



The Introduction To Artificial Intelligence

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2025-2026-1

The Introduction to Artificial Intelligence

- Teaching hours: 32 hours
- Assessment:

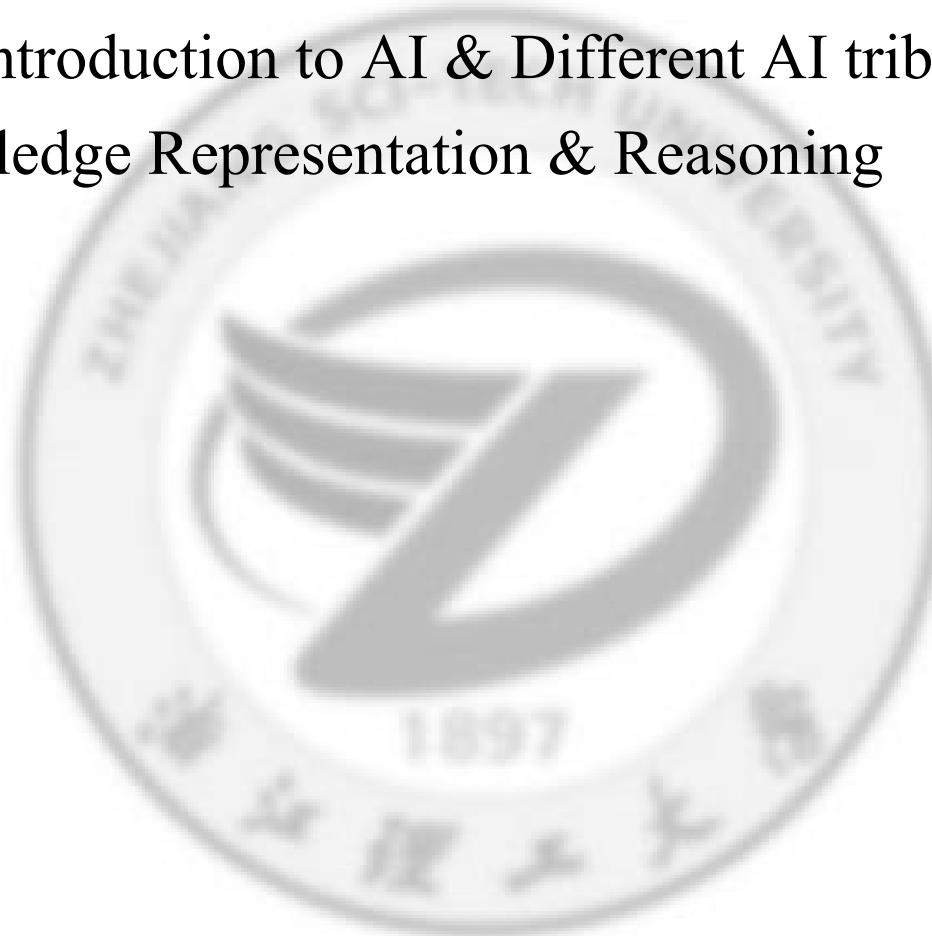
Assessment Methods	Assessment Requirements	Assessment Weighting
Homework	4-6 times (50%)	50% of grade
Literature review	Presentation, Q&A (50%)	
in-class tests	-	10% of grade
Final exam	Open-book examination	40% of grade

- No Copy! No Plagiarize!

The Introduction to Artificial Intelligence

- Part I Brief Introduction to AI & Different AI tribes

-  Part II Knowledge Representation & Reasoning



OUTLINE

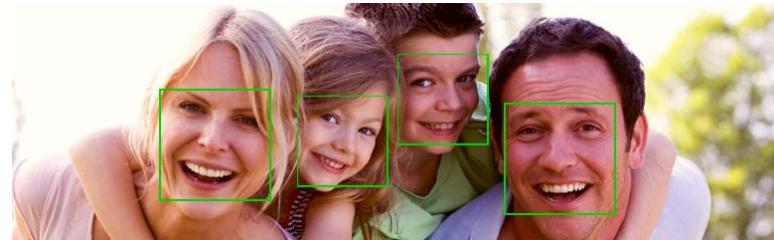
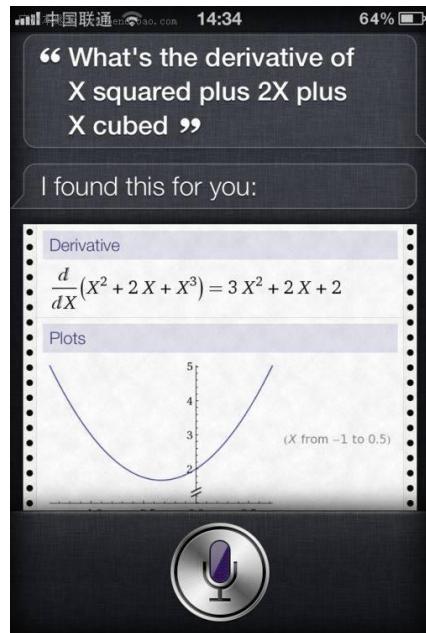
- *1.1 Brief Review*
- 1.2 Knowledge Representation & Reasoning



1.1 Brief Review

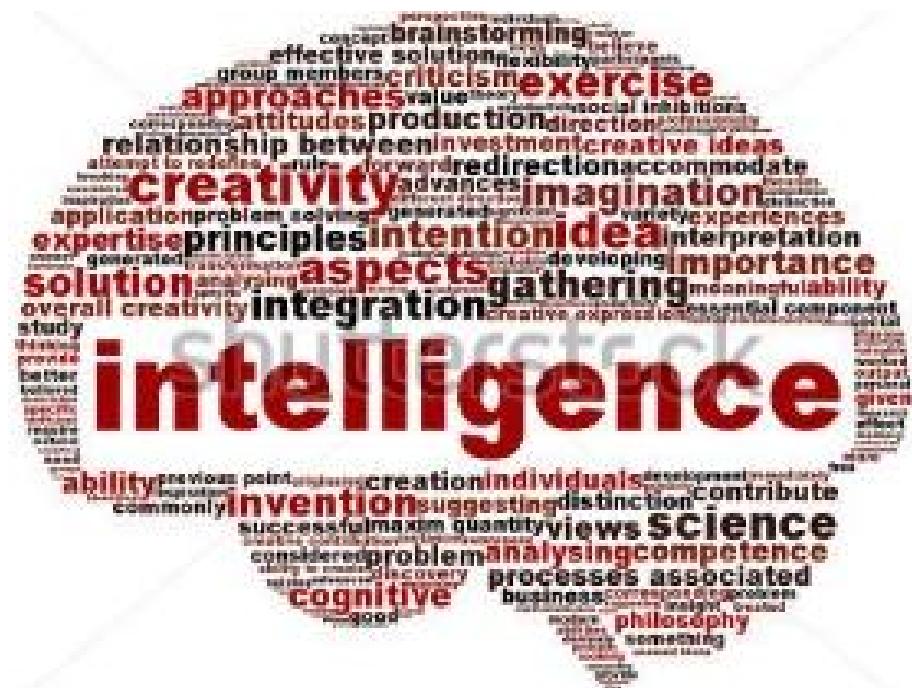
□ What Is Artificial Intelligence?

Actually,
artificial intelligence is **intelligence** exhibited by machines.



1.1 Brief Review

□ What Is Artificial Intelligence?

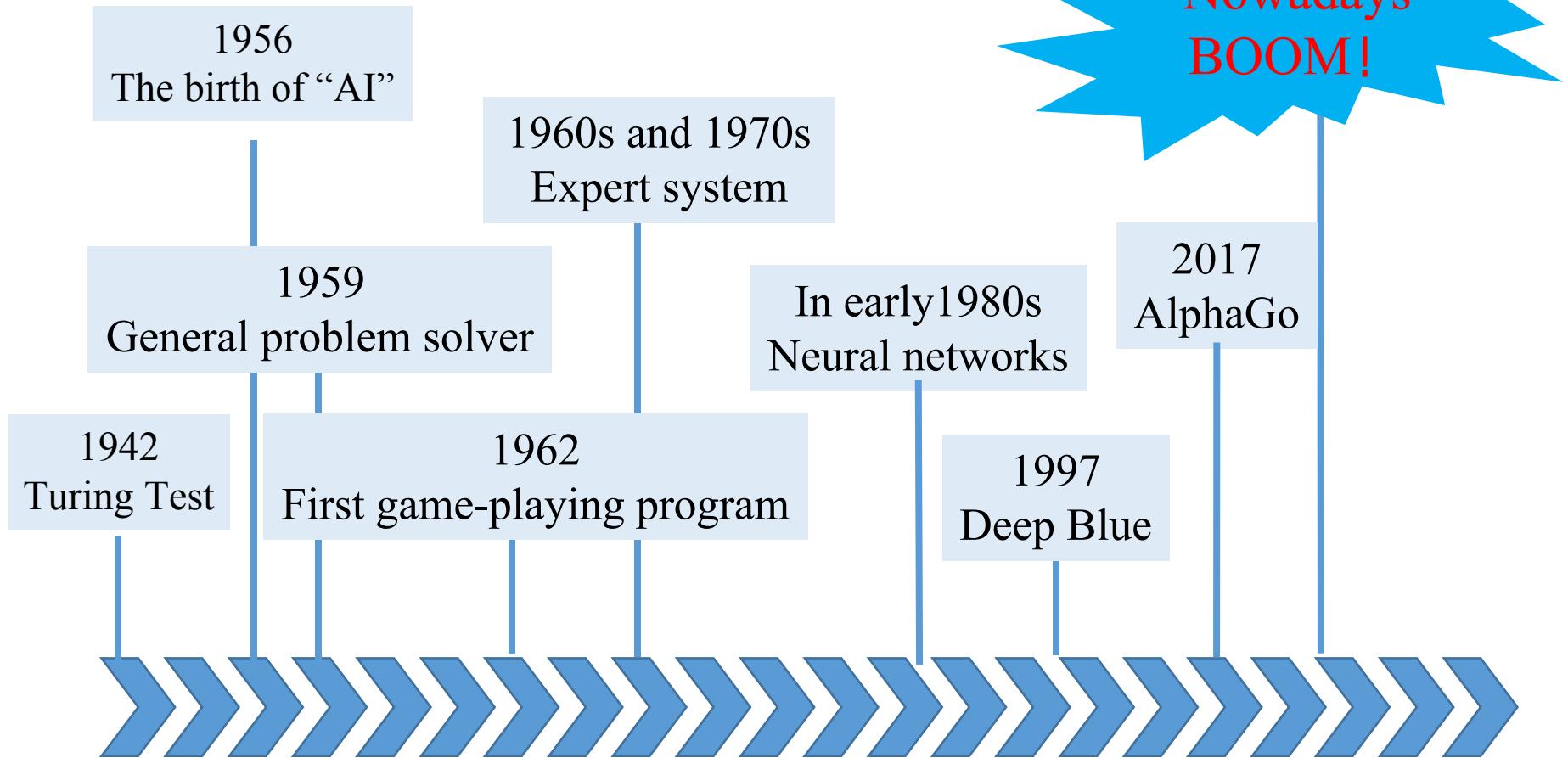


We agree with that:

Intelligence is the ability to learn or understand or to deal with new or trying situations; the ability to apply knowledge to manipulate one's environment or to think abstractly.

1.1 Brief Review

□ The History of AI



1.2 Different Tribes of AI

I.M. Intuition Machine

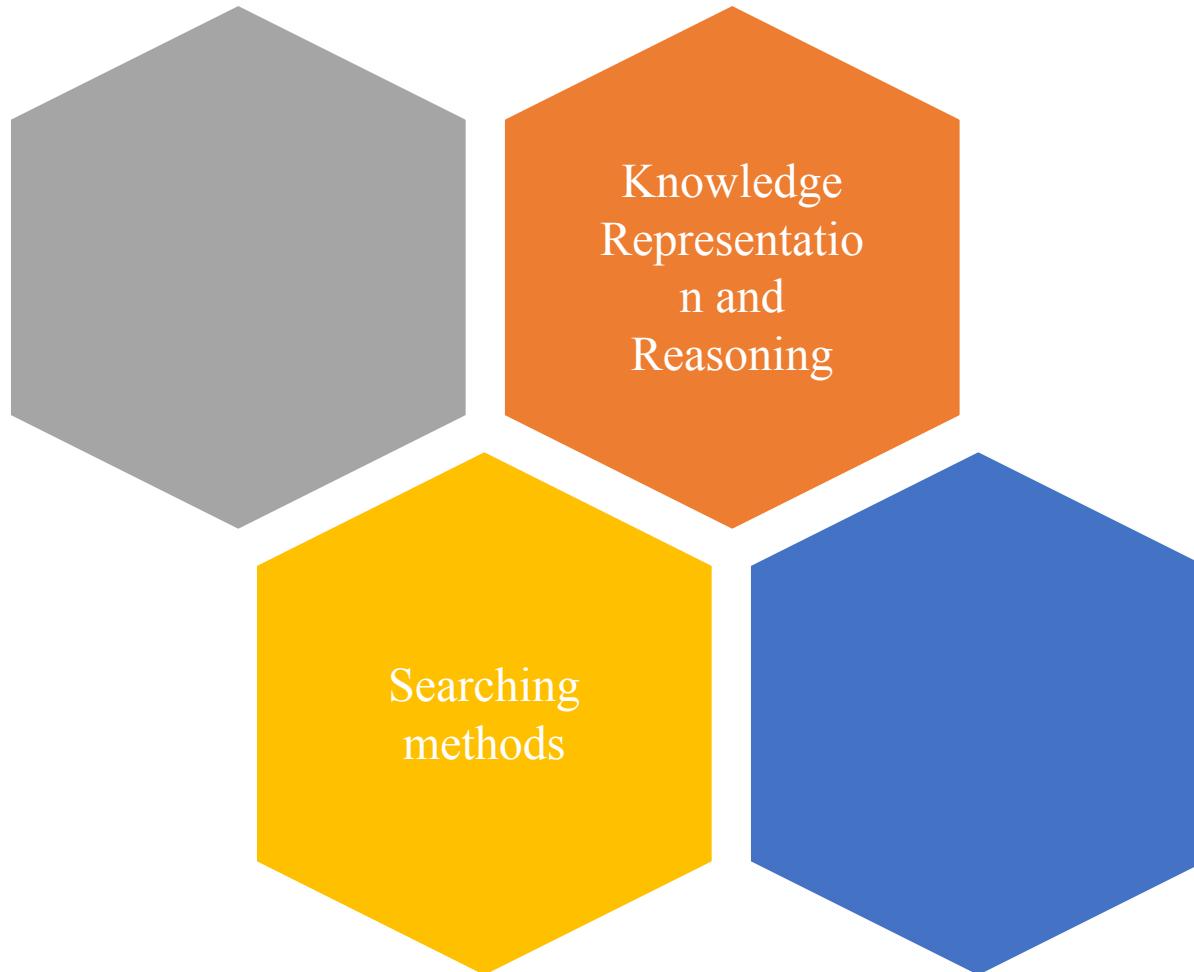
The Many Tribes of
ARTIFICIAL INTELLIGENCE



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1.2 Brief Review

□ Traditional AI Methods



OUTLINE

- 1.1 Brief Review
- 1.2 *Knowledge Representation & Reasoning*



1.2 Knowledge Representation & Reasoning

- *What is knowledge representation?*
- Propositional Logic
- Predicate Logic
- Production-rule System
- Frame-Based System
- State Space System
- Knowledge graph

1.3 Knowledge Representation & Reasoning

□ Puzzle Time

- An island has two kinds of inhabitants, *knaves*, who always tell the truth, and *knaves*, who always lie.
- You go to the island and meet A and B.
 - A says “B is a knight.”
 - B says “The two of us are of opposite types.”

Example: What are the types of A and B?



How to represent the problem and solve it by computer?



Knowledge representation and reasoning

1.3 Knowledge Representation & Reasoning

□ Knowledge

- The information, understanding and skills accumulated in long-term life and social practice, scientific research and experiments.
- Knowledge reflects the relationship between things in the objective world.

Example:

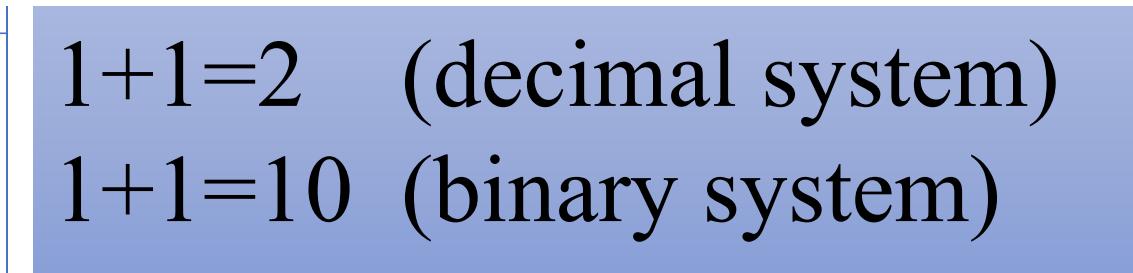
1. The snow is white. --- Facts
2. If you have a headache and a runny nose, you might have a cold. --- Rule

1.3 Knowledge Representation & Reasoning

□ Characteristic of Knowledge

➤ Relative correctness

Any knowledge is produced under certain conditions, and is correct under such conditions.

 $1+1=2$ (decimal system)
 $1+1=10$ (binary system)

1.3 Knowledge Representation & Reasoning

□ Characteristic of Knowledge

➤ Uncertainty: True, False, state between True and False

- Uncertainty caused by **randomness**
- Uncertainty caused by **ambiguity**
- Uncertainty caused by **experience**
- Uncertainty caused by **incompleteness**

Example:

1. If you have a headache and a runny nose, you **might** have a cold.
2. Li is **very high**.

1.3 Knowledge Representation & Reasoning

□ Characteristic of Knowledge

➤ Uncertainty caused by **ambiguity**



天气冷热



雨的大小



风的强弱



人的胖瘦



年龄大小



个子高低

--- A vague
concept

1.3 Knowledge Representation & Reasoning

- Characteristic of Knowledge
 - Representability and Exploitability

Representability of knowledge: Knowledge can be expressed in appropriate forms, such as language, writing, graphics, neural networks, etc.

Exploitability of Knowledge : Knowledge can be utilized

1.3 Knowledge Representation & Reasoning

□ Knowledge Representation

- Formalize or model human knowledge.
- a description of knowledge, or a set of conventions, a data structure that a computer can accept to describe knowledge.
- Principles for selecting knowledge representation methods:
 - Fully express domain knowledge. (充分表示领域知识。)
 - Conducive to the use of knowledge. (有利于对知识的利用。)
 - Easy to organize, maintain and manage. (便于对知识的组织、维护与管理。)
 - Easy to understand and implement. (便于理解与实现。)

1.3 Knowledge Representation & Reasoning

- What is knowledge representation?
- *Propositional Logic*
- Production-rule System
- Frame-Based System
- State Space System
- Knowledge graph

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

Proposition:

A := The street is wet.

B := It is raining.

A **proposition** is a statement that is either **true** or **false** but **not both**.

- Atomic formulas are denoted by letters A, B, C, etc.
- Each atomic formula is assigned a truth value: true (1) or false (0).

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

- A *proposition* is a **declarative sentence** that is either true or false.
- Examples of propositions:
 - a) The Moon is made of green cheese.
 - b) Beijing is the capital of China.
 - c) Hangzhou is the capital of Canada.
 - d) $1 + 0 = 1$
 - e) $0 + 0 = 2$
- Examples that are not propositions.
 - a) Sit down!
 - b) What time is it?
 - c) $x + 1 = 2$
 - d) $x + y = z$

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

- It is possible to determine whether any given statement is a proposition by prefixing it with
 - *It is true that . . .*
 - and seeing whether the result **makes grammatical sense.**
- What is the time?
- $2 + 3 = 5$
- “Phone” has five letters.
- $2 + 3 = 6$
- Oh dear!
- I like AI class.

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

- Have a try...

■ 您去电影院吗?	不是命题
■ $2 + 3 = 5$	命题
■ 看花去!	不是命题
■ 这句话是谎言。	不是命题
■ $X=2$	不是命题
■ 两个奇数之和是奇数。	命题
■ 李白要么擅长写诗，要么擅长喝酒。	命题

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

Proposition:

A := The street is wet.

B := It is raining.

“ Propositional logic is not the study of truth, but of the relationship between the truth of one statement and that of another”

——Hedman 2004



We can connect the two propositions A and B:

If it is raining, the street is wet.



Written more formally

It is raining. → The street is wet.

$$A \rightarrow B$$

1.3 Knowledge Representation & Reasoning

□ What is propositional logic

- Constructing Propositions

- Propositional Variables: p, q, r, s, \dots
- The proposition that is always true is denoted by **T** and the proposition that is always false is denoted by **F**.
- Compound Propositions: constructed from logical connectives and other propositions
 - Negation \neg (否定联结词)
 - Conjunction \wedge (合取联结词)
 - Disjunction \vee (析取联结词)
 - Implication \rightarrow (蕴涵联结词)
 - Biconditional \leftrightarrow (等价联结词)

1.3 Knowledge Representation & Reasoning

□ Syntax of logical connectives

- Conjunction (合取): And \wedge
- Disjunction (析取): Or \vee
- Negation (否定): Not \neg
- Implication (蕴涵): Implies \rightarrow (if... then...)
- Biconditional (等价): Iff \leftrightarrow (if and only if)

1.3 Knowledge Representation & Reasoning

□ AND (\wedge)

读作：“A并且B” “A与B”
称为：A与B的合取式
记作： $A \wedge B$

- The *conjunction* ' A **AND** B ', written $A \wedge B$, of two propositions is true **when both A and B are true**, false otherwise.

Translation of sentences to propositions

$A :=$ It's Monday.
 $B :=$ It's raining.

A	B	$A \wedge B$
t	t	t
t	f	f
f	t	f
f	f	f

It's Monday and it's raining.
It's Monday but it's raining.
It's Monday. It's raining.

} $A \wedge B$

1.3 Knowledge Representation & Reasoning

□ Translation of propositions to sentences

- In propositional logic :
- $A \wedge B$ and $B \wedge A$ should always have the same meaning.
- But...
- A:=She became sick .
- B:= she went to the doctor.

Logically the same!

- She became sick and she went to the doctor.
- and
- She went to the doctor and she became sick.

Different!

1.3 Knowledge Representation & Reasoning

□ OR (\vee)

- Also called *disjunction*.
- The disjunction “*A OR B*”, written $A \vee B$, of two propositions is true when *A* or *B* (or both) are true, false otherwise.

读作：“A或者B”
称为：A与B的析取式
记作： $A \vee B$

A	B	$A \vee B$
t	t	t
t	f	t
f	t	t
f	f	f

Translation of sentences to propositions

$A :=$ It's Monday.

$B :=$ It's raining.

It's Monday or it's raining. $A \vee B$

1.3 Knowledge Representation & Reasoning

□ NOT

- Also known as *negation*
- The negation “NOT A ” of a proposition (or $\neg A$) is true when A is false and is false otherwise.
- $\neg A$ may be read that it is
- false that A .

读作：“非A”
称为：A的否定式
记作： $\neg A$

A	$\neg A$
t	f
f	t

Translation of sentences to propositions

$A :=$ AI is easy.

It is false that AI is easy.

It is not the case that AI is easy. $\neg A$

AI is not easy.

1.3 Knowledge Representation & Reasoning

□ If . . . Then (\rightarrow)

- Also known as **implication**
- The implication “*A IMPLIES B*”, written $A \rightarrow B$, of two propositions is true when either *A* is false or *B* is true, and false otherwise.

读作：“如果A则B”
称为：A与B的蕴涵式
记作： $A \rightarrow B$

A:= I study hard.

B:= I get rich.

If I study hard then I get rich.

Whenever I study hard, I get rich.

That I study hard implies I get rich.

I get rich, if I study hard.

$A \rightarrow B$

A	B	$A \rightarrow B$
t	t	t
t	f	f
f	t	t
f	f	t

1.3 Knowledge Representation & Reasoning

□ If . . . Then (\rightarrow)

- In $p \rightarrow q$ there does not need to be any connection between the antecedent or the consequent.
- The “meaning” of $p \rightarrow q$ depends only on the truth values of p and q .
- These implications are perfectly fine, but would not be used in ordinary English.
 - “If the moon is made of green cheese, then I have more money than Bill Gates.”
 - “If the moon is made of green cheese then I’m on welfare.”

1.3 Knowledge Representation & Reasoning

□ Different Ways of Expressing $p \rightarrow q$

if p , then q

if p , q

q unless $\neg p$

q if p

q whenever p

q follows from p

p implies q

p only if q

q when p

p is sufficient for q

q is necessary for p

a necessary condition for p is q

a sufficient condition for q is p

1.3 Knowledge Representation & Reasoning

□ Converse, Contrapositive, and Inverse

- From $p \rightarrow q$ we can form new conditional statements .

- $q \rightarrow p$ is the **converse** of $p \rightarrow q$ (逆命题)
- $\neg q \rightarrow \neg p$ is the **contrapositive** of $p \rightarrow q$ (逆否命题)
- $\neg p \rightarrow \neg q$ is the **inverse** of $p \rightarrow q$ (否命题)

Example: Find the converse, inverse, and contrapositive of “It’s raining is a sufficient condition for my not going to town.”

Solution:

converse: ?

inverse: ?

contrapositive: ?

1.3 Knowledge Representation & Reasoning

□ Converse, Contrapositive, and Inverse

- From $p \rightarrow q$ we can form new conditional statements .
 - $q \rightarrow p$ is the **converse** of $p \rightarrow q$
 - $\neg q \rightarrow \neg p$ is the **contrapositive** of $p \rightarrow q$
 - $\neg p \rightarrow \neg q$ is the **inverse** of $p \rightarrow q$

Example: Find the converse, inverse, and contrapositive of “It’s raining is a sufficient condition for my not going to town.”

Solution: → “If it is raining, I do not go town.”

converse: If I do not go to town, then it is raining.

inverse: If it is not raining, then I will go to town.

contrapositive: If I go to town, then it is not raining.

1.3 Knowledge Representation & Reasoning

□ Biconditional

读作：“A当且仅当B”
称为：A与B的等价式
记作： $A \leftrightarrow B$

- Also known as iff or the biconditional.

The biconditional, written as $A \leftrightarrow B$, of two propositions is true when both A and B are true or when both A and B are false, and false otherwise.

A	B	$A \leftrightarrow B$
t	t	t
t	f	f
f	t	f
f	f	t

If p denotes “I am at home.” and q denotes “It is raining.” then $p \leftrightarrow q$ denotes “I am at home **if and only if** it is raining.”

1.3 Knowledge Representation & Reasoning

□ Semantics

- Example:

A := The street is wet.

B := It is raining.

If A is true, and B is true, then $A \wedge B$ is true.

A	B	$\neg A$	$A \wedge B$	$A \vee B$	$A \rightarrow B$	$A \leftrightarrow B$
t	t	f	t	t	t	t
t	f	f	f	t	f	f
f	t	t	f	t	t	f
f	f	t	f	f	t	t

Truth table

A	B	$A \wedge B$
t	t	t
t	f	f
f	t	f
f	f	f

Operator	precedence
\neg	1
\wedge, \vee	2
$\rightarrow, \leftrightarrow$	3

1.3 Knowledge Representation & Reasoning

□ Tautologies, Contradictions, and Contingencies

- A *tautology* (永真式) is a proposition which is always true.
 - Example: $p \vee \neg p$
- A *contradiction* (矛盾式) is a proposition which is always false.
 - Example: $p \wedge \neg p$
- A *contingency* is a proposition which is neither a tautology nor a contradiction, such as p

1.3 Knowledge Representation & Reasoning

□ Logical Equivalences

- Two compound propositions p and q are **logically equivalent** if $p \leftrightarrow q$ is a tautology.
- We write this as $p \Leftrightarrow q$ or as $p \equiv q$ where p and q are compound propositions.
- Two compound propositions p and q are equivalent if and only if the columns in a truth table giving their truth values agree.
- This truth table shows that $\neg p \vee q$ is equivalent to $p \rightarrow q$.

p	q	$\neg p$	$\neg p \vee q$	$p \rightarrow q$
T	T	F	T	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

1.3 Knowledge Representation & Reasoning

□ De Morgan's Laws

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$



Augustus De Morgan

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

1806-1871

This truth table shows that De Morgan's Second Law holds.

p	q	$\neg p$	$\neg q$	$(p \vee q)$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	F	T	F
F	F	T	T	F	T	T

1.3 Knowledge Representation & Reasoning

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p	q	$\neg p$	$\neg q$	$(p \vee q)$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	T	F	F
T	F	F	T	F	T	F
F	T	T	F	F	T	F
F	F	T	T	F	T	T

1.3 Knowledge Representation & Reasoning

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p	q	$\neg p$	$\neg q$	$(p \vee q)$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	T		
T	F	F	T	T		
F	T	T	F	T		
F	F	T	T	F		

1.3 Knowledge Representation & Reasoning

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p	q	$\neg p$	$\neg q$	$(p \vee q)$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	T	F	
T	F	F	T	T	F	
F	T	T	F	T	F	
F	F	T	T	F	T	

1.3 Knowledge Representation & Reasoning

□ De Morgan's Laws

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1806-1871

This truth table shows that De Morgan's Second Law holds.

p	q	$\neg p$	$\neg q$	$(p \vee q)$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	T	F	F
F	F	T	T	F	T	T

1.3 Knowledge Representation & Reasoning

□ Key Logical Equivalences

- Identity Laws:

(统一律)

$$p \wedge T \equiv p \quad p \vee F \equiv p$$

- Domination Laws:

(支配律, 零律)

$$p \vee T \equiv T \quad p \wedge F \equiv F$$

- Idempotent laws:

(等幂律)

$$p \vee p \equiv p \quad p \wedge p \equiv p$$

- Double Negation Law:

(对合律)

$$\neg(\neg p) \equiv p$$

- Negation Laws:

(否定律)

$$p \vee \neg p \equiv T \quad p \wedge \neg p \equiv F$$

1.3 Knowledge Representation & Reasoning

□ Key Logical Equivalences

- Commutative Laws: $p \vee q \equiv q \vee p$, $p \wedge q \equiv q \wedge p$
(交换律)
- Associative Laws: $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
(结合律)
 $(p \vee q) \vee r \equiv p \vee (q \vee r)$
- Distributive Laws:
(分配律)
 $(p \vee (q \wedge r)) \equiv (p \vee q) \wedge (p \vee r)$
 $(p \wedge (q \vee r)) \equiv (p \wedge q) \vee (p \wedge r)$
- Absorption Laws:
(吸收律)
 $p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$

1.3 Knowledge Representation & Reasoning

□ Logical Equivalences

- *The operations \wedge , \vee are commutative and associative, and the following equivalences are generally valid:*
- $\neg A \vee B \equiv A \rightarrow B$ (implication)
- $A \rightarrow B \equiv \neg B \rightarrow \neg A$ (contraposition)
- $(A \rightarrow B) \wedge (B \rightarrow A) \equiv (A \leftrightarrow B)$ (equivalence)
- $\neg(A \wedge B) \equiv \neg A \vee \neg B$ (De Morgan's law)
- $\neg(A \vee B) \equiv \neg A \wedge \neg B$
- $A \vee (B \wedge C) \equiv (A \vee B) \wedge (A \vee C)$ (distributive law)
- $A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$
- $A \vee \neg A \equiv t$ (tautology)
- $A \wedge \neg A \equiv f$ (contradiction)
- $A \vee f \equiv A$
- $A \vee t \equiv t$
- $A \wedge f \equiv f$
- $A \wedge t \equiv A$

1.3 Knowledge Representation & Reasoning

□ Logical Equivalences

- The operations \wedge, \vee are commutative and associative, and the
 - following equivalences are generally valid:
- | | | |
|-------------------|----------|----------------------------------------------|
| $\neg A \vee B$ | \equiv | $A \rightarrow B$ (implication) |
| $A \rightarrow B$ | \equiv | $\neg B \rightarrow \neg A$ (contraposition) |

A	B	$\neg A$	$\neg A \vee B$	$A \rightarrow B$
t	t	f	t	t
t	f	f	f	f
f	t	t	t	t
f	f	t	t	t

1.3 Knowledge Representation & Reasoning

□ Equivalence Proofs

Example: Show that $\neg(p \vee (\neg p \wedge q))$
is logically equivalent to $\neg p \wedge \neg q$

Solution:

1.3 Knowledge Representation & Reasoning

□ Equivalence Proofs

Example: Show that $\neg(p \vee (\neg p \wedge q))$
is logically equivalent to $\neg p \wedge \neg q$

Solution:

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) \\ &\equiv \neg p \wedge [\neg(\neg p) \vee \neg q] \\ &\equiv \neg p \wedge (p \vee \neg q) \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) \\ &\equiv F \vee (\neg p \wedge \neg q) \\ &\equiv (\neg p \wedge \neg q) \vee F \\ &\equiv (\neg p \wedge \neg q)\end{aligned}$$

by the second De Morgan law
by the first De Morgan law
by the double negation law
by the second distributive law
because $\neg p \wedge p \equiv F$
by the commutative law
for disjunction
by the identity law for **F**

1.3 Knowledge Representation & Reasoning

□ Equivalence Proofs

Example: Show that $(p \wedge q) \rightarrow (p \vee q)$

is a tautology.

Solution:

1.3 Knowledge Representation & Reasoning

□ Equivalence Proofs

Example: Show that $(p \wedge q) \rightarrow (p \vee q)$

is a tautology.

Solution:

$$\begin{aligned}(p \wedge q) \rightarrow (p \vee q) &\equiv \neg(p \wedge q) \vee (p \vee q) && \text{by truth table for } \rightarrow \\&\equiv (\neg p \vee \neg q) \vee (p \vee q) && \text{by the first De Morgan law} \\&\equiv (\neg p \vee p) \vee (\neg q \vee q) && \text{by associative and} \\&&& \text{commutative laws} \\&\equiv T \vee T && \text{laws for disjunction} \\&\equiv T && \text{by truth tables} \\&&& \text{by the domination law}\end{aligned}$$

1.3 Knowledge Representation & Reasoning

□ Now, we have learned...

- Three basic elements in proposition logic: propositions, operations, and the truth values.
- Logical equivalences

1.3 Knowledge Representation & Reasoning

□ Applications

- 1. Translate English Sentences
- 2. System Specifications
- 3. Logic Puzzles

1.3 Knowledge Representation & Reasoning

□ Example

Problem: Translate the following sentence into propositional logic:
“You can access the Internet from campus **if** you are a computer science major **or** you are **not** a freshman.”

Atomic propositions:

- A:= You can access the Internet from campus.
- B:= You are a computer science major .
- C:= You are a freshman.

$$(B \vee \neg C) \rightarrow A$$

1.3 Knowledge Representation & Reasoning

□ Example

Problem:

- (1) 除非你很努力，否则你将失败。
- (2) 张三或者李四都可以做这件事。

Atomic propositions:

- (1) A:= 你努力。 B:= 你失败。
- (2) A := 张三能做事。 B:= 李四可以做这件事。

- (1) $\neg A \rightarrow B$
- (2) $A \wedge B$

1.3 Knowledge Representation & Reasoning

□ Consistent System Specifications (系统规范说明)

- Definition: A list of propositions is consistent if it is possible to assign truth values to the proposition variables so that each proposition in the list is true.

1.3 Knowledge Representation & Reasoning

□ Consistent System Specifications

Example: Is this list of propositions consistent?

- “The diagnostic message is stored in the buffer or it is retransmitted.”
- “The diagnostic message is not stored in the buffer.”
- “If the diagnostic message is stored in the buffer, then it is retransmitted.”

1.3 Knowledge Representation & Reasoning

□ Consistent System Specifications

Example: Is this list of propositions consistent?

- “The diagnostic message is stored in the buffer or it is retransmitted.”
- “The diagnostic message is not stored in the buffer.”
- “If the diagnostic message is stored in the buffer, then it is retransmitted.”

P:= “The diagnostic message is stored in the buffer.”
Q:= “The diagnostic message is retransmitted”

$$P \vee Q$$

1.3 Knowledge Representation & Reasoning

□ Consistent System Specifications

Example: Is this list of propositions consistent?

- “The diagnostic message is stored in the buffer or it is retransmitted.”
- “The diagnostic message is not stored in the buffer.”
- “If the diagnostic message is stored in the buffer, then it is retransmitted.”

P:= “The diagnostic message is stored in the buffer.”
Q:= “The diagnostic message is retransmitted”

$$\neg P$$

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□ Consistent System Specifications

Example: Is this list of propositions consistent?

■ “The diagnostic message is stored in the buffer or it is retransmitted.” $P \vee Q$

■ “The diagnostic message is not stored in the buffer.” $\neg P$

■ “If the diagnostic message is stored in the buffer, then it is retransmitted.” $P \rightarrow Q$

P	Q	$P \vee Q$	$\neg P$	$P \rightarrow Q$
t	t	t	f	t
t	f	t	f	t
f	t	t	t	t
f	f	f	t	t

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Example: Is this list of propositions consistent?

- “The diagnostic message is stored in the buffer or it is retransmitted.”
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When P is false and Q is true all three statements are true. So the list of propositions is **consistent**.

P	Q	$P \vee Q$	$\neg P$	$P \rightarrow Q$
t	t	t	f	t
t	f	t	f	f
f	t	t	t	t
f	f	f	t	t

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Example: Is this list of propositions consistent?

■ “The diagnostic message is stored in the buffer or it is retransmitted.” $P \vee Q$

■ “The diagnostic message is not stored in the buffer.” $\neg P$

■ “If the diagnostic message is stored in the buffer, then it is retransmitted.” $P \rightarrow Q$

■ The diagnostic message is not retransmitted. $\neg Q$

P	Q	$P \vee Q$	$\neg P$	$P \rightarrow Q$	$\neg Q$
t	t	t	f	t	
t	f	t	f	f	
f	t	t	t	t	
f	f	f	t	t	

1.3 Knowledge Representation & Reasoning

□ Consistent System Specifications

Example: Is this list of propositions consistent?

- “The diagnostic message is stored in the buffer or it is retransmitted.” $P \vee Q$
- “The diagnostic message is not stored in the buffer.” $\neg P$
- “If the diagnostic message is stored in the buffer, then it is retransmitted.” $P \rightarrow Q$
- The diagnostic message is not retransmitted. $\neg Q$

The list of propositions is
not consistent!

P	Q	$P \vee Q$	$\neg P$	$P \rightarrow Q$	$\neg Q$
t	t	t	f	t	f
t	f	t	f	f	t
f	t	t	t	t	f
f	f	f	t	t	t

1.3 Knowledge Representation & Reasoning

□ Logic Puzzles

Knights: t
Knaves: f

- An island has two kinds of inhabitants, *knaves*, who always tell the truth, and *knaves*, who always lie.
- You go to the island and meet A and B.
 - A says “B is a knight.”
 - B says “The two of us are of opposite types.”

Example: What are the types of A and B?

1.3 Knowledge Representation & Reasoning

□ Logic Puzzles

- **Solution:** Let p and q be the statements that A is a knight and B is a knight, respectively. So, then $\neg p$ represents the proposition that A is a knave and $\neg q$ that B is a knave.
 - If A is a knight, then p is true. Since knights tell the truth, q must also be true. Then $(p \wedge \neg q) \vee (\neg p \wedge q)$ would have to be true, but it is not. So, A is not a knight and therefore $\neg p$ must be true.
 - If A is a knave, then B must not be a knight since knaves always lie. So, then both $\neg p$ and $\neg q$ hold since both are knaves.

Homework-1

□ Logic Puzzles

Knights: t
Knaves: f

- An island has two kinds of inhabitants, *knights*, who always tell the truth, and *knaves*, who always lie.
- You go to the island and meet A and B.
 - A says “At least one of us is a knave.”
 - B says nothing.

Example: What are the types of A and B?