</pre>
59
60
         for(size_t j = 0; j < P[i].rPartSet.size(); ++j)</pre>
61
62
           int idx = find(VN.begin(), VN.end(), P[i].
63
     1Part) - VN.begin() + VT.size();
           int 1 = 0;
65
           for(size_t k = 0; k < P[i].rPartSet[j].size();</pre>
      ++k)
67
             auto ptr = find(VN.begin(), VN.end(), P[i].
     rPartSet[j][k]);
             if(ptr == VN.end() || !nullable[ptr - VN.
     begin()])
70
               break;
71
             }
```

```
else
73
            {
              ++1;
              int oldSize = first[idx].size();
77
              int newSize = 0;
              auto ptr = find(VN.begin(), VN.end(), P[i
     ].rPartSet[j][k]);
              setUnion(first[idx], first[ptr - VN.begin
80
     () + VT.size()]);
              newSize = first[idx].size();
81
              if(oldSize != newSize)
82
83
                updated = true;
84
              }
86
          }
87
88
          // cout << P[i].1Part << "->" << P[i].rPartSet
90
     [j] << " " << 1 << endl;
91
          int oldSize = first[idx].size();
92
          int newSize = 0;
93
          if(1 >= P[i].rPartSet[j].size())
94
95
            first[idx].insert('e');
            newSize = first[idx].size();
97
          }
98
          else
99
100
            auto ptr = find(VN.begin(), VN.end(), P[i].
101
     rPartSet[j][l]);
            if(ptr == VN.end())
              if(P[i].rPartSet[j][1] != 'e')
104
                setUnion(first[idx], first[find(VT.begin
     (), VT.end(), P[i].rPartSet[j][1]) - VT.begin()]);
                108
                // cout << i << " " << idx << " ";
                // ::display(first[find(VT.begin(), VT.
     end(), P[i].rPartSet[j][1]) - VT.begin()]); cout <<</pre>
      endl << endl;</pre>
111
```

```
else
112
                 {
113
                   setUnion(first[idx], set < char > ({ 'e'}));
114
              }
              else
117
118
                 setUnion(first[idx], first[ptr - VN.begin
119
      () + VT.size()]);
              }
121
              newSize = first[idx].size();
122
123
124
            if(oldSize != newSize)
125
126
              updated = true;
127
128
         }
129
130
     }
131
     // sth about $ and e
     follow[find(VN.begin(), VN.end(), S) - VN.begin()].
135
      insert('$');
     updated = true;
     while(updated)
138
       updated = false;
140
141
       for(size_t i = 0; i < P.size(); ++i)</pre>
142
143
          for(size_t j = 0; j < P[i].rPartSet.size(); ++j)</pre>
144
145
            int idx = find(VN.begin(), VN.end(), P[i].
146
      1Part) - VN.begin();
147
            for(size_t k = 0; k < P[i].rPartSet[j].size();</pre>
148
       ++k)
              bool allNullable1 = true;
              for(size_t cnt = k + 1; cnt < P[i].rPartSet[</pre>
151
      j].size(); ++cnt)
              {
152
```

```
if(P[i].rPartSet[j][cnt] != 'e')
153
154
                  auto cntPtr = find(VN.begin(), VN.end(),
       P[i].rPartSet[j][cnt]);
                  if(cntPtr == VN.end() || !nullable[
      cntPtr - VN.begin()])
                    allNullable1 = false;
                    break;
                  }
               }
               else
162
               {
                 break;
164
               }
             if(allNullable1)
167
168
               auto kPtr = find(VN.begin(), VN.end(), P[i
      ].rPartSet[j][k]);
               if(kPtr != VN.end())
                  int oldSize = follow[kPtr - VN.begin()].
      size();
                  setUnion(follow[kPtr - VN.begin()],
173
      follow[idx]);
                  int newSize = follow[kPtr - VN.begin()].
      size();
                  if(oldSize != newSize)
175
                    updated = true;
178
179
181
             for(size_t l = k + 1; l < P[i].rPartSet[j].</pre>
      size(); ++1)
               bool allNullable2 = true;
184
               for(int cnt = k + 1; cnt < l - 1; ++cnt)</pre>
185
186
                  if(P[i].rPartSet[j][cnt] != 'e')
188
                    auto cntPtr = find(VN.begin(), VN.end
189
      (), P[i].rPartSet[j][cnt]);
                    if(cntPtr == VN.end() || !nullable[
190
```

```
cntPtr - VN.begin()])
191
                       allNullable2 = false;
                       break;
                     }
194
                  }
195
                  else
196
                     break;
198
200
                if(allNullable2)
201
202
                  auto kPtr = find(VN.begin(), VN.end(), P
203
      [i].rPartSet[j][k]);
                  if(kPtr != VN.end())
204
                  {
205
                     auto lPtr = find(VN.begin(), VN.end(),
206
       P[i].rPartSet[j][1]);
                     int lIdx = (lPtr == VN.end())?
207
                       (find(VT.begin(), VT.end(), P[i].
      rPartSet[j][l]) - VT.begin()):
                       (lPtr - VN.begin() + VT.size());
                     int oldSize = follow[kPtr - VN.begin()
210
      ].size();
                     set < char > __first__ = first[lIdx];
211
                     __first__.erase('e');
                     setUnion(follow[kPtr - VN.begin()],
      __first__);
                     int newSize = follow[kPtr - VN.begin()
214
      ].size();
                     if(oldSize != newSize)
215
                     {
216
                       updated = true;
217
218
                  }
219
                }
220
             }
           }
222
         }
224
     }
225
226
     // display nullable
     cout << "nullable" << endl;</pre>
228
     for(size_t i = 0; i < nullable.size(); ++i)</pre>
```

```
230
        cout << VN[i] << ": " << ((nullable[i])? "nullable</pre>
231
      ": "not nullable") << endl;
     }
     cout << endl;</pre>
234
     // display first
     cout << "first" << endl;</pre>
     for(size_t i = 0; i < first.size(); ++i)</pre>
237
        if(i < VT.size())</pre>
239
        {
240
          cout << VT[i] << ": ";
241
          ::display(first[i]);
242
          cout << endl;</pre>
243
        }
244
        else
245
246
          cout << VN[i - VT.size()] << ": ";</pre>
          ::display(first[i]);
248
          cout << endl;</pre>
        }
     }
     cout << endl;</pre>
252
253
     // display follow
254
     cout << "follow" << endl;</pre>
     for(size_t i = 0; i < follow.size(); ++i)</pre>
256
257
        cout << VN[i] << ": ";
258
        ::display(follow[i]);
        cout << endl;</pre>
260
     }
261
     cout << endl;</pre>
262
263 }
```

Listing 13: calNullableFirstFollow 成员函数代码清单

递归的计算方法要求指定求值顺序这一问题,为了避开这一棘手的问题, 我们使用迭代至不懂点的方法求 nullable 集, first 集和 follow集。需要说 明的是,我们使用字符 e 表示空串,并对它进行不同于其他终结符的特殊处 理。